

Dijkstra's Algorithm: Runtime Analysis



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by Christine Alvarado, Mia Minnes, and Leo Porter, 2015.

Concept Challenge: Procedure

- **Pause** Try to solve the problem yourself
- **Discuss** with other learners (if you can)
- **Watch** the UC San Diego learners video
- **Answer** the question again
- **Confirm** your understanding with our explanation



Dijkstra: Algorithm

Dijkstra(S, G):

Initialize: Priority queue (PQ), visited HashSet,
parent HashMap, and distances to infinity

Enqueue $\{S, 0\}$ onto the PQ

while PQ is not empty:

 dequeue node curr from front of queue

 if(curr is not visited)

 add curr to visited set

 If curr == G return parent map

 for each of curr's neighbors, n, not in visited set:

 if path through curr to n is shorter

 update curr as n's parent in parent map

 enqueue $\{n, \text{distance}\}$ into the PQ

// If we get here then there's no path

A. What is the running time of this algorithm in terms of $|V|$ and $|E|$? (If more than one might be correct—which is tighter?)

A. $O(|V|^2)$

B. $O(|E| + |V|)$

C. $O(|E| * \log(|E|))$

D. $O(|E| * \log(|E|) + |V|)$

E. $O(|E| * |V|)$

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Walkthrough

Dijkstra(S, G):

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$O(|V|)$

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$O(V)$
$O(E)$

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$O(V)$
$O(E)$

~~$O(|V|+1)$~~

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Initialize: Priority queue (PQ), visited HashSet,
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$O(|V|)$

Enqueue $\{S, 0\}$ onto the PQ

$O(?)$

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$O(|V|)$

$O(|E|)$

**Most lines in the
loop are $O(1)$**

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$O(|V|)$

$O(|E|)$

$O(\log |E|)$

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$O(|V|)$

$O(|E|)$

$O(\log |E|)$

$O(\log |E|)$

**Insert and remove from
a Priority Queue with N
elements is $\log(N)$**

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$O(|V|)$

$O(|E|)$

$O(\log |E|)$

**But what about this
for loop?**

$O(\log |E|)$

**Insert and remove from
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$O(|V|)$

$O(|E|)$

$O(\log |E|)$

**It's already
accounted for!**

$O(\log |E|)$

**Insert and remove from
a Priority Queue with N
elements is $\log(N)$**

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$$O(|V|)$$

$$O(|E| \log |E|)$$

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$$O(|V|)$$

$$O(|E| \log |E|)$$

Overall:

$$O(|E| \log |E| + |V|)$$

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$$O(|V|)$$

$$O(|E| \log |E|)$$

Overall:

$$O(|E| \log |E| + |V|)$$

Because $E \leq |V|^2$ and
 $\log(|V|^2)$ is just $O(\log(|V|))$
we could tighten to:
 $O(|E| \log |V| + |V|)$