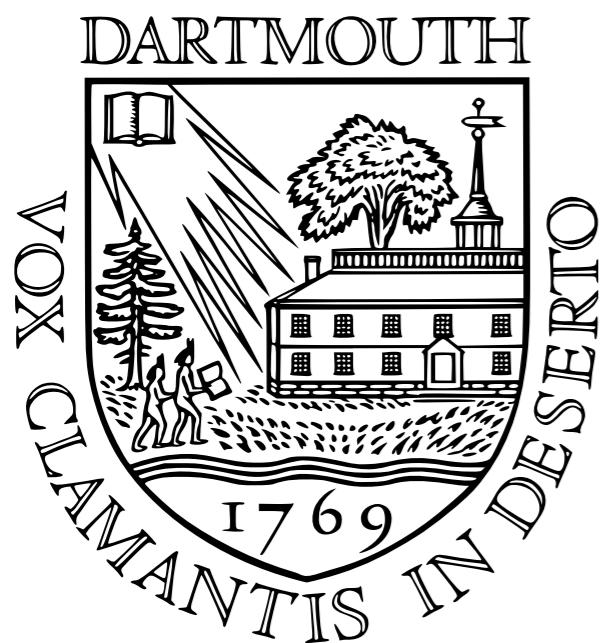




SPQR  
LAB RAT RY



# Mementos

**System Support for Long-Running Computations on RFID-Scale Devices**

Benjamin Ransford\*, UMass Amherst  
Jacob Sorber, Dartmouth College  
Kevin Fu, UMass Amherst

<http://spqr.cs.umass.edu/mementos>

ASPLOS XVI — March 8, 2011



CNS-0627529, CNS-0845874, CNS-0923313, Grad Res. Fel.

Any opinions, findings, and conclusions expressed in this material are those of the authors and do not necessarily reflect the views of the NSF.

# Ubiquitous Computing

---



*“... the most powerful things are those that are effectively **invisible** in use.”*

— Mark Weiser  
(PARC, 1988)

## Problem:

Batteryless **invisible** computer ⇒ constant reboots

# Baby Steps Toward Ubicomp

---

1. Take Emerging Platform

# Baby Steps Toward Ubicomp

---

1. Take Emerging Platform

2. Add Robustness Mechanism

# Baby Steps Toward Ubicomp

---

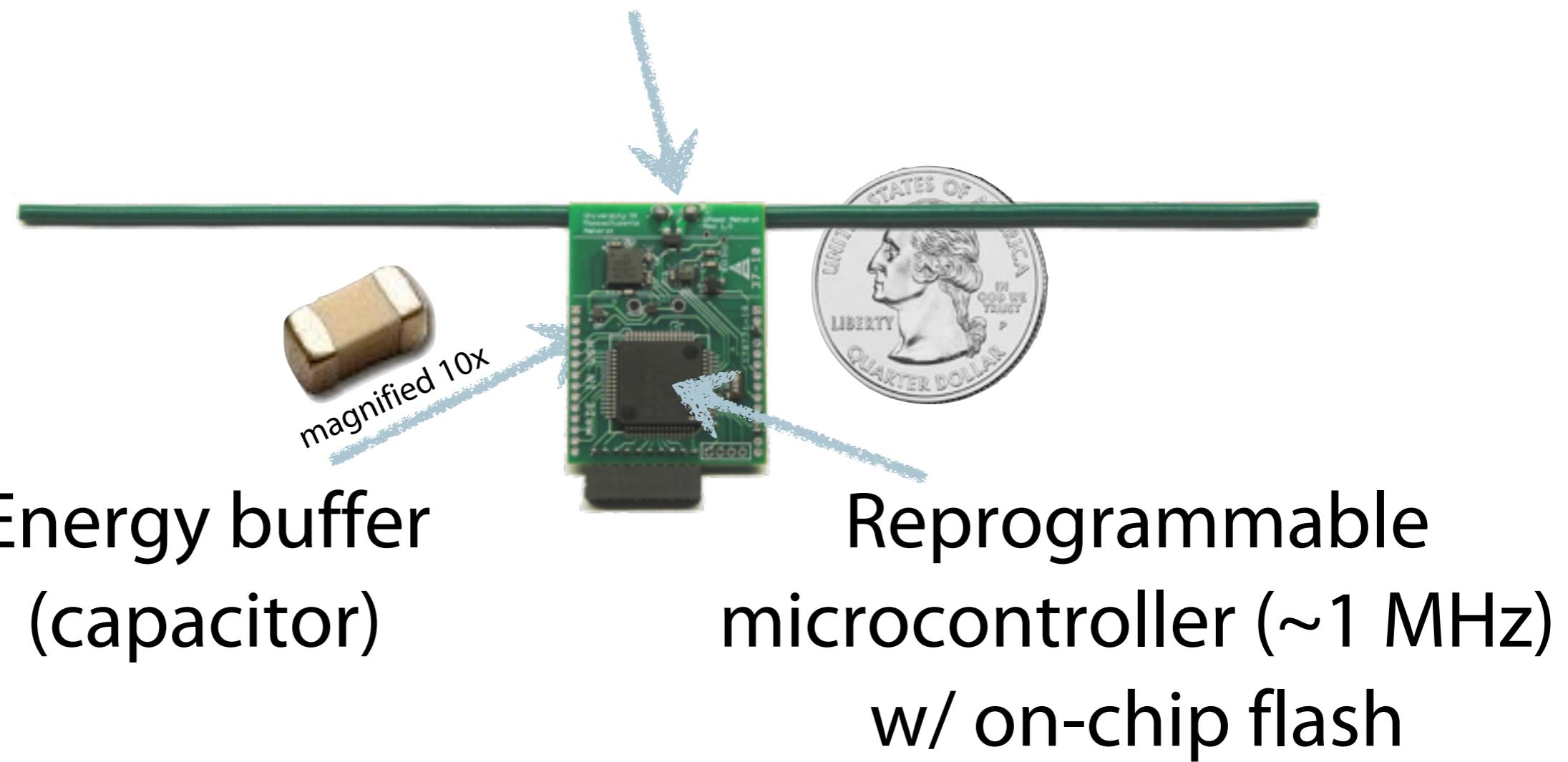
1. Take Emerging Platform

2. Add Robustness Mechanism

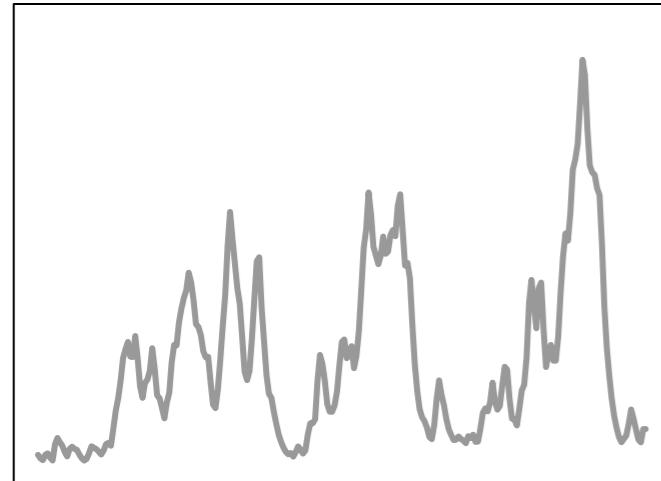
3. Provide Simulation Tools

# RFID-Scale Devices

Radio (RF) harvester



# RFID-Scale Devices



Radio (RF) harvester

Energy buffer  
(capacitor)

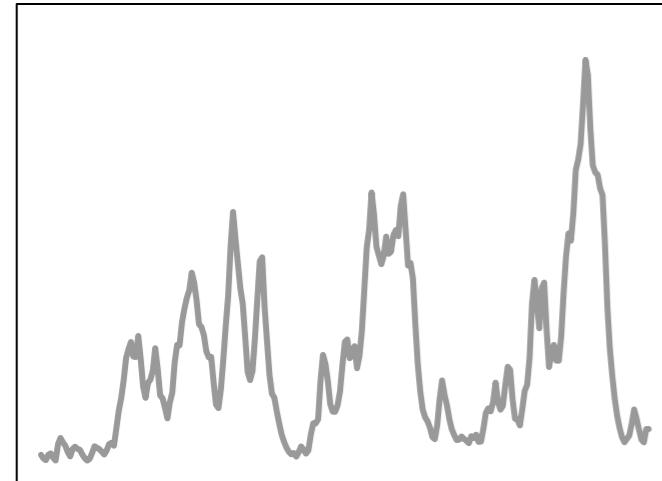


magnified 10x

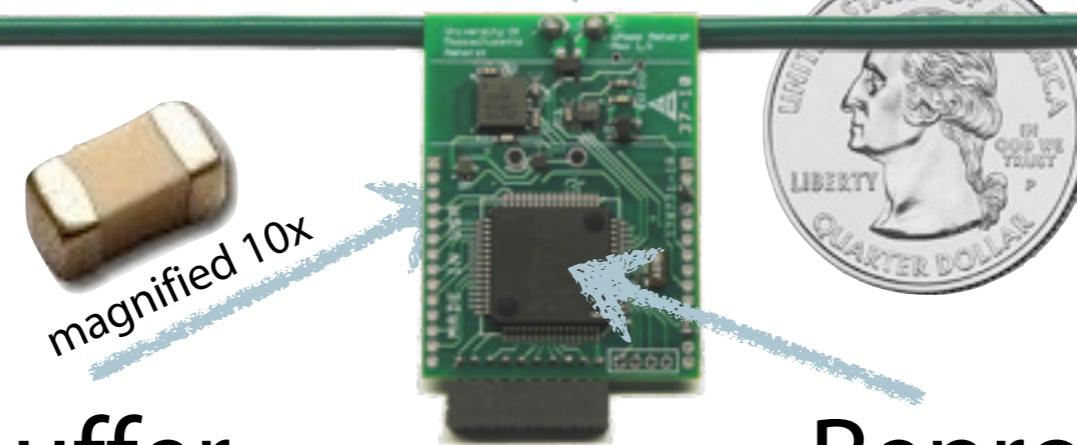


Reprogrammable  
microcontroller (~1 MHz)  
w/ on-chip flash

# RFID-Scale Devices



Radio (RF) harvester



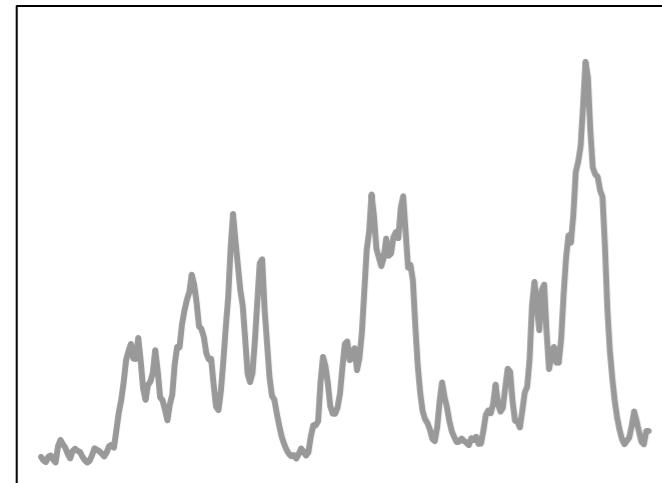
Energy buffer  
(capacitor)

Fills quickly,  
low capacity

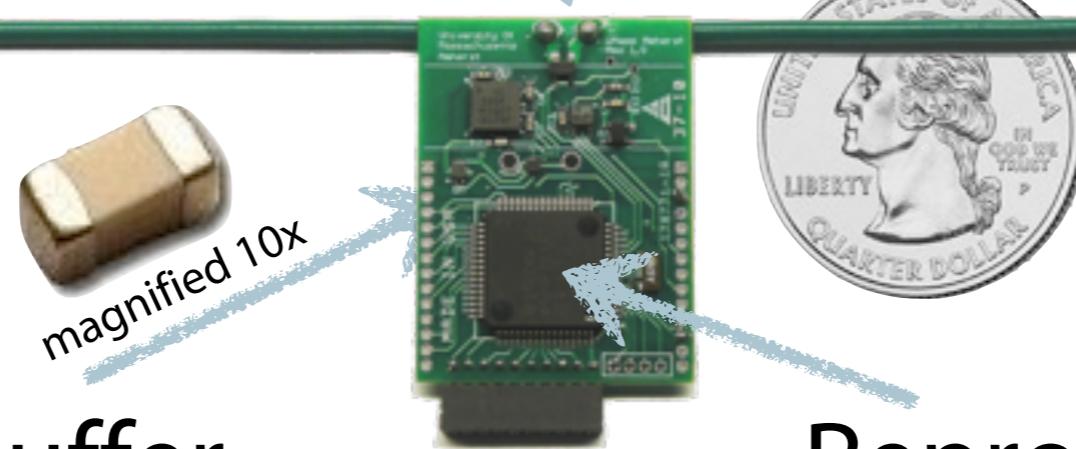


Reprogrammable  
microcontroller (~1 MHz)  
w/ on-chip flash

# RFID-Scale Devices



Radio (RF) harvester



Energy buffer  
(capacitor)

Fills quickly,  
low capacity



Reprogrammable  
microcontroller (~1 MHz)

w/ on-chip flash

Frequent reboots



ball: clipart.pierceinternet.com

# Mementos: System Support for Long-Running Computation on RFID-Scale Devices

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## Abstract

Transiently powered computing devices such as RFID tags, kinetic energy harvesters, and smart cards typically rely on programs that complete a task under tight time constraints before energy starvation leads to complete loss of volatile memory. *Mementos* is a software system that transforms general-purpose programs into interruptible computations that are protected from frequent power losses by automatic, energy-aware state checkpointing. *Mementos* comprises a collection of optimization passes for the LLVM compiler infrastructure and a linkable library that exercises hardware support for energy measurement while managing state checkpoints stored in nonvolatile memory. We evaluate *Mementos* against diverse test cases in a trace-driven simulator of transiently powered RFID-scale devices. Although *Mementos*'s energy checks increase run time when energy is plentiful, they allow *Mementos* to safely suspend execution when energy dwindles, effectively spreading computation across zero or more power failures. This paper's contributions are: a study of the runtime environment for programs on RFID-scale devices; an energy-aware state checkpointing system for these devices that is implemented for the MSP430 family of microcontrollers; and a trace-driven simulator of transiently powered RFID-scale devices.

*Categories and Subject Descriptors* C.3 [SPECIAL-PURPOSE AND APPLICATION-BASED SYSTEMS]: Real-time and embedded systems

*General Terms* Design, Experimentation

*Keywords* Mementos, RFID-Scale Devices, Computational RFID, Energy-Aware Checkpointing

## 1. Introduction

Demand for tiny, easily deployable computers has driven the development of general-purpose *remotely-powered computers* that lack both batteries and wired power, operating exclusively on energy harvested from remote supplies or the environment. Such devices range from *computational RFIDs* [16]—microcontroller-based devices that harvest RF from readers and communicate via RFID protocols—to general-purpose batteryless sensor devices [45].

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Computing under transient power conditions is a challenge. Transiently powered RFID tags use simple state machines instead of supporting general-purpose computation. Contactless smart cards perform more complicated special-purpose computations (e.g., cardholder authentication); however, they offer no execution guarantees, and instead rely on the user to provide the needed RF power for a sufficient period of time. When energy consumption outpaces energy harvesting, these computations fail and must restart from scratch, when adequate energy becomes available.

With ultra-low-power microcontrollers (MCUs), tiny programmable devices can perform computation and sensing under RFID-scale energy constraints; however, these MCUs consume more power than conventional RFID circuitry, and energy consumption can easily outpace harvesting, resulting in frequent power loss.

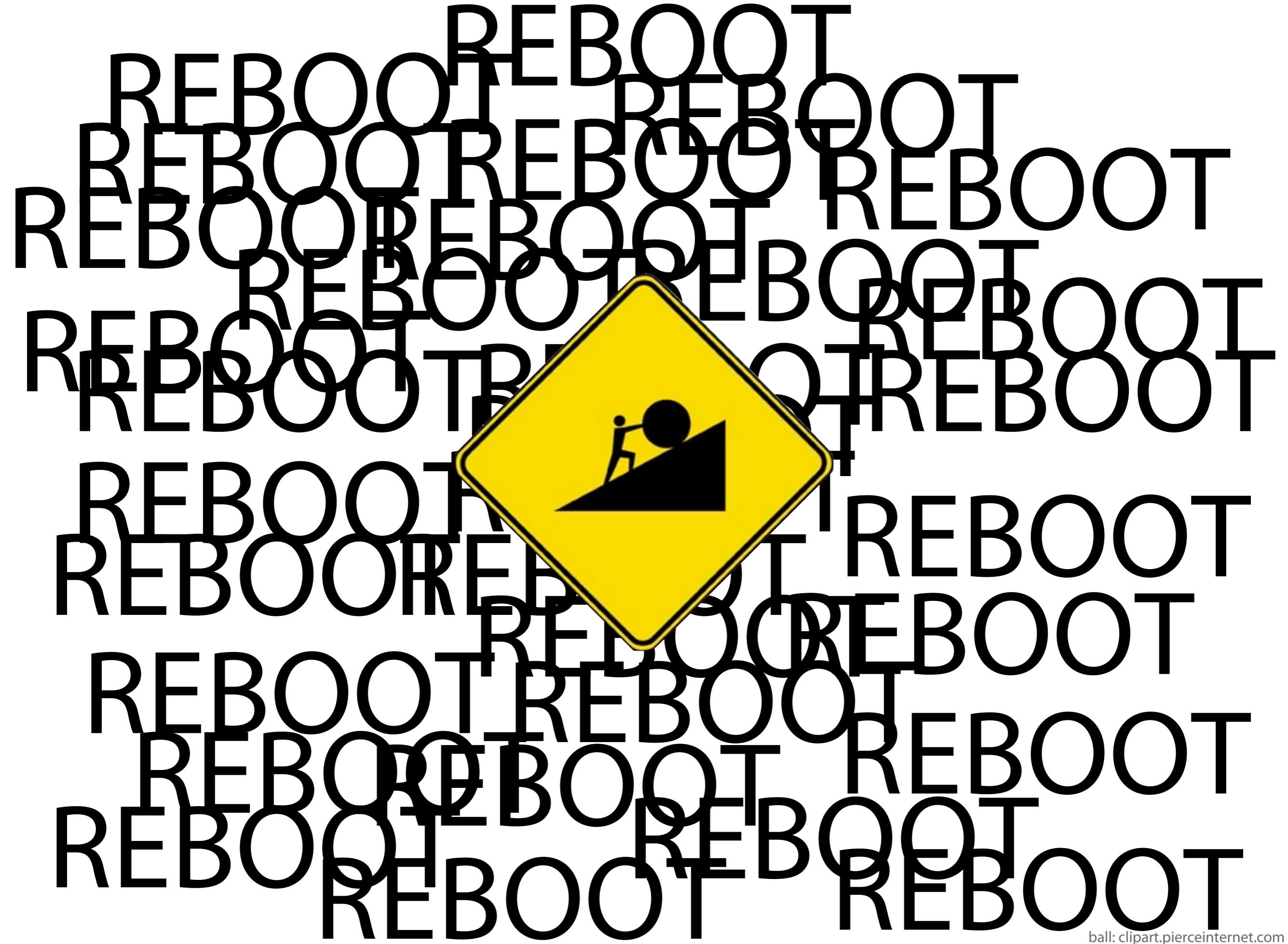
Today, programs that run CPU-intensive operations like cryptography on these devices are painstakingly and painstakingly hand-tuned to complete within a short time window (often under 100 ms) [7, 9]. The usefulness and power of RFID-scale devices can be dramatically improved if designers can confidently write programs without being limited by power failures.

*Mementos* is a software system that enables long-running computations to span power loss events by combining compile-time instrumentation and run-time energy-aware state checkpointing<sup>1</sup>. At compile time, *Mementos* inserts function calls that estimate available energy. At run time, *Mementos* predicts power losses and, when appropriate, saves program state to nonvolatile memory. After a failure, program state is restored and execution continues rather than restarting from scratch.

This paper makes the following contributions: (1) An energy-aware state checkpointing system that splits program execution across multiple lifecycles on transiently powered RFID-scale devices. The system is implemented for the MSP430 family of microcontrollers, requires no hardware modifications to existing devices, and operates automatically at run time without user intervention. (2) A suite of compile-time optimization passes that insert energy checks at control points in a program. The optimization passes employ three different instrumentation strategies that favor different program structures. (3) A trace-driven simulator to evaluate the behavior of programs on transiently powered RFID-scale devices. The simulator, modeled after a prototype hardware device with an off-the-shelf microcontroller, takes executable code as input and simulates power loss events during runs. We evaluate the simulator's accuracy and *Mementos*'s performance under simulation in Section 5.

<sup>1</sup>In the iMile file *Mementos*, the main character would unpredictably lose short-term memory, especially when sleeping. He checkpoints state with notes and tattoos in an attempt to execute long-running tasks.

ball: clipart.pierceinternet.com



ball: clipart.pierceinternet.com

# Robustness Under RF Harvesting

---

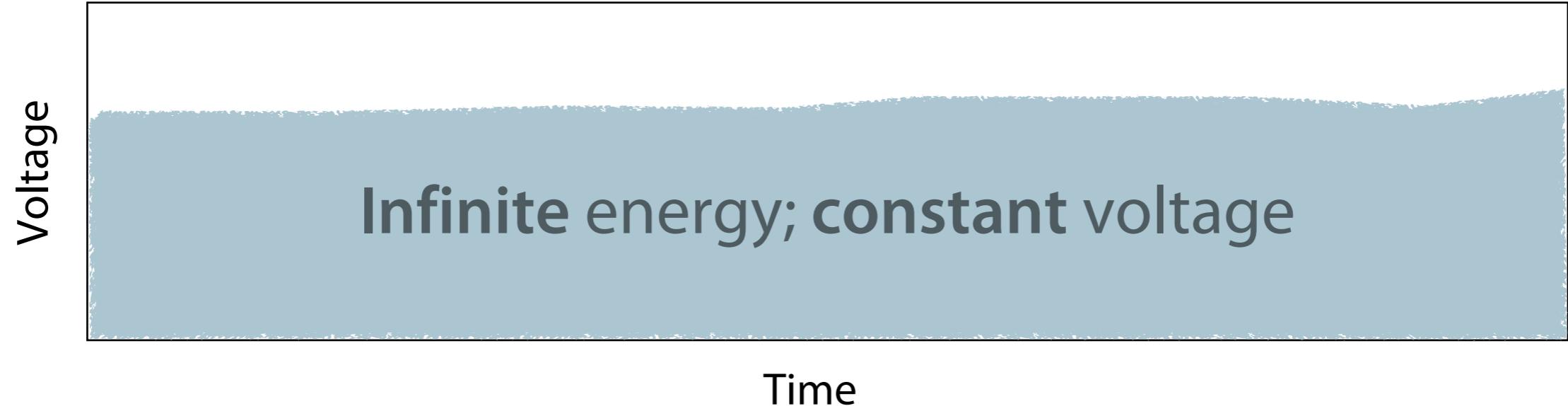


300 ms

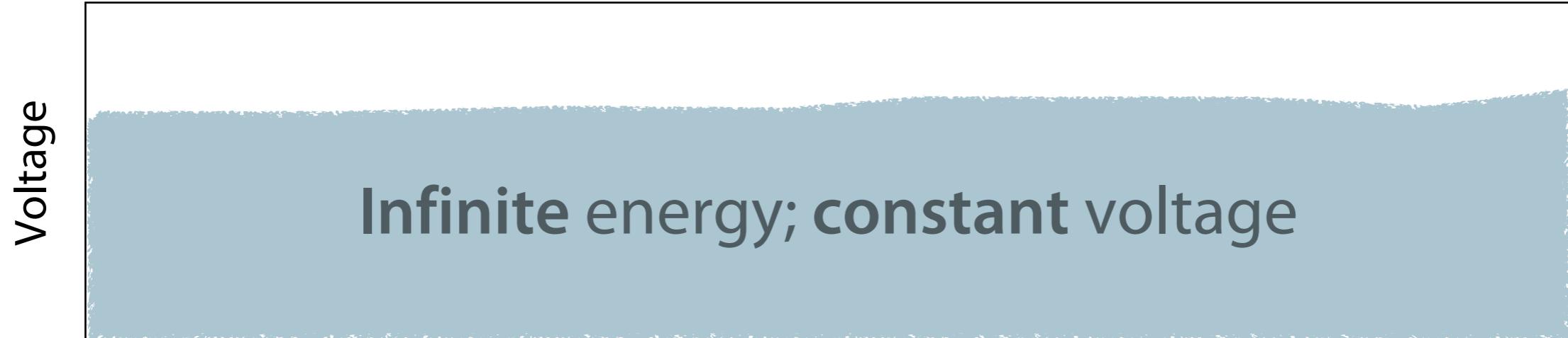
- Typical approach: constrain the problem
- **Mementos:** relax constraints to make general-purpose computation feasible

# Unpredictable Energy Morass

---

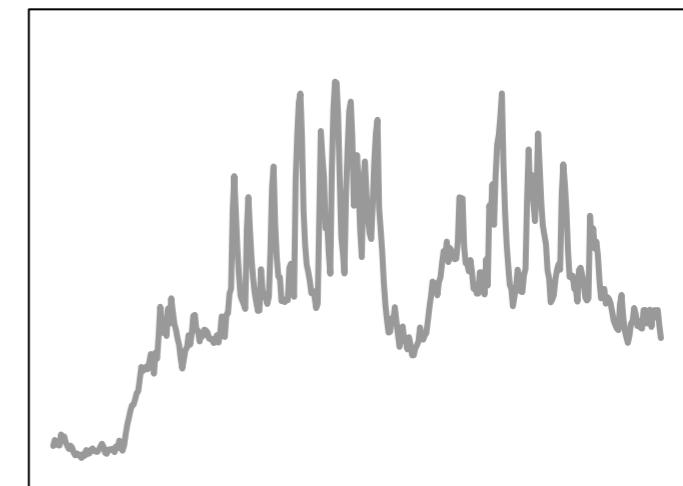
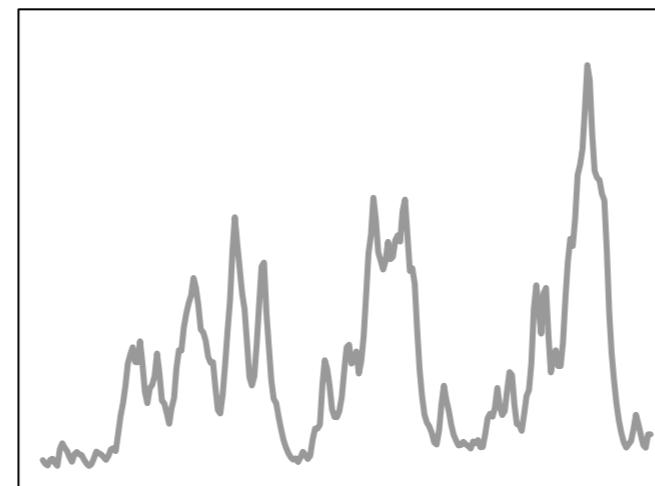
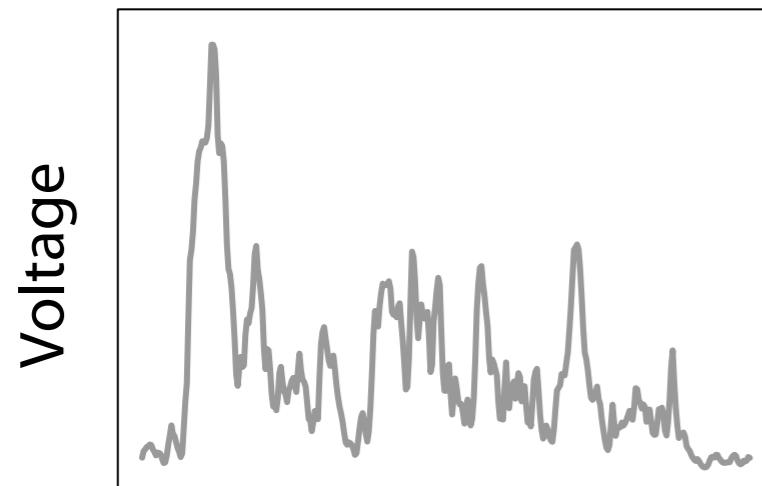


# Unpredictable Energy Morass

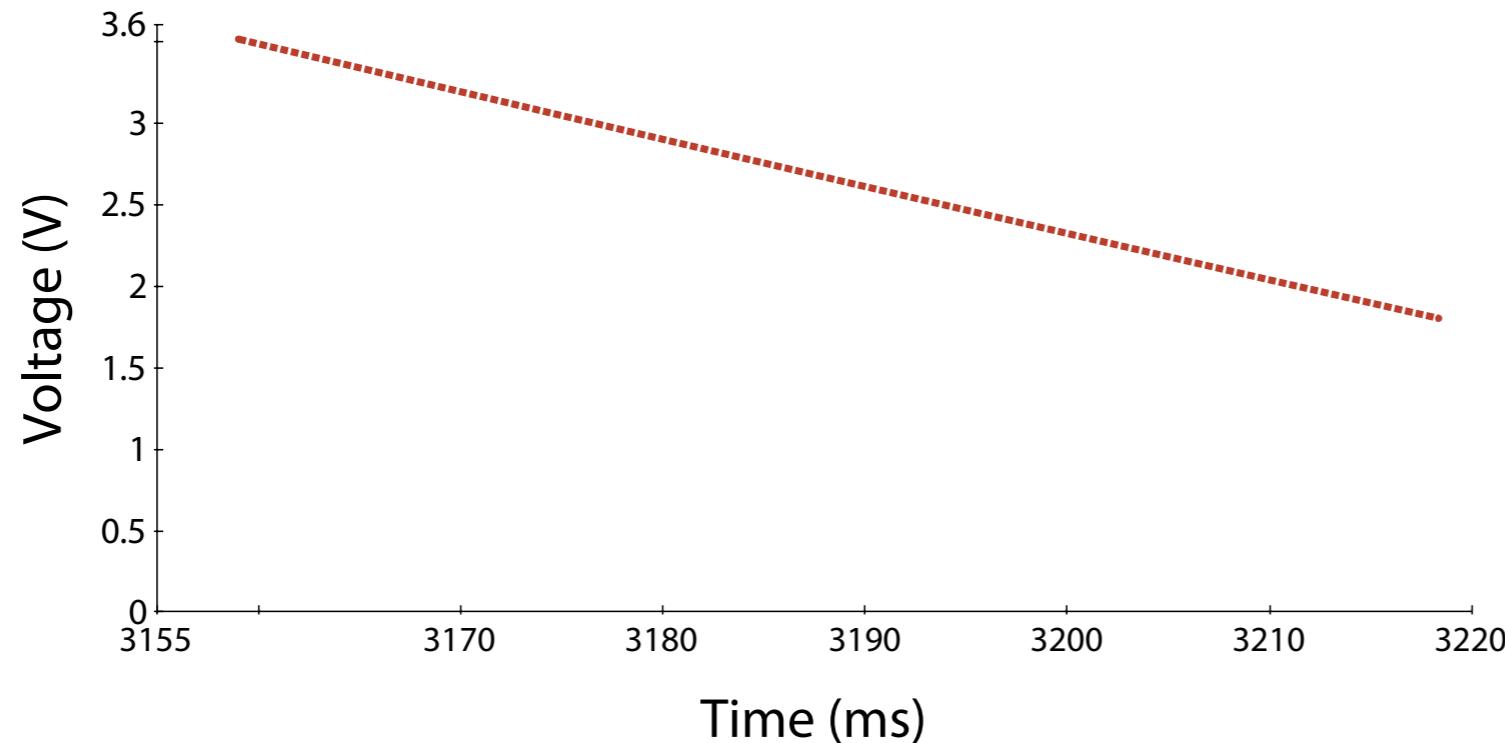


Time

vs.



# Mementos Approach

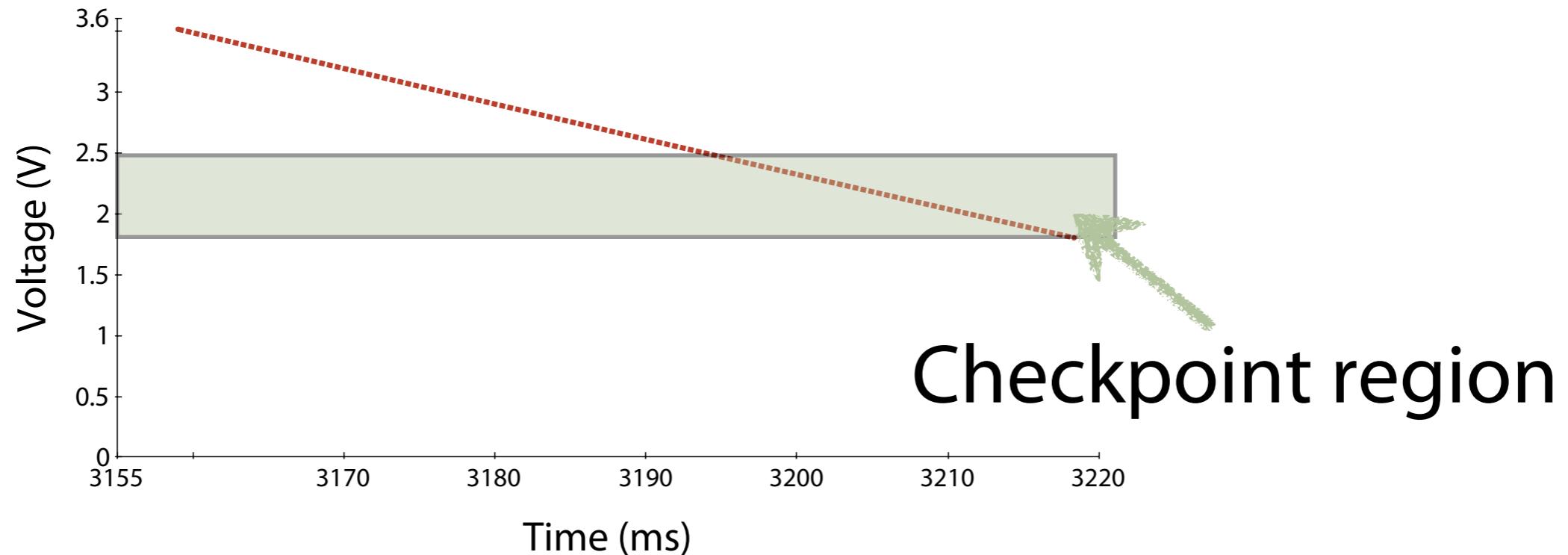


- Checkpoint when failure appears imminent
- Spread computation across reboots



Movie poster: publispain.com

# Mementos Approach

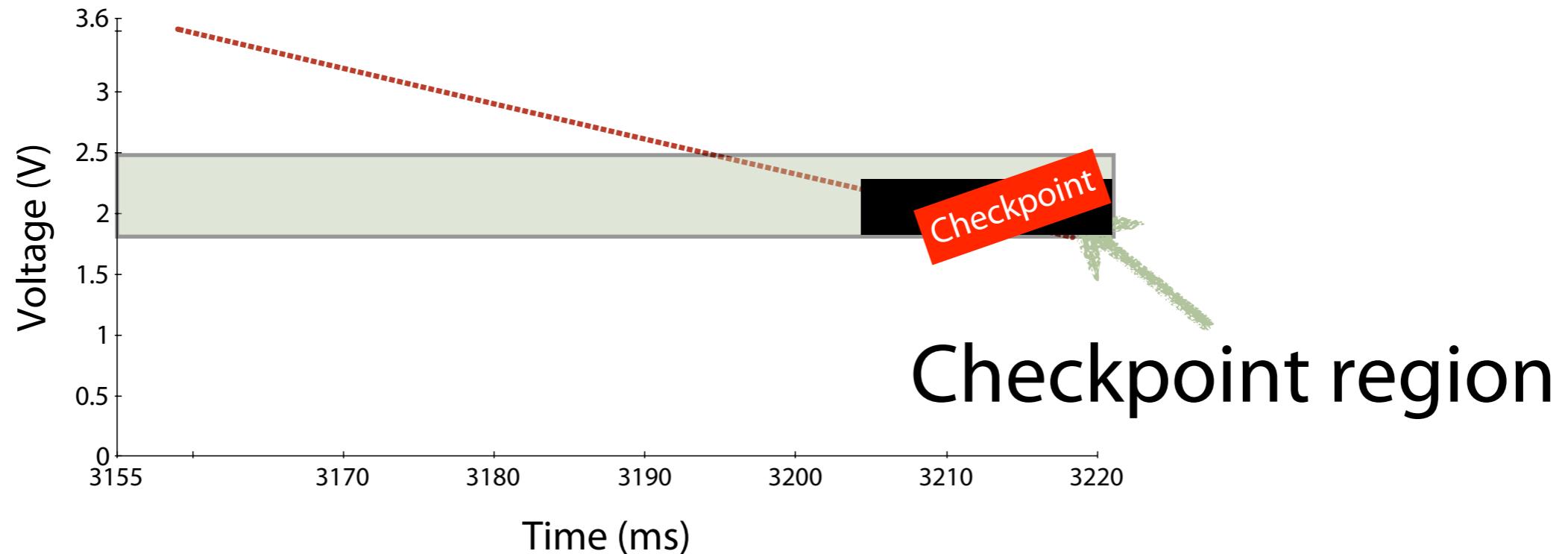


- Checkpoint when failure appears imminent
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Movie poster: publispain.com

# Mementos Approach

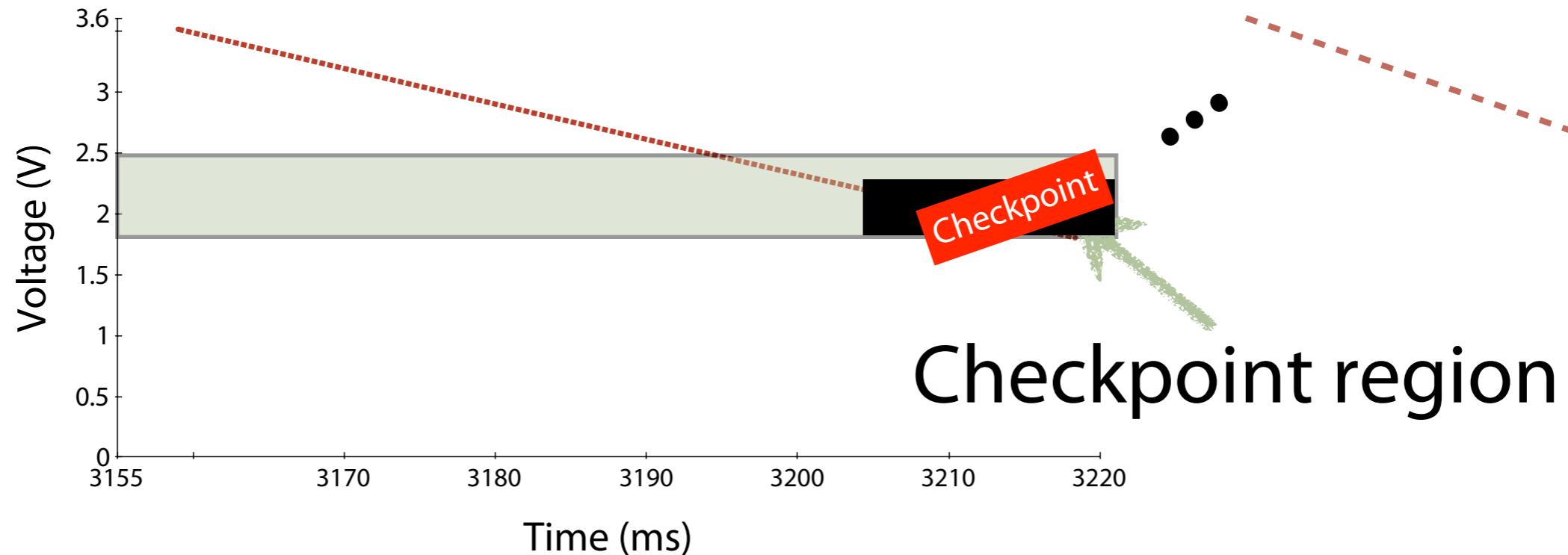


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Movie poster: publispain.com

# Mementos Approach

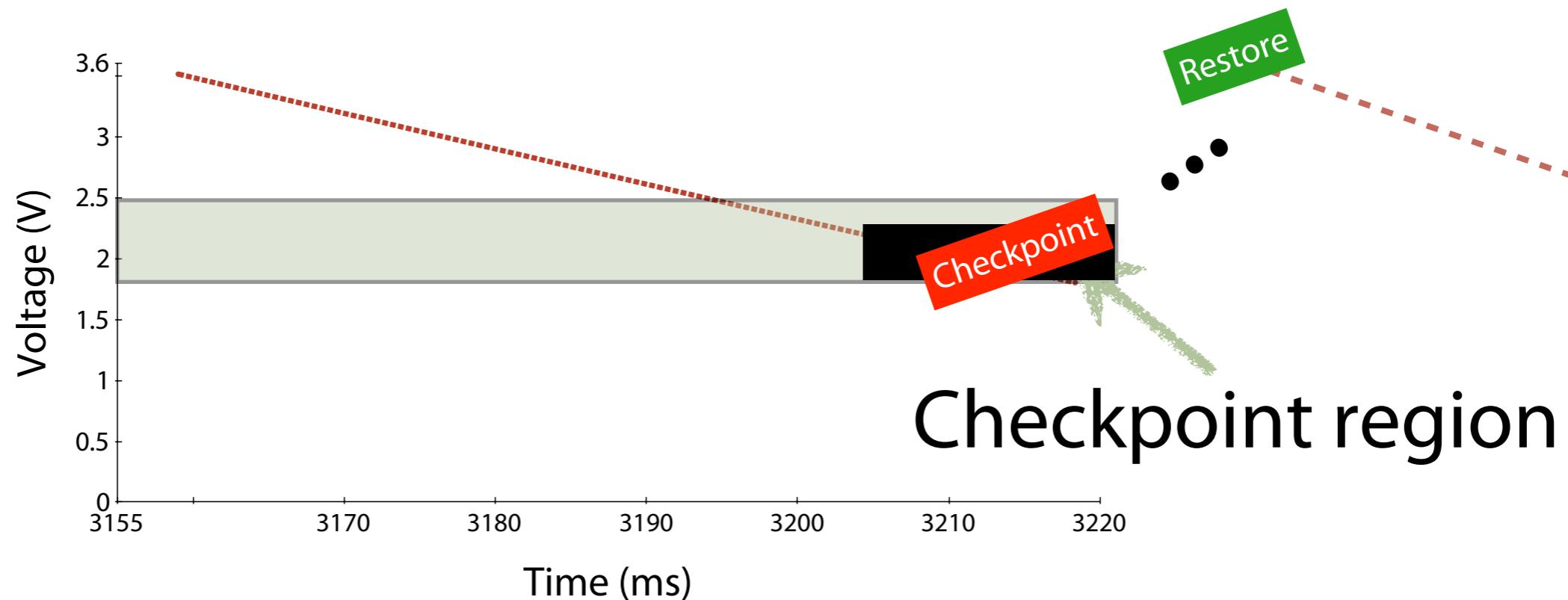


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Movie poster: publispain.com

# Mementos Approach



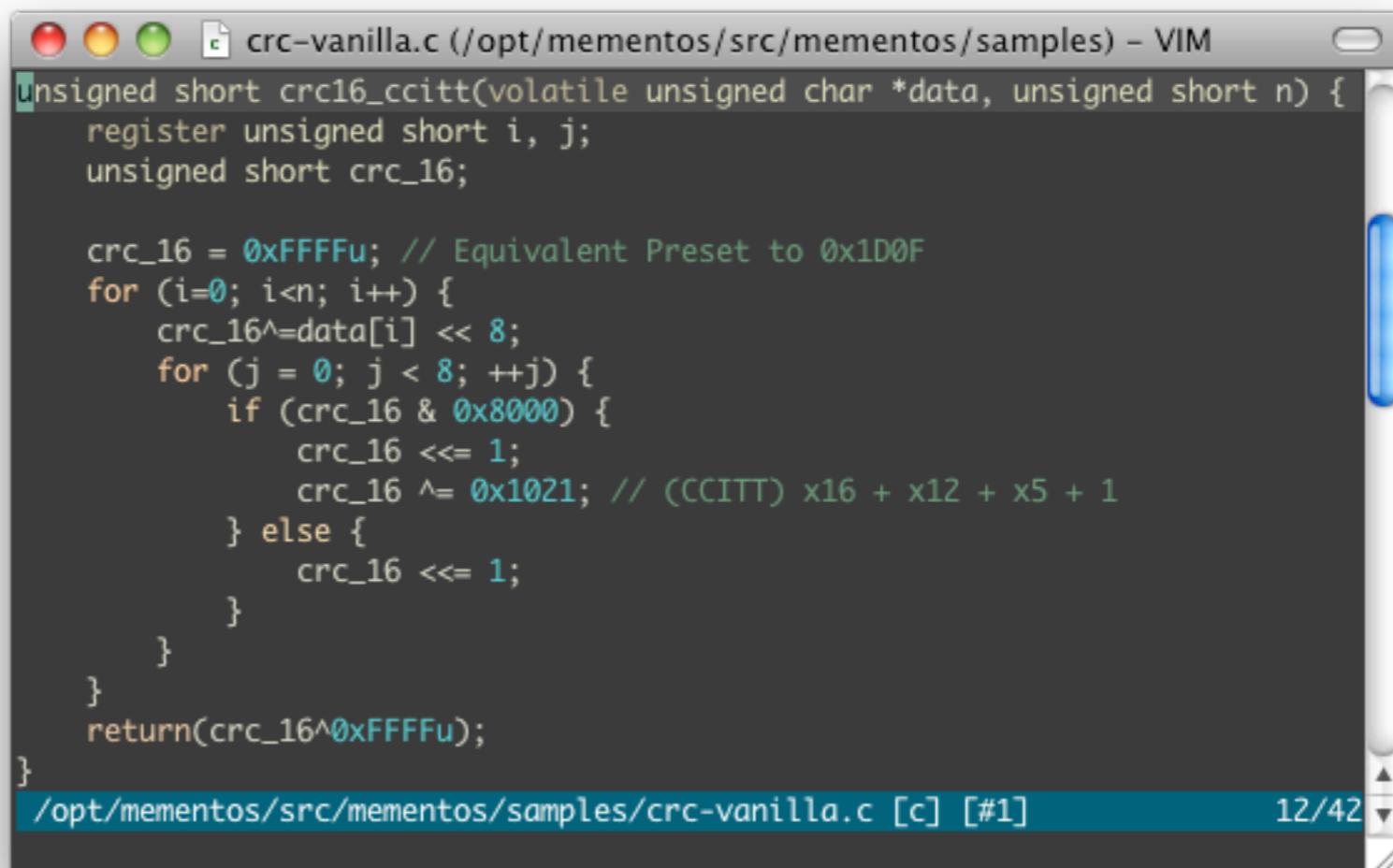
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Movie poster: publispain.com

# Running Example: CRC

- Compute CRC16-CCITT checksum over 2 KB data
- Tight nested loops
- 575,000 CPU cycles  $\sim$  575 ms



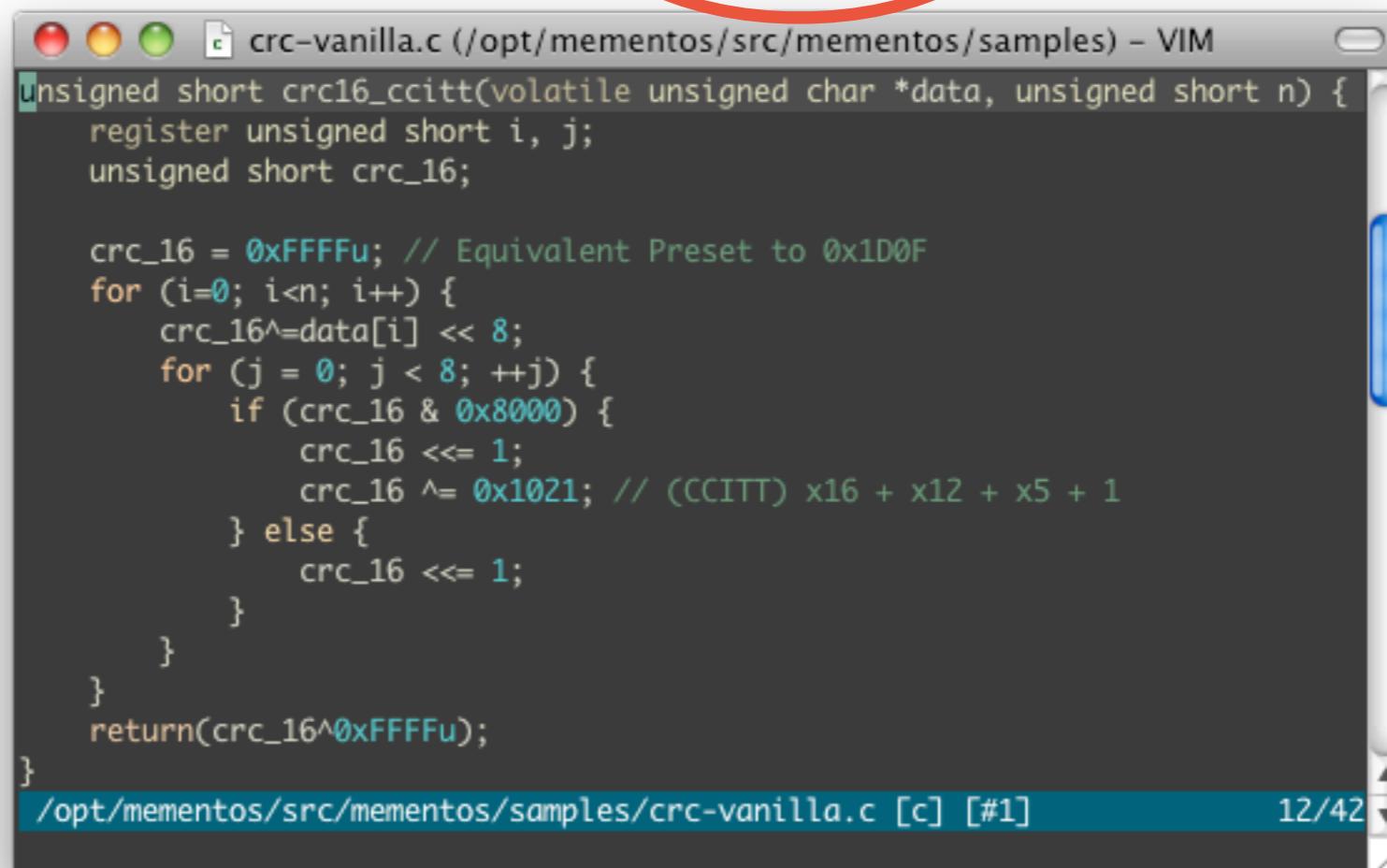
The screenshot shows a VIM editor window with the file `crc-vanilla.c` open. The code implements a CRC16-CCITT checksum function. It uses a tight nested loop structure to process data. The code includes comments explaining the equivalent preset value and the polynomial used.

```
unsigned short crc16_ccitt(volatile unsigned char *data, unsigned short n) {
    register unsigned short i, j;
    unsigned short crc_16;

    crc_16 = 0xFFFFu; // Equivalent Preset to 0x1D0F
    for (i=0; i<n; i++) {
        crc_16^=data[i] << 8;
        for (j = 0; j < 8; ++j) {
            if (crc_16 & 0x8000) {
                crc_16 <= 1;
                crc_16 ^= 0x1021; // (CCITT) x16 + x12 + x5 + 1
            } else {
                crc_16 <= 1;
            }
        }
    }
    return(crc_16^0xFFFFu);
}
```

# Running Example: CRC

- Compute CRC16-CCITT checksum over 2 KB data
- Tight nested loops Reboots every O(100) ms!
- 575,000 CPU cycles  $\sim 575$  ms



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```

/opt/mementos/src/mementos/samples/crc-vanilla.c [c] [#1] 12/42

# How to Use Mementos

---

Programmer

Mementos (our contributions)

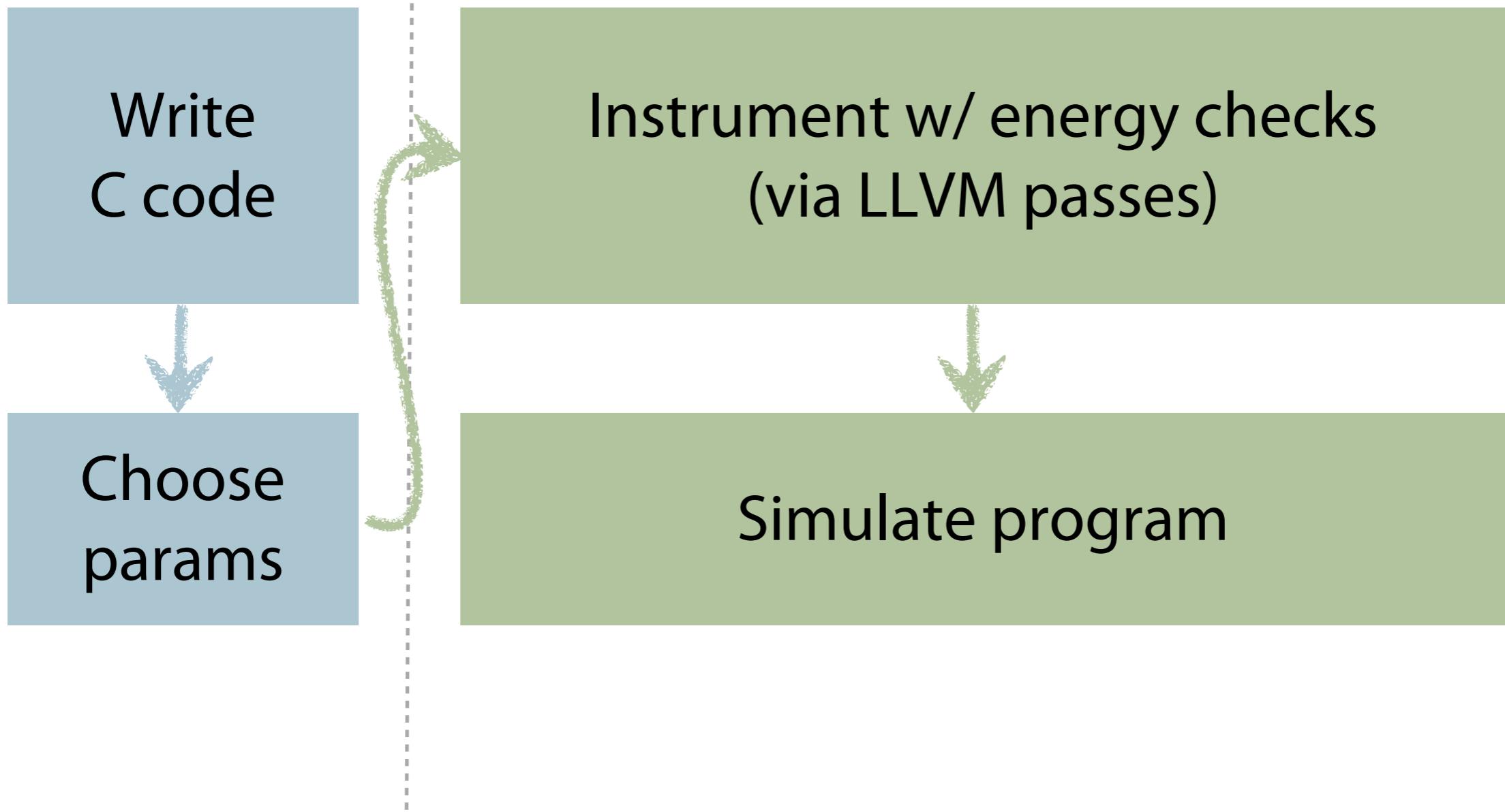
Write  
C code



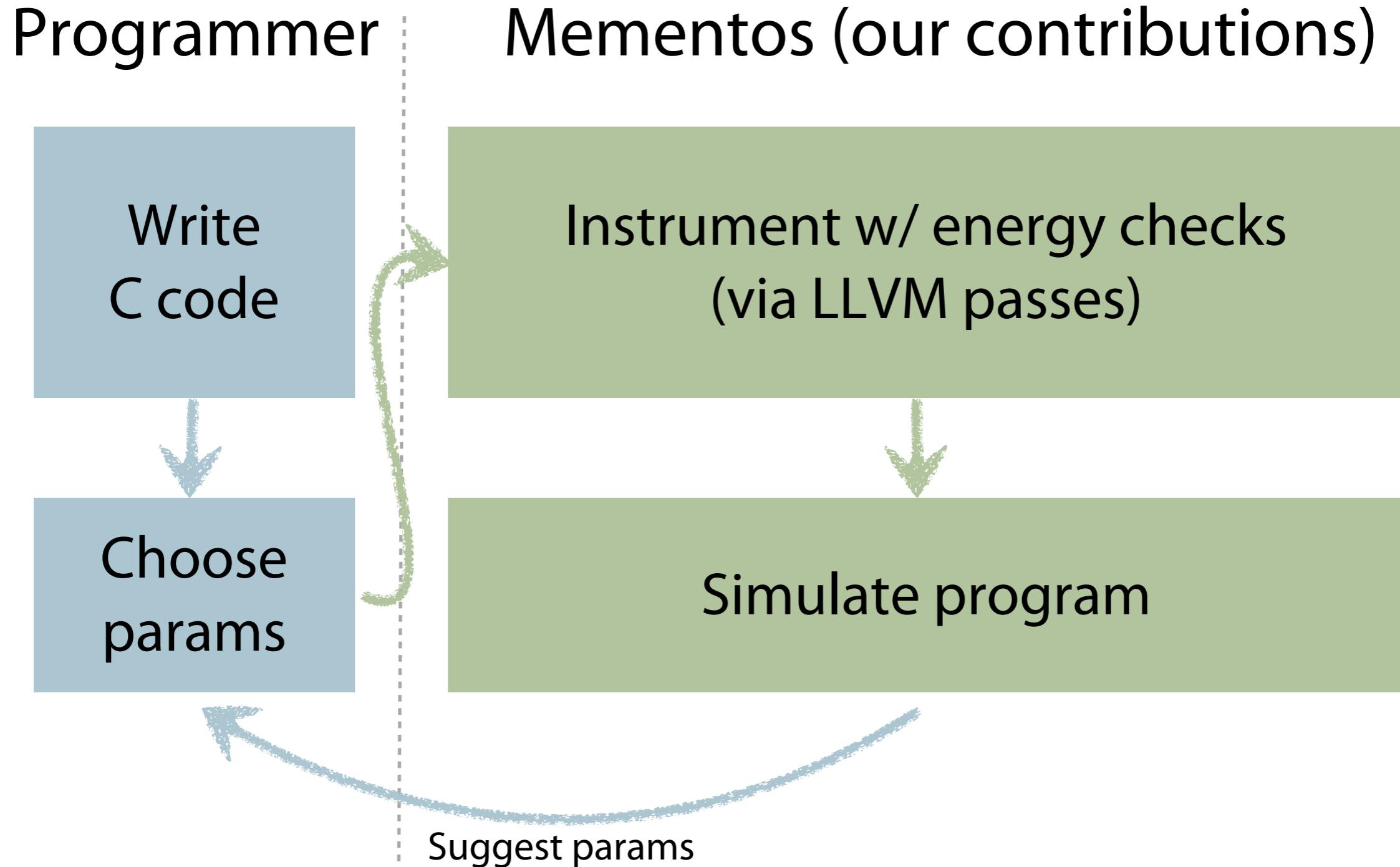
Choose  
params

# How to Use Mementos

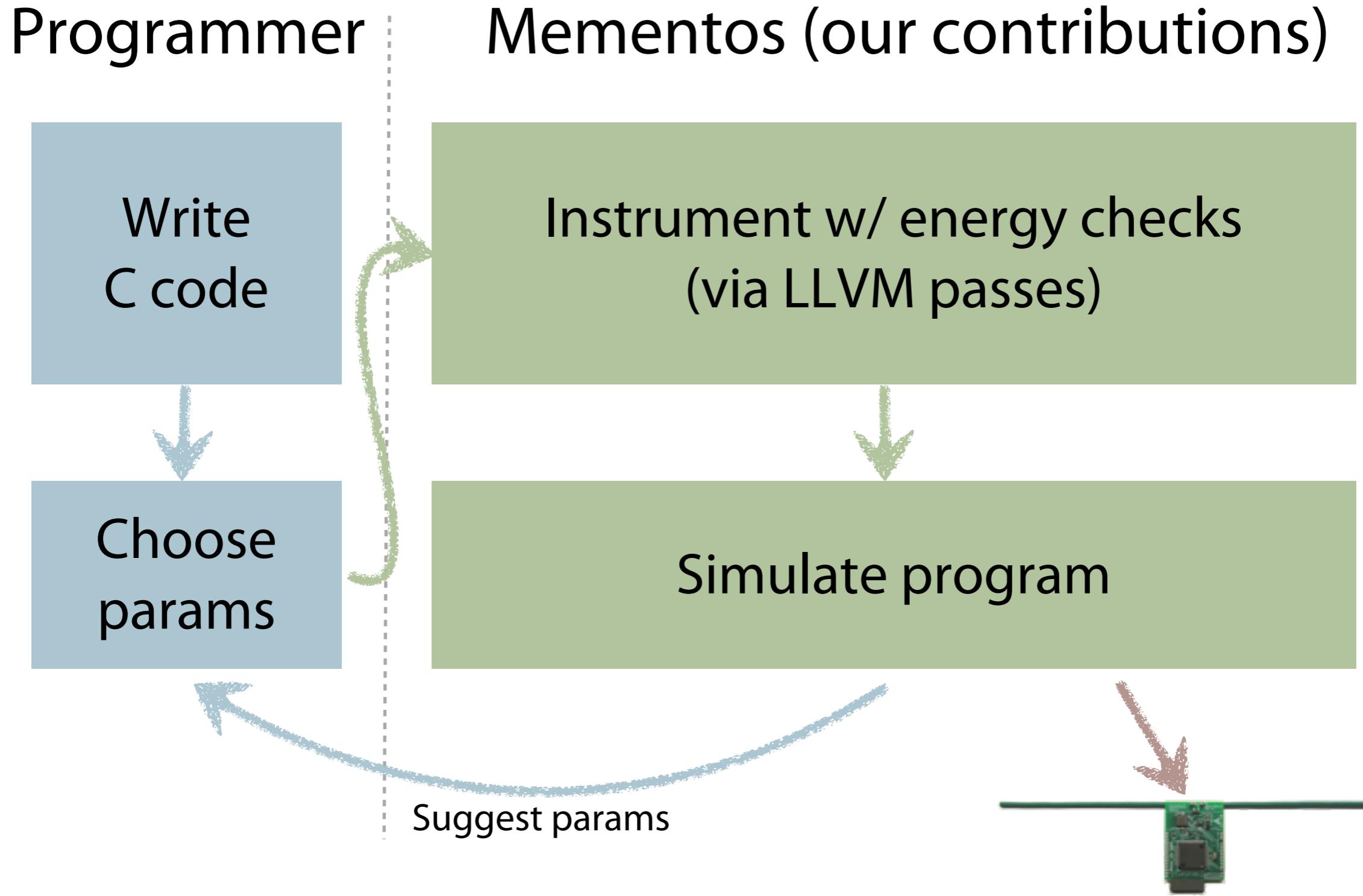
Programmer | Mementos (our contributions)



# How to Use Mementos



# How to Use Mementos



# Choosing Parameters (1/2)

---

Programmer

1) Instrumentation strategy

Write  
C code

Choose  
params



# Choosing Parameters (1/2)

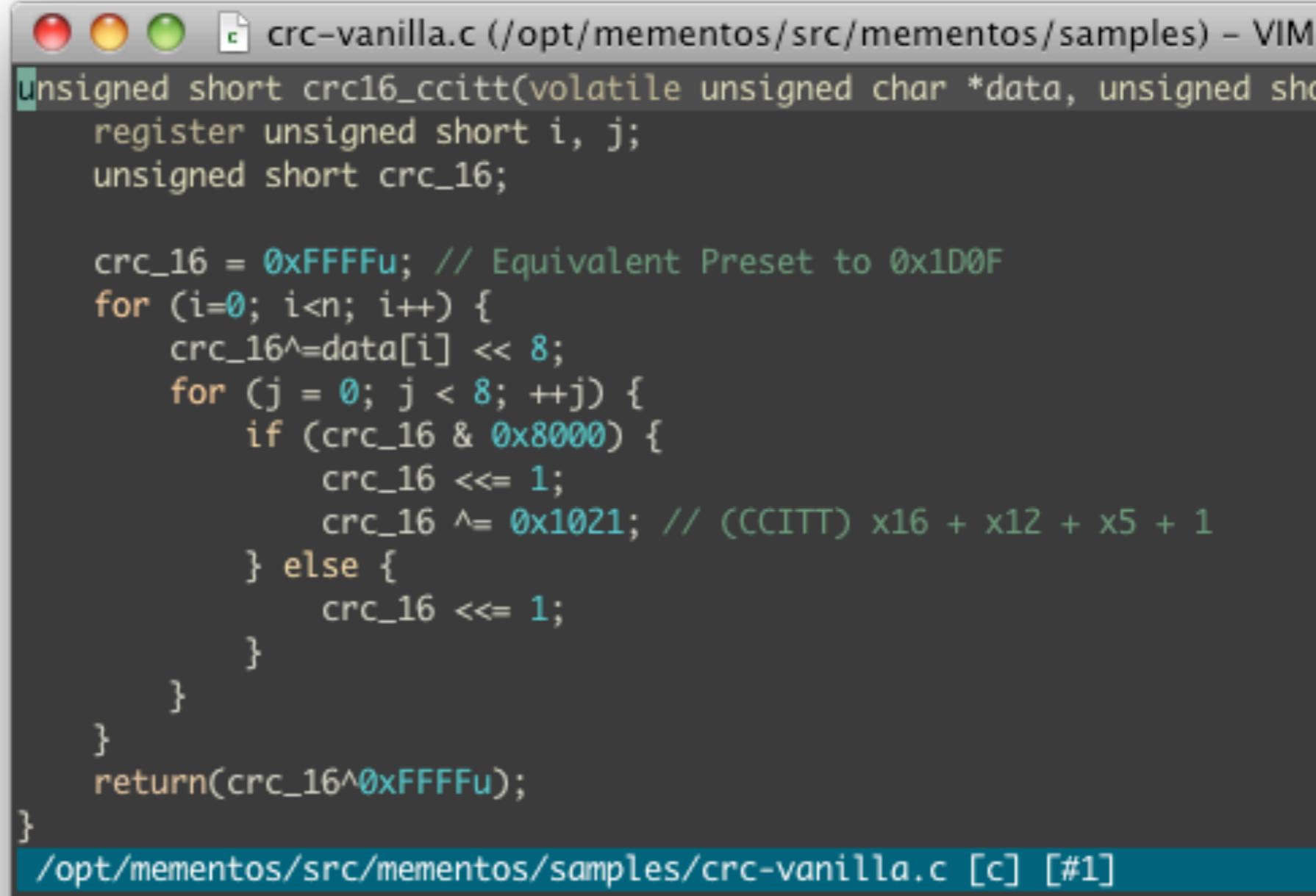
Programmer

Write  
C code



Choose  
params

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```

/opt/mementos/src/mementos/samples/crc-vanilla.c [c] [#1]

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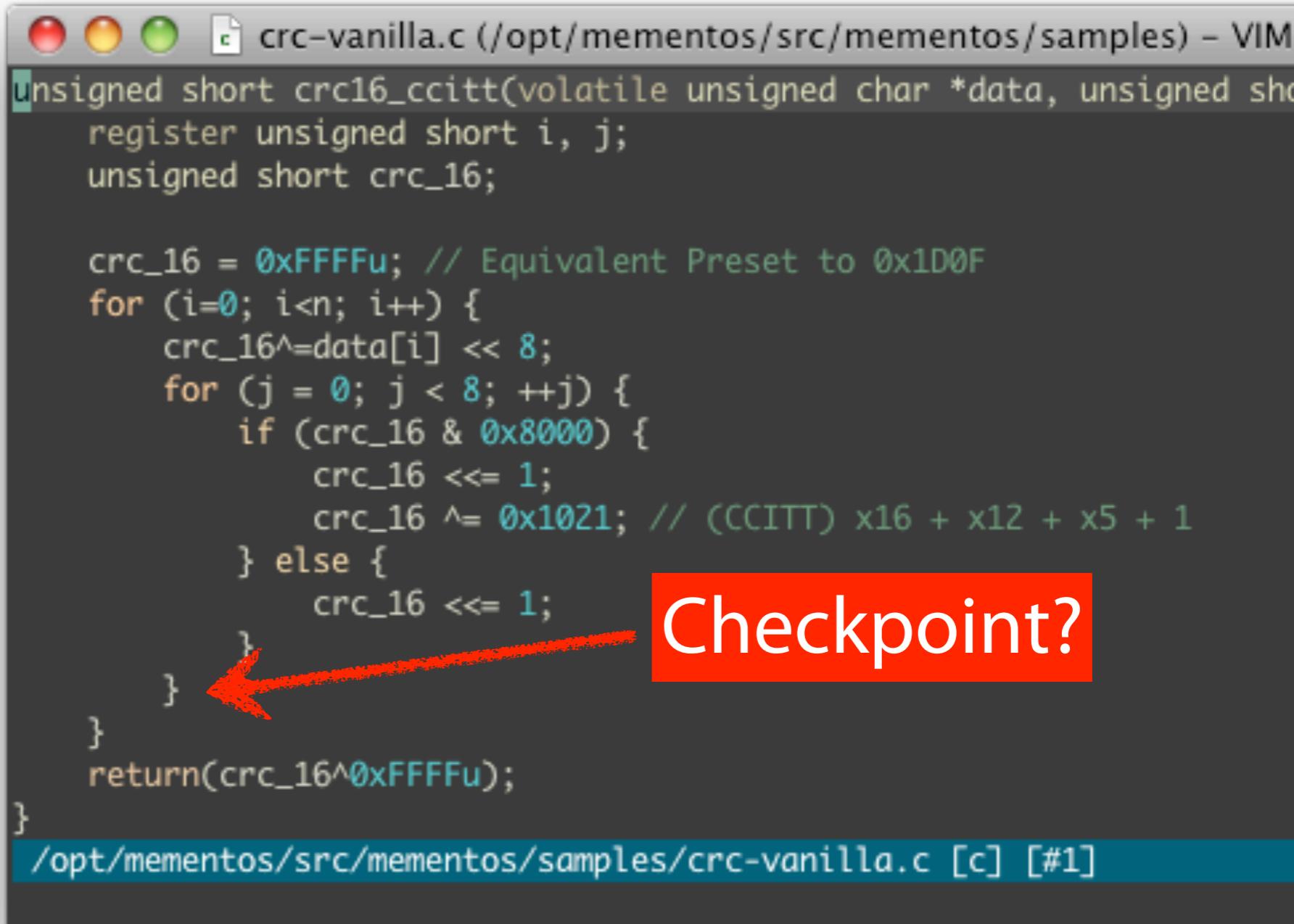
Programmer

Write  
C code



Choose  
params

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/opt/mementos/src/mementos/samples/crc-vanilla.c [c] [#1]

Checkpoint?

# Choosing Parameters (1/2)

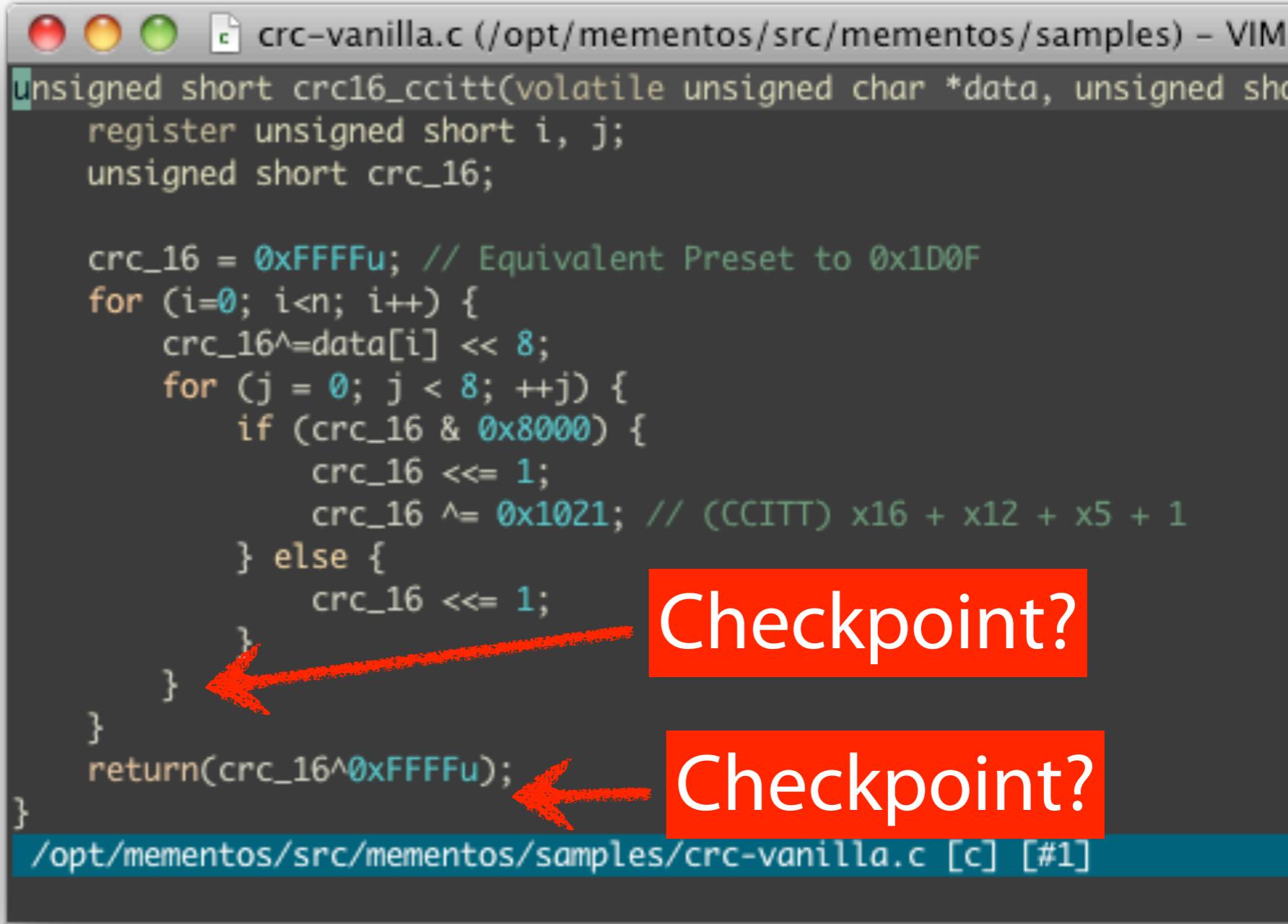
Programmer

Write  
C code



Choose  
params

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```

/opt/mementos/src/mementos/samples/crc-vanilla.c [c] [#1]

Checkpoint?

Checkpoint?

# Choosing Parameters (2/2)

---

Programmer

2) Checkpoint threshold  $V_{\text{thresh}}$

Write  
C code

Choose  
params



# Choosing Parameters (2/2)

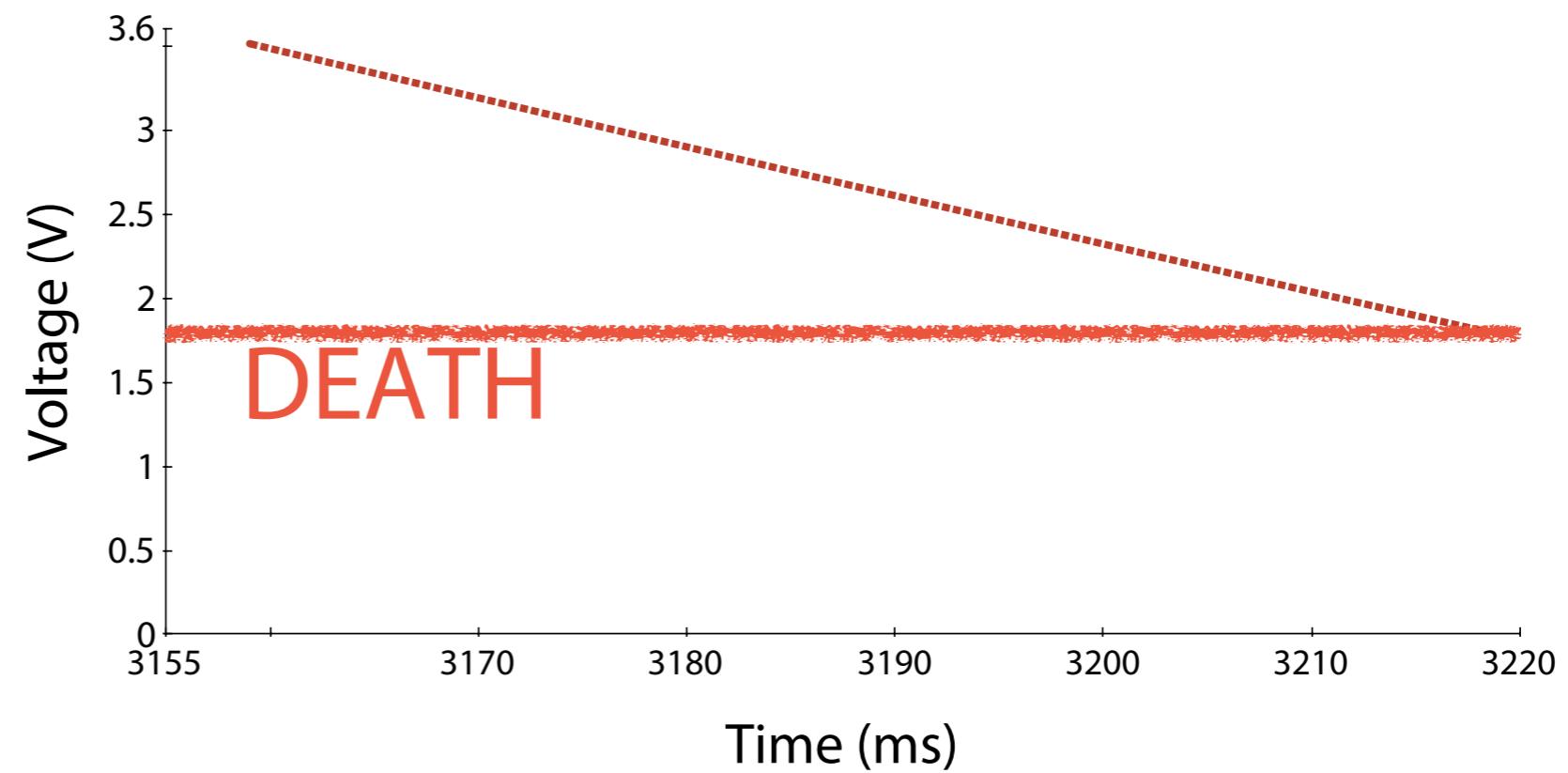
Programmer

Write  
C code



Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



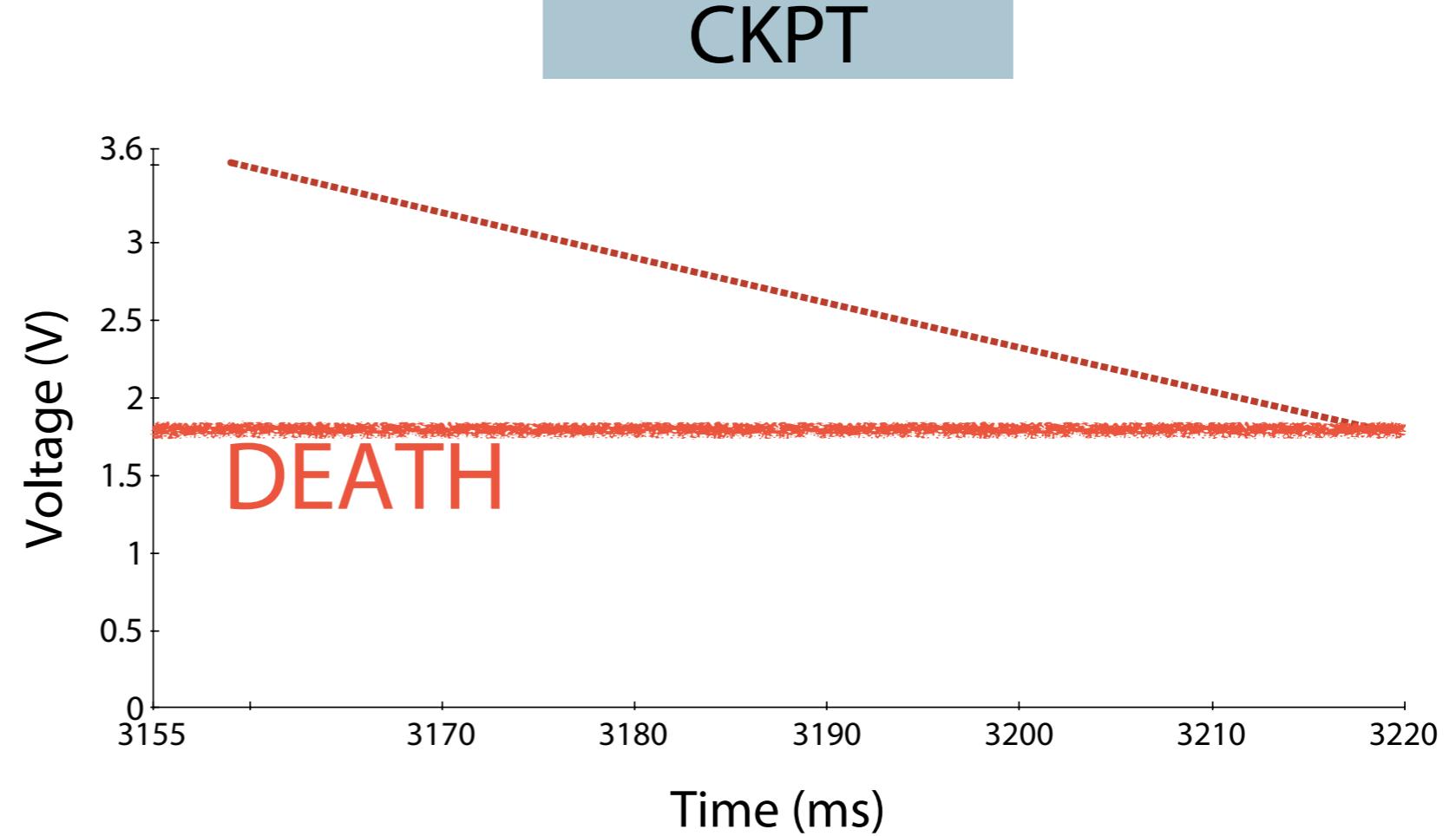
# Choosing Parameters (2/2)

Programmer

Write  
C code

Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



# Choosing Parameters (2/2)

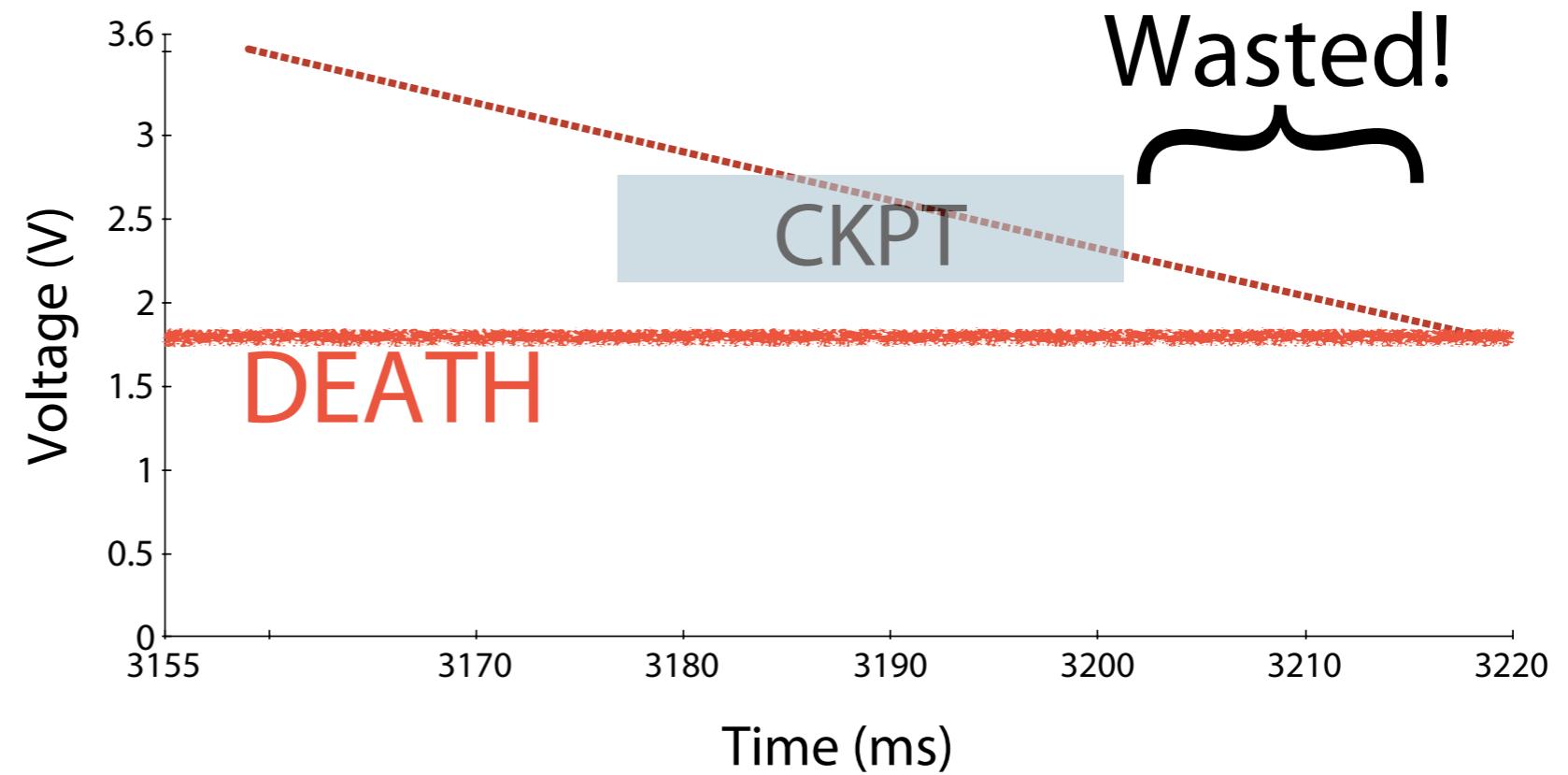
Programmer

Write  
C code



Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



# Choosing Parameters (2/2)

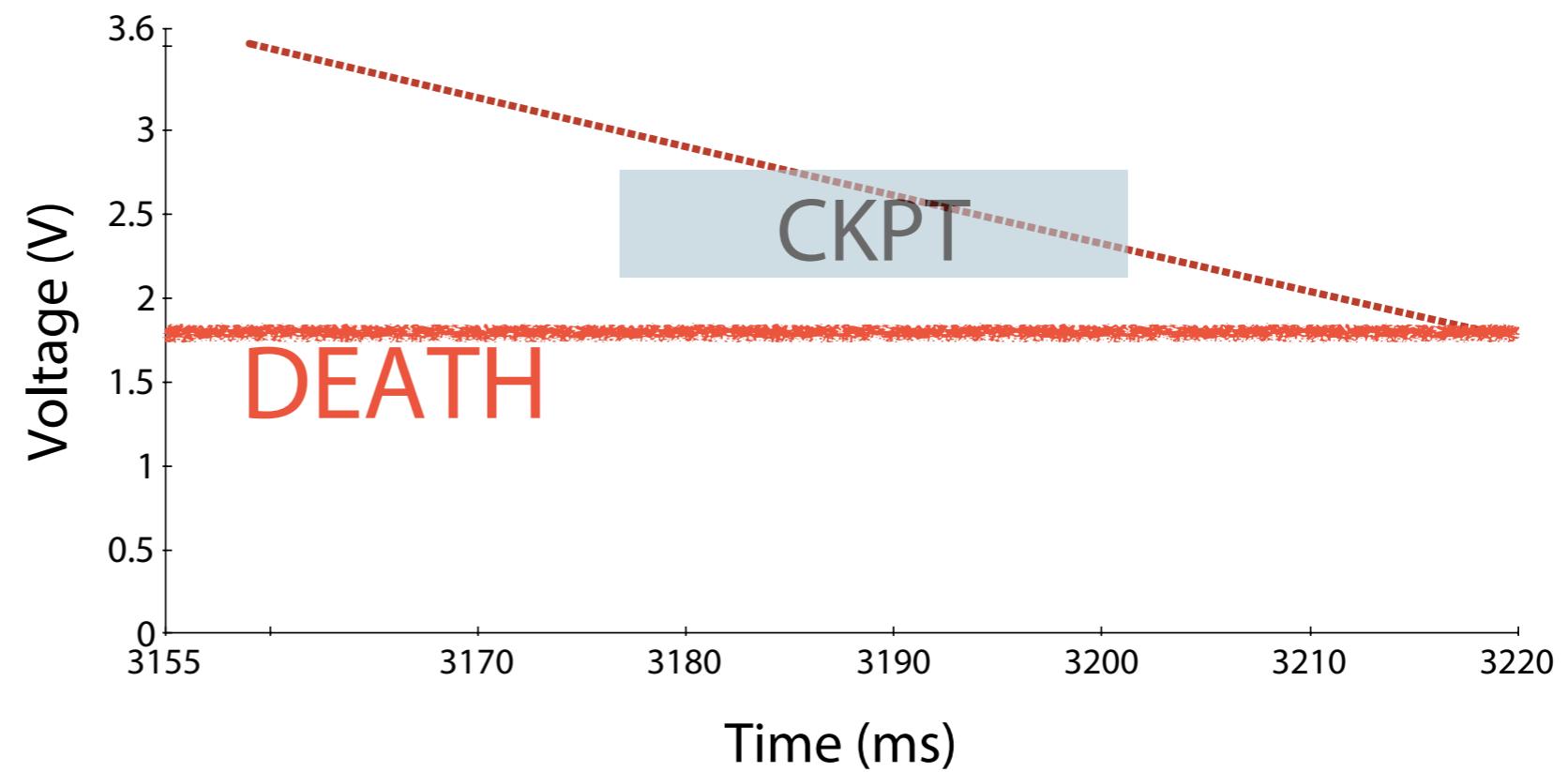
Programmer

Write  
C code



Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



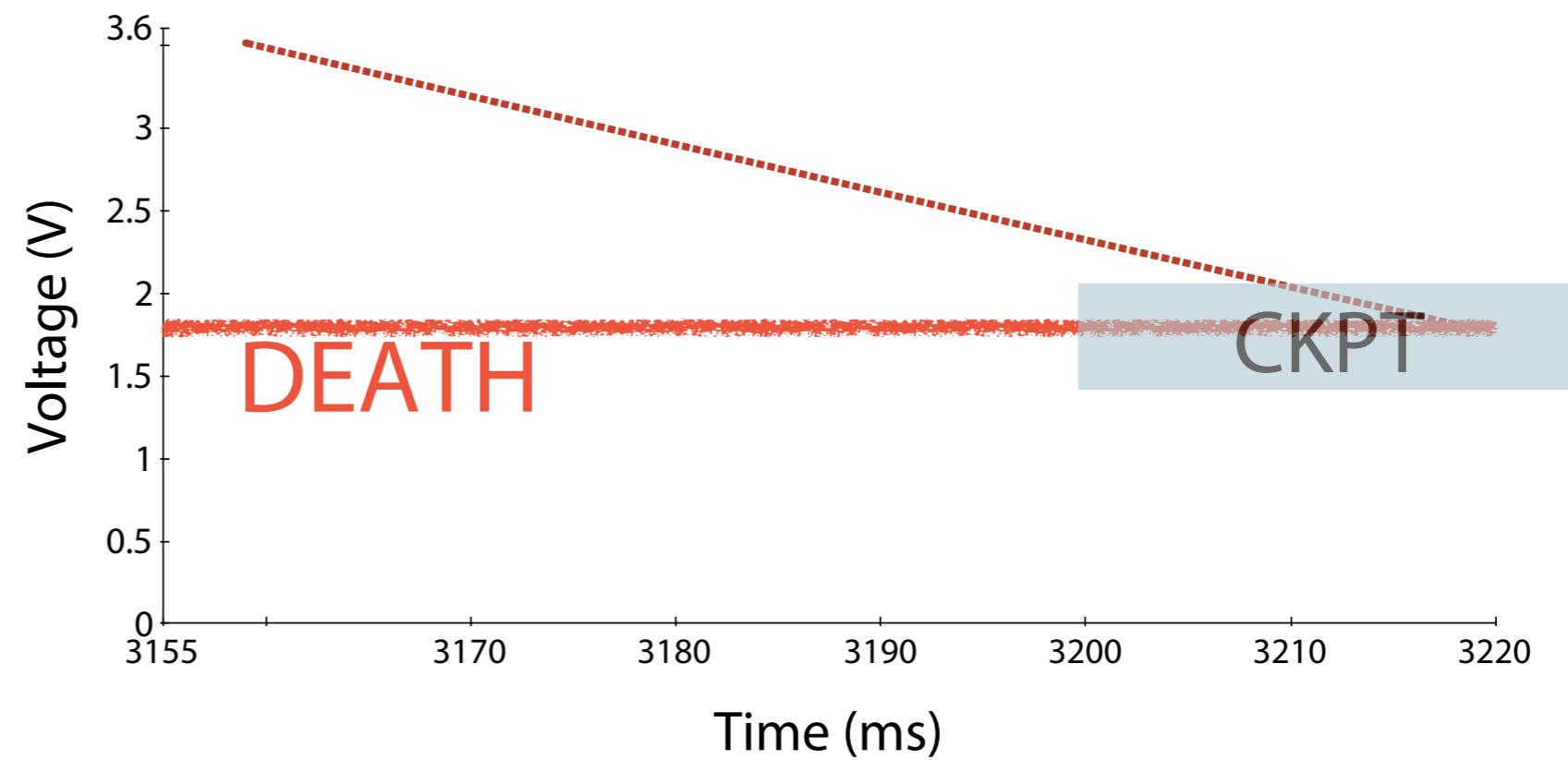
# Choosing Parameters (2/2)

Programmer

Write  
C code

Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



# Choosing Parameters (2/2)

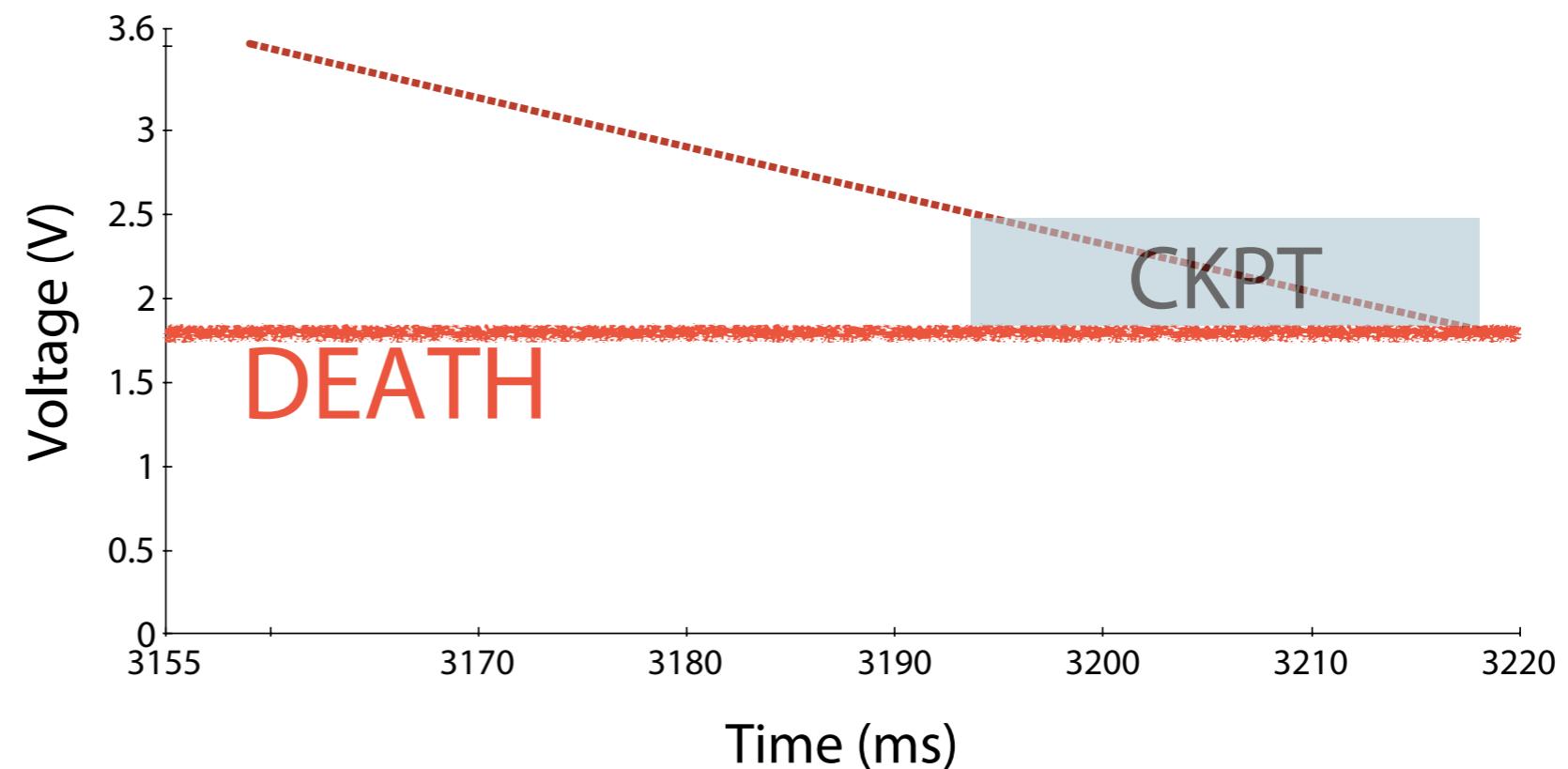
Programmer

Write  
C code



Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



# Choosing Parameters (2/2)

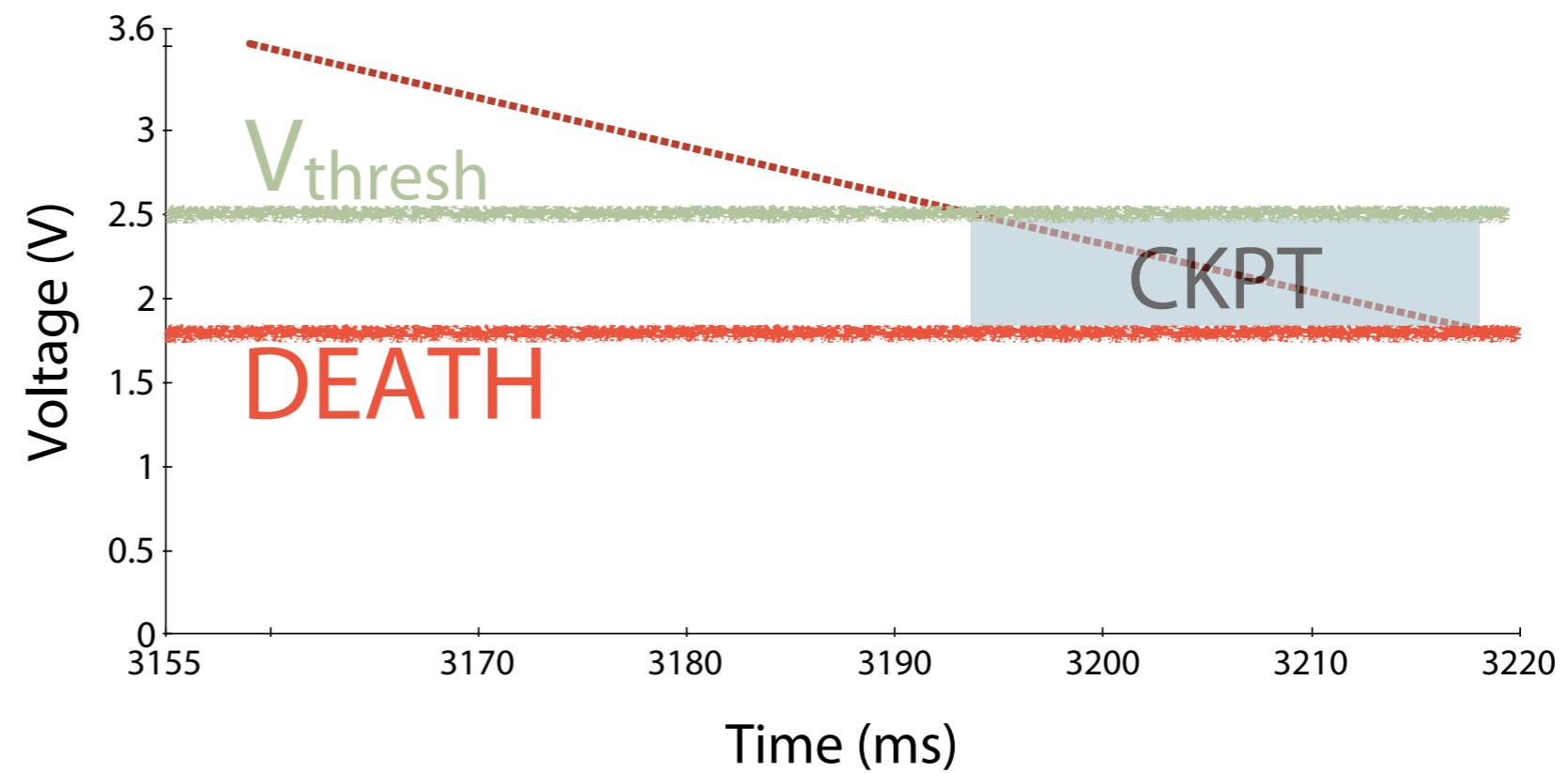
Programmer

Write  
C code



Choose  
params

2) Checkpoint threshold  $V_{\text{thresh}}$



# Assorted Challenges

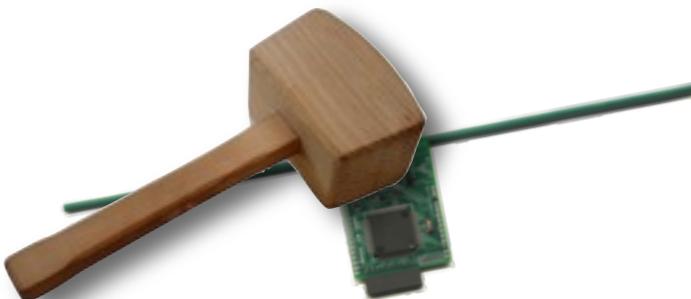
---

**Checkpointing isn't trivial in this context**

- No FTL; manage flash ourselves
- Can't overwrite arbitrary bit patterns in flash memory → tricky checkpoint maintenance

**Working on these devices is painful**

- Fickle harvesting → runs not reproducible
- Limited visibility into running hardware

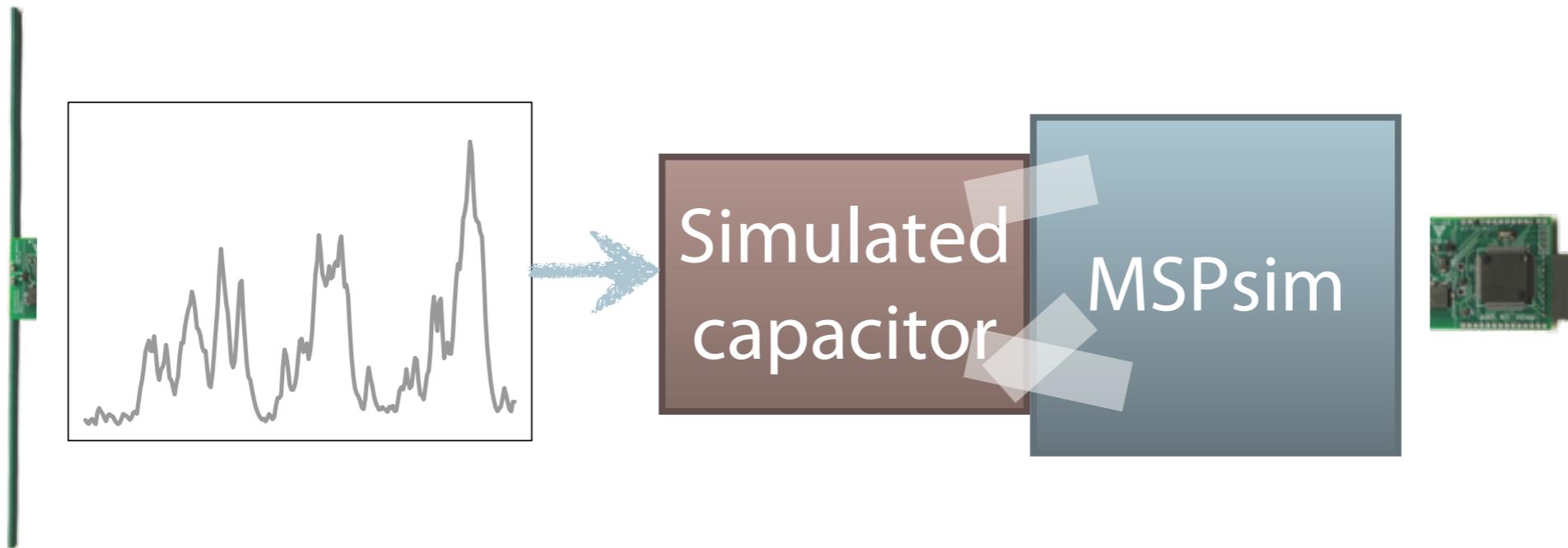


# Trace-Driven Simulator



- Based on MSPsim — cycle-accurate, open-source MSP430 simulator [EWSN '07]
- We augmented MSPsim with notions of energy (harvester, capacitor, power loss)

# Trace-Driven Simulator



- Based on MSPsim — cycle-accurate, open-source MSP430 simulator [EWSN '07]
- We augmented MSPsim with notions of energy (harvester, capacitor, power loss)

# Accurate Energy Simulation

---

- Simulated capacitor obeys capacitor equations to buffer incoming energy
- Validated with microbenchmarks (all chip modes, all instruction classes)
  - Measured MSP430 current down to  $\mu\text{A}$
  - Details in paper



# Straightforward Simulation

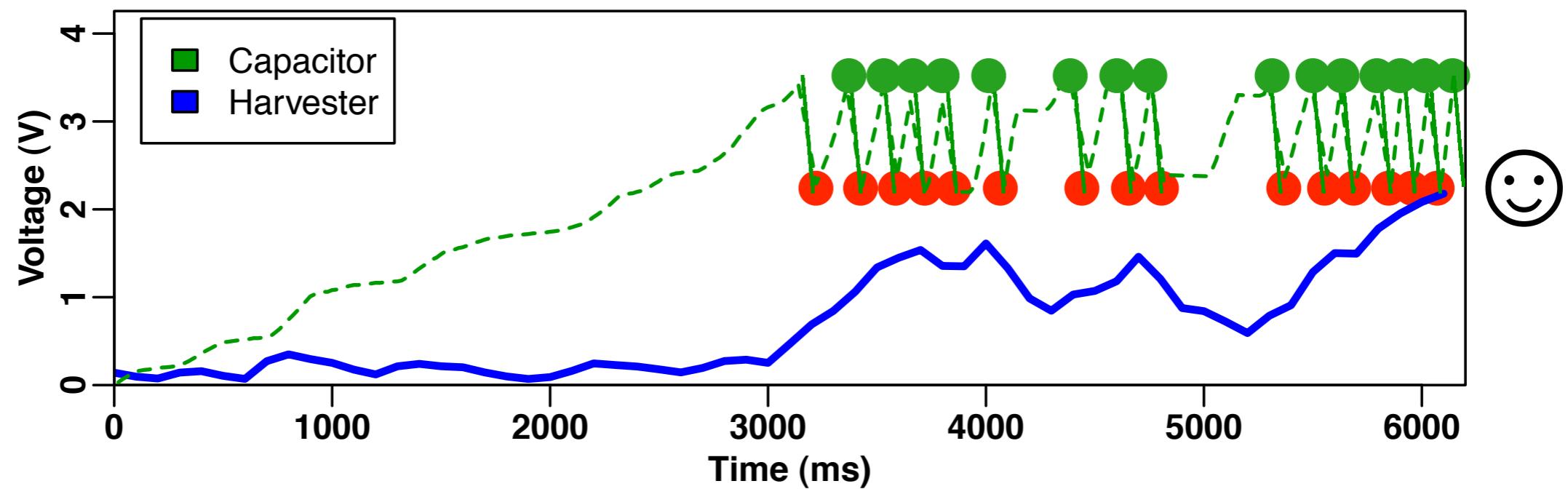
---

- Simulator input:  $\langle \text{executable}, \text{voltage trace} \rangle$
- Output:  $\langle \# \text{ reboots to completion}, \# \text{ CPU cycles}, \text{total time}, \text{execution trace} \rangle$

# Straightforward Simulation

---

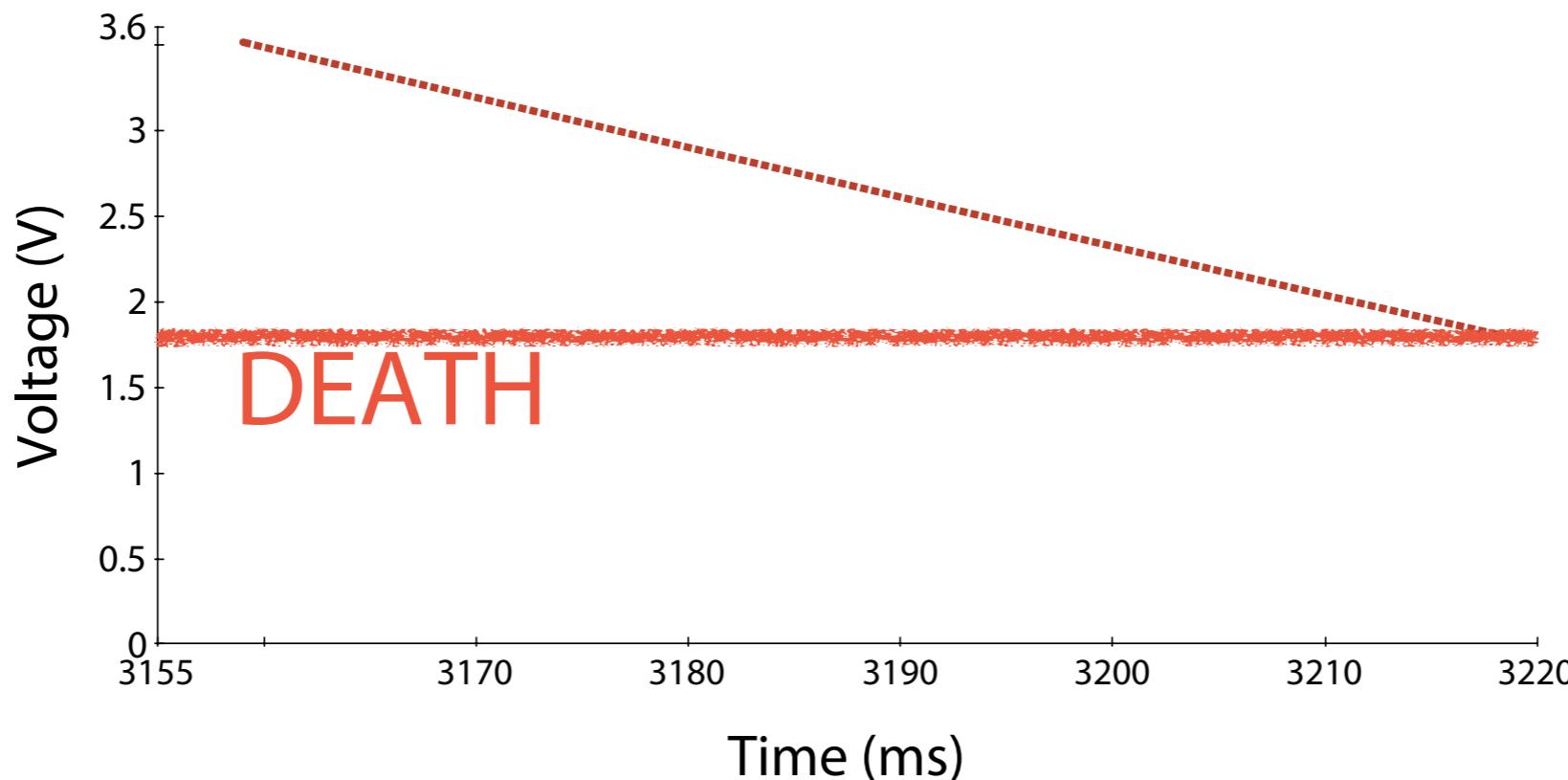
- Simulator input: <*executable, voltage trace*>
- Output: <*# reboots to completion, # CPU cycles, total time, execution trace*>



# Simulation with an Oracle

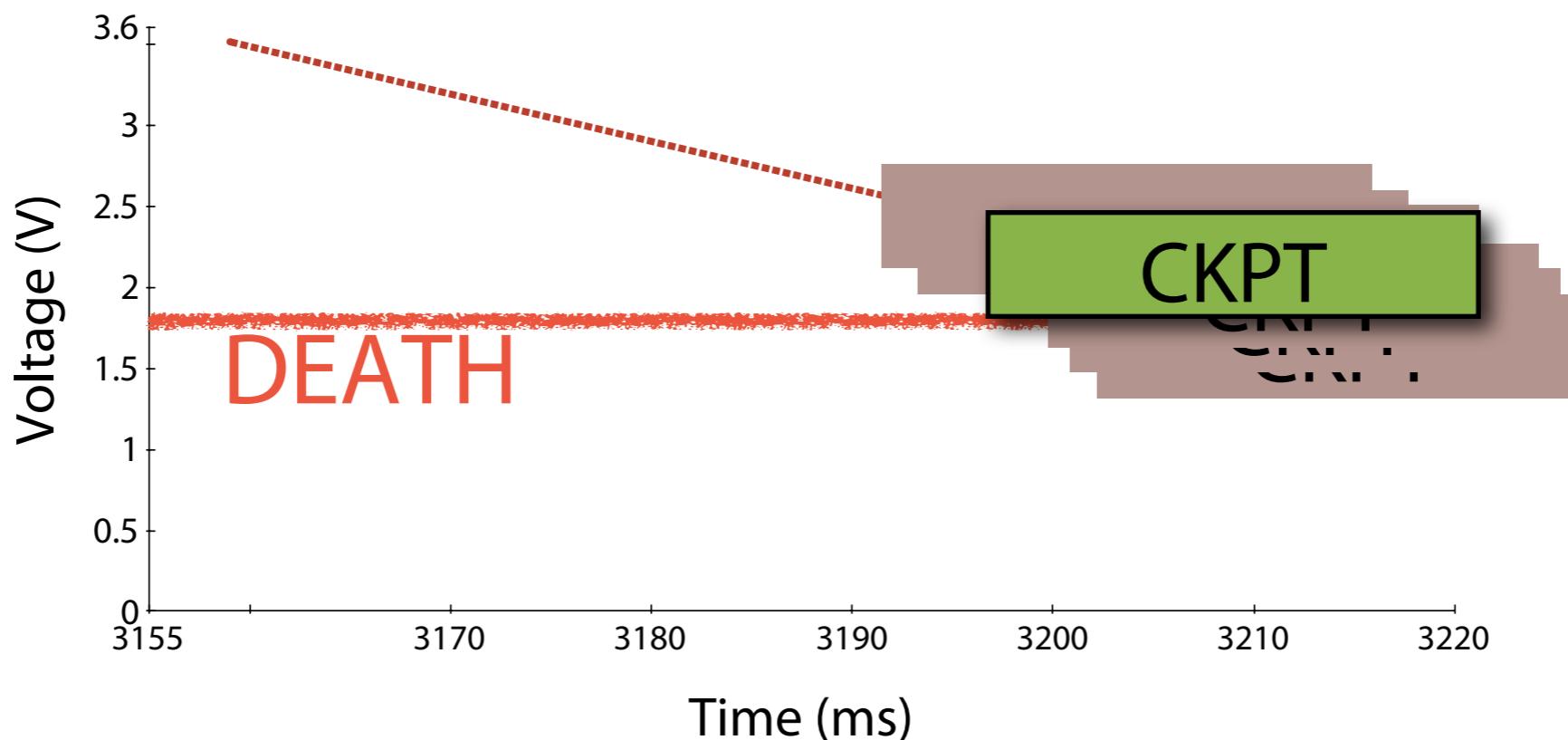
---

- Checkpoint oracle finds last practicable opportunity by binary search on  $V_{\text{thresh}}$ 
  - ▶ Uninstrumented code → best-case estimate
  - ▶ Final report: *lower bound* for  $V_{\text{thresh}}$



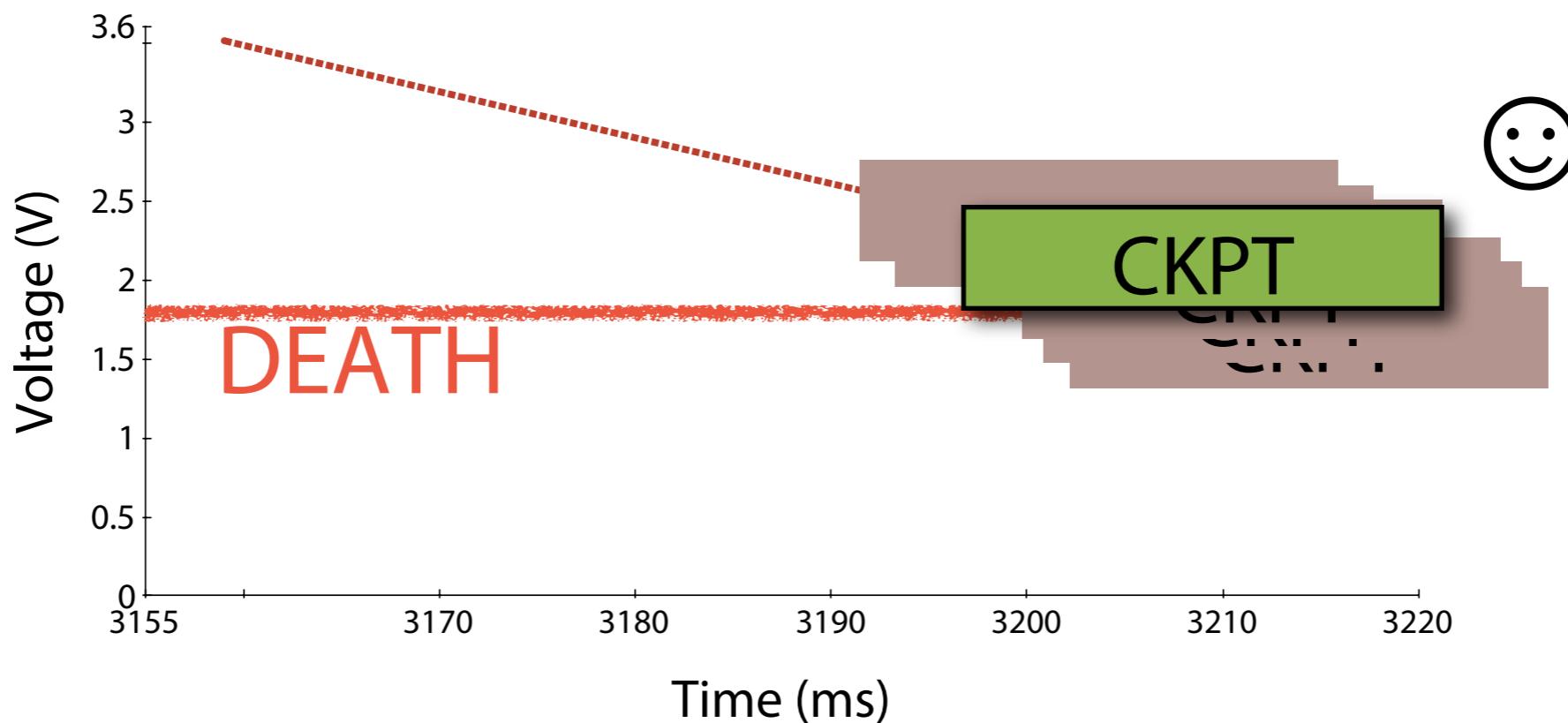
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# Simulation with an Oracle

- Checkpoint oracle finds last practicable opportunity by binary search on  $V_{\text{thresh}}$ 
  - ▶ Uninstrumented code → best-case estimate
  - ▶ Final report: *lower bound* for  $V_{\text{thresh}}$



# Evaluation

---

- **High-level:** Mementos splits execution in simulation and on hardware
- Focus on CRC example test case
- Baselines:
  - ▶ Execution without Mementos
  - ▶ Execution against checkpoint oracle

# Constant Part of Overhead

---

- Impact on code memory (NVRAM):
  - ▶ 2.4 KB for Mementos library
  - ▶ 1 KB reserved checkpoint storage
- Impact on run time:
  - ▶ ~0.1 ms per energy check (mostly ADC read)
  - ▶ CRC (46 bytes): checkpoint 4 ms, restore 2 ms

# Constant Part of Overhead

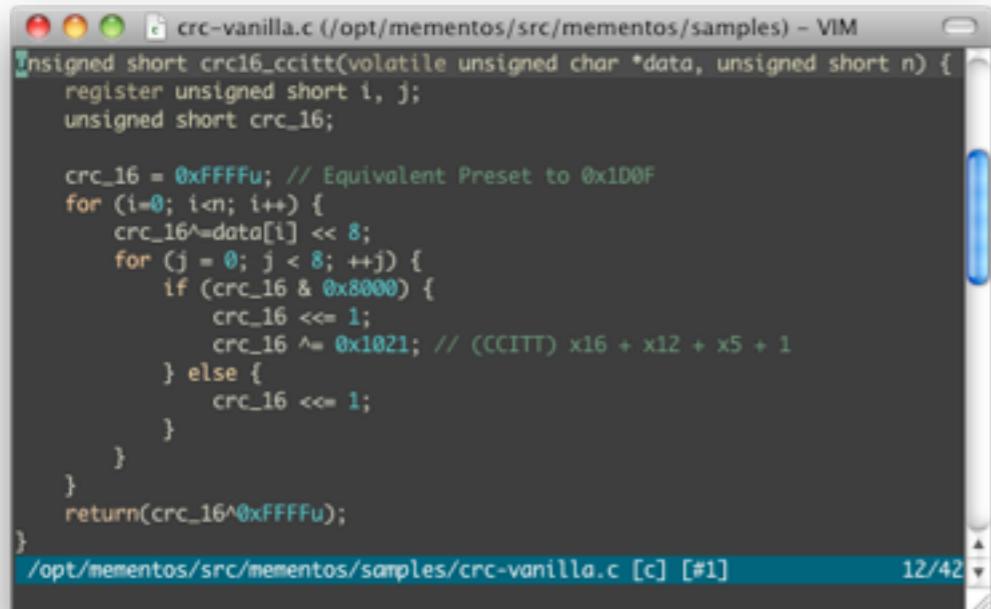
---

- Impact on code memory (NVRAM):
  - ▶ 2.4 KB for Mementos library
  - ▶ 1 KB reserved checkpoint storage
- Impact on run time:
  - ▶ ~0.1 ms per energy check (mostly ADC read)
  - ▶ CRC (46 bytes): checkpoint 4 ms, restore 2 ms

2 ms boot vs. TinyOS  $\geq 100$  ms



# CRC Test Case



A screenshot of a VIM editor window displaying a C program named `crc-vanilla.c`. The code implements a CRC16\_CCITT checksum algorithm. It starts with a function declaration:

```
unsigned short crc16_ccitt(volatile unsigned char *data, unsigned short n) {
```

It initializes the CRC register to `0xFFFFu` and processes the data in 8-bit chunks. For each byte, it shifts the current CRC value left by 8 bits and then performs a polynomial division using the CCITT polynomial  $x^{16} + x^{12} + x^5 + 1$ :

```
    crc_16 ^= data[i] << 8;
    for (j = 0; j < 8; ++j) {
        if (crc_16 & 0x8000) {
            crc_16 <= 1;
            crc_16 ^= 0x1021; // (CCITT) x16 + x12 + x5 + 1
        } else {
            crc_16 <= 1;
        }
    }
}
```

The function concludes by returning the final CRC value after applying a preset value of `0xFFFFu`:

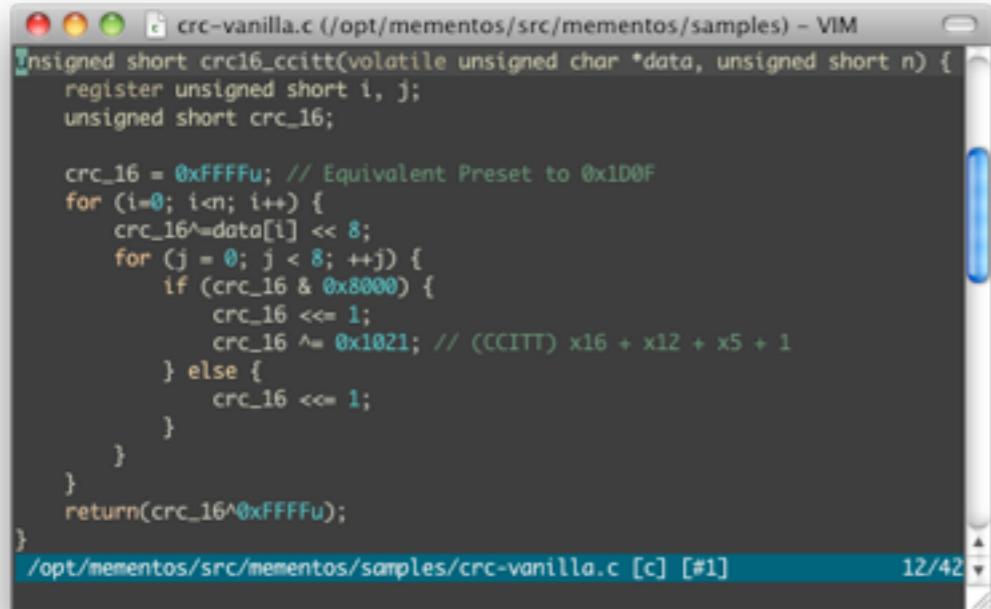
```
return(crc_16 ^ 0xFFFFu);
}
```

The status bar at the bottom of the VIM window shows the file path `/opt/mementos/src/mementos/samples/crc-vanilla.c [c] [#1]` and the line number `12/42`.

Uninstrumented,  
unlimited energy:

575,315 cycles  
575 ms

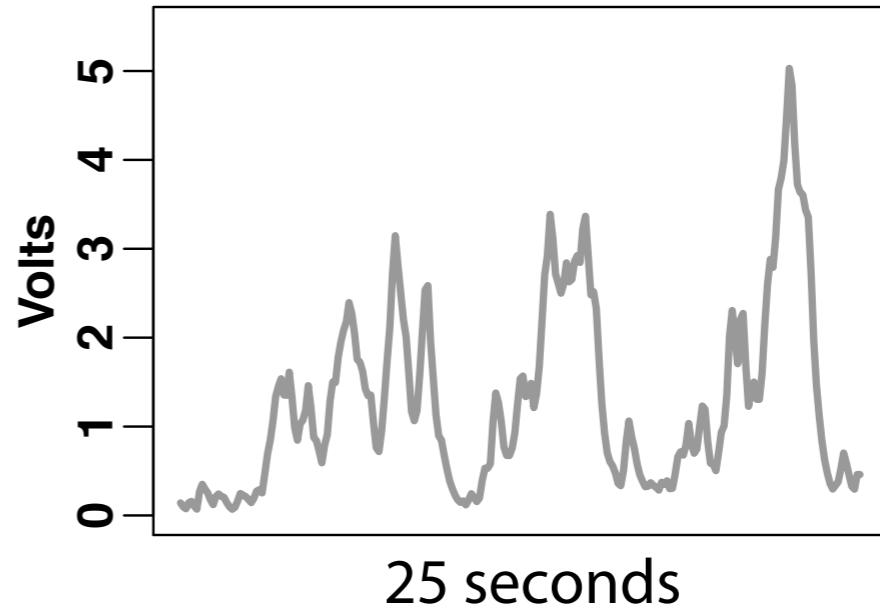
# CRC Test Case



A screenshot of a VIM editor window displaying the file `crc-vanilla.c`. The code implements a CRC16\_CCITT checksum calculation. It uses a preset value of `0xFFFFu` and iterates over the input data, shifting it 8 bits at a time and applying a polynomial of `0x1021` if the 16th bit is set. The code is annotated with comments explaining the logic.

```
unsigned short crc16_ccitt(volatile unsigned char *data, unsigned short n) {
    register unsigned short i, j;
    unsigned short crc_16;

    crc_16 = 0xFFFFu; // Equivalent Preset to 0x1D0F
    for (i=0; i<n; i++) {
        crc_16^=data[i] << 8;
        for (j = 0; j < 8; ++j) {
            if (crc_16 & 0x8000) {
                crc_16 <= 1;
                crc_16 ^= 0x1021; // (CCITT) x16 + x12 + x5 + 1
            } else {
                crc_16 <= 1;
            }
        }
    }
    return(crc_16^0xFFFFu);
}
```



Uninstrumented,  
unlimited energy:

575,315 cycles  
575 ms

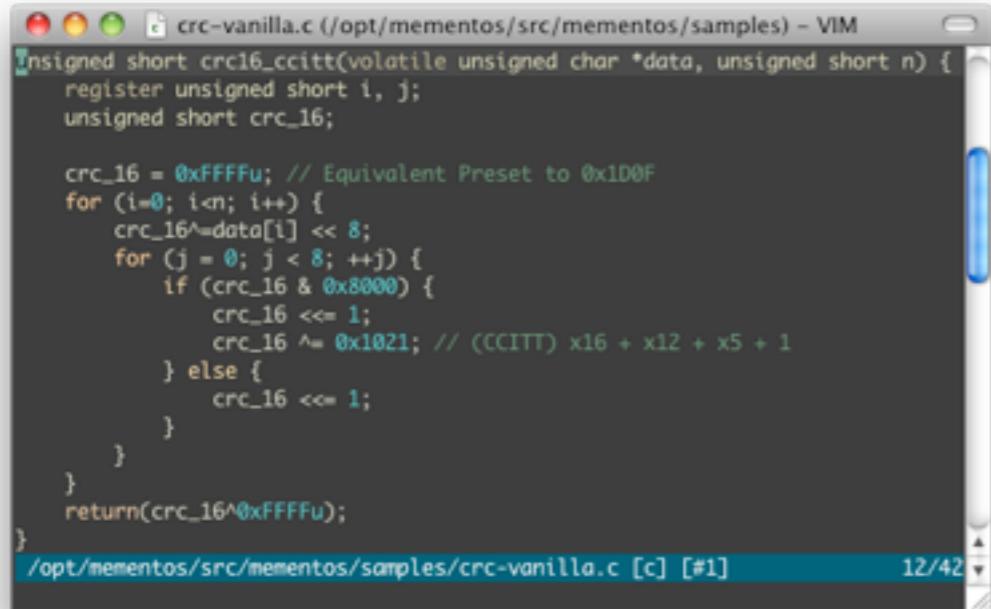
Oracle:

685,608 cycles  
~4,000 ms  
14 reboots  
 $V_{thresh} \geq 2.35 \text{ V}$

Best execution:

761,983 cycles  
6,145 ms  
16 reboots  
 $V_{thresh} = 2.6 \text{ V}$

# CRC Test Case



A screenshot of a VIM editor window displaying the file `crc-vanilla.c`. The code implements a CRC16\_CCITT checksum calculation. It uses a for loop to iterate over a buffer of length `n`, shifting the current CRC value left by 8 bits and applying a polynomial if the 16th bit is set. The final result is XORed with `0xFFFF`.

```
signed short crc16_ccitt(volatile unsigned char *data, unsigned short n) {
    register unsigned short i, j;
    unsigned short crc_16;

    crc_16 = 0xFFFF; // Equivalent Preset to 0x1D0F
    for (i=0; i<n; i++) {
        crc_16^=data[i] << 8;
        for (j = 0; j < 8; ++j) {
            if (crc_16 & 0x8000) {
                crc_16 <= 1;
                crc_16 ^= 0x1021; // (CCITT) x16 + x12 + x5 + 1
            } else {
                crc_16 <= 1;
            }
        }
    }
    return(crc_16^0xFFFF);
}
```

10x blowup is better than  
never finishing at all!

Uninstrumented,  
unlimited energy:

575,315 cycles  
575 ms

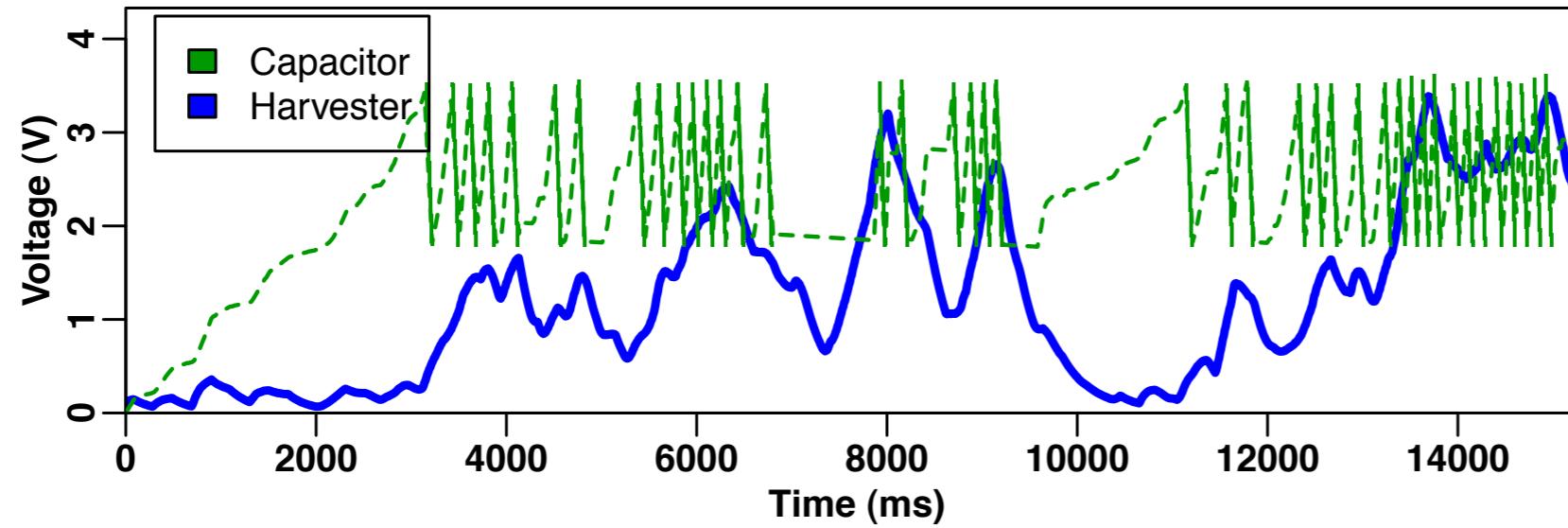
Oracle:

685,608 cycles  
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14 reboots  
 $V_{thresh} \geq 2.35 \text{ V}$

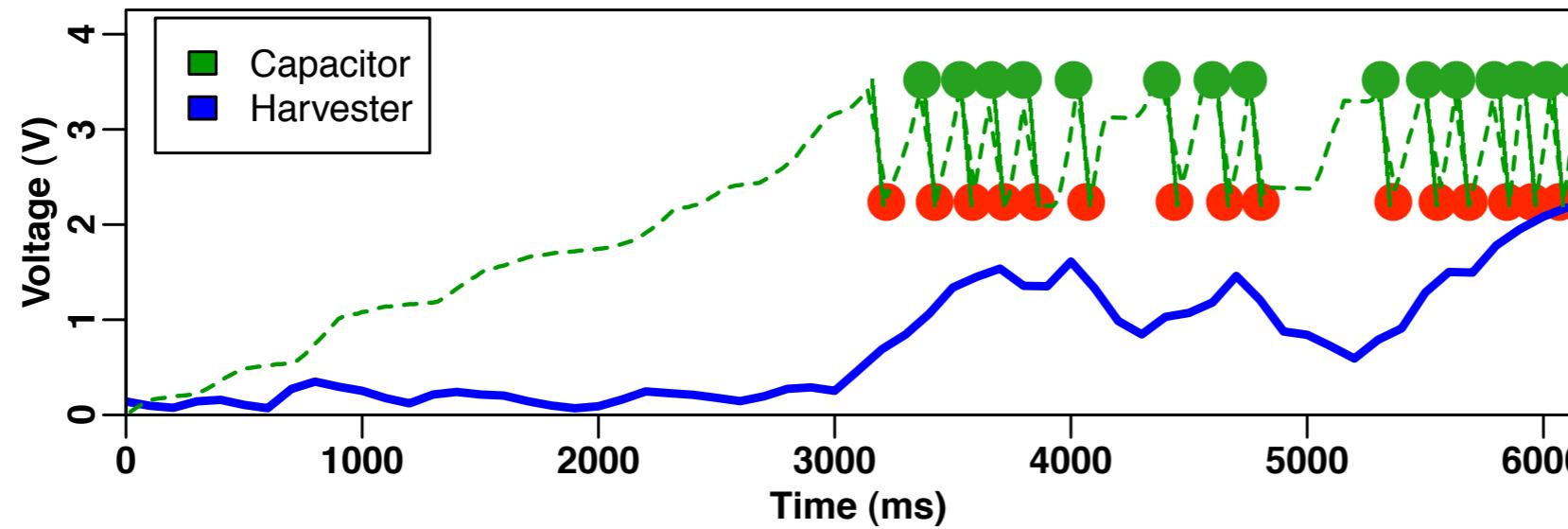
Best execution:

761,983 cycles  
6,145 ms  
16 reboots  
 $V_{thresh} = 2.6 \text{ V}$

# With and Without Mementos



CRC  
w/o Mementos:  
never finishes



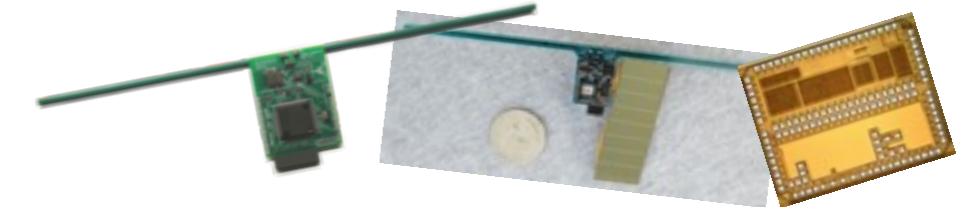
😊 CRC  
w/ Mementos:  
16 reboots

Oracle: 14 reboots

# Related Work

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- **RFID-scale devices**
  - ▶ Mementos workshop paper [HotPower '08]
  - ▶ Dewdrop scheduler for RFID-scale devices [NSDI 2011]
  - ▶ WISP [IEEE TIM '08] and friends
- **Checkpointing**
  - ▶ Sensornet checkpointing [EWSN '09]
  - ▶ Checkpointing for process migration (Condor [ICDCS '88], Porch [IEEE Micro '98])



# Extensions

---

- Dynamic or randomized  $V_{\text{thresh}}$  adaptation
- NVRAM technology (PCM? FeRAM?)
- Smarter checkpointing (incremental, LVA...)
- Integrate with asynchronous communications  
on upcoming RFID-scale prototype (June '11)



# Mementos Conclusion

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- Energy-aware checkpoints for computation on batteryless RFID-scale devices
- Tools available today; built on LLVM and MSPsim
- Applications: implantable devices, insect-scale tracking, infrastructure monitoring...

# Acknowledgements

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# Your Homework

---

Get Mementos, simulator, hardware:  
<http://spqr.cs.umass.edu/mementos>

- What should be moved off-chip to save energy?
- Right combo of RAM & NVRAM? Tiny off-chip NVRAM?
- HW/SW interface for detecting an impending failure?
- Conditional branches predicated on available energy?
- Task scheduling when failure is common case?
- Should we write a SuperTinyOS that expects failure?
- Compile-time optimizations for expected failure?
- How to not repeat non-idempotent actions?

HW

OS

PL

# Contingency Slides

# $V_{thresh}$ Subtleties

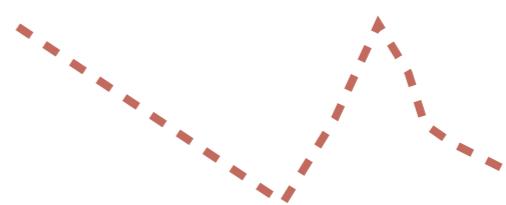
---

- Size, duration of CKPT are application dependent
- $V_{thresh}$  is a conservative estimate
  - ▶ More energy might arrive
  - ▶ Choose according to risk tolerance

# $V_{thresh}$ Subtleties

---

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  - ▶ More energy might arrive
  - ▶ Choose according to risk tolerance



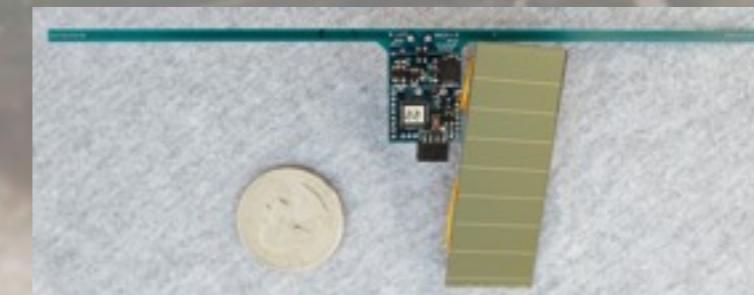
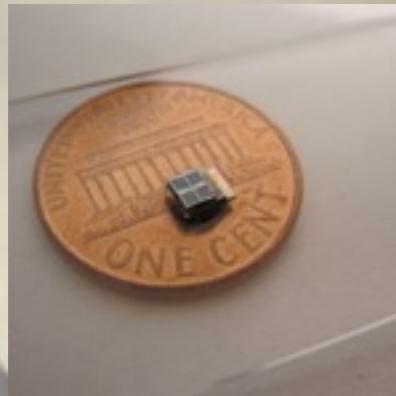
Spoiler: We help the programmer choose  $V_{thresh}$

# Low-Power Modes

---

- MSP430 has a variety of low-power modes (~1  $\mu$ A in LPM3/LPM4) that retain RAM
- Sleeping when energy is low is an optimistic strategy
- We don't know when or whether energy will return — should Mementos guess?

# Applications



Daeyeon Kim; Gyouho Kim; (SolarWISP) Shane Clark; Medtronic

# Thanks for asking about TinyOS

---

- Condor, Porch, libckpt — all depend on OS-level or out-of-band facilities
- Sensor OSes (e.g.,  TinyOS) designed to boot *infrequently* and sip from batteries
  - ▶ TinyOS boot:  $\geq 100$  ms (too slow)

# Thanks for asking about TinyOS

---

- Condor, Porch, libckpt — all depend on OS-level or out-of-band facilities
- Sensor OSes (e.g.,  TinyOS) designed to boot *infrequently* and sip from batteries
  - ▶ TinyOS boot:  $\geq 100$  ms (too slow)

Existing lightweight OSes *still* too heavy

# CRC Example: Overhead

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**How much CPU overhead for checks?**  
Consistently high voltage ( $V > V_{\text{thresh}}$ ):

Instrumentation	CPU Cycles	Mementos
Uninstrumented	575,315	0%
Loop latches	619,450	6.9%
Function returns	577,702	0.2%
Timer + latches	598,171	3.4%

# Why Not Just Add...

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- Thin-film battery?
- Deeper charge pump (higher voltage)?
- Tiny dedicated NVRAM?
- Hardware “low energy” interrupt support?

# Other Use Cases

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Programmer can:

- Disable instrumentation at a function level
- Manually call Mementos routines