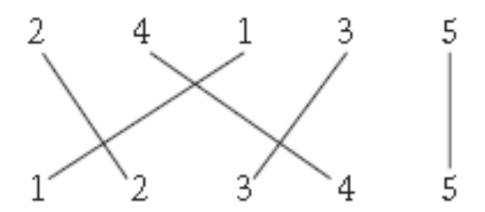
Introduction

- We are given a sequence of n numbers a1,...,an.
- We will assume that all the numbers are distinct.
- We want to define a measure that tells us how far this list is from being in ascending order.
- The value of the measure should be 0 if a1< a2 < . . . < an, and should increase as the numbers become more scrambled.

- Counting the Number of Inversions
- The two indices i < j form an inversion if ai > aj, that is, if the two elements ai and aj are "out of order."
- Counting the number of inversions is to determine the number of inversions in the sequence a1, . . . 3,

- Example
- Sequence is 2, 4, 1, 3, 5.
- There are three inversions in this sequence: (2, 1), (4, 1), and (4, 3).



- Example
- If the sequence is in descending order, then every pair forms an inversion, and so there are **n(n-1)/2** pairs.

```
{5,4}, {5,3}, {5,2}, {5,1}
{4,3}, {4,2}, {4,1}
{3,2}, {3,1}
{2,1}
```

Algorithm

- Look at every pair of numbers (ai, aj) and determine whether they constitute an inversion.
- This would take O(n^2) time.
- The basic idea is to follow the strategy of divide and conquer.
- We set m = n/2 and divide the list into the two pieces a1, . . . , am and am+1, . . . , an.
- We first count the number of inversions in each of these two halves separately.
- Then we count the number of inversions (ai, aj), where the two numbers belong to different halves.

Algorithm

- Note that these first-half/second-half inversions have a particularly nice form: they are precisely the pairs (ai, aj), where ai is in the first half, aj is in the second half.
- Suppose we have recursively sorted the first and second halves of the list and counted the inversions in each.
- We now have two sorted lists A and B, containing the first and second halves, respectively.
- We want to produce a single sorted list C from their union, while also counting the number of pairs (a, b) with $a \in A$, $b \in B$, and \dot{a}

Algorithm: Counting Inversions

```
//Purpose: To count the inversions for a given list L(a1,a2,.....an)
//Input: An unsorted list of distinct numbers L(a1,a2,.....an)
//Output: The number of inversions r for the list L(a1,a2,.....an)
    Sort-and-Count(L)
      If the list has one element then
           there are no inversions
      Else
           Divide the list into two halves:
              A contains the first \lceil n/2 \rceil elements.
               B contains the remaining \lfloor n/2 \rfloor elements.
           (r_A, A)=Sort-and-Count(A)
           (r_B, B)=Sort-and-Count(B)
           (r, L)=Merge-and-Count(A, B)
        Endif
       Return r = r_A + r_B + r, and the sorted list L
```

Merge-and-Count(A,B)

Maintain a Current pointer into each list, initialized to point to the front elements

Maintain a variable Count for the number of inversions, initialized to 0

While both lists are nonempty:

Let ai and bj be the elements pointed to by the Current pointer Append the smaller of these two to the output list

If bj is the smaller element then

Increment Count by the number of elements remaining in A

Endif

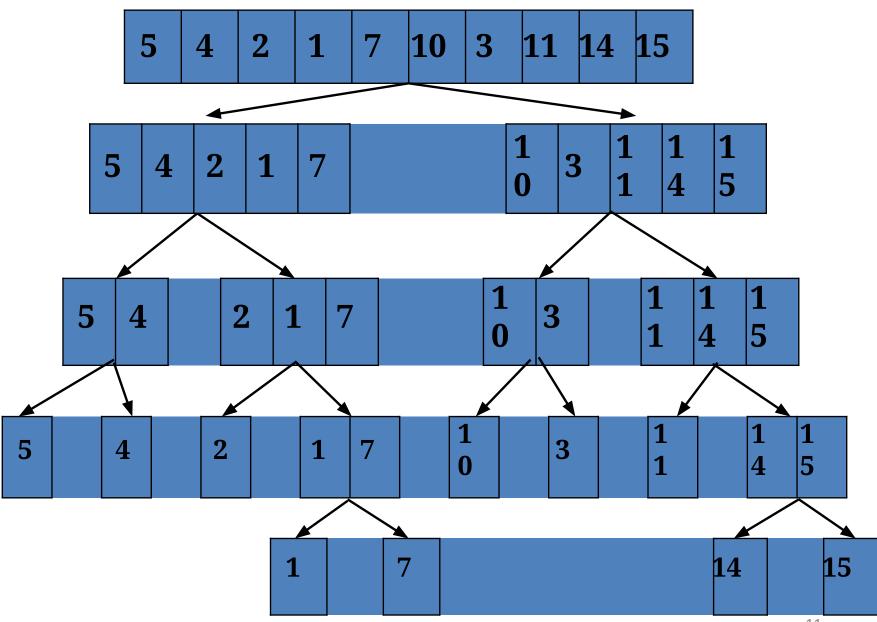
Advance the Current pointer in the list from which the smaller element was selected.

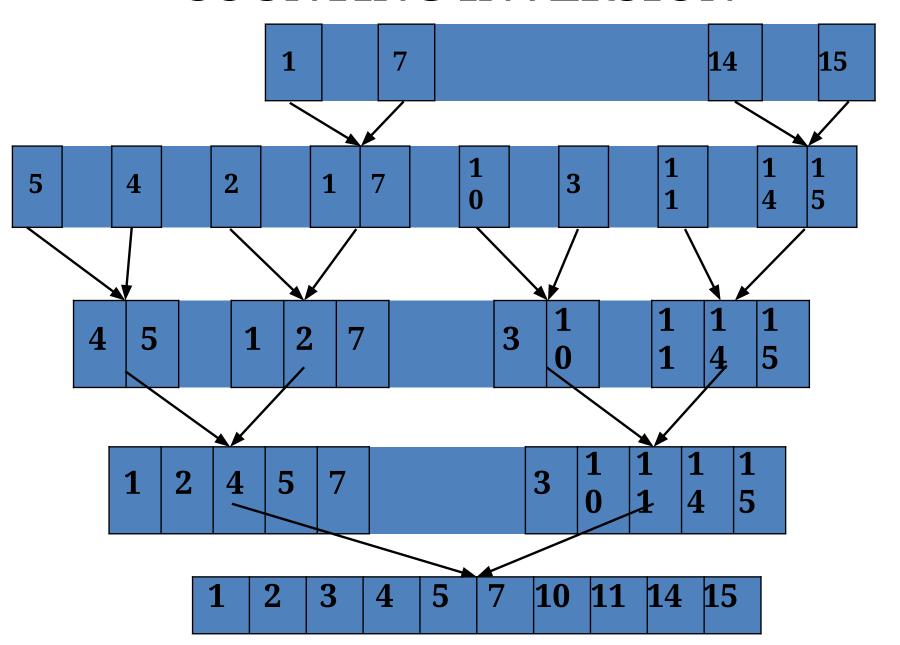
EndWhile

Once one list is empty, append the remainder of the other list to the output Return Count and the merged list

EXAMPLE

Total Inversions: 10





THANK YOU