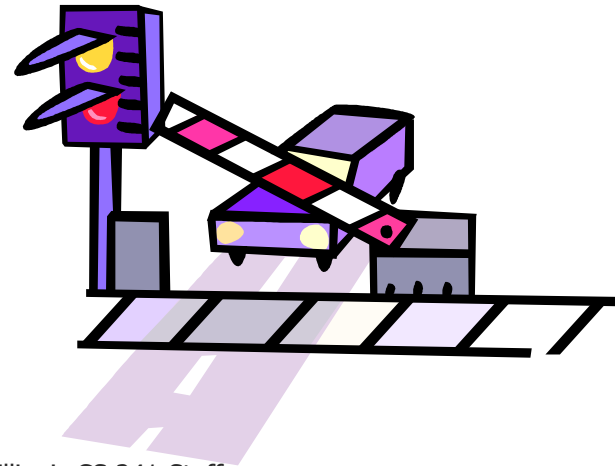


Synchronization and Semaphores



[Synchronization Primitives]

- Counting Semaphores
 - Permit a limited number of threads to execute a section of the code
- Binary Semaphores - Mutexes
 - Permit only one thread to execute a section of the code
- Condition Variables
 - Communicate information about the state of shared data



[POSIX Semaphores]

■ Named Semaphores

- Provides synchronization between unrelated process and related process as well as between threads
- Kernel persistence
- System-wide and limited in number
- Uses `sem_open`

→ ■ Unnamed Semaphores

- Provides synchronization between threads and between related processes
- Thread-shared or process-shared
- Uses `sem_init`



[POSIX Semaphores]

- Data type

- Semaphore is a variable of type `sem_t`

- Include `<semaphore.h>`

- Atomic Operations

```
int sem_init(sem_t *sem, int pshared, unsigned value);
```

```
int sem_destroy(sem_t *sem);
```

```
int sem_post(sem_t *sem);
```

```
int sem_trywait(sem_t *sem);
```

```
int sem_wait(sem_t *sem);
```



Unnamed Semaphores

```
#include <semaphore.h>
```

```
int sem_init(sem_t *sem, int pshared, unsigned value);
```

- Initialize an unnamed semaphore

- Returns

- 0 on success
- -1 on failure, sets **errno**

- Parameters

- **sem:**

- Target semaphore

- **pshared:**

- 0: only threads of the creating process can use the semaphore
 - Non-0: other processes can use the semaphore

- **value:**

- Initial value of the semaphore

You cannot make a copy of a semaphore variable!!!



[Sharing Semaphores]

- Sharing semaphores between threads within a process is easy, use **pshared==0**
- A non-zero **pshared** allows any process that can access the semaphore to use it
 - Places the semaphore in the global (OS) environment
 - Forking a process creates copies of any semaphore it has
 - Note: unnamed semaphores are not shared across unrelated processes



[`sem_init` can fail]

- On failure
 - `sem_init` returns -1 and sets `errno`

<code>errno</code>	cause
<code>EINVAL</code>	<code>Value > sem_value_max</code>
<code>ENOSPC</code>	Resources exhausted
<code>EPERM</code>	Insufficient privileges

```
sem_t semA;
```

```
if (sem_init(&semA, 0, 1) == -1)
    perror("Failed to initialize semaphore semA");
```



Semaphore Operations

```
#include <semaphore.h>
```

```
int sem_destroy(sem_t *sem) ;
```

- Destroy an semaphore

- Returns

- 0 on success
- -1 on failure, sets **errno**

- Parameters

- **sem**:
 - Target semaphore

- Notes

- Can destroy a **sem_t** only once
- Destroying a destroyed semaphore gives undefined results
- Destroying a semaphore on which a thread is blocked gives undefined results



[Semaphore Operations]

```
#include <semaphore.h>
```

```
int sem_post(sem_t *sem);
```

- Unlock a semaphore - same as signal
- Returns
 - 0 on success
 - -1 on failure, sets **errno** (**== EINVAL** if semaphore doesn't exist)
- Parameters
 - **sem**:
 - Target semaphore
 - $\text{sem} > 0$: no threads were blocked on this semaphore, the semaphore value is incremented
 - $\text{sem} == 0$: one blocked thread will be allowed to run



Semaphore Operations

```
#include <semaphore.h>
```

```
int sem_wait(sem_t *sem);
```

- Lock a semaphore
 - Blocks if semaphore value is zero
- Returns
 - 0 on success
 - -1 on failure, sets **errno** (**== EINTR** if interrupted by a signal)
- Parameters
 - **sem**:
 - Target semaphore
 - $\text{sem} > 0$: thread acquires lock
 - $\text{sem} == 0$: thread blocks



Semaphore Operations

```
#include <semaphore.h>
```

```
int sem_trywait(sem_t *sem);
```

- Test a semaphore's current condition
 - Does not block
- Returns
 - 0 on success
 - -1 on failure, sets **errno** (**== AGAIN** if semaphore already locked)
- Parameters
 - **sem**:
 - Target semaphore
 - $\text{sem} > 0$: thread acquires lock
 - $\text{sem} == 0$: thread returns



[Example: bank balance

- Protect shared variable **balance** with a semaphore when used in:
 - **decshared**
 - Decrements current value of **balance**
 - **incshared**
 - increments the **balance**



[Example: bank balance]

```
int decshared() {  
    while (sem_wait(&balance_sem) == -1)  
        if (errno != EINTR)  
            return -1;  
    balance--;  
    return sem_post(&balance_sem);  
}
```

```
int incshared() {  
    while (sem_wait(&balance_sem) == -1)  
        if (errno != EINTR)  
            return -1;  
    balance++;  
    return sem_post(&balance_sem);  
}
```





[Example: bank balance]

```
#include <errno.h>
#include <semaphore.h>
```

```
static int balance = 0;
static sem_t bal_sem;
```

```
int initshared(int val) {
    if (sem_init(&bal_sem, 0, 1) == -1)
        return -1;
    balance = val;
    return 0;
}
```

pshared

value



[Example: bank balance]

```
int decshared() {  
    while (sem_wait(&bal_sem)  
           == -1)  
        if (errno != EINTR)  
            return -1;  
    balance--;  
    return sem_post(&bal_sem);  
}
```

```
int incshared() {  
    while (sem_wait(&bal_sem)  
           == -1)  
        if (errno != EINTR)  
            return -1;  
    balance++;  
    return sem_post(&bal_sem);  
}
```

Which one is going first?



[Advanced Semaphores]

```
int semget(key_t key, int nsems, int semflg);
```

- Get set of semaphores

```
int semop(int semid, struct sembuf *sops,  
          unsigned int nsops);
```

- Atomically perform a user-defined array of semaphore operations on the set of semaphores



[Pthread Synchronization]

- Two primitives
 - Mutex
 - Semaphore with maximum value 1
 - Condition variable
 - Provides a shared signal
 - Combined with a mutex for synchronization



[Pthread Mutex]

■ States

- Locked
 - Some thread holds the mutex
- Unlocked
 - No thread holds the mutex

■ When several threads compete

- One wins
- The rest block
 - Queue of blocked threads



[Mutex Variables]

- A typical sequence in the use of a mutex
 1. Create and initialize **mutex**
 2. Several threads attempt to lock **mutex**
 3. Only one succeeds and now owns **mutex**
 4. The owner performs some set of actions
 5. The owner unlocks **mutex**
 6. Another thread acquires **mutex** and repeats the process
 7. Finally **mutex** is destroyed



[Creating a mutex]

```
#include <pthread.h>
```

```
int pthread_mutex_init(pthread_mutex_t *mutex,  
    const pthread_mutexattr_t *attr);
```

- Initialize a pthread mutex: the mutex is initially unlocked
- Returns
 - 0 on success
 - Error number on failure
 - **EAGAIN**: The system lacked the necessary resources; **ENOMEM**: Insufficient memory; **EPERM**: Caller does not have privileges; **EBUSY**: An attempt to re-initialise a mutex; **EINVAL**: The value specified by attr is invalid
- Parameters
 - **mutex**: Target mutex
 - **attr**:
 - NULL: the default mutex attributes are used
 - Non-NULL: initializes with specified attributes



[Creating a mutex]

- Default attributes
 - Use **PTHREAD_MUTEX_INITIALIZER**
 - Statically allocated
 - Equivalent to dynamic initialization by a call to **pthread_mutex_init()** with parameter **attr** specified as NULL
 - No error checks are performed



Destroying a mutex

```
#include <pthread.h>
```

```
int pthread_mutex_destroy(pthread_mutex_t  
    *mutex);
```

- Destroy a pthread mutex

- Returns

- 0 on success

- Error number on failure

- **EBUSY**: An attempt to re-initialise a mutex; **EINVAL**: The value specified by attr is invalid

- Parameters

- **mutex**: Target mutex



[Locking/unlocking a mutex]

```
#include <pthread.h>
int pthread_mutex_lock(pthread_mutex_t *mutex);
int pthread_mutex_trylock(pthread_mutex_t
    *mutex);
int pthread_mutex_unlock(pthread_mutex_t *mutex);
```

■ Returns

- 0 on success
- Error number on failure
 - **EBUSY**: already locked; **EINVAL**: Not an initialised mutex; **EDEADLK**: The current thread already owns the mutex; **EPERM**: The current thread does not own the mutex



[Simple Example]

```
#include <pthread.h>
#include <stdio.h>
#include <stdlib.h>
```

```
static pthread_mutex_t my_lock =
    PTHREAD_MUTEX_INITIALIZER;
```

```
void *mythread(void *ptr) {
    long int i,j;
    while (1) {
```

```
        pthread_mutex_lock (&my_lock);
```

```
        for (i=0; i<10; i++) {
            printf ("Thread %d\n", (int) ptr);
            for (j=0; j<500000000; j++);
        }
```

```
        pthread_mutex_unlock (&my_lock);
        for (j=0; j<500000000; j++);
```

```
    }
```

```
}
```

```
int main (int argc, char *argv[]) {
    pthread_t thread[2];
```

```
    pthread_create(&thread[0], NULL,
        mythread, (void *)0);
```

```
    pthread_create(&thread[1], NULL,
        mythread, (void *)1);
```

```
    getchar();
```

```
}
```



[Condition Variables]

- Used to communicate information about the state of shared data
 - Execution of code depends on the state of
 - A data structure or
 - Another running thread
- Allows threads to synchronize based upon the actual value of data
- Without condition variables
 - Threads continually poll to check if the condition is met



[Condition Variables]

- Signaling, not mutual exclusion
 - A mutex is needed to synchronize access to the shared data
- Each condition variable is associated with a single mutex
 - Wait atomically unlocks the mutex and blocks the thread
 - Signal awakens a blocked thread



[Creating a Condition Variable]

- Similar to pthread mutexes

```
int pthread_cond_init(pthread_cond_t *cond, const  
    pthread_condattr_t *attr);
```

```
int pthread_cond_destroy(pthread_cond_t *cond);
```

```
pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
```



Using a Condition Variable

■ Waiting

- Block on a condition variable.
- Called with **mutex** locked by the calling thread
- Atomically release **mutex** and cause the calling thread to block on the condition variable
- On return, **mutex** is locked again

```
int pthread_cond_wait(pthread_cond_t *cond,  
    pthread_mutex_t *mutex);
```

```
int pthread_cond_timedwait(pthread_cond_t *cond,  
    pthread_mutex_t *mutex, const struct timespec  
    *abstime);
```



[Using a Condition Variable]

■ Signaling

```
int pthread_cond_signal(pthread_cond_t *cond);
```

- unblocks at least one of the blocked threads

```
int pthread_cond_broadcast(pthread_cond_t *cond);
```

- unblocks all of the blocked threads

■ Signals are not saved

- Must have a thread waiting for the signal or it will be lost



Condition Variable: Why do we need the mutex?

```
pthread_mutex_lock(&mutex);           /* lock mutex */
while (!predicate) {                  /* check predicate */

    pthread_cond_wait(&condvar, &mutex);

    /* go to sleep - recheck
       pred on awakening */
}
pthread_mutex_unlock(&mutex);          /* unlock mutex */
```

```
pthread_mutex_lock(&mutex);           /* lock the mutex      */
predicate=1;                          /* set the predicate   */
pthread_cond_broadcast(&condvar);      /* wake everyone up    */
pthread_mutex_unlock(&mutex);          /* unlock the mutex    */
```



[Condition Variable: No mutex!]

```
pthread_mutex_lock(&mutex);           /* lock mutex */
while (!predicate) {                  /* check predicate */
    pthread_mutex_unlock(&mutex);      /* unlock mutex */
    pthread_cond_wait(&condvar);       /* go to sleep - recheck
                                         pred on awakening */
    pthread_mutex_lock(&mutex);        /* lock mutex */
}
pthread_mutex_unlock(&mutex);          /* unlock mutex */
```

————— What can happen here? —————

```
pthread_mutex_lock(&mutex);           /* lock the mutex */
predicate=1;                          /* set the predicate */
pthread_cond_broadcast(&condvar);     /* wake everyone up */
pthread_mutex_unlock(&mutex);          /* unlock the mutex */
```



Condition Variable: Why do we need the mutex?

- Separating the condition variable from the mutex
 - Thread goes to sleep when it shouldn't
 - Problem
 - `pthread_mutex_unlock()` and `pthread_cond_wait()` are not guaranteed to be atomic
- Joining condition variable and mutex
 - Call to `pthread_cond_wait()` unlocks the mutex
 - UNIX kernel can guarantee that the calling thread will not miss the broadcast



Using a Condition Variable: Challenges

- Call `pthread_cond_signal()` before calling `pthread_cond_wait()`
 - Logical error – waiting thread will not catch the signal
- Fail to lock the mutex before calling `pthread_cond_wait()`
 - May cause it NOT to block
- Fail to unlock the mutex after calling `pthread_cond_signal()`
 - May not allow a matching `pthread_cond_wait()` routine to complete (it will remain blocked).



Example without Condition Variables

```
int data_avail = 0;
pthread_mutex_t data_mutex =
    PTHREAD_MUTEX_INITIALIZER;

void *producer(void *) {
    pthread_mutex_lock(&data_mutex);

    <Produce data>
    <Insert data into queue;>
    data_avail=1;

    pthread_mutex_unlock(&data_mutex);
}
```



Example without Condition Variables

```
void *consumer(void *) {
```

```
    while( !data_avail ); /* do nothing */
```

```
    pthread_mutex_lock(&data_mutex);
```

```
    <Extract data from queue;>
```

```
    if (queue is empty)
```

```
        data_avail = 0;
```

```
    pthread_mutex_unlock(&data_mutex);
```

```
    <Consume Data>
```

```
}
```

Busy Waiting!



Example with Condition Variables

```
int data_avail = 0;
pthread_mutex_t data_mutex = PTHREAD_MUTEX_INITIALIZER;
pthread_cond_t data_cond = PTHREAD_COND_INITIALIZER;

void *producer(void *) {
    pthread_mutex_lock(&data_mutex);
    <Produce data>
    <Insert data into queue;>
    data_avail = 1;

    pthread_cond_signal(&data_cond);
    pthread_mutex_unlock(&data_mutex);
}
```



Example with Condition Variables

Mutex solution

```
while( !data_avail ); /* do nothing */
```

```
void *consumer(void *) {  
    pthread_mutex_lock(&data_mutex);  
    while( !data_avail ) {  
        /* sleep on condition variable*/  
        pthread_cond_wait(&data_cond, &data_mutex);  
    }  
    /* woken up */  
    <Extract data from queue;>  
    if (queue is empty)  
        data_avail = 0;  
    pthread_mutex_unlock(&data_mutex);  
    <Consume Data>  
}
```

No Busy Waiting!



[More Complex Example]

- Master thread
 - Spawns a number of concurrent slaves
 - Waits until all of the slaves have finished to exit
 - Tracks current number of slaves executing
- A mutex is associated with count and a condition variable with the mutex



[Example]

```
#include <stdio.h>
#include <pthread.h>

#define NO_OF_PROCS 4

typedef struct _SharedType {
    int count;                /* number of active slaves */
    pthread_mutex_t lock;     /* mutex for count */
    pthread_cond_t done;      /* sig. by finished slave */
} SharedType, *SharedType_ptr;

SharedType_ptr shared_data;
```





[Example: Main]

```
main(int argc, char **argv) {
    int res;

    /* allocate shared data */
    if ((sh_data = (SharedType *)
        malloc(sizeof(SharedType))) ==
        NULL) {
        exit(1);
    }
    sh_data->count = 0;

    /* allocate mutex */
    if ((res =
        pthread_mutex_init(&sh_data-
            >lock, NULL)) != 0) {
        exit(1);
    }
}
```

```
/* allocate condition var */
if ((res =
    pthread_cond_init(&sh_data-
        >done, NULL)) != 0) {
    exit(1);
}
```

```
/* generate number of slaves
to create */
srandom(0);
/* create up to 15 slaves */
master((int) random()%16);
}
```





[Example: Main]

```
main(int argc, char **argv) {
    int res;

    /* allocate shared data */
    if ((sh_data = (SharedType *)
        malloc(sizeof(SharedType))) ==
        NULL) {
        exit(1);
    }
    sh_data->count = 0;
```

```
pthread_mutex_t data_mutex =
    PTHREAD_MUTEX_INITIALIZER;
```

```
pthread_cond_t data_cond =
    PTHREAD_COND_INITIALIZER;
```

```
/* generate number of slaves
   to create */
srandom(0);
/* create up to 15 slaves */
master((int) random()%16);
}
```



[Example: Master]

```
master(int nslaves) {
    int i;
    pthread_t id;
    for (i = 1; i <= nslaves; i +=
        1) {
        pthread_mutex_lock(&sh_data-
            >lock);
        /* start slave and detach */
        shared_data->count += 1;
        pthread_create(&id, NULL,
            (void* (*)(void*))slave,
            (void *)sh_data);
        pthread_mutex_unlock(&sh_data-
            >lock);
    }
```

```
pthread_mutex_lock(&sh_data-
    >lock);
```

```
while (sh_data->count != 0)
    pthread_cond_wait(&sh_data-
        >done, &sh_data->lock);
```

```
pthread_mutex_unlock(&sh_data-
    >lock);
```

```
printf("All %d slaves have
    finished.\n", nslaves);
pthread_exit(0);
}
```





[Example: Slave]

```
void slave(void *shared) {
    int i, n;
    sh_data = shared;
    printf("Slave.\n", n);
    n = random() % 1000;

    for (i = 0; i < n; i+= 1)
        Sleep(10);
```

```
/* mutex for shared data */
pthread_mutex_lock(&sh_data-
    >lock);
```

```
/* dec number of slaves */
sh_data->count -= 1;
```

```
/* done running */
printf("Slave finished %d
    cycles.\n", n);

/* signal that you are done
    working */
pthread_cond_signal(&sh_data-
    >done);

/* release mutex for shared
    data */
pthread_mutex_unlock(&sh_data-
    >lock);
}
```



Semaphores vs. CVs

Semaphore

- Integer value (≥ 0)
- Wait does not always block
- Signal either releases thread or inc's counter
- If signal releases thread, both threads continue afterwards

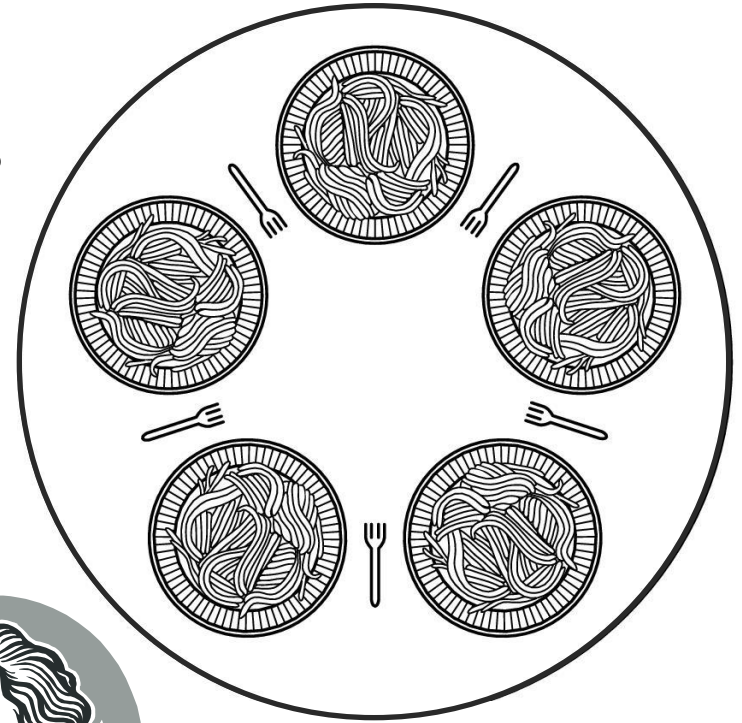
Condition Variables

- No integer value
- Wait always blocks
- Signal either releases thread or is lost
- If signal releases thread, only one of them continue



Dining Philosophers

- N philosophers and N forks
 - Philosophers eat/think
 - Eating needs 2 forks
 - Pick one fork at a time



Descartes Aristotle Plato Melville Kennedy