

Introduction

Ch 1 Introduction

Womp... Womp.

Ch 2 Analog vs Digital communication

• Fourier Transform: $X(f) = \int_{-\infty}^{\infty} x(t) \exp(-j2\pi f t) dt$ forward

$x(t) = \int_{-\infty}^{\infty} X(f) \exp(j2\pi f t) df$ inverse.

linearity: $F\{ax_1(t) + bx_2(t)\} = aX_1(f) + bX_2(f)$

Convolution: $F\{x_1(t) * x_2(t)\} = X_1(f) \cdot X_2(f)$.

Real signals: $X(f) = X^*(-f)$ if $x(t) \in \mathbb{R}$

Linear-Time Invariant (LTI) Systems

impulse response $h(t)$: $y(t) = x(t) * h(t)$; freq domain: $H(f) = X(f) \cdot H(f)$.

• Parseval's Theorem: Energy: $\int_{-\infty}^{\infty} |g(t)|^2 dt = \int_{-\infty}^{\infty} |G(f)|^2 df$.

• $M \rightarrow \text{set.}$; Bayes Rule: $\Pr\{Y=y | X=x\} = \frac{\Pr\{X=x | Y=y\} \Pr\{Y=y\}}{\Pr\{X=x\}}$

• PMF: $\Pr\{X=x\}$; Law Total $\Pr\{Y=y\} = \sum_x \Pr\{X=x, Y=y\}$

• PDF: $\int_{-\infty}^{\infty} p_X(r) dr = 1$

White Gaussian Noise (WGN)

- power spectral density (PSD): $S_{Nw}(f) = N_0/2$

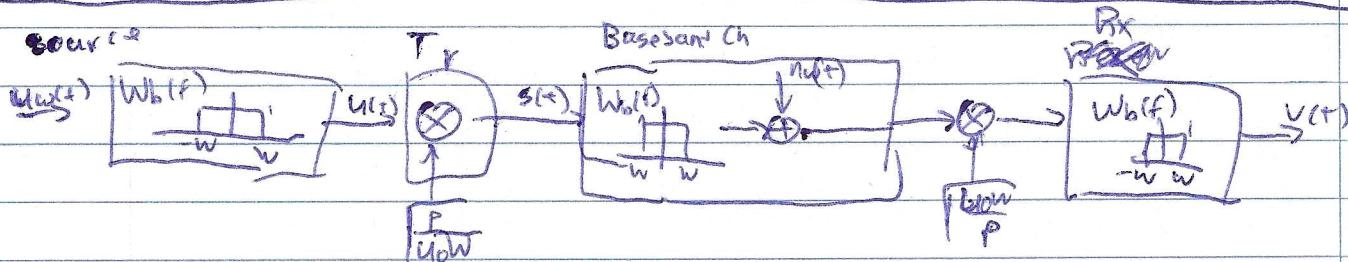
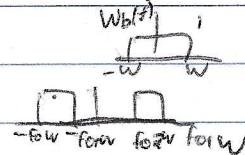
- zero mean: $E[Nw(t)] = 0$

- autocorrelation: $R_{Nw}(t, s) = N_0/2 \delta(t-s)$

• Wideband AWGN: $v(t) = s(t) + n_w(t)$

• Baseband AWGN: $v(t) = s(t) + n_b(t)$

• Passband AWGN: $v(t) = s(t) + n_p(t)$



$$\text{Power distribution } d(t) \triangleq v(t) - u(t) = [U_{ow}/P \cdot u(t) + N_{ow}(t)] \otimes W_b(t)$$

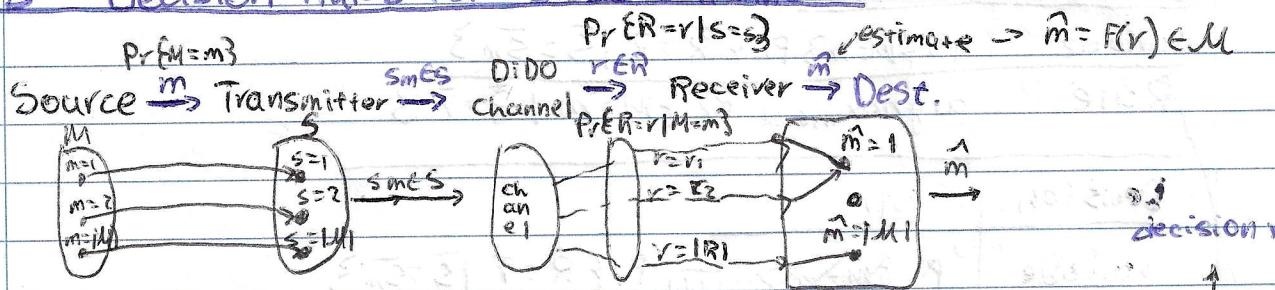
$$\text{Variance distortion } E[d^2(t)] = U_{ow}^2 \frac{W_b}{2} \cdot \sigma_w^2 = U_{ow} \frac{N_{ow}}{P}$$

$$\text{Mean Square Distortion } D = E[d^2(t)]$$

$$SDR \triangleq \frac{E[u^2(t)]}{E[d^2(t)]} = P/N_{ow} \quad \text{for DSB-SC}$$

Channel Processing Receivers.

Ch 3 - Decision Rules for DIDD Channels.



$$P_e \triangleq \Pr\{\hat{M} \neq M\} \rightarrow \text{optimal receiver minimizes } \rightarrow P_c \triangleq \Pr\{\hat{M} = M\} = 1 - P_e \rightarrow f(r)$$

$$P_c = \Pr\{M = \hat{M}\} \rightarrow \hat{M} = f(R) \rightarrow P_c = \sum_{m \in M} \sum_{r \in R, f(r)=m} \Pr\{M=m, R=r\}$$

$$= \dots \Pr\{f(R=r | M=m)\} \Pr\{M=m\} = \Pr\{M=m, R=r\}$$

ML

Example Tx signals: $S \in S = \{S_1, S_2\}$ ($|M|=2$); Rx signals: $R \in R = \{a, b, c\}$

$$\Pr\{M=m\} \cdot \text{A-priori: } \Pr\{M=1\} = 0.4, \quad \Pr\{M=2\} = 0.6$$

$$\Pr\{S_1, S_2\} \cdot \text{Conditional: } \Pr\{R=a | S_1\} = 0.5, \quad \Pr\{R=b | S_1\} = 0.4, \quad \Pr\{R=c | S_1\} = 0.1$$

$$\Pr\{R=a | S=S_2\} = 0.1, \quad \Pr\{R=b | S=S_2\} = 0.3, \quad \Pr\{R=c | S=S_2\} = 0.6$$

decision probability • **Decision:** $r \in \{a, b\} \rightarrow m=1$ ($f(r)=1$); $r \in \{c\} \rightarrow m=2$ ($f(r)=2$).

decision variables • **Joint PMF (PMF):** $\Pr\{M=m, R=r\} = \Pr\{M=m\} \cdot \Pr\{R=r | S=S_m\}$

m	$\Pr\{M=m, R=a\}$	$\Pr\{M=m, R=b\}$	$\Pr\{M=m, R=c\}$
1	0.2	0.16	0.04
2	0.06	0.18	0.36

$$P_c = 0.2 + 0.16 + 0.36 = 0.72; \quad P_e = 0.28$$

• Max Likelihood detection: $P_c = 0.2 + 0.18 + 0.36 = 0.74$

$$0.74 > 0.72$$

MAP

$$\hat{m}^{\text{MAP}}(r) \triangleq \arg \max_{m \in M} \Pr\{M=m | R=r\} \rightarrow \max \Pr$$

PMF
probability Mass Func

$$\text{A-posteriori Probability: } \Pr\{M=m | R=r\} = \Pr\{M=m\} \cdot \Pr\{R=r | S=S_m\}$$

Bayes' Rule \rightarrow Pr of M transmitted based on received R.

$$\Pr\{R=a\} = 0.2 + 0.06 = 0.26 \rightarrow \Pr\{M=1, R=a\} + \Pr\{M=2, R=a\}$$

$$\Pr\{R=b\} = 0.34; \quad \Pr\{R=c\} = 0.4$$

$$\Pr\{M=1, R=a\} / \Pr\{R=a\} = 0.2 / 0.26 = 0.77 = 77/26$$

m	$\Pr\{M=m R=a\}$	$\Pr\{M=m R=b\}$	$\Pr\{M=m R=c\}$
1	20/26	16/34	4/10
2	6/26	18/34	36/40

$$(\text{MAP}) \rightarrow P_c = 20/26 + 18/34 + 36/40$$

A-Priori: $\Pr\{M=m\}$; A-Posteriori: $\Pr\{M=m | R=r\}$

• uniform a-priori: $\Pr\{M=m\} = 1/|M| \rightarrow \Pr\{M=m, R=r\} = \frac{1}{|M|} \Pr\{R=r | S=S_m\}$

$$\hat{m}^{\text{ML}}(r) \triangleq \arg \max_{m \in M} \Pr\{R=r | M=m\} \rightarrow \max \Pr$$

conditional.

ML

$$\text{If } \Pr\{M=m\} = 1/2 \rightarrow \text{equal } P_c^{\text{ML}} = 1/2 (0.5 + 0.4) + 1/2 (0.6) = 0.75$$

MAP vs ML Summary

Decision | MAP

Variable $\Pr\{M=m \mid \bar{R}=\bar{r} \mid \bar{S}=\bar{s}_m\}$

Rule $\arg\max_{m \in \mathcal{M}} \Pr\{\bar{E} \mid M=m \mid \bar{R}=\bar{r}\}$

Decision | ML

Variable $\Pr\{\bar{E} \mid M=m \mid \bar{R}=\bar{r} \mid \bar{S}=\bar{s}_m\}$

Rule $\arg\max_{m \in \mathcal{M}} \Pr\{\bar{E} \mid \bar{R}=\bar{r} \mid M=m\}$

ML

1. A-Priori:	m	$\Pr\{M=m\}$
	1	0.4
	2	0.6

conditional Pr

2.	m	$\Pr\{\bar{E} \mid \bar{R}=a \mid \bar{S}=\bar{s}_m\}$	$\Pr\{\bar{E} \mid \bar{R}=b \mid \bar{S}=\bar{s}_m\}$	$\Pr\{\bar{E} \mid \bar{R}=c \mid \bar{S}=\bar{s}_m\}$
	1	0.5	0.4	0.1
	2	0.1	0.3	0.6

$$3. f(r) = 1 \quad r \in \{a, b\} \quad f(r)=2 \quad r \in \{c\}$$

4. Joint Pr	m	$\Pr\{\bar{E} \mid M=m, R=a\}$	$\Pr\{\bar{E} \mid M=m, R=b\}$	$\Pr\{\bar{E} \mid M=m, R=c\}$
	1	0.2	0.16	0.04
	2	0.6	0.18	0.36

$$5. P_C = 0.72 //$$

MAP → have to know output → harder & to implement even though better.

$$1. \Pr\{\bar{E} \mid R=a\} = \Pr\{\bar{E} \mid M=1, R=a\} + \Pr\{\bar{E} \mid M=2, R=a\} = 0.26.$$

$$\Pr\{\bar{E} \mid R=b\} = 0.34 ; \quad \Pr\{\bar{E} \mid R=c\} = 0.4.$$

$$2. \text{A-Posteriori Pr:}$$

$$\Pr\{\bar{E} \mid M=m, \bar{R}=\bar{r}\}$$

$$\text{Bayes eq: } \Pr\{\bar{E} \mid M=m \mid \bar{R}=\bar{r}\} = \frac{\Pr\{\bar{E} \mid M=m\}}{\Pr\{\bar{E} \mid \bar{R}=\bar{r}\}}$$

$$\Pr\{\bar{E} \mid M=1 \mid \bar{R}=\bar{r}\} = 0.2 / 0.26 =$$

m	$\Pr\{\bar{E} \mid M=m \mid R=a\}$	$\Pr\{\bar{E} \mid M=m \mid R=b\}$	$\Pr\{\bar{E} \mid M=m \mid R=c\}$
1	0.16	0.18	0.04
2	0.6	0.36	0.4

$$3. P_C = 0.2 + 0.18 + 0.36 = 0.74 //$$

Equal A-Priori → ML

$$1. \Pr\{\bar{E} \mid M=m\} = 0.5 \quad 2. f(r)=1 \quad r \in \{a, b\} \quad f(r)=2 \quad r \in \{c\}$$

$$3. P_C = \frac{1}{2}(0.5+0.4) + \frac{1}{2}(0.6) = 0.75$$

$$P_{\text{r}}\{M=m | R=r\} = \frac{P_{\text{r}}\{M=m\} P_R(r | M=m)}{P_R(r)} \rightarrow 1$$

Ch 9 Decision Rules for DICO Channels

- Decision Variables DICO: $P(M=m, R=r) = P(M=m) \cdot P_R(r | S=s_m)$
 $= P(M=m) \cdot P_R(r | M=m)$

- Max A-Posteriori (MAP) Decision Rule: $\hat{m}_{\text{MAP}}(r) = \arg \max_{m \in M} P\{M=m | R=r\}$

Bayes rule: $= \arg \max_{m \in M} P\{M=m\} \cdot P_R(r | M=m)$

- Max Likelihood (ML) Decision Rule: $\hat{m}_{\text{ML}} = \arg \max_{m \in M} P_R(r | M=m) \rightarrow \text{all equal a-priori}$

- Scalar Additive Noise Channel AGN: $r = s + n$

$$\text{Gauss distribution } p_N = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{r^2}{2\sigma^2}\right) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(r-s_m)^2}{2\sigma^2}\right)$$

- Conditional Probability Density AGN channel: $P_R(r | S=s_m) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(r-s_m)^2}{2\sigma^2}\right)$

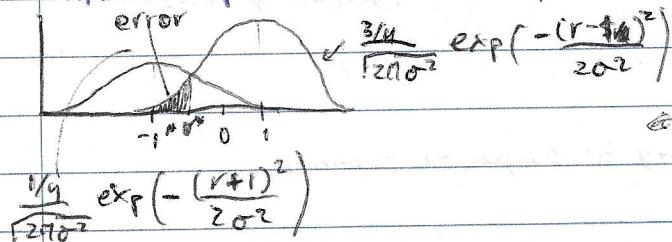
- MAP rule for Binary-input AGN (bi-AGN) channel: $M \in \{1, 2\}$

$$\hat{m} = 1 \text{ if: } P(M=1) \cdot P_R(r | S=s_1) \geq P(M=2) \cdot P_R(r | S=s_2)$$

- MAP Threshold: $r^* = \frac{\sigma^2}{(s_1 - s_2)} \ln\left(\frac{P(M=2)}{P(M=1)}\right) + \frac{s_1 + s_2}{2} \quad \hat{m} = \begin{cases} 1 & r \geq r^* \\ 2 & \text{else.} \end{cases}$

- ML Thr, equiprobable: $P(M=1) = P(M=2) = 0.5 ; r^* = \frac{s_1 + s_2}{2}$

- MAP Example: $P\{M=1\} = 3/4, s_1 = +1, P\{M=2\} = 1/4, s_2 = -1 \rightarrow r^* = \frac{s_1 + s_2}{2} = -0.5 \approx -0.546$



Similar for ML but both

$$\frac{1}{\sqrt{2\pi\sigma^2}} \text{ even if } P\{M=1\} = P\{M=2\} = 1/2$$

- If $r^* = 0 \rightarrow P_e = \int_{-\infty}^{r^*} \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(r-r^*)^2}{2\sigma^2}\right) dr.$

- Q-function: tail (Pr) of Gaussian Random: $Q(x) = \int_x^{\infty} \frac{1}{\sqrt{2\pi}} \exp\left(-\frac{a^2}{2}\right) da$

$$Q(x) = 1 - Q(-x) \quad \text{and} \quad Q(x) + Q(-x) = 1 \quad \Pr\{R < r^* | M=1\} \quad \Pr\{R > r^* | M=2\}$$

- Error (Pr) for bi-AGN channel: $P_e = P(M=1) Q\left(\frac{s_1 - r^*}{\sigma}\right) + P(M=2) Q\left(\frac{s_2 - r^*}{\sigma}\right)$

$$\text{if Equally likely: } P(M=1) = P(M=2) = 0.5 \rightarrow P_e = Q\left(\frac{s_1 - s_2}{2\sigma}\right)$$

- Decision Region: $I_m = \{r \in \mathbb{R}^N : f(r) = m\} \quad \mathbb{R}^N \rightarrow \text{Observation space}$

Example MAP vs ML bi-AGN

$$s_1 = 1, s_2 = -1, \sigma^2 = 1, P\{M=1\} = 3/4, P\{M=2\} = 1/4$$

$$P_e = \frac{3}{4} Q\left(1 + \frac{1}{2}\right) + \frac{1}{4} Q\left(-1 + \frac{1}{2}\right) = 0.1270 // \text{MAP}$$

$$- \text{ML: } Q\left(\frac{1}{2}\right) = 0.1587 //$$

- PDF of Noise Vector: $p_N(n) = \frac{1}{(2\pi\sigma^2)^{N/2}} \exp\left(-\frac{\|n\|^2}{2\sigma^2}\right)$

- PDF of Vector: $R = S_m + n \rightarrow P_R(r | S=s_m) = \frac{1}{(2\pi\sigma^2)^{N/2}} \exp\left(-\frac{\|r - s_m\|^2}{2\sigma^2}\right)$

decision rule based on distance from hyper plane

$$\text{if } \|r - s_m\| = A, s_1 = A \rightarrow \text{it's same} //$$

squared Euclidean distance only gives optimal if a-priori msgs are equiprobable.

$$d = \sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$$

Decision Rules for AGN Vector Channels:

- MAP: $\hat{m}_{MAP}(\bar{r}) = \arg \min_{m \in M} \{ \| \bar{r} - \bar{s}_m \|^2 - 2\sigma^2 \ln \Pr\{M=m\}\}$
- ML: $\hat{m}_{ML}(\bar{r}) = \arg \min_{m \in M} \{ \| \bar{r} - \bar{s}_m \|^2 \}$
- Error (P_e) in AGN vector channel: $P_e = Q\left(\frac{d_{min}}{\sigma}\right)$ distance from signal point to decision hyperplane

Upper Bound on Error Probability (ML)

$$P_e \leq \sum_{m \in M} \frac{1}{|M|} \sum_{m' \in M, m' \neq m} Q\left(\frac{d_{min}}{\sigma}\right) \quad d_{min} = \frac{\| \bar{s}_m - \bar{s}_{m'} \|^2 / ML}{\sigma^2}$$

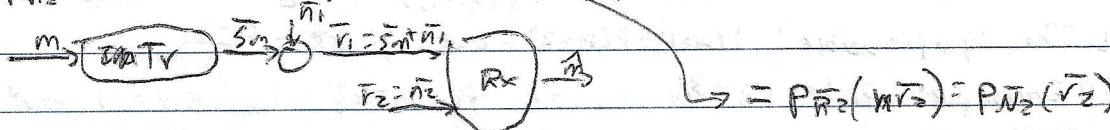
Can be used for MAP as well

σ^2 = Per dimension noise variance
distance to hyper plane

Theorem of Irrelevance:

Output \bar{r}_2 of multivector channel irrelevant (not effect P_e) if for all \bar{r}_1, \bar{r}_2

$$\Pr_{\bar{r}_2}(\bar{r}_2 | \bar{s}_m = \bar{s}_1, \bar{r}_1 = \bar{r}_1) \text{ not depend on } M$$

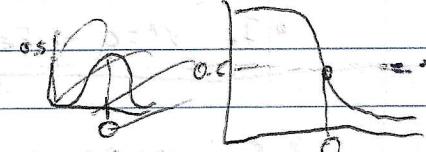


Theorem of Reversability:

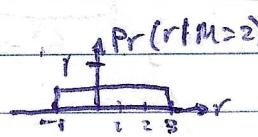
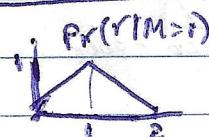
Min P_e not effected by reversability of output of channel



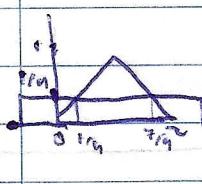
Random decision: $P_e = 1/2 \rightarrow 50\%$ $\rightarrow Q(0) =$



$$\text{MAP or ML} \Rightarrow \frac{\Pr(r_1 | s_c)}{\Pr(r_1 | s_i)} \geq 1$$



$$\begin{cases} \hat{m} \geq 2 & r_1 \geq 2 \\ M=1 & 1 \leq r_1 < 2 \\ \hat{m}=2 & r_1 > 2 \end{cases}$$



$$\Pr\{M=1\} = P_1 \Pr\{r_1 \leq 1 | r_1 \geq 1, M=1\} + \Pr\{M=2\} \Pr\{r_1 \leq 1 | r_1 \geq 1, M=2\}$$

$$\textcircled{1}: 2 \cdot \int_0^{1/4} x dx = 1/16 \quad \textcircled{2}: 1/4 \cdot \int_{1/4}^{3/4} dx = 3/8$$

$$1/16 \cdot \frac{1}{2} + 3/8 \cdot 1/2 = 7/32 = P_{eff}$$

(MAP)

SNR

Noise spectral density:
average noise power per unit frequency (W/Hz)
Power distribution of noise across diff frequencies

Channel Processing and Receivers

Ch5: Waveform Channels

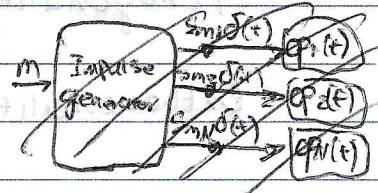
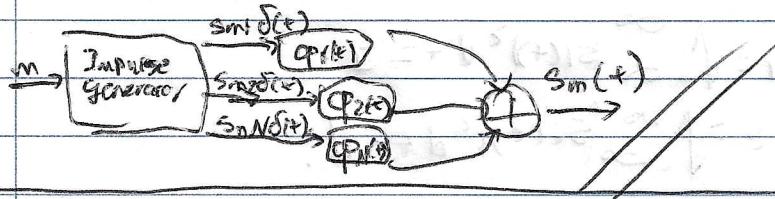
• Waveform energy: $E_x = \int_{-\infty}^{\infty} x^2(t) dt$

• Orthogonality: $\int_{-\infty}^{\infty} \varphi_i(t) \varphi_j(t) dt = \begin{cases} E_i, & i=j \\ 0, & i \neq j \end{cases}$

• Orthonormality: $E_i = 1$; waveforms \rightarrow orthonormal.

• Signal Synthesis: $s_m(t) = \sum_{i=1}^N s_{mi} \varphi_i(t)$, $m \in M = \{1, 2, \dots, M\}$.
where $\varphi_i(t)$ orthonormal basis function.

↳ building block waveforms



Gram-Schmidt \rightarrow arbitrary signals to orthonormal basis.

$$1) \varphi_1(t) = s_1(t) / \sqrt{E_1}, \quad S_{11} = \sqrt{E_1}$$

$$2) \varphi_2(t) = s_2(t) - S_{21} \varphi_1(t), \quad S_{21} = \int_{-\infty}^{\infty} s_2(t) \varphi_1(t) dt.$$

$$3) \varphi_3(t) = \varphi_3(t) / \sqrt{E_{22}}, \quad S_{22} = \int_{-\infty}^{\infty} s_3(t) \varphi_2(t) dt.$$

$$4) \varphi_m(t) = s_m(t) - \sum_{i=1}^{m-1} s_{mi} \varphi_i(t)$$

$$S_{mi} = \int_{-\infty}^{\infty} s_m(t) \varphi_i(t) dt; \quad i = 1, \dots, m-1.$$

$$P_o = d \left(\frac{d}{t_0} \right)^2$$

$$5a) \text{ IF } \varphi_m(t) = 0 \text{ stop or } \rightarrow \sum_{i=1}^{m-1} s_{mi} \varphi_i(t).$$

$$5b) \text{ if } \varphi_m(t) \neq 0 \rightarrow \varphi_m(t) = \varphi_m(t) / \sqrt{E_{mm}} \rightarrow \text{step 4.}$$

• Signal in vector space:

$$s_m(t) = \sum_{i=1}^N s_{mi} \varphi_i(t) \rightarrow \bar{s}_m = (s_{m1}, s_{m2}, \dots, s_{mN})$$

↳ each waveform \rightarrow corresponds to a vector with N coefficients

$$d = \sqrt{(s_{m1} - s_{21})^2 + (s_{m2} - s_{22})^2}$$

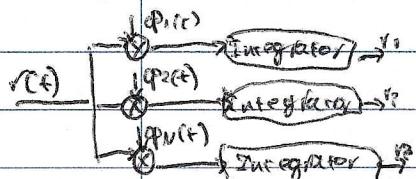
• Signal Space: $\{\bar{s}_1, \bar{s}_2, \bar{s}_M\} \rightarrow$ signal constellation, set of vectors.

$$\bullet \text{No noise } r(t) = s_m(t) \rightarrow \int_{-\infty}^{\infty} s_m(t) \varphi_i(t) dt = \int_{-\infty}^{\infty} \left(\sum_{j=1}^N s_{mj} \varphi_j(t) \right) \varphi_i(t) dt. \quad r_i = \sum_{j=1}^N s_{mj} r_j$$

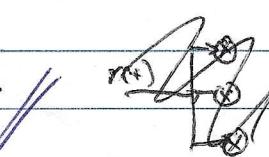
$$\bullet \text{noise } r(t) = s_m(t) + n_w(t) \rightarrow r_i = \int_{-\infty}^{\infty} r(t) \varphi_i(t) dt = \int_{-\infty}^{\infty} s_m(t) \varphi_i(t) dt + \int_{-\infty}^{\infty} n_w(t) \varphi_i(t) dt$$

$$\bullet \text{Noise signal space: Joint PDF of relevant noise: } p_{\bar{n}}(\bar{n}) = \frac{1}{(nW)^{N/2}} \exp\left(-\frac{|\bar{n}|^2}{nW}\right) \quad \text{symmetric}$$

• Irrelevant noise: noise is signal space relevant, all other dimensions not.



$$\bullet \text{SNR in vector space} = \frac{\|s\|^2}{nW}$$



dimensionality = 1 \rightarrow when all signals are scalar
1 building block if $\phi_1, \phi_2 \rightarrow$ dimensionality = 2

1 dimensional mean (\bar{x}) not (\bar{xy})

so no S_{11}, S_{21} only S_{11}

$$\text{Orthogonality} \rightarrow \int_{-\infty}^{\infty} S_1(t) S_2(t) dt = 0$$

$$\text{Orthonormality} \rightarrow E_1 = \int_{-\infty}^{\infty} |S_1(t)|^2 dt = 1$$

$$E_2 = \int_{-\infty}^{\infty} |S_2(t)|^2 dt = 1.$$

$$P_{\text{avg}} = \sum_{m \in M} \Pr\{M=m\} \|S_m - \bar{S}\|^2 = E[\|\bar{S} - \bar{S}\|^2]$$

basically the euclidean mean distance

sum \bar{S} is P_{avg} .

antipodal minimize P_{avg}

Channel Processing Receivers

Ch 6 - Receiver Implementation, Matched Filters

vector representation of signals

- original representation: $f(t) = \sum_{i=1}^N f_i(t)$, $g(t) = \sum_{i=1}^N g_i(t)$
- vector product: $(f \cdot g) = \sum_{i=1}^N f_i g_i$
- Norm and Distance: $\|f\|^2 = (f \cdot f) = \sum_{i=1}^N f_i^2$; $\|f-g\|^2 = (f-f)^2 = \|f\|^2 + \|g\|^2 - 2(f \cdot g)$
- correlation as dot prod: $\int_{-\infty}^{\infty} f(t)g(t)dt = (f \cdot g)$

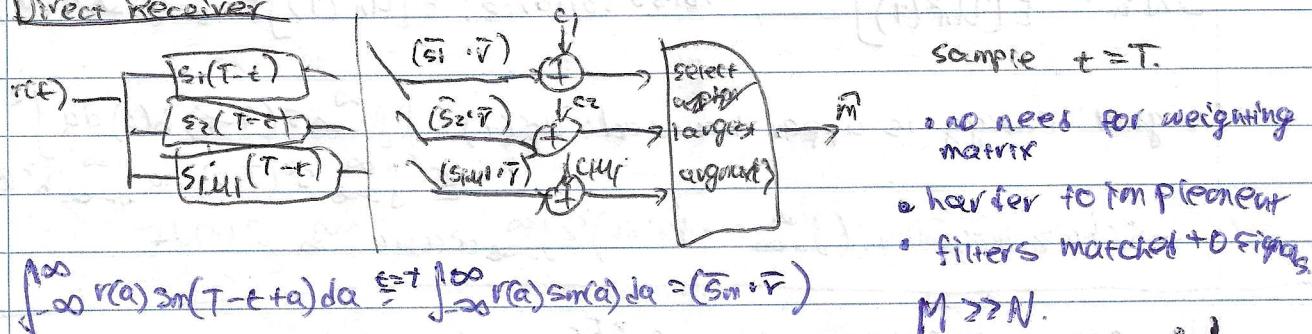
Optimum Receiver

$$\hat{m}_{MAP}(r) = \arg \max_{m \in M} \{(\bar{r} \cdot \bar{s}_m) + c_m\}; c_m = N_0 / \epsilon \ln P_r \{M=m\} - E_m / \epsilon$$

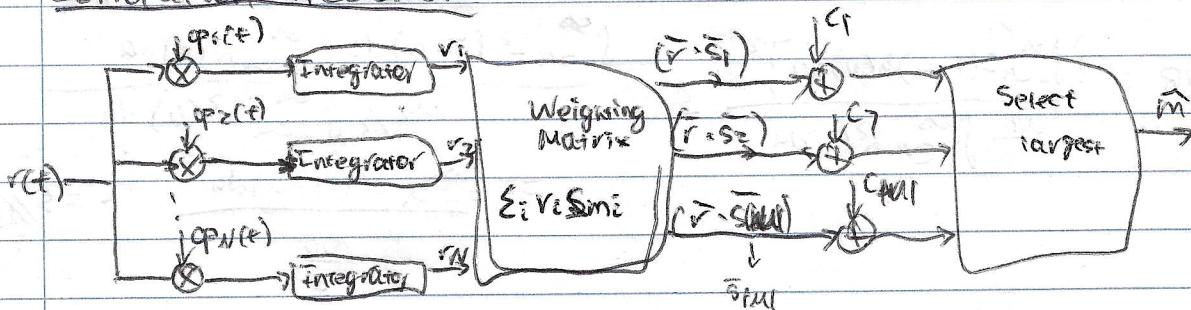
$$E_m = \int_{-\infty}^{\infty} S_m^2(t) dt.$$

- Decision Criterion: minimize metric: $\|\bar{r} - \bar{s}_m\|^2 - N_0 \ln P_r \{M=m\}$
- $$\Rightarrow \|\bar{r}\|^2 - 2(\bar{r} \cdot \bar{s}_m) + \|\bar{s}_m\|^2$$

Direct Receiver



Correlation Receiver



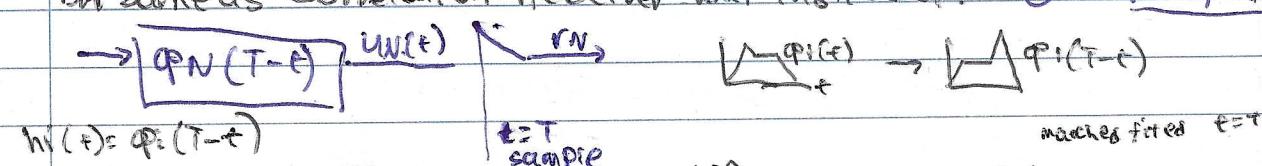
Recovered Signal Components

$$r_i = \int_{-\infty}^{\infty} r(a) q_i(t) da; (\bar{r} \cdot \bar{s}_m) = \sum_i r_i s_{mi} \rightarrow \text{matrix rep of decisions}$$

$$\begin{bmatrix} (r \cdot s_1) \\ (r \cdot s_2) \\ \vdots \\ (r \cdot s_M) \end{bmatrix} = \begin{bmatrix} s_{11} & s_{12} & \cdots & s_{1N} \\ s_{21} & s_{22} & \cdots & s_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ s_{M1} & s_{M2} & \cdots & s_{MN} \end{bmatrix} \begin{bmatrix} r_1 \\ r_2 \\ \vdots \\ r_N \end{bmatrix}$$

Matched Filter Receiver

BN Same as Correlation Receiver but instead of:



$$v_i(t) = r(t) * h_i(t) = \int_{-\infty}^{\infty} r(a) h_i(t-a) da = \int_{-\infty}^{\infty} r(a) q_i(T-t+a) da =$$

$$\text{for } t=T \rightarrow v_i(T) = \int_{-\infty}^{\infty} r(a) q_i(a) da = v_i //$$

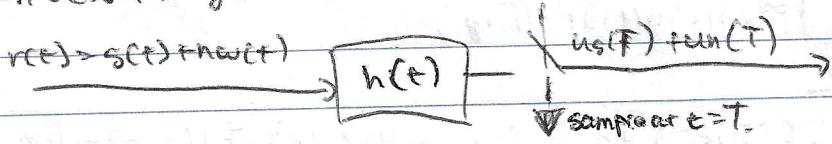
Instead of multipliers/integrators, samplers → easier practically for analog signals.

Matched Filter Receiver

Same as correlation receiver but instead of

Signal to Noise Ratio

ANALYING FILTER



$$u(T) = \int_{-\infty}^{\infty} r(T-a) h(a) da = u_s(T) + u_n(T)$$

$$u_s(T) \triangleq \int_{-\infty}^{\infty} s(T-a) h(a) da; u_n(T) \triangleq \int_{-\infty}^{\infty} n_w(T-a) h(a) da.$$

$$SNR \triangleq \frac{u_s^2(T)}{E[u_n^2(T)]}; \text{ noise variance: } E[u_n^2(T)] = \frac{N_0}{2} \int_{-\infty}^{\infty} h^2(a) da.$$

signal energy is upper bounded: $u_s^2(T) = [\int_{-\infty}^{\infty} s(T-a) h(a) da]^2 \leq \int_{-\infty}^{\infty} s^2(T-a) da \int_{-\infty}^{\infty} h^2(a) da$

Schwarz inequality: $(\int_{-\infty}^{\infty} a(t)b(t) dt)^2 \leq \int_{-\infty}^{\infty} a^2(t) dt \int_{-\infty}^{\infty} b^2(t) dt$

Max attained SNR: $SNR \leq \frac{2E_s}{N_0}$! Matched Filter Receiver

minimizes P_e , max SNR

$$SNR = \frac{[\int_{-\infty}^{\infty} s(T-a) h(a) da]^2}{\frac{N_0}{2} \int_{-\infty}^{\infty} h^2(a) da} \leq \frac{\int_{-\infty}^{\infty} s^2(T-a) da \int_{-\infty}^{\infty} h^2(a) da}{N_0 / 2 \int_{-\infty}^{\infty} h^2(a) da} = \frac{\int_{-\infty}^{\infty} s^2(T-a) da}{\frac{N_0}{2}} = \frac{E_s}{N_0 / 2}$$

Channel Processing:

Transmitters

Ch 7 Signal Energy Considerations, Orthogonal Signals

Binary Signaling

• energy/avg energy: $E_{sm} = \int_0^T S_m^2(t) dt = \|S_m\|^2$

$$\lim_{M \rightarrow \infty} E[S] = 0$$

$$E_{av} = \sum_m P(M=m) E_{sm} = E[\|S\|^2]$$

• Binary Orthogonal Signaling: signal vectors for orthogonal signals.

$$S_1 = (\sqrt{E_s}, 0), \quad S_2 = (0, \sqrt{E_s}), \quad P_e = Q(\frac{\sqrt{E_s}}{\sqrt{N_0}})$$

$$P_e^{orth} = Q\left(\frac{\sqrt{E_s}}{\sqrt{N_0}}\right) \quad \text{or} \quad N_0/2 \Rightarrow 0 = \sqrt{\frac{N_0}{2}}$$

$$\text{AWGN: } Q\left(\frac{d}{2\sigma}\right) \quad \text{or}$$

• Binary antipodal signaling: signal vectors for antipodal signals.

$$S_1 = (\sqrt{E_s}, 0), \quad S_2 = (-\sqrt{E_s}, 0); \quad P_e^{antip} = Q\left(\frac{\sqrt{2E_s}}{\sqrt{N_0}}\right)$$

Non-Binary Orthogonal Signaling / capacity results

• Orthogonal Signal Sets:

- signal representation: $S_m = \sqrt{E_s} \vec{\varphi}_m \quad m \in M$

$$\vec{\varphi}_1 = (1, 0, 0), \quad \vec{\varphi}_2 = (0, 1, 0), \quad \vec{\varphi}_3 = (0, 0, 1) \quad M=3$$

• Optimum Receiver for orthogonal signals: decision rule:

$$\hat{m} = \operatorname{argmax}_{m \in M} (\vec{r} \cdot \vec{s}_m) \Rightarrow \operatorname{argmin}_{m \in M} \{\|\vec{r} - \vec{s}_m\|^2\} = \operatorname{argmin}_{m \in M} \{r_m\}$$

$$P_c = \int_{-\infty}^{\infty} P_R(a | M=1) P(\hat{M}=1 | M=1, R_1=a) da.$$

$$P_e \cancel{=} \left\{ \begin{array}{l} 2 \exp(-[E_b/N_0 - 1/\ln 2]^2 \log_2 M) \text{ if } E_b/N_0 < 4/\ln 2 \\ 2 \exp(-[E_b/(2N_0) - 1/\ln 2]^2 \log_2 M) \text{ if } E_b/N_0 \geq 4/\ln 2 \end{array} \right.$$

• Channel Capacity: wideband capacity:

$$R = P_s / E_b \leq P_s / N_0 \ln 2 \quad \text{bits/sec} \quad P_s \rightarrow \text{transmitter power}$$

$$\cdot SDR = P/N_0 w \quad , \quad \text{Shannon capacity: } C = B \log_2 (1 + SNR)$$

lets say
you want
0.001 Pe

a. Achieve same Pe \rightarrow orthogonal signaling requires $(\times 2)$ energy

of antipodal signaling ($\approx 3\text{dB}$ disadvantage).

b. orthogonal: $s_1 = (\sqrt{E_s}, 0) \rightarrow s_2 = (0, \sqrt{E_s}) \rightarrow \text{Euclidean dist: } \sqrt{2E_s} \rightarrow Q\left(\frac{\sqrt{2E_s}}{\sqrt{N_0}}\right)$

c. antipodal: $s_1 = (-\sqrt{E_s}, 0), \quad s_2 = (\sqrt{E_s}, 0) \rightarrow \text{euclidean dist: } \sqrt{2E_s}$

d. For AWGN channels \rightarrow Pe depends on Euclidean distance.

$$10 \log_{10}(2) = 3\text{dB}.$$

to achieve same Pe orthogonal $(\times 2)$

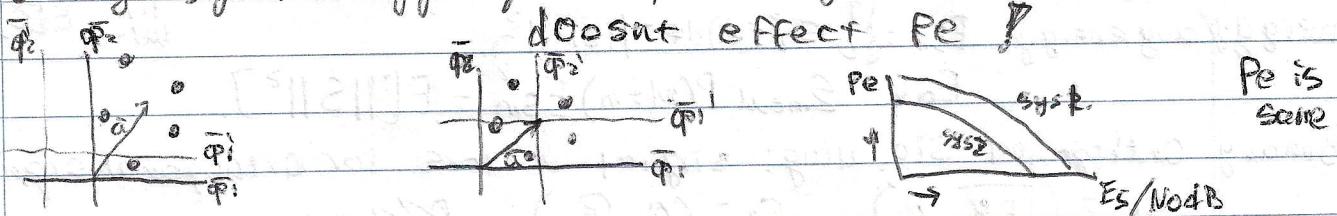
but orthogonal easier to implement on
coherent detection \rightarrow since receiver doesn't
need signal phase (better for practical systems)

$$Q\left(\frac{\sqrt{2E_s}}{\sqrt{N_0}}\right) \rightarrow$$

$$Q\left(\frac{\sqrt{2E_s}}{\sqrt{2N_0}}\right) = Q\left(\frac{\sqrt{E_s}}{\sqrt{N_0}}\right)$$

$$Q\left(\frac{\sqrt{E_s}}{\sqrt{N_0}}\right)$$

- Translating signal structure closer to origin minimizes average signal energy required, keeping decision regions same



- Orthogonal (FSK) waveforms ($M=2^1, 2^2$, $\Pr\{M=1\} = \Pr\{M=2\} = 1/2$)

$$S_1(t) = \sqrt{2E_s} \sin(2\pi f_1 t) \quad 0 \leq t < T$$

$$\bar{s}_1 = (\sqrt{E_s}, 0)$$

$\Pr\{M=1\}$ WUR

$$S_2(t) = \sqrt{2E_s} \sin(2\pi f_2 t) \quad 0 \leq t < T$$

$$\bar{s}_2 = (0, \sqrt{E_s})$$

$\Pr\{M=2\}$

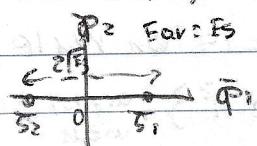
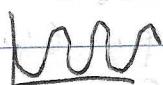
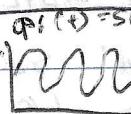
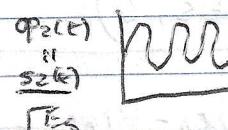
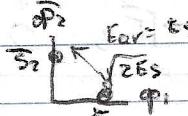
- Antipodal (PSK) waveforms

$$|M|=2^1, 2^2 \quad \Pr\{M=1\} = \Pr\{M=2\} = 1/2$$

$$s_1(t) = \sqrt{2E_s} \sin(10\pi t) \quad 0 \leq t < T$$

$$s_2(t) = -\sqrt{2E_s} \sin(10\pi t) \quad 0 \leq t < T$$

$$\bar{s}_1 = (\sqrt{E_s}, 0), \quad \bar{s}_2 = (-\sqrt{E_s}, 0)$$



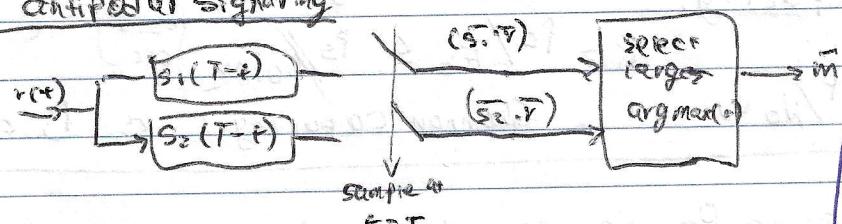
AGN Vector Channel \rightarrow Pr that noise pushes signal to wrong hyperplane $P_I = Q(\frac{\Delta}{\sigma})$ $\Delta \rightarrow$ dist point + hyperplane

Pe arch \rightarrow Pe doesn't change but for antipodal \rightarrow same Pe for lower energy can be achieved

rotations \rightarrow no change $Pe \rightarrow$ ~~but~~ save energy.

Receivers of antipodal signaling

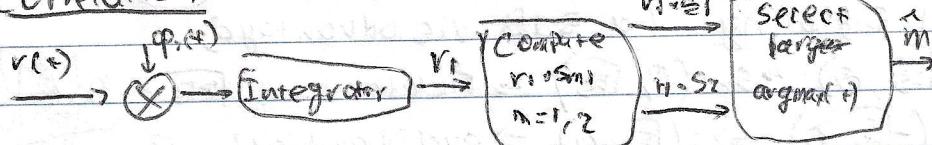
Direct



$$E_b \triangleq \frac{E_s}{\log(M)}$$

energy per transmitter bit of info.

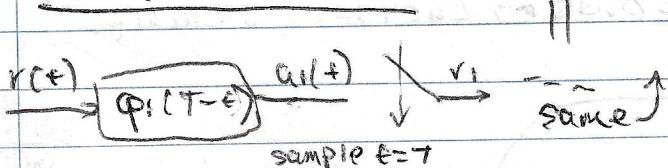
Correlation



$$E_{av} = \sum_{i=1}^N \Pr\{M=i\} \|s_i\|^2$$

$$= 1/2 (f_1)^2 + 1/2 (f_2)^2 = 2$$

Matched Filter Receiver



assuming $\bar{x} = 0$. origin.

antipodal minimize
Pe

Channel Processing: Transmitters.

$$A = \frac{h \cdot b}{2} \quad \text{Area of triangle}$$

Ch 8 Message Sequence, Bandwidth

Bit-by-bit signaling

$$E_{avg} = E_s = P_s T \quad \text{Joule}$$

• Transmission Rate: $R = \frac{\log_2 M}{T} \quad \text{bits/second}$

$$T = \frac{T}{2^K}$$

• Energy per bit: $E_b = \frac{E_s}{R} = \frac{P_s}{\log_2 M} \quad \text{Joule/bit}$

• Signal representation: transmitted signal: $s(t) = \sum_{i=1}^K (-1)^{b_i} p(t-(i-1)\tau)$

• Signal representation: normalized waveforms: $\phi_i(t) = p(t-(i-1)\tau)$

• Signal as vector: $s_m = \sqrt{E_b} ((-1)^{b_1}, (-1)^{b_2}, \dots, (-1)^{b_K})$

$$\sqrt{E_b}$$

signal dimensionality: $N = K$

• $P_c = (1 - Q(\sqrt{\frac{2E_b}{N_0}}))^K \quad ; \quad P_e = 1 - P_c$

$$Q(\frac{t}{\sqrt{E_b}})$$

Block Orthogonal Signaling

$$E_b/N_0 = (H/G)^2 \ln 2.$$

$R = \frac{K}{T} \quad ; \quad \text{orthogonal signal: } s_m(t) = \sqrt{E_s} \phi((t-(m-1)\tau), m = 1, \dots, K)$

$P_e \leq 2 \exp(-[\sqrt{\frac{E_b}{N_0}} - \sqrt{\ln 2}]^2 \log_2 M) = 2 \exp(-E^2 R T / \ln 2); \quad E = \text{energy efficiency}$

• Channel Capacity: $C_{ch} = P_s / (N_0 \ln 2) \quad \text{bit/sec.}$

• If $R < C_{ch}$ achievable with low Pe increasing T.

• Dimensionality / bandwidth: number of orthogonal dimensions: $N \leq W T$, where W is bandwidth and T, time duration.

• Stream of messages $\rightarrow s_m(t) \neq 0$ OLTCT \rightarrow equiprobable

• bit/bit signaling: $\rightarrow K$ binary digits: $(b_1, b_2, b_3, \dots, b_K)$ at int., $M = 2^K, R = \frac{\log_2 M}{T} = \frac{K}{T}$

geometric: $s(t) = \sum_{i=1}^K (-1)^{b_i} \sqrt{E_s} \phi_i(t).$

$$s_m = \sqrt{E_s} ((-1)^{b_1}, (-1)^{b_2}, \dots, (-1)^{b_K})$$

$$S = \{s_1(t), s_2(t), \dots\}$$

• optimum receiver: $A_i = 1$ if $b_i < 0$.

$$K = N = 1 \quad M = 2$$

no error if $r_i = -\sqrt{E_b} + n_i \quad (0 \text{ for all } i = 1, \dots, K)$

$$\underbrace{\sqrt{E_s}}_{2\sqrt{E_s}} \underbrace{\phi_i}_{\sqrt{E_b}} \underbrace{n_i}_{\sqrt{N_0}}$$

$$11010 \dots$$

• $K = RT$ and $E_b = P_s / R$

$$P_e = 1 - \left(1 - Q\left(\sqrt{\frac{2P_s}{RN_0}}\right)\right)^{RT}$$

$$RT = 1/R \Rightarrow K = 1 \Rightarrow P_e = Q\left(\sqrt{\frac{2P_s}{RN_0}}\right)$$

• If Pe decreased by increasing P_s decreasing R

$$R = \frac{1}{\sqrt{1+e^P} N_0 \ln 2}$$

$$P_e \leq 2 \exp(-CRT \ln 2)$$

• reliability can't increase by inc T,

! reliability increase by increasing
 P_s and decreasing R (bit by bit)
increasing code word T.

for bit by bit signaling

→ But can with orthogonal

Signaling	Dimensions per block	Dimensions per sec.
Bit by Bit	$K = RT$	$K/T = R$
Block Orthogonal	$2^K = 2^{RT}$	$2^K/T = 2^{RT}/T$

• Block Orth. signaling: Pros / cons.

- large reliability in $C/I(T)$ (not bit-by-bit)

- num dimensions explode $\propto T$ per inc in T .

- Channel with finite B can't accommodate all dimensions

Dimensionality theorem

$$1. \int_{-W}^{+W} |\mathcal{P}(\text{CCDF}(f))| dt \leq 1.$$

$$2. P_1(t) = 0 \text{ for all } t \in [0, T] \quad *$$

max number of orthogonal waveforms

upper bounded by $2WT$.

$$\cdot 2^K/T \rightarrow (2RT)^1/T.$$

• Increasing P_s , lowering $R \rightarrow$ bit/bit and Orthogonality

• Increasing codeword $T \rightarrow$ orthogonal (reliability)

- Only up to $C/I \rightarrow$ achieve reliability

in non-freq limited spectrum

• here time limited not freq-limited

Ch10 \rightarrow freq limited

Channel Processing: Transmitters.

Ch 9 - Capacity of the Baseband and Wideband Channels

Channel Capacity of AWGN Vector Channel.

- Rate per dimension: $R_N = \frac{\log_2 M}{N} \text{ bits/dimension}$ $\rightarrow C_N = \frac{1}{2} \log_2 \left(\frac{EN + N_0/2}{N_0/2} \right) \text{ bits/sec}$
- $P_e \rightarrow 0$ if $R_N < C_N$.

- Channel capacity of baseband AWGN: $C = W \log_2 \left(1 + \frac{P_s}{W N_0} \right) \text{ bits/sec}$

→ scales linearly with bandwidth.

- Power limited capacity: $C_{\text{PS}} = \frac{P_s}{N_0 \ln 2}$

- bandwidth limited: $C = W \log_2 (\text{SNR}) \rightarrow C = \frac{P_s}{N_0 \ln 2}$

- Sphere hardening: mean/variance of square norm of normalized Gaussian vector: $E[(|G'|^2)] = \sigma_g^2$, $\text{var}[|G'|^2] = \frac{2\sigma_g^4}{N}$

- Error Prob in random code Gen: $R_N = C_N - \delta$

$$\lim_{N \rightarrow \infty} P_e \leq \lim_{N \rightarrow \infty} 2^{-\delta N} = 0$$

- Digital Comm: capacity binary symmetric channels (BSC): $C = 1 - H(p)$

Shannon Theorem reliable comm: $R_C \leq C \Rightarrow P_e \rightarrow 0$

→ sufficiently long codes //

- Module Objective: Reliable transmission $\rightarrow (P_e \rightarrow 0)$

at certain rate (R) increasing T . when (W) bandwidth
(P_s) transmission power \rightarrow (Fixed)

• block Orthogonal signalling \rightarrow reliable comm rates

$$\text{if } R < \frac{P_s}{N_0 \ln 2}$$

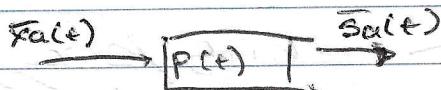
Channel Processing:
Transmitters.

K k

Chapter 8 \rightarrow time limited \rightarrow bandwidth unlimited

Ch 10 - Pulse Transmission.

[where $\bar{\alpha} = (\alpha_0, \alpha_1, \dots, \alpha_{K-1})$]



$$x(t) = \sum_{k=0}^{K-1} \alpha_k \delta(t - kT)$$

$$\bar{x}(t) = \sum_{k=0}^{K-1} \alpha_k \delta(t - kT)$$

$$\bar{S}(t) = \sum_{k=0}^{K-1} \alpha_k \Phi(t - kT)$$

$$S(t) = \sum_{k=0}^{K-1} \alpha_k \Phi(t - kT)$$

$$T_s > \frac{1}{R_b}$$

\Rightarrow linear Filter.

- building blocks \rightarrow impulses \rightarrow non-time-limited
- Finite bandwidth \rightarrow ISI created. \rightarrow infinite-time.

- Can have binary/non-binary (Pulse Amplitude Modulation) PAM.

$$\begin{array}{c} 0 \\ \vdots \\ i \\ \vdots \\ m \\ \hline 1 \\ \vdots \\ 00 \\ \vdots \\ 01 \\ \vdots \\ 11 \\ \vdots \\ 10 \\ \hline \end{array} \quad R = m = \log_2(M) \rightarrow M = 4 \rightarrow \log_2(M) = 2$$

trans rate

- ISI $\xrightarrow{\text{soln}}$ Nyquist criteria.

$$h(t) = p(t) * p(-t) \rightarrow H(f) = P(f)P^*(f) = |P(f)|^2$$

$|P(f)|^2 \rightarrow$ because $p(t) \rightarrow$ real

$$\int_{-\infty}^{\infty} p(t - kT) p(t - lT) dt = \begin{cases} 1 & k=l \\ 0 & k \neq l \end{cases}$$

$$= p(t) * p(-t) \Big|_{t=kT} = h(kT) = \begin{cases} 1 & k=0 \\ 0 & k \neq 0 \end{cases}$$

\circ Time Domain Restriction on $p(t) \rightarrow$

so nyquist in f-domain = $\frac{1}{T} \sum_{m=-\infty}^{\infty} |P(f + m/T)|^2$ zero forcing ($\geq F$)

$$|P(f + m/T)|^2 = H(f + m/T)$$

$$= 1 \quad \text{for all } f$$

Orthonormal pulses/Bandwidth

- $W < 1/(2T)$ $H(f) \neq |P(f)|^2$

Bandwidth: $W=3$

- $W > 1/(2T)$ \circ excess bandwidth \rightarrow ISI \rightarrow higher transmission capacity

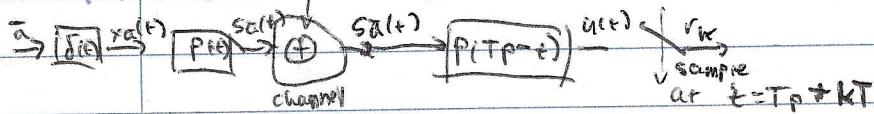
- $W > 1/(2T)$ \circ orthonormal but to min ISI \rightarrow sqrt raised cosine.

$$\circ W = \frac{1}{2\pi f_c} \quad \begin{array}{c} \text{Graph of a sinc-like pulse} \\ \text{from } -\frac{1}{2T} \text{ to } \frac{1}{2T} \end{array} \quad \circ P(t) = \frac{1}{\pi f_c} \sin(\pi f_c t / f_c); \quad T=1 \rightarrow \text{sinc pulses.}$$

$$\circ P(f) = \begin{cases} \pi f_c & |f| < 1/(2T) \\ 0 & |f| > 1/(2T) \end{cases}$$

- \circ orthogonal \rightarrow no ISI \rightarrow unrealistic --

PAM Transceiver. $nw(t) \rightarrow$ potential



$$x(t) \rightarrow [P(t - T_p - t)] \rightarrow u(t)$$

$$u(t) = \int_{-\infty}^{\infty} x(\tau) P(t - \tau + T_p) d\tau$$

$$rx = \sum_{k=0}^{K-1} \alpha_k \int_{-\infty}^{\infty} p(t - kT) p(-kT + \tau) d\tau + \int_{-\infty}^{\infty} nw(\tau) p(-kT + \tau) d\tau$$

\downarrow
 α_k
 \circ unless $k=0$

\downarrow
 nw

freq
limited
time
unlimited

bits / symbol $\rightarrow \log_2(N)$

$$R = r \cdot t \cdot e = \frac{\text{bits/symbol}}{T} = \text{bits/sec}$$

$$\text{Nyquist bandwidth: } \frac{1}{2T} = B$$

Power of the noise that will remain after

$$P_{\text{nw}} = \frac{N_0}{2} W$$

sampling receiver easier than match filter receiver

output noise power $\int_{-\infty}^{\infty} \frac{N_0}{2} |P(f)|^2 df \rightarrow \text{band limited}$

$$\int_{-\infty}^{\infty} \frac{N_0}{2} |P(f)|^2 df \stackrel{T/4 \text{ inserted}}{\longrightarrow} \left(\frac{N_0}{2}\right)^2 \int_0^{T/2} P(f)^2 df = \frac{N_0}{2} \text{ watts}$$

your question about noise reduction seems

and how to get maximum bandwidth

and still have a good SNR

Channel Processing: Transmitters

Pass Ch II - Pass-Band Channels

Passband Transmission / Capacity:

bandwidth of baseband: $BW \geq \frac{1}{2T}$. $T \rightarrow$ symbol period.

Passband channel capacity: $C_N = \frac{1}{2} \log_2 \left(1 + \frac{P_S}{N_0 W} \right)$ bits/dimension.

Passband Channel Capacity:

$$C_N = \frac{1}{2} \log_2 \left(1 + \frac{P_S}{N_0 W} \right) \text{ bits/dimension.}$$

$$C = 2W \log_2 \left(1 + \frac{P_S}{2N_0 W} \right) \text{ bits/sec.}$$

Baseband channel capacity:

$$C = W \log_2 \left(1 + \frac{P_S}{N_0 W} \right)$$

Quadrature Multiplexing.

Transmitted signal: $s_m(t) = s_c(t) \sqrt{2} \cos(2\pi f_c t) + s_s(t) \sqrt{2} \sin(2\pi f_c t)$

Orthogonality proof: $\int_{-\infty}^{\infty} q_{c,i}(t) q_{c,j}(t) dt = \delta_{i,j}$

$$\int_{-\infty}^{\infty} q_{s,i}(t) q_{s,j}(t) dt = \delta_{i,j}$$

$$\boxed{\int_{-\infty}^{\infty} q_{c,i}(t) q_{s,j}(t) dt = 0}$$

Matched / Correlation Receiver:

$$r_i = \int_{-\infty}^{\infty} r(t) q_i(t) dt$$

$$\hat{m} = \arg \max_{m \in M} \{ r \cdot s_m \} + c_m \}$$

$$C_m = \frac{N_0}{2} \ln P_r (M = m) = \frac{E_m}{2}$$

Quadrature multiplexing model

$$s_m(t) = s_c(t) \sqrt{2} \cos(2\pi f_c t) + s_s(t) \sqrt{2} \sin(2\pi f_c t)$$

$\hookrightarrow q_i(t), q_j(t) \rightarrow$ building block waveforms

$$s_m(t) = \sum_{i=1}^{N_c} s_{mi} q_i(t) \quad \text{and} \quad s_m(t) = \sum_{j=1}^{N_s} s_{mj} q_j(t)$$

\hookrightarrow transmitted waveform is addition of the two.

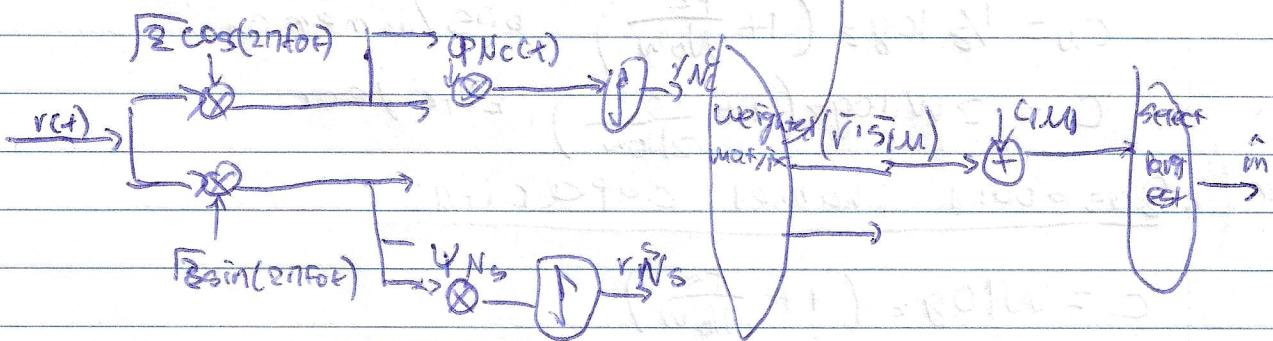
Quadrature optimum receiver.

$$\int_{-\infty}^{\infty} r(t) q_{ci}(t) dt = \int_{-\infty}^{\infty} r(t) \sqrt{2} \cos(2\pi f_0 t) q_{ci}(t) dt = r_i^c, \quad i=1, 2, \dots, N_c.$$

$$\int_{-\infty}^{\infty} r(t) q_{sj}(t) dt = \int_{-\infty}^{\infty} r(t) \sqrt{2} \sin(2\pi f_0 t) q_{sj}(t) dt = r_j^s, \quad i=1, 2, \dots, N_s$$

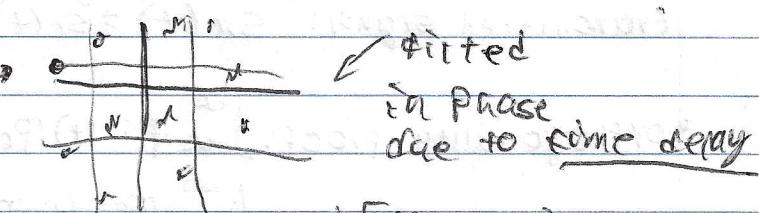
$$(\bar{r}, \bar{s}_m) = \sum_{i=1}^{N_c} r_i^c s_m^i + \sum_{j=1}^{N_s} r_j^s s_m^j$$

Correlation RRC Quadrature Receiver.

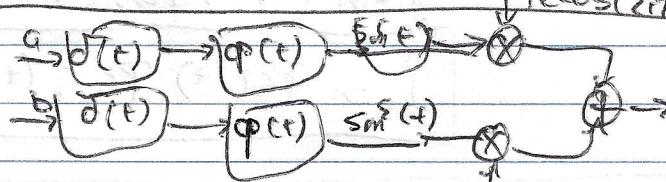


$$N_c > N_s = N \rightarrow q_i(t) = v_i(t) \rightarrow s_m(t) = \sum_{i=1}^{N_c} q_i(t) (s_m^c, F_2 \cos(2\pi f_0 t), s_m^s, F_2 \sin(2\pi f_0 t))$$

if $s_m^c \rightarrow 1 + n_w(t)$

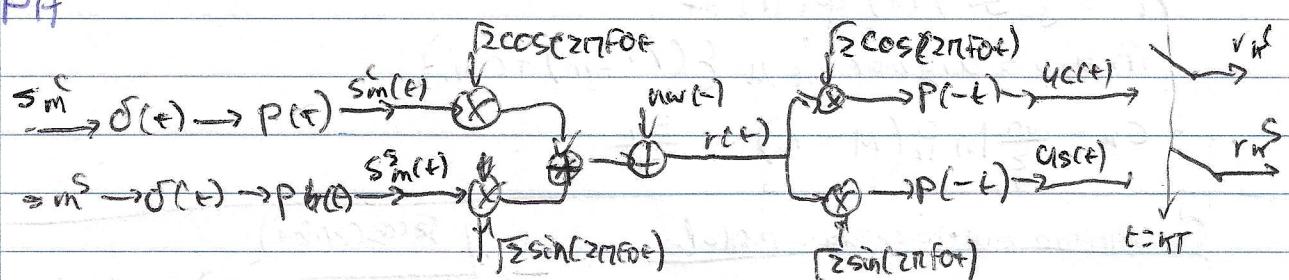


Normal QAM:



Serial QAM: $(p(t) = p_f)$

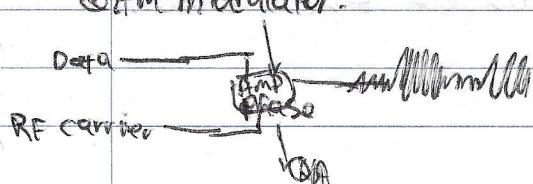
PA



! See Separate Page on Quadrature

64QAM & 16QAM \rightarrow noise tolerance.

QAM modulator.



QPSK = 01 symbols

16QAM = 0001 symbols

64QAM = 001011 symbols

Channel Processing : Transmitters.

Ch 12 - Random Carrier-Phase

• Coherent vs Incoherent Communication

• Incoherent \rightarrow no knowledge of θ phase means there is only cosine but can be decomposed to have sine component:

$$\cos(a+b) = \cos a \cos b + \sin a \sin b$$

• Antipodal event work for incoherent reception.

\rightarrow best but unrealistic: coherent antipodal.

Schreiber Signal Model: Quadrature-multiplexed Transmitted Waveform
N or $2N$ dimension.

$$s_m(t) = s_m^b(t)\sqrt{2} \cos(2\pi f_0 t - \theta)$$

$$= s_m^b(t) \cos(\theta) \sqrt{2} \cos(2\pi f_0 t) + s_m^b(t) \sin(\theta) \sqrt{2} \sin(2\pi f_0 t)$$

$$= \sum_{i=1}^N s_m^c \varphi_i(t) \sqrt{2} \cos(2\pi f_0 t) + \sum_{i=1}^N \varphi_i(t) \sqrt{2} \sin(2\pi f_0 t)$$

$$\text{where } \bar{s}_m = (s_{m1}, s_{m2}, \dots, s_{mN}) \cos(\theta) = \bar{s}_m \cos(\theta)$$

$$\bar{s}_m^s = (s_{m1}, s_{m2}, \dots, s_{mN}) \sin(\theta) = \bar{s}_m \sin(\theta)$$

Incoherent Receiver:

• decision rule: $\hat{m}_{MAP}(r) = \arg \max_{m \in M} \left\{ I_0 \left(\frac{E_m}{N_0} \right) \exp \left(-\frac{E_m}{N_0} \right) \right\}$

$$X_m = \sqrt{(r_c \cdot s_m)^2 + (r_s \cdot s_m)^2} ; I_0(x) \rightarrow \text{zero}$$

Order modified Bessel funct

E_m = message energy

• Equal energy Signals: $\hat{m}_{MAP}(r) = \arg \max_{m \in M} \left\{ (r_c \cdot s_m)^2 + (r_s \cdot s_m)^2 \right\}$

• Error Prob \rightarrow

$$\text{Incoherent: } P_e^{in} = 1/2 \exp \left(-\frac{E_s}{2N_0} \right)$$

$$\text{Coherent: } P_e^{coh} = Q \left(\sqrt{\frac{E_s}{N_0}} \right)$$

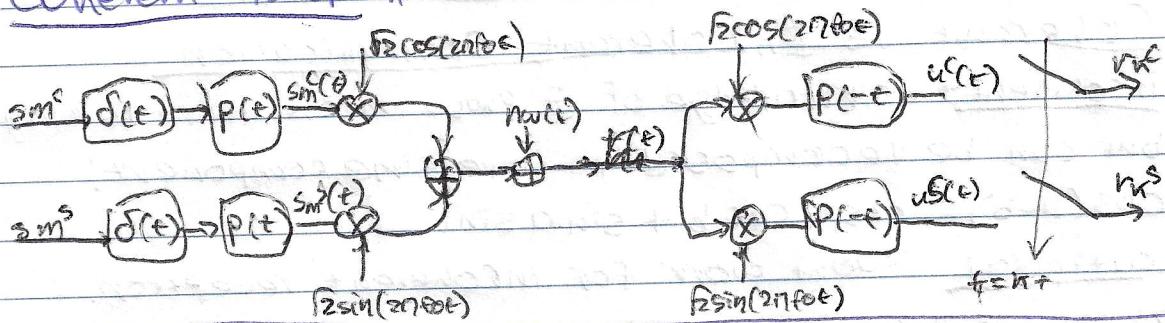
• Envelope detection: $X_m = \sqrt{(u_c^m(T))^2 + (u_s^m(T))^2}$

$u_c^m(T)$ correlation of $r(t)$ with $s_m(\omega_0 t)$ $\hookrightarrow u_s^m(T)$ correlation of $r(t)$ with $s_m(\omega_0 t)$

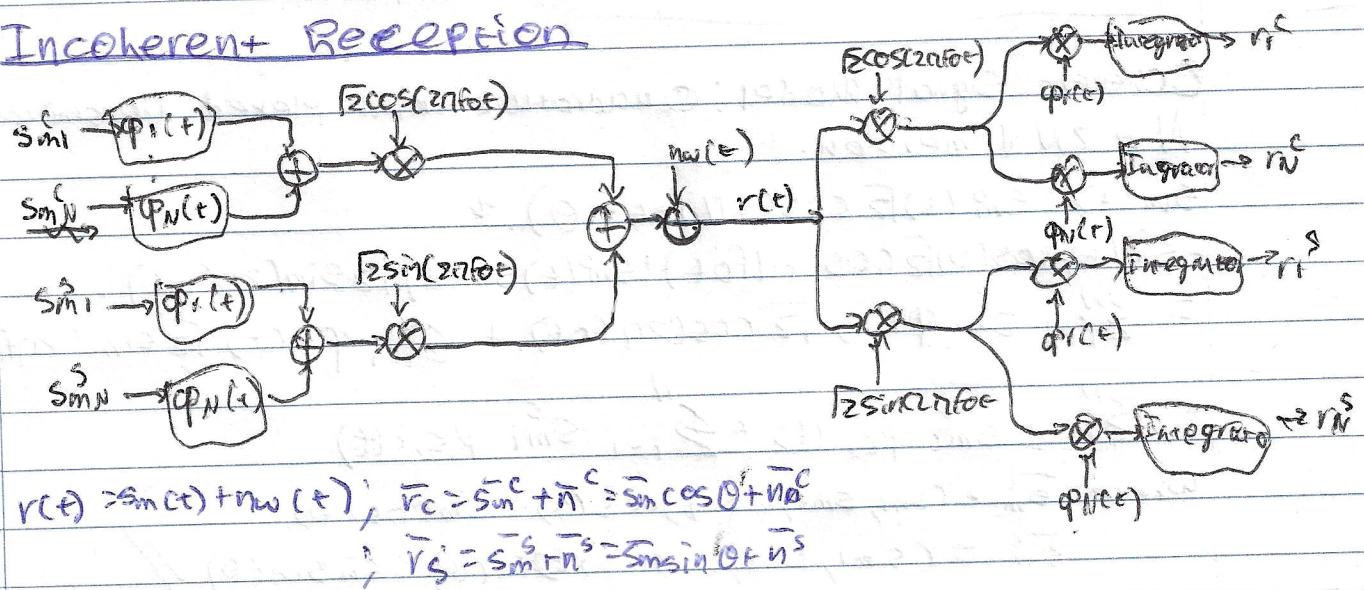
Wideband AWGN: $r(t) = s(t) + n_w(t)$.

Baseband AWGN: $r(t) = s(t) * w_b(t) + n_w(t)$.

Coherent Reception.

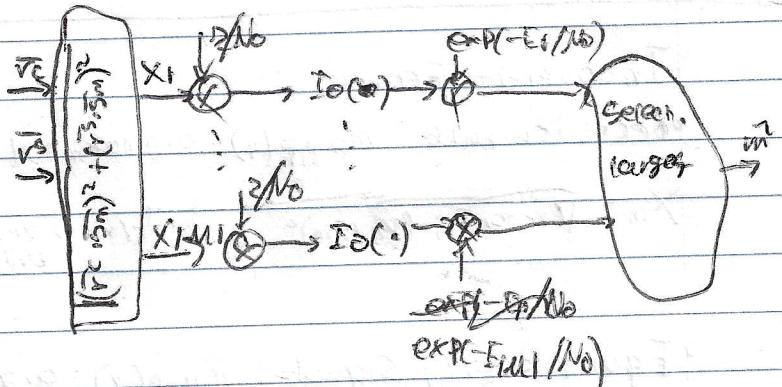


Incoherent Reception



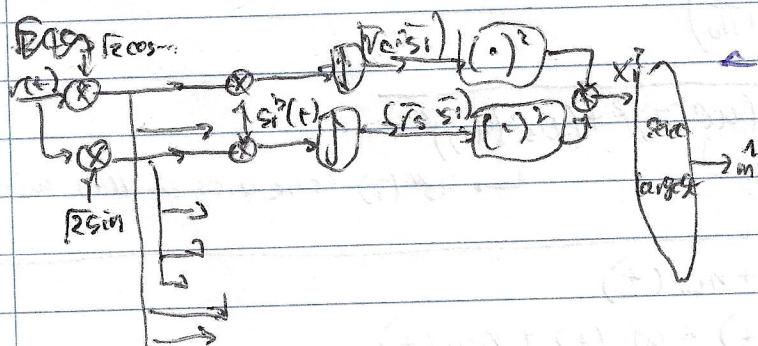
General Structure of Incoherent Receiver.

$$\hat{m}^{MAP}(r) = \underset{M \in M}{\operatorname{argmax}} \left\{ I_0 \left(\frac{2r}{N_0} \right) \exp \left(- \frac{E_m}{N_0} \right) \right\}$$



If equal energy

$$\underset{M \in M}{\operatorname{argmax}} \left\{ (r_c \cdot s_m)^2 + (r_s \cdot s_m)^2 \right\} \rightarrow \text{much simpler} //$$



The equal energy

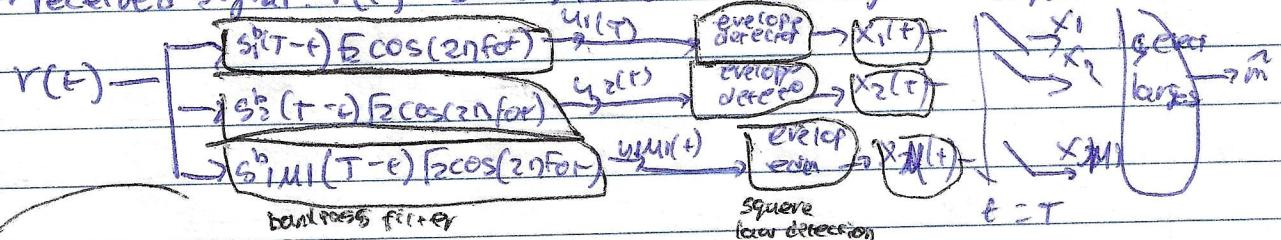
receiver is also easier
computation/Direct
Random phase receiver

Ch 12 - Cont - Random Carrier Receiver

Envelope Detection: Receiver structure, same energy:

$$\tilde{m}_{\text{map}}(\tau) = \arg \max_m \{x_m^2\} \rightarrow x_m = \sqrt{(r^2 \cdot s_m)^2 + (\tilde{r} \cdot s_m)^2}$$

Received signal: $r(t) = s_m^b(t) \sqrt{2} \cos(2\pi f_0 t - \theta) + n_w(t)$.



Output of filters: $u_m^c(t) = u_m^h(t) \cos(2\pi f_0 t) + u_m^s(t) \sin(2\pi f_0 t)$

baseband

$$\int_{-\infty}^{\infty} r(a) F_2 \cos(2\pi f_0 a) \overline{s_m^b(T-t+a)} da$$

baseband

$$\int_{-\infty}^{\infty} r(a) F_2 \sin(2\pi f_0 a) \overline{s_m^b(T-t+a)} da$$

baseband

$$x_m(t) = \sqrt{(u_m^h(t))^2 + (u_m^s(t))^2}$$

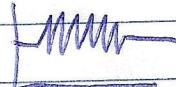
$s_m^b(t)$

$$s_m(t) = s_m^b(t) \sqrt{2} \sin(2\pi f_0 t)$$

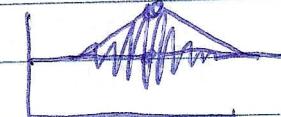
$$\theta = \pi/2, T=1$$



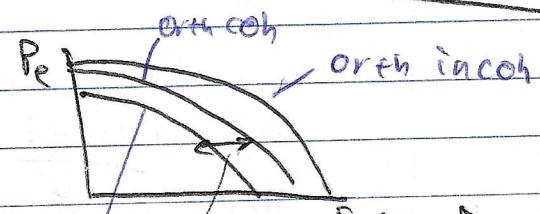
$u_m(t)$



$$u_m(t) = s_m^b(t) \sqrt{2} \cos(2\pi f_0 t)$$



$$P_e^{\text{in/ant}} = 1/2 \exp\left(-\frac{E_s}{2N_0}\right)$$



$$P_e^{\text{orth/coh}} = Q\left(\sqrt{\frac{E_s}{N_0}}\right) \leq 1/2 \exp\left(-\frac{E_s}{2N_0}\right)$$

$$P_e^{\text{anti/coh}} = Q\left(\sqrt{\frac{\sum E_s}{N_0}}\right)$$

anti
anti-
coh

incoherent cheaper easier but worse SNR
3dB loss

Invited Lecture - Ch 13

FSO : (Free Space optics) vs RF (Radio Frequency Communication)

- FSO \rightarrow 75 dB more gain for same antenna size, power, distance.

$$\text{• Receive power, } P_R = P_{RF} \eta P_T \left(\frac{D_1}{\lambda} \right)^2 \left(\frac{D_2}{\lambda} \right)^2 \left(\frac{1}{4\pi R} \right)^2$$

$P_{\text{Receive}} \propto 1/f^2 \propto f^{-2}$

- Due to high transmission gains: high directivity : FSO: LFD, LPI
Low Probability Detection, Interception.

$$\text{• Diffraction-limited spot size: } w(R) = w_0 \sqrt{1 + \left(\frac{R}{w_0^2}\right)^2} = \frac{d_R}{MUB}$$

- high throughput security, requires less energy per bit, license free

- DSP design, hard availability (water limited), location (directive sign), eye safety.

Moonshot FSO link to moon: 0.01nw Avg power received; single photon detector.

NIR wavelength 1550nm; given 256 ppm; proton efficiency?

$\text{ph } E = h \cdot F$; such that we have $\text{fix } f(h \cdot F) \text{ ph/sec.}$

- M-PPM \rightarrow M timeslots per symbol : send $\log_2(M)$ info bits per symbol.

- $$1 \text{ ph per symbol} \rightarrow \log_2(N) \text{ bits per symbol} \rightarrow \lceil \log_2(M) \rceil \text{ ph/bit}$$

- $$P_{\text{Tx}} = 0.01 \text{ mW} \rightarrow 7.8 \times 10^7 \text{ pW/sec.} \rightarrow M = 256 \rightarrow 623 \text{ Mbps} \rightarrow 0.125 \text{ ppb}$$

Calculate throughput.:

A: num of exanpls offical channels

B : Bandrate per channel

C : Modulation Format.

multiply. X

D: FEC & coding rate.

E: Polarizations

F : Transmission success factor.

G : Over head, header and protocols

- #### Solutions for fading and diversity methods.

- Temporal, ~~div~~ spatial, site frequency diversifying //

resend
message

multiple
receivers

↓
diff
place

↓
diff
freq

free space optical power rating than fibre optic since traverses through air.

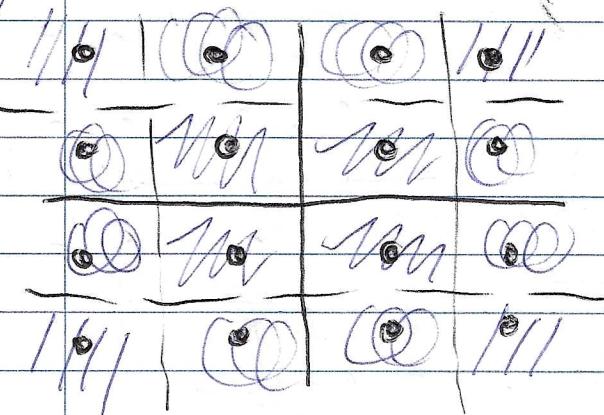
$$P_c = (1-2Q)(1-2Q)$$

$$= 1 - 4Q + 4Q^2$$

$$P_e = 1 - P_c$$

$$= 4Q - 4Q^2$$

Understanding Quadrature



Union Bound

neighbors

$$\text{I } P_e \leq 4Q\left(\frac{d}{20}\right)$$

$$\text{Q } P_e \leq 3Q\left(\frac{d}{20}\right)$$

$$\text{M } P_e \leq 2Q\left(\frac{d}{20}\right)$$

Exact P_e

Exact P_e

always
the same

$$\text{I } P_e = 4Q\left(\frac{d}{20}\right) - 4\left(\frac{d}{20}\right)^2$$

$$\text{Q } P_e = 3Q\left(\frac{d}{20}\right) - 2\left(\frac{d}{20}\right)^2$$

$$\text{M } P_e = 4Q\left(\frac{d}{20}\right) - 4Q^2\left(\frac{d}{20}\right)$$

$$\text{Q } P_e = 3Q\left(\frac{d}{20}\right) - 2Q^2\left(\frac{d}{20}\right)$$

$$\text{M } P_e = 2Q\left(\frac{d}{20}\right) - Q^2\left(\frac{d}{20}\right)$$

not the correct way to
look at it see 2024

Exam paper: → and ←

Serial QAM Upper Bound total P_e :

$$P_e \leq \frac{1}{16} \cdot 2Q\left(\frac{d}{20}\right) + \frac{8}{16} \cdot 3Q\left(\frac{d}{20}\right) + \frac{4}{16} \cdot 4Q\left(\frac{d}{20}\right) = \frac{48}{16}Q\left(\frac{d}{20}\right)$$

$$\text{Think of it this way: } \sum_{m \in \text{M}} \frac{1}{|m|} \sum_{m' \in \text{M}, m' \neq m} Q\left(\frac{d}{20}\right)$$

$$P_e = \frac{1}{16} (4(4Q\left(\frac{d}{20}\right)) + 4(2Q\left(\frac{d}{20}\right)) + 8(3Q\left(\frac{d}{20}\right)))$$

4 symbols | that have
4 neighbors | 4 symbols | that have
that have 2 neighbors

Exact Error Prob

- Center exact $P_e \rightarrow 4$ neighbors, (4 useful neighbors (II))

$$\hookrightarrow 4Q\left(\frac{d}{20}\right) - 4Q^2\left(\frac{d}{20}\right)$$

dist far → $\frac{d}{2}$

Upper bound base on union bound

$$P_e = \frac{1}{16} (1(4Q\left(\frac{d}{20}\right)) + 4(1Q\left(\frac{d}{20}\right) + 2Q\left(\frac{d}{20}\right)))$$

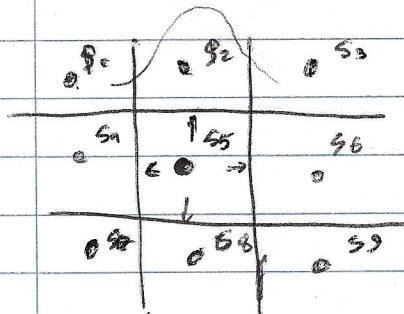
$B = A$

$$\hookrightarrow P_e = \frac{1}{16} (8Q\left(\frac{d}{20}\right) + 8Q\left(\frac{d}{20}\right)) //$$

union bound of average error probability.

Exact Pe Calculation

Exact Probability of Error



$$S5 \rightarrow P_{Pe} = (1-2Q)(1-2Q)$$

$$P_c = 4Q^2 - 4Q + 1$$

$$\begin{aligned} P_{Pe} &= 1 - P_c = 4Q^2 - 4Q \\ &= 1 - (4Q^2 - 4Q + 1) \\ &= 4Q - 4Q^2 \end{aligned}$$

The distances are relative to up/down in this case both same, both $\sqrt{2}$. $O(\sqrt{2})$.