# HW5

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### 1 Problem 1

First prove the base case. When n=1, 1 = 1<sup>2</sup>, the statement holds. Then assume the statement holds for n=k, namely  $\sum_{1}^{k}(2n-1)=k^2, k\geq 1$ . When n=k+1,  $\sum_{1}^{k+1}(2n-1)=\sum_{1}^{k}(2n-1)+2k+1=k^2+2k+1=(k+1)^2$  By induction, the statement is TRUE Q.E.D.

### 2 Problem 2

First consider the base cases. When n=1, n= $F_1$ ; when n=2, n= $F_3$ , when n=3, n= $F_4$ 

Assume the statement holds  $\forall n \leq k, k \geq 3$ 

When n=k+1, if there is a fibonacci number  $F_p$  that is equal to n, n =  $F_p$ . Else, there must be a  $0 < F_t < n < F_{t+1}$ . Let m = n- $F_t$ . Since the statement holds  $\forall n \leq k$  and m < n, m could be expressed as the sum of distinctive non-consecutive Fibonacci numbers  $\implies n = F_t + m$ 

Meanwhile,  $F_{t-1}$  is not in m since  $m = n - F_t < F_{t+1} - F_t = F_{t-1} \implies$  n could be represented as the sum of distinctive non-consecutive Fibonacci numbers.

By induction, the statement is true Q.E.D.

### 3 Problem 3

We prove it using induction. We first prove the base case. When n=1, i=j=k=1, we do not need to cut the cake at all, which means the minimum cut is 1-1=0. The statement holds.

Then we assume that for all  $n \leq k, k \geq 1$ , the statement holds. When n=k+1, we let the longest edge be i, and divide the rectangle into 2 smaller rectangles, Denote them as  $R_1, R_2$ . Since the rectangle has integer sides, the volumes  $r_1$  and  $r_2$  of  $R_1$  and  $R_2$  are smaller than n  $\implies$  it took  $r_1 - 1 + r_2 - 1$  in total to cut the two slices.

Since  $r_1 + r_2 = n$ , and it took 1 cut to divide the rectangle into 2 pieces, so in total is n - 1 cuts.

By induction, the statement holds.

Q.E.D.