

CS 4820, Spring 2017 Homework 10, Problem 3

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Solution:

1.) Prove  $m(s)$  tends to infinity as  $s$  tends to infinity

For  $s_1 < s_2$ , suppose  $m(s_1) > m(s_2)$ , this is a contradiction as if  $m(s_1) > m(s_2)$  according to the definition of  $m(s)$ ,  $s_2$  has to  $< s_1$ . So  $m(s_1) \leq m(s_2)$ . Suppose  $m(s_1) = m(s_2)$ , for  $Y = \{\text{input } y \text{ with length } = m(s_1)\}$ , let  $s$  be the max halting step for any  $y$  in  $Y$ . If  $s_2 > s$ ,  $m(s_1) = m(s_2)$  is a contradiction. So  $m(s_1) < m(s_2)$ , since the length of  $y$  increases by an integer, as  $s$  goes to infinity,  $m(s)$  goes to infinity.

2.) Reduce halting problem to this problem

Suppose such a  $M$  exists, For a halting problem with input:  $N;x$ , denote the binary representation of length of  $(N;x)$  by  $l$ , iteratively increase  $s$  from 1 and run  $M$  on  $x$  corresponding to  $s$  until the output of  $M(x)$   $z > l$ , find  $m(s)$  for this  $z$ , now run universal turing machine  $U$  on  $(N;x)$  for  $s$  steps. If  $U$  halts within  $s$  steps,  $N$  halts on  $x$ . It is a YES instance. Otherwise  $N$  never halts on  $x$ , it is a NO instance. To prove this, suppose  $U$  halts at  $s_3 > s$ , according to the definition of  $m(s)$ ,  $z$  for this  $m(s) \leq s$  which is a contradiction.

We now have a algorithm that decides the halting problem which is a contradiction as halting problem is undecidable. Therefore, there does not exist such a  $M$ .