

# Outline

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  - Significance
- 2 Section: Computational Model
  - Quantum Information
  - Quantum Transformations
  - Measurement
- 3 Section: Quantum Algorithms
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# Founding Fathers

Time: early 1980-s

People: Paul Benioff, Richard Feynman, and Yuri Manin



Key Tool: **Quantum Parallelism** (tricky to employ)

# The Pearl: Shor's Algorithm (1994)

Problem: Integer Factorization

Best "Classical" Solution:  $O\left(e^{1.9(\log N)^{1/3}(\log \log N)^{2/3}}\right)$

Shor's Algorithm:  $O((\log N)^2(\log \log N)(\log \log \log N))$

# State Space Postulate

## Postulate 1 (State Space Postulate)

The state of a system is described by a unit vector in a Hilbert space  $\mathcal{H}$ .

### "Systems"

A piece of physical reality used to encode information akin to transistors, e.g.

- electron and its spin,
- photon and its polarization,
- spins of other particles

### Hilbert space

Complex **vector space** with a **scalar product**  
(to measure angles and lengths).

Once you pick a basis, its basically  $\mathbb{C}^n$ .

Example: 2 dimensions, fixed basis:  $|0\rangle, |1\rangle \Rightarrow$  **qubit**

$$\alpha_0 |0\rangle + \alpha_1 |1\rangle, \quad \alpha_i \in \mathbb{C} : |\alpha_0|^2 + |\alpha_1|^2 = 1.$$

# Composite Systems

## Composition of Systems Postulate

If one system is in the state  $|\varphi_1\rangle$  and the second system in the state  $|\varphi_2\rangle$ , then the state of the combined system is described by the *tensor product*:

$$|\varphi_1\rangle \otimes |\varphi_2\rangle \in \mathcal{H}_1 \otimes \mathcal{H}_2, \quad \text{if } |\varphi_i\rangle \in \mathcal{H}_i.$$

**Notation** Instead of  $|\varphi_1\rangle \otimes |\varphi_2\rangle$  we write  $|\varphi_1\rangle |\varphi_2\rangle$  or even  $|\varphi_1\varphi_2\rangle$ .

Tensor Product in a Nutshell: bilinear pairing operation.

## Example: two qubits

Given

$$|\varphi_1\rangle = 1/\sqrt{2}(|0\rangle + |1\rangle), \quad |\varphi_2\rangle = 1/\sqrt{2}(|0\rangle - |1\rangle).$$

their composite is described with:

What about dimension?

$$|\varphi_1\rangle |\varphi_2\rangle = 1/2(|0\rangle + |1\rangle)(|0\rangle - |1\rangle) = 1/2(|00\rangle - |01\rangle + |10\rangle + |11\rangle).$$

# Entangled States

Can we always un-tensor states of two qubits?

No (...t at all):

$$|\varphi\rangle = 1/\sqrt{2}(|00\rangle + |11\rangle)$$

Fancy name: **EPR pair** (for Einstein, Podolsky, and Rosen)

# The Way To Compute

What to require from a computational tool when work with *unit vectors*?

## Evolution Postulate

The time-evolution of the state of a closed quantum system is described by a **unitary operator**:  $\exists U: |\varphi_{t0+x}\rangle = U|\varphi_{t0}\rangle$ .

## Unitary Operators

$U$  unitary  $\Leftrightarrow (Ux, Uy) = (x, y)$ .

Using adjoint operator:

$$(Ux, Uy) = (x, U^* Uy),$$

an equivalent formulation:

$$U \text{ unitary} \Leftrightarrow U^* = U^{-1}.$$

## Case of Real-Valued Matrices

First,  $U^*$  is simply  $U^T$ .

$$U \text{ unitary} \Leftrightarrow U^{-1} = U^T.$$

Second, if an operator is self-inverse ( $U^{-1} = U$ ), then

$$U \text{ unitary} \Leftrightarrow U = U^T.$$

# The Cats Of It

## Problem: Can't Access The Amplitudes

A “classical” observer can only get to basis states, e.g.  $|0\rangle$ ,  $|1\rangle$ .  
So what do superposition states mean?

cat:  $1/\sqrt{2}$  dead +  $1/\sqrt{2}$  alive...

## Measurement Postulate

Consider a system  $A$  and its state space  $\mathcal{H}_A$ .

For an orthonormal basis  $B = \{|\varphi_i\rangle\}$  in  $\mathcal{H}_A$ , there exists a description of  $A$ :

$$|\varphi\rangle = \sum_i \alpha_i |\varphi_i\rangle, \quad \sum_i |\alpha_i|^2 = 1.$$

It is possible to perform a (*Von Neumann*) *measurement* on system  $A$  with respect to the basis  $B$ . With probability  $|\alpha_i|^2$ , this act outputs a label  $i$  and leaves the system in state  $|\varphi_i\rangle$ .



# Measurement Examples

$$|\varphi\rangle = \sqrt{1/11} |00\rangle + \sqrt{5/11} |01\rangle + \sqrt{2/11} |10\rangle + \sqrt{3/11} |11\rangle$$

## Measure both qubits

What are the possible outcomes of measuring the pair of qubits?

Here they are:

- 00 — with probability  $1/11$ ; the system turns into  $|00\rangle$
- 01 — with probability  $5/11$ ; the system turns into  $|01\rangle$  ...

## Measure the first qubit

*Note from above:*

the probability of getting 0 for the first qubit is  $1/11 + 5/11 = 6/11$ .

So, we factor out  $\sqrt{6/11}$  first to get:

$$|\varphi\rangle = \sqrt{6/11} |0\rangle \left( \sqrt{1/6} |0\rangle + \sqrt{5/6} |1\rangle \right) + \sqrt{5/11} |1\rangle \left( \sqrt{2/5} |0\rangle + \sqrt{3/5} |1\rangle \right).$$

When output=0 the system becomes:  $|0\rangle (\sqrt{1/6} |0\rangle + \sqrt{5/6} |1\rangle)$ .

# The Road to Shor

## Idea 1 (a simple one)

We only need to be able to find one factor of the input  $N$ .

## Idea 2 (a technical one)

Take random  $A < N$  coprime to  $N$ . An integer  $r$  is called the **order** of  $A$  modulo  $N$  if  $A^r \equiv 1 \pmod{N}$ .

Finding a factor of  $N$  can be efficiently reduced to the order-finding problem.

## The Insight About Periodicity

The order-finding problem can be seen as the problem of finding the period of the following mapping:

$$\exp_{A,N}: \quad b \mapsto A^b \pmod{N}, \quad \text{where } b \in \mathbb{Z}.$$

# Periodic Functions: What and Why

## Definition

$f: G \rightarrow X$  defined on a group  $G$  is *periodic with period*  $r \in G$  ( $r \neq e$ ) if:

$$\forall n \in \mathbb{Z} \forall g \in G: f(g + rn) = f(g).$$

## Shor's Problem

Find the period of a particular function ( $\exp_{A,N}$ ) defined on the group  $\mathbb{Z}$ .

*Note:* In fact, on a smaller group,  $\mathbb{Z}_{\varphi(N)}$ , but we don't know it in advance, hence the other name: the *hidden subgroup problem*.

We are in rush, so we tackle an easier instance of the same problem.

## Simon's Problem

Find the period of a function defined on the group  $\mathbb{Z}_2^n$ .

# Simon's Problem

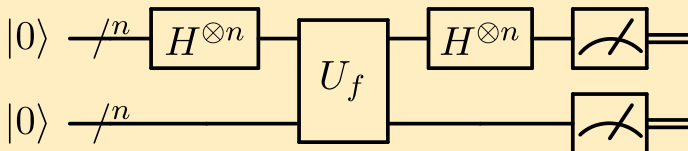
## Formulation

**Input:** A black-box for computing an unknown function  $f: \mathbb{Z}_2^n \rightarrow X$ .

**Promise:** There exists  $\bar{s}$ , s.t.:  $f(\bar{x}) = f(\bar{y})$  iff  $\bar{x} = \bar{y}$  or  $\bar{x} = \bar{y} \oplus \bar{s}$ .

**Problem:** Determine  $\bar{s}$  by making queries to  $f$ .

## The Circuit



$$|0^{2n}\rangle \xrightarrow{(1)} \sum_{\bar{x} \in \mathbb{Z}_2^n} |\bar{x}\rangle |0^n\rangle \xrightarrow{(2)} \sum_{\bar{x} \in \mathbb{Z}_2^n} |\bar{x}\rangle |f(\bar{x})\rangle \xrightarrow{(3)} (|\bar{x}\rangle + |\bar{x} \oplus \bar{s}\rangle) |f(\bar{x})\rangle \xrightarrow{(4)} \dots$$

# Why The Circuit Works (And What it Outputs)

Hadamard on a coset

$$H^{\oplus n}(|\bar{x}\rangle + |\bar{x} \oplus \bar{s}\rangle) = \sum_{\bar{z} \in \bar{s}^{\perp}} (-1)^{\bar{x} \cdot \bar{z}} |\bar{z}\rangle .$$

Output of The Circuit

Claim: the output on the first  $n$  wires of the circuit is a vector  $\bar{z} \in \bar{s}^{\perp}$ .

# Quantum Computing

- Massively data-parallel model with probabilistic outcomes;
- good only for certain classes of tasks,  
esp. when searching for global properties of functions;
- some of the applications are very important (e.g. in cryptography);
- practise lags behind theory.