

Free Theorems for Nested Types

ANONYMOUS AUTHOR(S)

1 FREE THEOREM FOR TYPE OF POLYMORPHIC IDENTITY

Suppose $\vdash g : \text{Nat}^\alpha \alpha \alpha$, let $G^{\text{Set}} = \llbracket \vdash g : \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Set}}$, and let $G^{\text{Rel}} = \llbracket \vdash g : \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Rel}}$. By Theorem ??, $(G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) = G^{\text{Rel}} \rho$. Thus, for all $\rho \in \text{RelEnv}$ and any $(a, b) \in \llbracket \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$, eliding the only possible instantiations of a and b gives that

$$\begin{aligned} (G^{\text{Set}}, G^{\text{Set}}) &= (G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) \in \llbracket \vdash \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Rel}} \rho \\ &= \{\eta : id \Rightarrow id\} \\ &= \{(\eta_1 : id \Rightarrow id, \eta_2 : id \Rightarrow id)\} \end{aligned}$$

That is, G^{Set} is a natural transformation from the identity functor on Set to itself.

Now let S be any set. If $S = \emptyset$, then there is exactly one morphism $id_S : S \rightarrow S$, so $G_S^{\text{Set}} : S \rightarrow S$ must be id_S . If $S \neq \emptyset$, then if a is any element of S and $K_a : S \rightarrow S$ is the constantly a -valued morphism on S , then instantiating the naturality square implied by the above equality gives that $G_S^{\text{Set}} \circ K_a = K_a \circ G_S^{\text{Set}}$, i.e., $G_S^{\text{Set}} a = a$, i.e., $G_S^{\text{Set}} = id_S$. Putting these two cases together we have that for every $S : \text{Set}$, $G_S^{\text{Set}} = id_S$, i.e., G^{Set} is the identity natural transformation for the identity functor on Set . So every closed term g of closed type $\text{Nat}^\alpha \alpha \alpha$ always denotes the identity natural transformation for the identity functor on Set , i.e., every closed term g of type $\text{Nat}^\alpha \alpha \alpha$ denotes the polymorphic identity function.