

# Concept Challenge: DFS and BFS



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by Christine Alvarado, Mia Minnes, and Leo Porter, 2015.

# Concept Challenge: Procedure

- **Pause** Try to solve the problem yourself
- **Discuss** with other learners (if you can)
- **Watch** the UC San Diego learners video
- **Answer** the question again
- **Confirm** your understanding with our explanation



**DFS(S, G):**

Initialization of structures

Push S on stack and add to visited

while stack is not empty:

    pop node curr from top of stack

    if curr == G return parent map

    for each of curr's unvisited neighbors, n:

        add n to visited set

        add curr as n's parent in parent map

        push n to top of stack

// If we get here then there's no path

**BFS(S, G):**

Initialization of structures

Enqueue S in queue and add to visited

while queue is not empty:

    dequeue node curr from front of queue

    if curr == G return parent map

    for each of curr's unvisited neighbors, n:

        add n to visited set

        add curr as n's parent in parent map

        enqueue n to back of queue

// If we get here then there's no path

**Which algorithm is asymptotically faster in the worst case?**

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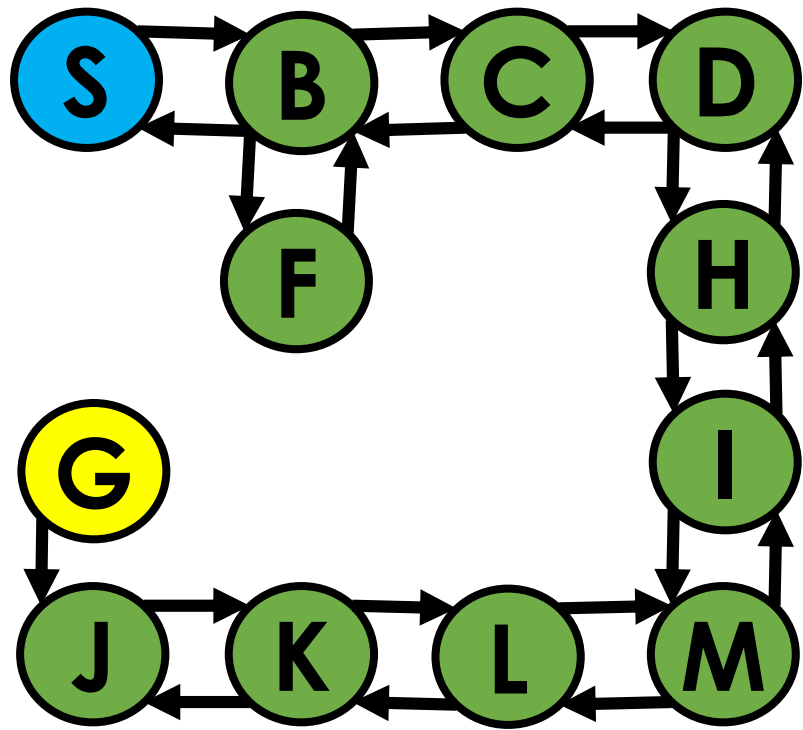
**Which algorithm is asymptotically faster in the worst case?**

A. BFS is faster

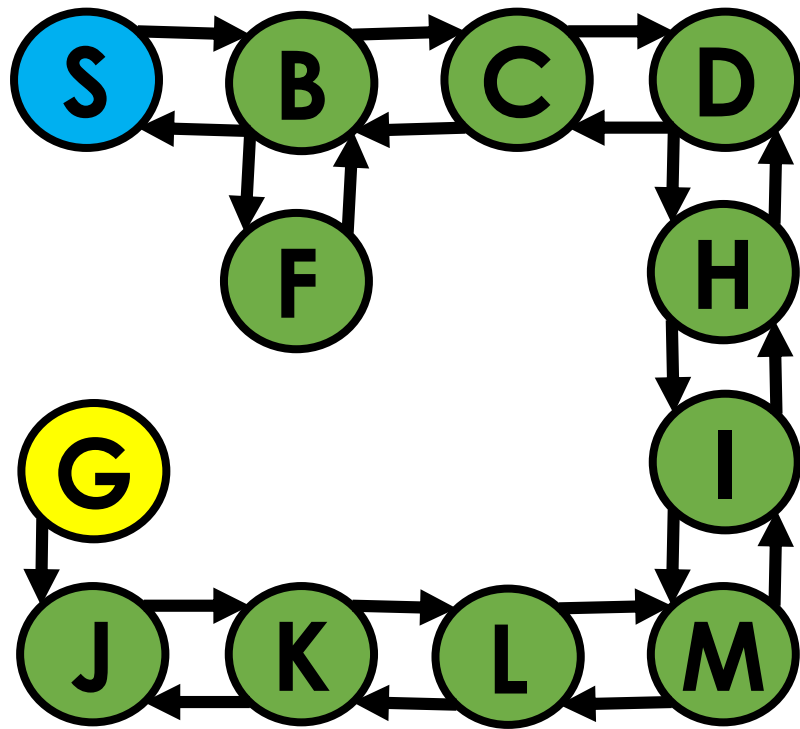
B. DFS is faster

C. They are both the same

Give tight big O bounds for each algorithm.



What does the worst case look like?



**How many nodes will each algorithm visit?**

Search(S, G):

Initialization of structures

Add S to [stack/queue] and add to visited

while [stack/queue] is not empty:

remove node curr

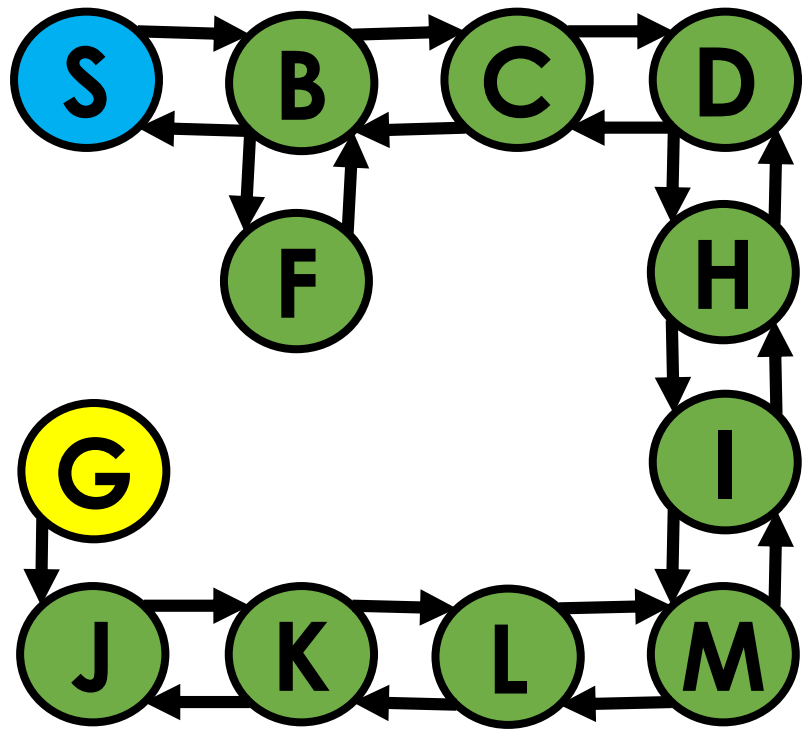
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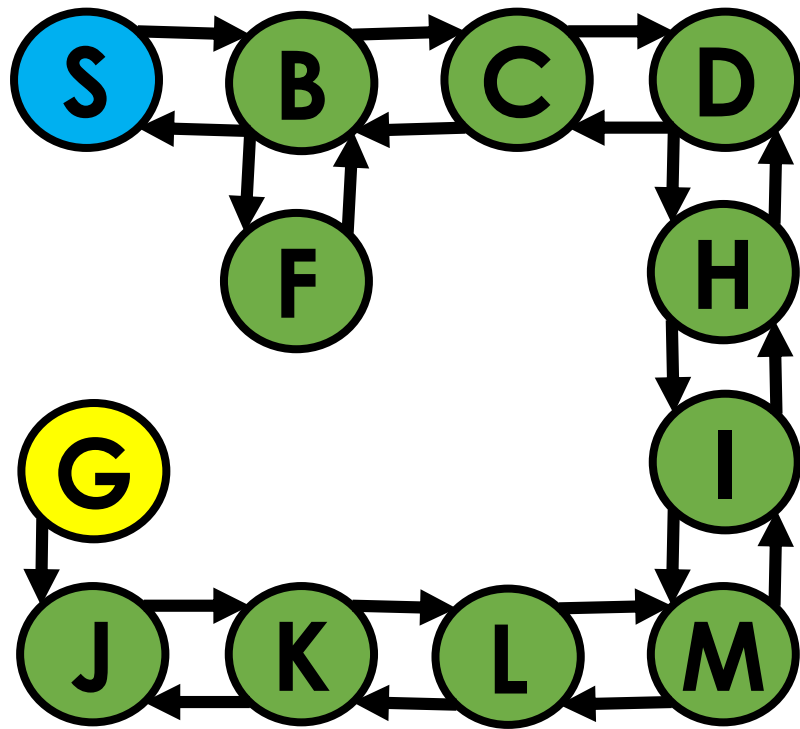
add n to [stack/queue]



How many nodes will each algorithm visit?

$$|V| - 1 = O(|V|)$$

So is that it?



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Initialization of structures

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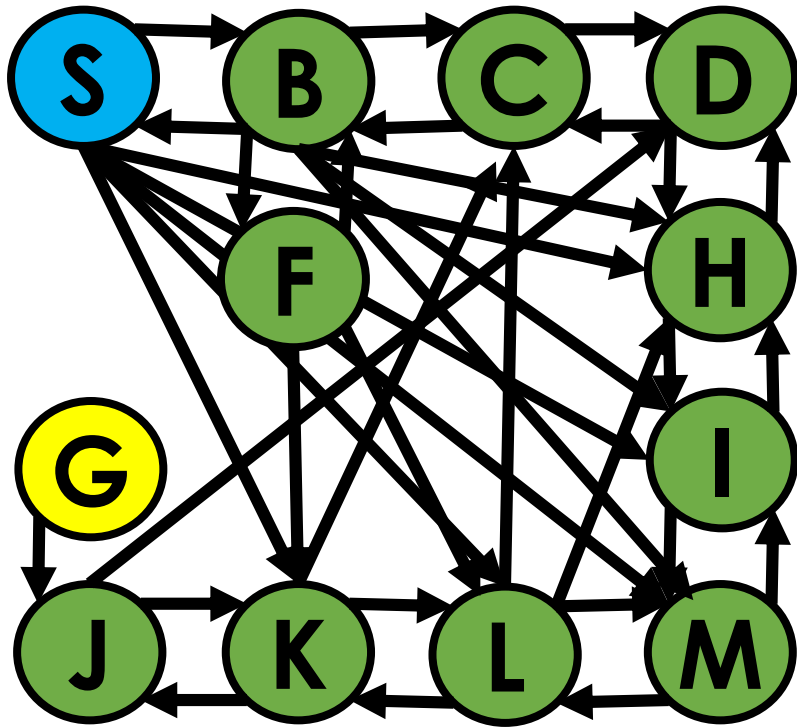
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add curr as n's parent in parent map

add n to [stack/queue]





Search(S, G):

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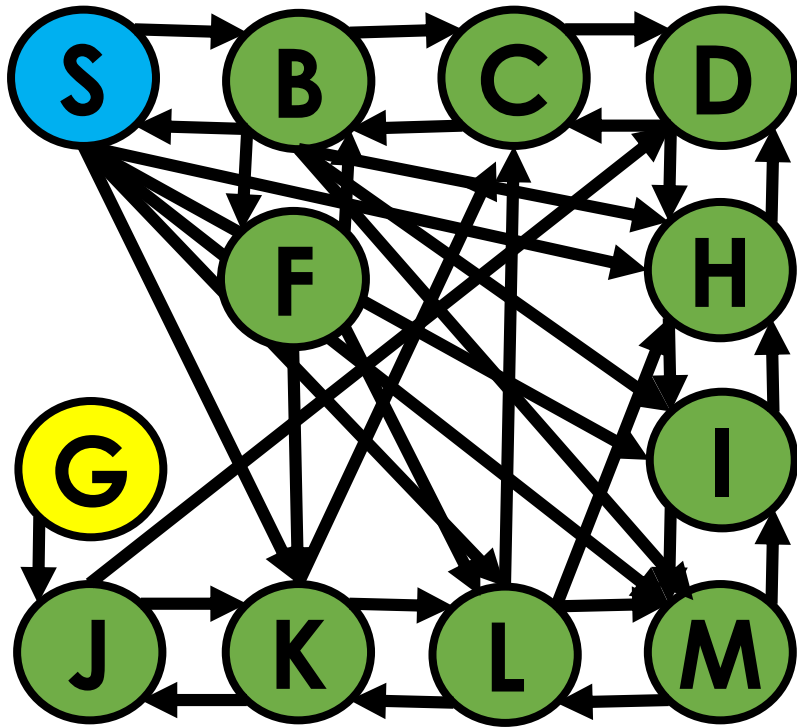
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add n to [stack/queue]

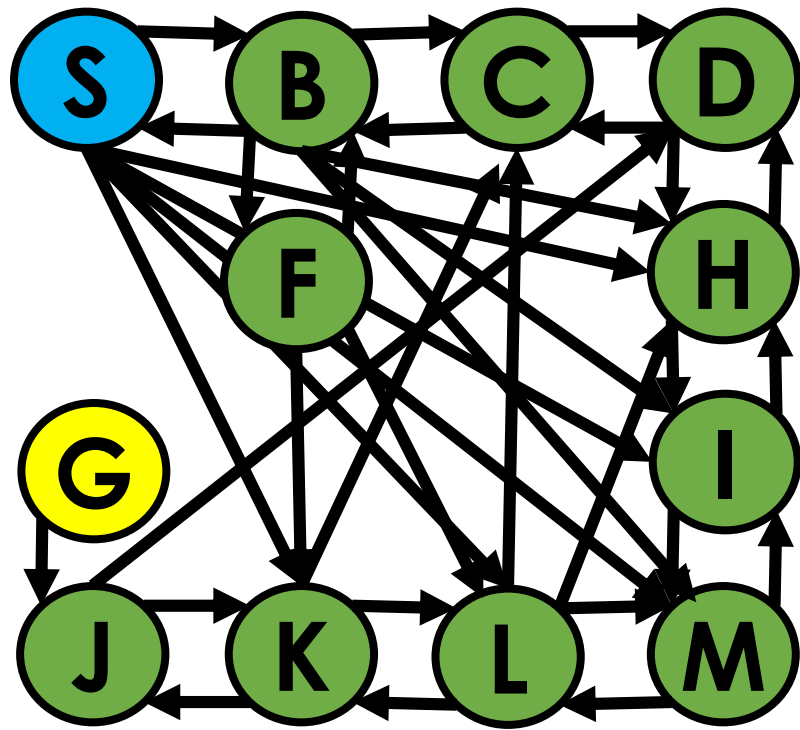


How many nodes will each algorithm visit?  $O(|V|)$

But DFS and BFS will traverse each edge too!

How many edges can there be in a graph?

$|E|$  can be up to  $|V|^2$



How many nodes will each algorithm visit?  $O(|V|)$

How many edges will each algorithm traverse?  
 $O(|E|)$

Worst case running time for both  $\rightarrow O(|V| + |E|)$