

Free Theorems for Nested Types

ANONYMOUS AUTHOR(S)

1 FREE THEOREM FOR TYPE OF POLYMORPHIC BOTTOM

Suppose $\vdash g : \text{Nat}^\alpha \mathbb{1} \alpha$, let $G^{\text{Set}} = \llbracket \vdash g : \text{Nat}^\alpha \mathbb{1} \alpha \rrbracket^{\text{Set}}$, and let $G^{\text{Rel}} = \llbracket \vdash g : \text{Nat}^\alpha \mathbb{1} \alpha \rrbracket^{\text{Rel}}$. By Theorem ??, $(G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) = G^{\text{Rel}} \rho$. Thus, for all $\rho \in \text{RelEnv}$ and any $(a, b) \in \llbracket \phi, \alpha; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$, eliding the only possible instantiations of a and b gives that

$$\begin{aligned} (G^{\text{Set}}, G^{\text{Set}}) &= (G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) \in \llbracket \vdash \text{Nat}^\alpha \mathbb{1} \alpha \rrbracket^{\text{Rel}} \rho \\ &= \{\eta : K_1 \Rightarrow id\} \\ &= \{(\eta_1 : K_1 \Rightarrow id, \eta_2 : K_1 \Rightarrow id)\} \end{aligned}$$

That is, G^{Set} is a natural transformation from the constantly 1-valued functor to the identity functor in Set . In particular, for every $S : \text{Set}$, $G_S^{\text{Set}} : 1 \rightarrow S$. Note, however, that if $S = \emptyset$, then there can be no such morphism, so no such natural transformation can exist in Set , and thus no term $\vdash g : \text{Nat}^\alpha \mathbb{1} \alpha$ can exist in our calculus. That is, our calculus does not admit any terms with the closed type $\text{Nat}^\alpha \mathbb{1} \alpha$ of the polymorphic bottom.