

Practical Parametricity for GADTs

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Abstract goes here

Maybe develop our theory for *any* $\lambda \geq \omega_1$, and then specialize to ω_1 when discussing GADTs? Can we do that? It seems we really use properties of ω CPO to get that interpretations of Nat-types are well-defined.

1 THE CALCULUS

1.1 Types

For each $k \geq 0$, we assume countable sets \mathbb{T}^k of *type constructor variables of arity k* and \mathbb{F}^k of *functorial variables of arity k* , all mutually disjoint. The sets of all type constructor variables and functorial variables are $\mathbb{T} = \bigcup_{k \geq 0} \mathbb{T}^k$ and $\mathbb{F} = \bigcup_{k \geq 0} \mathbb{F}^k$, respectively, and a *type variable* is any element of $\mathbb{T} \cup \mathbb{F}$. We use lower case Greek letters for type variables, writing ϕ^k to indicate that $\phi \in \mathbb{T}^k \cup \mathbb{F}^k$, and omitting the arity indicator k when convenient, unimportant, or clear from context. We reserve letters from the beginning of the alphabet to denote type variables of arity 0, i.e., elements of $\mathbb{T}^0 \cup \mathbb{F}^0$. We write $\bar{\zeta}$ for either a set $\{\zeta_1, \dots, \zeta_n\}$ of type constructor variables or a set of functorial variables when the cardinality n of the set is unimportant or clear from context. If P is a set of type variables we write $P, \bar{\phi}$ for $P \cup \bar{\phi}$ when $P \cap \bar{\phi} = \emptyset$. We omit the vector notation for a singleton set, thus writing ϕ , instead of $\bar{\phi}$, for $\{\phi\}$.

DEFINITION 1. Let V be a finite subset of \mathbb{T} , P be a finite subset of \mathbb{F} , $\bar{\alpha}$ be a finite subset of \mathbb{F}^0 disjoint from P , and $\phi^k \in \mathbb{F}^k \setminus P$. The set $\mathcal{F}^P(V)$ of functorial expressions over P and V are given by

$$\begin{aligned} \mathcal{F}^P(V) ::= & \quad 0 \mid 1 \mid \text{Nat}^P \mathcal{F}^P(V) \mathcal{F}^P(V) \mid \overline{P \mathcal{F}^P(V)} \mid V \overline{\mathcal{F}^P(V)} \mid \mathcal{F}^P(V) + \mathcal{F}^P(V) \\ & \mid \mathcal{F}^P(V) \times \mathcal{F}^P(V) \mid \left(\mu \phi^k . \lambda \alpha_1 \dots \alpha_k . \mathcal{F}^{P, \alpha_1, \dots, \alpha_k, \phi}(V) \right) \overline{\mathcal{F}^P(V)} \\ & \mid (\text{Lan}_{\mathcal{F}^{\bar{\alpha}}}^{\bar{\alpha}} \mathcal{F}^{P, \bar{\alpha}}) \overline{\mathcal{F}^P} \end{aligned}$$

A *type* over P and V is any element of $\mathcal{F}^P(V)$. The difference with [Johann et al. 2020] here lies solely in the incorporation of functorial expressions constructed from Lan.

The notation for types entails that an application $FF_1 \dots F_k$ is allowed only when F is a type variable of arity k , or F is a subexpression of the form $\mu \phi^k . \lambda \alpha_1 \dots \alpha_k . F'$ or $\text{Lan}_{\bar{K}}^{\bar{\alpha}} F'$. Moreover, if F has arity k then F must be applied to exactly k arguments. Accordingly, an overbar indicates a sequence of subexpressions whose length matches the arity of the type applied to it. The fact that types are always in *η -long normal form* avoids having to consider β -conversion at the level of types. In a subexpression $\text{Nat}^\Phi F G$, the Nat operator binds all occurrences of the variables in Φ in F and G . Note that, by contrast with [Johann et al. 2020], variables of arity greater than 0 are allowed in Φ ; this is necessary to construct well-typed terms of Lan types. In a subexpression $\mu \phi^k . \lambda \bar{\alpha} . F$, the μ operator binds all occurrences of the variable ϕ , and the λ operator binds all occurrences of the variables in $\bar{\alpha}$, in the body F . And in a subexpression $(\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A}$, the Lan operator binds all occurrences of the variables in $\bar{\alpha}$ in every element of \bar{K} , as well as in F .

A *type constructor context* is a finite set Γ of type constructor variables, and a *functorial context* is a finite set Φ of functorial variables. In Definition 2, a judgment of the form $\Gamma; \Phi \vdash F$ indicates that the type F is intended to be functorial in the variables in Φ but not necessarily in those in Γ .

DEFINITION 2. The formation rules for the set $\mathcal{F} \subseteq \bigcup_{V \subseteq \mathbb{T}, P \subseteq \mathbb{F}} \mathcal{F}^P(V)$ of well-formed types are

$$\begin{array}{c}
\frac{}{\Gamma; \Phi \vdash 0} \quad \frac{}{\Gamma; \Phi \vdash 1} \\
\\
\frac{\Gamma; \Phi \vdash F \quad \Gamma; \Phi \vdash G}{\Gamma; \Phi \vdash F + G} \quad \frac{\Gamma; \Phi \vdash F \quad \Gamma; \Phi \vdash G}{\Gamma; \Phi \vdash F \times G} \\
\\
\frac{\Gamma; \Phi \vdash F \quad \Gamma; \Phi \vdash G}{\Gamma; \emptyset \vdash \text{Nat}^\Phi F G} \\
\\
\frac{\phi^k \in \Gamma \cup \Phi \quad \overline{\Gamma; \Phi \vdash F}}{\Gamma; \Phi \vdash \phi^k \overline{F}} \\
\\
\frac{\Gamma; \overline{\gamma^0}, \overline{\alpha^0}, \phi^k \vdash F \quad \overline{\Gamma; \Phi, \overline{\gamma^0} \vdash G}}{\Gamma; \Phi, \overline{\gamma^0} \vdash (\mu \phi^k. \lambda \overline{\alpha^0}. F) \overline{G}} \\
\\
\frac{\Gamma; \Phi, \overline{\alpha^0} \vdash F \quad \overline{\Gamma; \overline{\alpha^0} \vdash K} \quad \overline{\Gamma; \Phi \vdash A}}{\Gamma; \Phi \vdash (\text{Lan}_{\overline{K}}^{\overline{\alpha^0}} F) \overline{A}}
\end{array}$$

In addition to textual replacement, we also have a proper substitution operation on types. If F is a type over P and V , if P and V contain only type variables of arity 0, and if $k = 0$ for every occurrence of ϕ^k bound by μ in F , then we say that F is *first-order*; otherwise we say that F is *second-order*. Substitution for first-order types is the usual capture-avoiding textual substitution. We write $F[\alpha := \sigma]$ for the result of substituting σ for α in F , and $F[\alpha_1 := F_1, \dots, \alpha_k := F_k]$, or $F[\overline{\alpha} := \overline{F}]$ when convenient, for $F[\alpha_1 := F_1][\alpha_2 := F_2, \dots, \alpha_k := F_k]$. Substitution for second-order types is defined below, where we adopt a similar notational convention for vectors of types. Note that it is not correct to substitute along non-functorial variables.

DEFINITION 3. If $\Gamma; \Phi, \phi^k \vdash H$ and if $\Gamma; \Phi, \overline{\alpha} \vdash F$ with $|\overline{\alpha}| = k$, then $\Gamma; \Phi \vdash H[\phi :=_{\overline{\alpha}} F]$. Similarly, if $\Gamma, \phi^k; \Phi \vdash H$, and if $\Gamma; \overline{\psi}, \overline{\alpha} \vdash F$ with $|\overline{\alpha}| = k$ and $\Phi \cap \overline{\psi} = \emptyset$, then $\Gamma, \overline{\psi}; \Phi \vdash H[\phi :=_{\overline{\alpha}} F[\overline{\psi} := \overline{\psi'}]]$. Here, the operation $(\cdot)[\phi :=_{\overline{\alpha}} F]$ of second-order type substitution along $\overline{\alpha}$ is defined by:

$$\begin{array}{ll}
0[\phi :=_{\overline{\alpha}} F] & = 0 \\
1[\phi :=_{\overline{\alpha}} F] & = 1 \\
(\text{Nat}^{\overline{\beta}} G K)[\phi :=_{\overline{\alpha}} F] & = \text{Nat}^{\overline{\beta}} (G[\phi :=_{\overline{\alpha}} F]) (K[\phi :=_{\overline{\alpha}} F]) \\
(\psi \overline{G})[\phi :=_{\overline{\alpha}} F] & = \begin{cases} \psi \overline{G[\phi :=_{\overline{\alpha}} F]} & \text{if } \psi \neq \phi \\ F[\alpha := G[\phi :=_{\overline{\alpha}} F]] & \text{if } \psi = \phi \end{cases} \\
(G + K)[\phi :=_{\overline{\alpha}} F] & = G[\phi :=_{\overline{\alpha}} F] + K[\phi :=_{\overline{\alpha}} F] \\
(G \times K)[\phi :=_{\overline{\alpha}} F] & = G[\phi :=_{\overline{\alpha}} F] \times K[\phi :=_{\overline{\alpha}} F] \\
((\mu \psi. \lambda \overline{\beta}. G) \overline{K})[\phi :=_{\overline{\alpha}} F] & = (\mu \psi. \lambda \overline{\beta}. G[\phi :=_{\overline{\alpha}} F]) \overline{K[\phi :=_{\overline{\alpha}} F]} \\
((\text{Lan}_{\overline{H}}^{\overline{\beta}} G) \overline{K})[\phi :=_{\overline{\alpha}} F] & = (\text{Lan}_{\overline{H}}^{\overline{\beta}} G[\phi :=_{\overline{\alpha}} F]) \overline{K[\phi :=_{\overline{\alpha}} F]}
\end{array}$$

We note that $(\cdot)[\phi^0 :=_{\emptyset} F]$ coincides with first-order substitution. We also omit $\overline{\alpha}$ when convenient.

1.2 Terms

We now define our term calculus. To do so we assume an infinite set \mathcal{V} of term variables disjoint from \mathbb{T} and \mathbb{F} . If Γ is a type constructor context and Φ is a functorial context, then a *term context* for Γ and Φ is a finite set of bindings of the form $x : F$, where $x \in \mathcal{V}$ and $\Gamma; \Phi \vdash F$. We adopt the same conventions for denoting disjoint unions and for vectors in term contexts as for type constructor contexts and functorial contexts.

DEFINITION 4. *Let Δ be a term context for Γ and Φ . The formation rules for the set of well-formed terms over Δ are*

$$\begin{array}{c}
\frac{\Gamma; \Phi \vdash F}{\Gamma; \Phi \mid \Delta, x : F \vdash x : F} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : \mathbb{0} \quad \Gamma; \Phi \vdash F}{\Gamma; \Phi \mid \Delta \vdash \perp_F t : F} \quad \frac{}{\Gamma; \Phi \mid \Delta \vdash \top : \mathbb{1}} \\
\\
\frac{\Gamma; \Phi \mid \Delta \vdash s : F}{\Gamma; \Phi \mid \Delta \vdash \text{inl } s : F + G} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : G}{\Gamma; \Phi \mid \Delta \vdash \text{inr } t : F + G} \\
\\
\frac{\Gamma; \Phi \vdash F, G \quad \Gamma; \Phi \mid \Delta \vdash t : F + G \quad \Gamma; \Phi \mid \Delta, x : F \vdash l : K \quad \Gamma; \Phi \mid \Delta, y : G \vdash r : K}{\Gamma; \Phi \mid \Delta \vdash \text{case } t \text{ of } \{x \mapsto l; y \mapsto r\} : K} \\
\\
\frac{\Gamma; \Phi \mid \Delta \vdash s : F \quad \Gamma; \Phi \mid \Delta \vdash t : G}{\Gamma; \Phi \mid \Delta \vdash (s, t) : F \times G} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : F \times G}{\Gamma; \Phi \mid \Delta \vdash \pi_1 t : F} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : F \times G}{\Gamma; \Phi \mid \Delta \vdash \pi_2 t : G} \\
\\
\frac{\Gamma; \Phi \vdash F \quad \Gamma; \Phi \vdash G \quad \Gamma; \Phi \mid \Delta, x : F \vdash t : G}{\Gamma; \emptyset \mid \Delta \vdash L_{\Phi} x. t : \text{Nat}^{\Phi} F G} \\
\\
\frac{\overline{\Gamma; \Phi, \bar{\beta} \vdash K} \quad \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\psi}} F G \quad \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\psi} :=_{\bar{\beta}} \bar{K}]}{\Gamma; \Phi \mid \Delta \vdash t_{\bar{K}} s : G[\bar{\psi} :=_{\bar{\beta}} \bar{K}]} \\
\\
\frac{\Gamma; \Phi, \bar{\phi} \vdash H \quad \overline{\Gamma; \Phi, \bar{\beta} \vdash F} \quad \overline{\Gamma; \Phi, \bar{\beta} \vdash G}}{\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{H}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\Phi, \bar{\beta}} F G) (\text{Nat}^{\Phi} H[\bar{\phi} :=_{\bar{\beta}} \bar{F}] H[\bar{\phi} :=_{\bar{\beta}} \bar{G}])} \\
\\
\frac{\Gamma; \Phi, \phi, \bar{\alpha} \vdash H}{\Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\Phi, \bar{\beta}} H[\phi :=_{\bar{\beta}} (\mu\phi. \lambda\bar{\alpha}. H)\bar{\beta}][\bar{\alpha} := \bar{\beta}]} (\mu\phi. \lambda\bar{\alpha}. H)\bar{\beta} \\
\\
\frac{\Gamma; \phi, \Phi, \bar{\alpha} \vdash H \quad \Gamma; \Phi, \bar{\beta} \vdash F}{\Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\Phi, \bar{\beta}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\Phi, \bar{\beta}} (\mu\phi. \lambda\bar{\alpha}. H)\bar{\beta} F)} \\
\\
\frac{\Gamma; \Phi, \bar{\alpha} \vdash F \quad \overline{\Gamma; \bar{\alpha} \vdash K} \quad \overline{\Gamma; \Phi \vdash A} \quad \Gamma; \Phi \mid \Delta \vdash t : F[\bar{\alpha} := \bar{A}]}{\Gamma; \Phi \mid \Delta \vdash \int_{\bar{K}, F}^{\bar{\alpha}} t : (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) K[\bar{\alpha} := \bar{A}]} \\
\\
\frac{\Gamma; \emptyset \mid \Delta \vdash \eta : \text{Nat}^{\Phi, \bar{\alpha}} F G[\bar{\beta} := \bar{K}] \quad \overline{\Gamma; \Phi \vdash A} \quad \Gamma; \Phi \mid \Delta \vdash t : (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A}}{\Gamma; \Phi \mid \Delta \vdash \partial_F^{G, \bar{K}} \eta t : G[\bar{\beta} := \bar{A}]}
\end{array}$$

In the rule for $L_{\bar{\alpha}}x.t$, the L operator binds all occurrences of the type variables in $\bar{\alpha}$ in the type of the term variable x and in the body t , as well as all occurrences of x in t . In the rule for $t_{\bar{K}}s$ there is one functorial expression in \bar{K} for every functorial variable in $\bar{\alpha}$. In the rule for $\text{map}_{\bar{H}}^{\bar{F}, \bar{G}}$ there is one functorial expression F and one functorial expression G for each functorial variable in $\bar{\phi}$. Moreover, for each ϕ^k in $\bar{\phi}$ the number of functorial variables in $\bar{\beta}$ in the judgments for its corresponding functorial expressions F and G is k . In the rules for in_H and fold_H^F , the functorial variables in $\bar{\beta}$ are fresh with respect to H , and there is one β for every α . (Recall from above that, in order for the types of in_H and fold_H^F to be well-formed, the length of α must equal the arity of ϕ .) In the rule for $\int_{\bar{K}, H} t$, there is one functorial expression A for every functorial variable in $\bar{\alpha}$, and in the rule for $\partial_F^{G, \bar{K}} \eta t$, there is one functorial expression A for every functorial expression in \bar{K} (and hence for every functorial variable in $\bar{\beta}$).

Substitution for terms is the obvious extension of the usual capture-avoiding textual substitution, and Definition 4 ensures that the expected weakening rules for well-formed terms hold.

Sum and product intro and elim rules should be annotated with constituent types for consistency?

We should have a computation rule along the lines of: If $\eta : \text{Nat}^{\bar{\alpha}F} G[\bar{\beta} := \bar{K}]$ then

$$\begin{aligned} & (\partial_F^{G, \bar{K}} \eta)_{\overline{K[\alpha := A]}} \circ (\int_{K, F})_{\bar{A}} \rightarrow \eta_{\bar{A}} \\ & : F[\bar{\alpha} := \bar{A}] \rightarrow G[\bar{\beta} := \overline{K[\alpha := A]}] \\ & = F[\bar{\alpha} := \bar{A}] \rightarrow G[\bar{\beta} := \bar{K}][\bar{\alpha} := \bar{A}] \end{aligned}$$

This will appear as a computational property of the term interpretations.

2 INTERPRETING TYPES

The fundamental idea underlying Reynolds' parametricity is to give each type $F(\alpha)$ with one free variable α both an *object interpretation* F_0 taking sets to sets and a *relational interpretation* F_1 taking relations $R : \text{Rel}(A, B)$ to relations $F_1(R) : \text{Rel}(F_0(A), F_0(B))$, and to interpret each term $t(\alpha, x) : F(\alpha)$ with one free term variable $x : G(\alpha)$ as a map t_0 associating to each set A a homomorphism $t_0(A) : G_0(A) \rightarrow F_0(A)$, and to each relation R a morphism $t_1(R) : G_1(R) \rightarrow F_1(R)$. These interpretations are to be given inductively on the structures of F and t in such a way that they imply two fundamental theorems. The first is an *Identity Extension Lemma*, which states that $F_1(\text{Eq}_A) = \text{Eq}_{F_0(A)}$, and is the essential property that makes a model relationally parametric rather than just induced by a logical relation. The second is an *Abstraction Theorem*, which states that, for any $R : \text{Rel}(A, B)$, $(t_0(A), t_0(B))$ is a morphism in Rel from $(G_0(A), G_0(B), G_1(R))$ to $(F_0(A), F_0(B), F_1(R))$. The Identity Extension Lemma is similar to the Abstraction Theorem except that it holds for *all* elements of a type's interpretation, not just those that are interpretations of terms. Similar theorems are expected to hold for types and terms with any number of free variables.

To accommodate GADTs, we will need to transition Reynolds' approach from a Set-based semantics to a semantics based on ω -complete partial orders. We denote the category of ω -complete partial orders (ω CPOs) and their sup-preserving morphisms by ωCPO . The underlying set of an ωCPO A is denoted $|A|$. The category ωCPOrel of relations on ωCPOs has as its objects triples (A, B, R) , where $A, B : \omega\text{CPO}$ and $R : \text{Rel}(|A|, |B|)$, and has as its morphisms from (A, B, R) to (A', B', R') pairs $(f : A \rightarrow A', g : B \rightarrow B')$ of morphisms in ωCPO such that $(f a, g b) \in R'$

whenever $(a, b) \in R$. We write $R : \omega\text{CPOrel}(A, B)$ in place of $(A, B, R) : \omega\text{CPOrel}$ when convenient. If $R : \omega\text{CPOrel}(A, B)$ then we write $\pi_1 R$ and $\pi_2 R$ for the *domain* A of R and the *codomain* B of R , respectively. If $A : \omega\text{CPO}$, then we write $\text{Eq}_A = (A, A, \text{Eq}_{|A|})$ for the *equality relation* on A .

To adapt Reynolds' approach, we first inductively define, for each type, an object interpretation in ωCPO and a relational interpretation in ωCPOrel . Next, we show that these interpretations satisfy both an Identity Extension Lemma (Theorem ??) and an Abstraction Theorem (Theorem ??) appropriate to the ωCPO setting. The key to proving our Identity Extension Lemma is a familiar "cutting down" of the interpretations of universally quantified types to include only the "parametric" elements; as in [Johann et al. 2020], the relevant types of the calculus defined above are the (now richer) Nat -types. The requisite cutting down requires that the object interpretations of our types in ωCPO are defined simultaneously with their relational interpretations in ωCPOrel . We give the object interpretations for our types in Section 2.1 and give their relational interpretations in Section 2.2. While the former are relatively straightforward, the latter are less so, mainly because of the cocontinuity conditions, adapted from the Set-based setting of [Johann et al. 2020], that must hold if they are to be well-defined. We develop these conditions in Section 2.2, which separates Definitions 8 and 18 in space, but otherwise has no impact on the fact that they are given by mutual induction.

2.1 Object Interpretations of Types

The object interpretations of the types in our calculus will be ω_1 -cocontinuous functors between categories of ω_1 -cocontinuous functors on categories constructed from the locally ω_1 -presentable category ωCPO . We therefore begin by recording some important facts about locally ω_1 -presentable categories and functors on them, and verifying the properties needed to interpret our syntax.

2.1.1 Preliminaries. Perhaps have as preliminaries to entire paper. Do everything for λCPOs Define these; investigate their properties.

A category is *small* if its collection of morphisms is a set. It is *locally small* if, for any two objects A and B , the collection of morphisms from A to B is a set. A *small (co)limit* in a category \mathcal{C} is a (co)limit of a diagram $F : \mathcal{A} \rightarrow \mathcal{C}$, where \mathcal{A} is a small category. A category \mathcal{C} is *(co)complete* if it has all small (co)limits.

A poset $\mathcal{D} = (D, \leq)$ is ω_1 -directed if every countable subset of D has a supremum. When \mathcal{D} is considered as a category, we write $d \in \mathcal{D}$ to indicate that d is an object of \mathcal{D} (i.e., $d \in D$). An ω_1 -directed *colimit* in a category \mathcal{C} is a colimit of a diagram $F : \mathcal{D} \rightarrow \mathcal{C}$, where \mathcal{D} is an ω_1 -directed poset. A category \mathcal{C} is ω_1 -cocomplete if it has all ω_1 -directed colimits; a *cocomplete* category is one that has all colimits.

If \mathcal{A} and \mathcal{C} are ω_1 -cocomplete, then the functor $F : \mathcal{A} \rightarrow \mathcal{C}$ is ω_1 -cocontinuous if it preserves ω_1 -directed colimits. We write $[\mathcal{A}, \mathcal{C}]_{\omega_1}$ for the category of ω_1 -cocontinuous functors from \mathcal{A} to \mathcal{C} , and $\mathcal{C}^{\mathcal{A}}$ for the category of *all* functors from \mathcal{A} to \mathcal{C} . Since (co)limits of functors are computed pointwise, $\mathcal{C}^{\mathcal{A}}$ has all (co)limits that \mathcal{C} has, and (co)limits of (co)continuous functors are again (co)continuous. It follows that $[\mathcal{A}, \mathcal{C}]_{\omega_1}$ is $(\omega_1 1)(\text{co})$ complete whenever \mathcal{C} is.

If \mathcal{A} is locally small, then an object A of \mathcal{A} is ω_1 -presentable if the functor $\text{Hom}_{\mathcal{A}}(A, -) : \mathcal{A} \rightarrow \text{Set}$ preserves ω_1 -directed colimits, i.e., if for every ω_1 -directed poset \mathcal{D} and every functor $F : \mathcal{D} \rightarrow \mathcal{C}$, there is a canonical isomorphism $\varinjlim_{d \in \mathcal{D}} \text{Hom}_{\mathcal{A}}(A, Fd) \simeq \text{Hom}_{\mathcal{A}}(A, \varinjlim_{d \in \mathcal{D}} Fd)$. A locally small category \mathcal{A} is ω_1 -accessible if it is ω_1 -cocomplete and has a set \mathcal{A}_0 of ω_1 -presentable objects such that every object is an ω_1 -directed colimit of objects in \mathcal{A}_0 ; a locally small category is *locally ω_1 -presentable* if it is ω_1 -accessible and cocomplete.

The category ωCPO is locally ω_1 -presentable (but not locally finitely presentable); see Examples 1.18(2) of [Adámek and Rosický 1994]. Its ω_1 -presentable objects are precisely the ωCPOs

cardinality less than ω_1 , i.e., the countable ω CPOs. In the next subsection we will interpret type variables in $\mathbb{T}^k \cup \mathbb{F}^k$ as elements of $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$; the following special cases of standard results (see, e.g., [Adámek and Rosický 1994]) will therefore be critical to deducing important properties of our object interpretations of types:

- PROPOSITION 5. (1) If C_1, \dots, C_n are locally ω_1 -presentable categories then so is $C_1 \times \dots \times C_n$. Moreover, the presentable objects of $C_1 \times \dots \times C_n$ are exactly the tuples of the form (P_1, \dots, P_n) , where, for each $i = 1, \dots, n$, the object P_i is presentable in C_i .
- (2) If \mathcal{A} is ω_1 -accessible and C is λ -cocomplete, then the category $[\mathcal{A}, C]_{\omega_1}$ is naturally equivalent to the category $C^{\mathcal{A}_0}$.
- (3) If C is locally ω_1 -presentable and \mathcal{A}_0 is essentially small, then $C^{\mathcal{A}_0}$ is locally ω_1 -presentable.

Together, the statements in Proposition 5 give that if \mathcal{A} and C are locally ω_1 -presentable, then $[\mathcal{A}, C]_{\omega_1}$ is naturally equivalent to $C^{\mathcal{A}_0}$, and thus is ω_1 -presentable. Thus, for all $k_1, \dots, k_n \in \mathbb{N}^n$, $[\mathcal{A}, C]_{\omega_1}^{k_1} \times \dots \times [\mathcal{A}, C]_{\omega_1}^{k_n}$ is locally ω_1 -presentable, and therefore $[[\mathcal{A}, C]_{\omega_1}^{k_1} \times \dots \times [\mathcal{A}, C]_{\omega_1}^{k_n}, C]_{\omega_1}$ is as well. Taking both \mathcal{A} and C to be ω CPO — as we will to ensure that the fixpoints interpreting μ -types in ω CPO exist — we have

PROPOSITION 6. For all $k_1, \dots, k_n \in \mathbb{N}^n$,

$$[[\omega\text{CPO}, \omega\text{CPO}]_{\omega_1}^{k_1} \times \dots \times [\omega\text{CPO}, \omega\text{CPO}]_{\omega_1}^{k_n}, \omega\text{CPO}]_{\omega_1}$$

is locally ω_1 -presentable.

2.1.2 *Object Interpretations.* To define the object interpretations of the types in Definition 2 we must first interpret their variables. We have:

DEFINITION 7. A ω CPO environment maps each type variable in $\mathbb{T}^k \cup \mathbb{F}^k$ to an element of $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$. A morphism $f : \rho \rightarrow \rho'$ for set environments ρ and ρ' with $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$ maps each type constructor variable $\psi^k \in \mathbb{T}$ to the identity natural transformation on $\rho\psi^k = \rho'\psi^k$ and each functorial variable $\phi^k \in \mathbb{F}$ to a natural transformation from the k -ary functor $\rho\phi^k$ on ω CPO to the k -ary functor $\rho'\phi^k$ on ω CPO. Composition of morphisms on set environments is given componentwise, with the identity morphism mapping each set environment to itself. This gives a category of set environments and morphisms between them, which we denote ωCPOEnv .

When convenient we identify a functor in $[\omega\text{CPO}^0, \omega\text{CPO}]_{\omega_1}$ with its value on $*$ and consider a ω CPO environment to map a type variable of arity 0 to an ω_1 -cocontinuous functor from ωCPO^0 to ωCPO , i.e., to an ω CPO. If $\Phi = \{\phi_1^{k_1}, \dots, \phi_n^{k_n}\}$ and $\bar{K} = \{K_1, \dots, K_n\}$ where $K_i : [\omega\text{CPO}^{k_i}, \omega\text{CPO}]_{\omega_1}$ for $i = 1, \dots, n$, then we write $\rho[\Phi := \bar{K}]$ for the ω CPO environment ρ' such that $\rho'\phi_i = K_i$ for $i = 1, \dots, n$ and $\rho'\phi = \rho\phi$ if $\phi \notin \Phi$. If ρ is an ω CPO environment, we write Eq_ρ for the ωCPOrel environment (see Definition 16) such that $\text{Eq}_\rho v = \text{Eq}_{\rho v}$ for every type variable v . The categories ωCPORT_k and relational interpretations appearing in the third clause of Definition 8 are given in full in Section ??.

Need to describe ω_1 -directed colimits in ωCPOEnv for the type of $[\![\cdot]\!]^{\omega\text{CPO}}$ to make sense. We do it pointwise.

DEFINITION 8. The object interpretation $\llbracket \cdot \rrbracket^{\omega\text{CPO}} : \mathcal{F} \rightarrow [\omega\text{CPOEnv}, \omega\text{CPO}]_{\omega_1}$ is defined by

$$\begin{aligned}
& \llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\omega\text{CPO}} \rho = 0 \\
& \llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\omega\text{CPO}} \rho = 1 \\
& \llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPO}} \rho = \{ \eta : \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}] \Rightarrow \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}] \\
& \quad \mid \forall K = (K_0, K_1, K^*) : \omega\text{CPORT}_K. \\
& \quad (\eta_{\bar{K}_0}, \eta_{\bar{K}_1}) : \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \text{Eq}_\rho [\bar{\Phi} := \bar{K}] \rightarrow \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \text{Eq}_\rho [\bar{\Phi} := \bar{K}] \\
& \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPO}} \rho = (\rho \phi) \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho \\
& \llbracket \Gamma; \Phi \vdash F + G \rrbracket^{\omega\text{CPO}} \rho = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho + \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} \rho \\
& \llbracket \Gamma; \Phi \vdash F \times G \rrbracket^{\omega\text{CPO}} \rho = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho \times \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} \rho \\
& \llbracket \Gamma; \Phi, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{G} \rrbracket^{\omega\text{CPO}} \rho = (\mu T_{H, \rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho \\
& \quad \text{where } T_{H, \rho}^{\omega\text{CPO}} F = \lambda \bar{A}. \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}} \rho [\phi := F] [\bar{\alpha} := \bar{A}] \\
& \quad \text{and } T_{H, \rho}^{\omega\text{CPO}} \eta = \lambda \bar{A}. \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}} \text{id}_\rho [\phi := \eta] [\bar{\alpha} := \text{id}_\rho \bar{A}] \\
& \llbracket \Gamma; \Phi \vdash (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A} \rrbracket^{\omega\text{CPO}} \rho = (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho [\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} \rho [\bar{\alpha} := _]) \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPO}} \rho
\end{aligned}$$

If $\rho \in \omega\text{CPOEnv}$ and $\vdash F$ then we write $\llbracket \vdash F \rrbracket^{\omega\text{CPO}}$ instead of $\llbracket \vdash F \rrbracket^{\omega\text{CPO}} \rho$ since the environment is immaterial. The third clause of Definition 8 does indeed define an ωCPO . First, we have

LEMMA 9. The collection of all natural transformations

$$\eta : \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}] \Rightarrow \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$$

defines a set.

PROOF. We first note that $\lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$ and $\lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$ are both in $[[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}, \omega\text{CPO}]_{\omega_1}$. By Proposition 6, $[[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}, \omega\text{CPO}]_{\omega_1}$ is locally ω_1 -presentable. It is therefore locally small, so there are only Set-many morphisms (i.e., natural transformations) between any two functors in $[[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}, \omega\text{CPO}]_{\omega_1}$. In particular, there are only Set-many natural transformations from $\lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$ to $\lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$. \square

Need to modify to take equality conditions into account. If C is any category, then given functors $F, G : C \rightarrow \omega\text{CPO}$ and natural transformations $\eta, \eta' : F \rightarrow G$, we define $\eta \leq \eta'$ iff $\eta_c \leq \eta'_c$ for all $c \in C$, i.e., iff $\eta_c x \leq \eta'_c x$ in Gc for all $c \in C$ and $x \in Fc$. If $(\eta_i)_{i < \omega}$ is a chain in $\{\eta : F \rightarrow G\}$, then the family of morphisms $(\bigvee_{i < \omega} \eta_i)_c = \lambda x. \bigvee_{i < \omega} (\eta_i)_c x : Fc \rightarrow Gc$ defines a natural transformation $\bigvee_{i < \omega} \eta_i : F \rightarrow G$, and this natural transformation is clearly the supremum of $(\eta_i)_{i < \omega}$ in $\{\eta : F \rightarrow G\}$. We therefore have that $\{\eta : F \rightarrow G\}$ is itself an ωCPO . Letting $F = \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$ and $G = \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\text{Set}} \rho [\bar{\Phi} := \bar{K}]$, we have that $\llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPO}} \rho$ is indeed an ωCPO . **Does this use functorial actions of interps? I don't think so, but double-check with Enrico. That each of the above interpretations is ω_1 -cocontinuous follows from Corollary 12 of [Johann and Polonsky 2019] if we APPROPRIATELY RESTRICT THE SUBSCRIPTS OF Lans. For Nat-types, we note that $\llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPO}}$ is an ω_1 -cocontinuous functor because it is constant on ω_1 -directed sets. Interpretations of Nat-types ensure that $\llbracket \Gamma \vdash F \rightarrow G \rrbracket^{\omega\text{CPO}}$ and $\llbracket \Gamma \vdash \forall \bar{\alpha}. F \rrbracket^{\omega\text{CPO}}$ are as expected in any parametric model.**

To make sense of the next-to-last clause in Definition 8, we need to know that, for each $\rho \in \omega\text{CPOEnv}$, $T_{H,\rho}^{\omega\text{CPO}}$ is an ω_1 -cocontinuous endofunctor on $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$, and thus admits a fixpoint. Since $T_{H,\rho}^{\omega\text{CPO}}$ is defined in terms of $\llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}}$, this means that interpretations of types must be such functors, which in turn means that the actions of set interpretations of types on objects and on morphisms in ωCPOEnv are intertwined. Fortunately, we know from [Johann and Polonsky 2019] that, for every $\Gamma; \bar{\alpha} \vdash F$, $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}}$ is actually in $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$, where $k = |\bar{\alpha}|$. Therefore, for each $\llbracket \Gamma; \bar{\gamma}, \phi^k, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}}$, the corresponding operator $T_{H,\rho}^{\omega\text{CPO}}$ can be extended to a functor from ωCPOEnv to $[[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}, [\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}]_{\omega_1}$. The action of $T_H^{\omega\text{CPO}}$ on an object $\rho \in \omega\text{CPOEnv}$ is given by the higher-order functor $T_{H,\rho}^{\omega\text{CPO}}$, whose actions on objects (functors in $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$) and on morphisms (natural transformations between such functors) are given in Definition 8. The action of $T_H^{\omega\text{CPO}}$ on a morphism $f : \rho \rightarrow \rho'$ is the higher-order natural transformation $T_{H,f}^{\omega\text{CPO}} : T_{H,\rho}^{\omega\text{CPO}} \rightarrow T_{H,\rho'}^{\omega\text{CPO}}$ whose action on $F : [\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$ is the natural transformation $T_{H,f}^{\omega\text{CPO}} F : T_{H,\rho}^{\omega\text{CPO}} F \rightarrow T_{H,\rho'}^{\omega\text{CPO}} F$ whose component at \bar{A} is $(T_{H,f}^{\omega\text{CPO}} F)_{\bar{A}} = \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}} f[\phi := id_F][\alpha := id_A]$.

In addition, for each \bar{K} , we have that $\text{Lan}_{\bar{K}}$ is itself a (higher-order) functor. Specifically, given functors $F, F' : \mathcal{C} \rightarrow \mathcal{D}$, a sequence of functors $\bar{K} = K_1, \dots, K_n$ with $K_i : \mathcal{C} \rightarrow \mathcal{C}_i$ for $i = 1, \dots, n$, and a natural transformation $\alpha : F \rightarrow G$, the functorial action $\text{Lan}_{\bar{K}}\alpha : \text{Lan}_{\bar{K}}F \rightarrow \text{Lan}_{\bar{K}}F'$ of $\text{Lan}_{\bar{K}}$ on α is defined to be the unique natural transformation such that $((\text{Lan}_{\bar{K}}\alpha) \circ \langle K_1, \dots, K_n \rangle) \circ \eta_F = \eta_{F'} \circ \alpha$. Here, $\eta_F : F \rightarrow (\text{Lan}_{\bar{K}}F) \circ \langle K_1, \dots, K_n \rangle$ and $\eta_{F'} : F' \rightarrow (\text{Lan}_{\bar{K}}F') \circ \langle K_1, \dots, K_n \rangle$ are the natural transformations associated with the functors $\text{Lan}_{\bar{K}}F$ and $\text{Lan}_{\bar{K}}F'$ from $\prod_{i \in \{1, \dots, n\}} \mathcal{C}_i$ to \mathcal{D} , respectively. It is not hard to see that $\text{Lan}_{\bar{K}}$ is a (higher-order) functor under this definition.

The next definition uses the functors $T_H^{\omega\text{CPO}}$ and $\text{Lan}_{\bar{K}}$ to define the actions of functors interpreting types on morphisms between set environments.

DEFINITION 10. Let $f : \rho \rightarrow \rho'$ be a morphism between ωCPO environments ρ and ρ' (so that $\rho|_{\top} = \rho'|_{\top}$). The action $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f$ of $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}}$ on the morphism f is given as follows:

- If $\Gamma; \Phi \vdash \perp$ then $\llbracket \Gamma; \Phi \vdash \perp \rrbracket^{\omega\text{CPO}} f = id_{\perp}$
- If $\Gamma; \Phi \vdash \perp$ then $\llbracket \Gamma; \Phi \vdash \perp \rrbracket^{\omega\text{CPO}} f = id_{\perp}$
- If $\Gamma; \emptyset \vdash \text{Nat}^{\Phi} F G$ then $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\Phi} F G \rrbracket^{\omega\text{CPO}} f = id_{\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\Phi} F G \rrbracket^{\omega\text{CPO}} \rho}$
- If $\Gamma; \Phi \vdash \phi \bar{F}$ then

$$\begin{aligned} \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPO}} f : \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPO}} \rho &\rightarrow \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPO}} \rho' \\ &= (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho} \rightarrow (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho'} \end{aligned}$$

is defined by

$$\begin{aligned} \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPO}} f &= (f\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho'} \circ (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f} \\ &= (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f} \circ (f\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho} \end{aligned}$$

The latter equality holds because $\rho\phi$ and $\rho'\phi$ are functors and $f\phi : \rho\phi \rightarrow \rho'\phi$ is a natural transformation, so the following naturality square commutes:

$$\begin{array}{ccc} (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho} & \xrightarrow{(f\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho}} & (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho} \\ (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f} \downarrow & & (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f} \downarrow \\ (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho'} & \xrightarrow{(f\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho'}} & (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho'} \end{array} \quad (1)$$

- If $\Gamma; \Phi \vdash F + G$ then $\llbracket \Gamma; \Phi \vdash F + G \rrbracket^{\omega\text{CPO}} f$ is defined by

$$\llbracket \Gamma; \Phi \vdash F + G \rrbracket^{\omega\text{CPO}} f(\text{inl } x) = \text{inl } (\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f x)$$

$$\llbracket \Gamma; \Phi \vdash F + G \rrbracket^{\omega\text{CPO}} f(\text{inr } y) = \text{inr } (\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} f y)$$

- If $\Gamma; \Phi \vdash F \times G$ then $\llbracket \Gamma; \Phi \vdash F \times G \rrbracket^{\omega\text{CPO}} f = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} f \times \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} f$
- If $\Gamma; \Phi, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{G}$ then

$$\begin{aligned} \llbracket \Gamma; \Phi, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{G} \rrbracket^{\omega\text{CPO}} f & : \llbracket \Gamma; \Phi, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{G} \rrbracket^{\omega\text{CPO}} \rho \rightarrow \llbracket \Gamma; \Phi, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{G} \rrbracket^{\omega\text{CPO}} \rho' \\ & = (\mu T_{H,\rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho \rightarrow (\mu T_{H,\rho'}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho' \end{aligned}$$

is defined by

$$\begin{aligned} & (\mu T_{H,f}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho' \circ (\mu T_{H,\rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} f \\ & = (\mu T_{H,\rho'}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} f \circ (\mu T_{H,f}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho \end{aligned}$$

The latter equality holds because $\mu T_{H,\rho}^{\omega\text{CPO}}$ and $\mu T_{H,\rho'}^{\omega\text{CPO}}$ are functors and $\mu T_{H,\rho}^{\omega\text{CPO}} : \mu T_{H,\rho'}^{\omega\text{CPO}} \rightarrow \mu T_{H,\rho}^{\omega\text{CPO}}$ is a natural transformation, so the following naturality square commutes:

$$\begin{array}{ccc} (\mu T_{H,\rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho & \xrightarrow{(\mu T_{H,f}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho} & (\mu T_{H,\rho'}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho \\ (\mu T_{H,\rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} f \downarrow & & (\mu T_{H,\rho'}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} f \downarrow \\ (\mu T_{H,\rho}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho' & \xrightarrow{(\mu T_{H,f}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho'} & (\mu T_{H,\rho'}^{\omega\text{CPO}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPO}} \rho' \end{array} \quad (2)$$

- If $\Gamma; \Phi \vdash (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A}$ then

$$\llbracket \Gamma; \Phi \vdash (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A} \rrbracket^{\omega\text{CPO}} f : \llbracket \Gamma; \Phi \vdash (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A} \rrbracket^{\omega\text{CPO}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash (\text{Lan}_{\bar{K}}^{\bar{\alpha}} F) \bar{A} \rrbracket^{\omega\text{CPO}} \rho'$$

is defined by

$$\begin{aligned} & (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} f[\bar{\alpha} := \text{id}__] \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPO}} \rho' \\ & \circ (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _] \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPO}} f \\ & = (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} \rho'[\bar{\alpha} := _] \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPO}} f \\ & \circ (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} f[\bar{\alpha} := \text{id}__] \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPO}} \rho \end{aligned}$$

where the above equality holds by naturality of $\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPO}} \rho[\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPO}} f[\bar{\alpha} := \text{id}__]$.

Definitions 8 and 10 respect weakening, i.e., ensure that a type and its weakenings have the same set interpretations.

2.2 Relational Interpretations of Types

DEFINITION 11. A k -ary ωCPO relation transformer F is a triple (F^1, F^2, F^*) , where $F^1, F^2 : [\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$ and $F^* : [\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$ are such that if $R_1 : \omega\text{CPOrel}(A_1, B_1), \dots, R_k : \omega\text{CPOrel}(A_k, B_k)$ then $F^* \bar{R} : \omega\text{CPOrel}(F^1 \bar{A}, F^2 \bar{B})$, and if $(\alpha_1, \beta_1) \in \text{Hom}_{\omega\text{CPOrel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\omega\text{CPOrel}}(R_k, S_k)$ then $F^*(\alpha, \beta) = (F^1 \bar{\alpha}, F^2 \bar{\beta})$. We define $F \bar{R}$ to be $F^* \bar{R}$ and $F(\alpha, \beta)$ to be $F^*(\alpha, \beta)$.

The last condition of the first sentence of Definition 11 expands to: if $(a, b) \in R$ implies $(\alpha a, \beta b) \in S$ then $(c, d) \in F^* \bar{R}$ implies $(F^1 \bar{\alpha} c, F^2 \bar{\beta} d) \in F^* \bar{S}$. When convenient we identify a 0-ary ωCPO relation transformer (A, B, R) with $R : \omega\text{CPOrel}(A, B)$. We may also write $\pi_1 F$ for F^1 and $\pi_2 F$ for F^2 . We

extend these conventions to ω CPO relation environments, introduced in Definition 16 below, in the obvious way.

DEFINITION 12. The category ωCPORT_k of k -ary ω CPO relation transformers is given by the following data:

- An object of ωCPORT_k is a k -ary ω CPO relation transformer.
- A morphism $\delta : (G^1, G^2, G^*) \rightarrow (H^1, H^2, H^*)$ in ωCPORT_k is a pair of natural transformations (δ^1, δ^2) , where $\delta^1 : G^1 \rightarrow H^1$ and $\delta^2 : G^2 \rightarrow H^2$ are such that, for all $R : \omega\text{CPOrel}(A, B)$, if $(x, y) \in G^* \bar{R}$ then $(\delta_A^1 x, \delta_B^2 y) \in H^* \bar{R}$.
- Identity morphisms and composition are inherited from the category of functors on ωCPO .

DEFINITION 13. An endofunctor H on ωCPORT_k is a triple $H = (H^1, H^2, H^*)$, where

- H^1 and H^2 are functors from $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$ to $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$
- H^* is a functor from ωCPORT_k to $[\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$
- for all $R : \omega\text{CPOrel}(A, B)$, $\pi_1((H^*(\delta^1, \delta^2))_{\bar{R}}) = (H^1 \delta_A^1)_{\bar{A}}$ and $\pi_2((H^*(\delta^1, \delta^2))_{\bar{R}}) = (H^2 \delta_B^2)_{\bar{B}}$
- The action of H on objects is given by $H(F^1, F^2, F^*) = (H^1 F^1, H^2 F^2, H^*(F^1, F^2, F^*))$
- The action of H on morphisms is given by $H(\delta^1, \delta^2) = (H^1 \delta^1, H^2 \delta^2)$ for $(\delta^1, \delta^2) : (F^1, F^2, F^*) \rightarrow (G^1, G^2, G^*)$

Since the results of applying an endofunctor H to k -ary ω CPO relation transformers and morphisms between them must again be k -ary ω CPO relation transformers and morphisms between them, respectively, Definition 13 implicitly requires that the following three conditions hold: i) if $R_1 : \omega\text{CPOrel}(A_1, B_1), \dots, R_k : \omega\text{CPOrel}(A_k, B_k)$, then $H^*(F^1, F^2, F^*) \bar{R} : \omega\text{CPOrel}(H^1 F^1 \bar{A}, H^2 F^2 \bar{B})$; ii) if $(\alpha_1, \beta_1) \in \text{Hom}_{\omega\text{CPOrel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\omega\text{CPOrel}}(R_k, S_k)$, then $H^*(F^1, F^2, F^*) (\bar{\alpha}, \bar{\beta}) = (H^1 F^1 \bar{\alpha}, H^2 F^2 \bar{\beta})$; and iii) if $(\delta^1, \delta^2) : (F^1, F^2, F^*) \rightarrow (G^1, G^2, G^*)$ and $R_1 : \omega\text{CPOrel}(A_1, B_1), \dots, R_k : \omega\text{CPOrel}(A_k, B_k)$, then $((H^1 \delta_A^1)_{\bar{A}} x, (H^2 \delta_B^2)_{\bar{B}} y) \in H^*(G^1, G^2, G^*) \bar{R}$ whenever $(x, y) \in H^*(F^1, F^2, F^*) \bar{R}$. Note, however, that this last condition is automatically satisfied because it is implied by the third bullet point of Definition 13.

DEFINITION 14. If H and K are endofunctors on ωCPORT_k , then a natural transformation $\sigma : H \rightarrow K$ is a pair $\sigma = (\sigma^1, \sigma^2)$, where $\sigma^1 : H^1 \rightarrow K^1$ and $\sigma^2 : H^2 \rightarrow K^2$ are natural transformations between endofunctors on $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$ and the component of σ at $F \in \omega\text{CPORT}_k$ is given by $\sigma_F = (\sigma_{F^1}^1, \sigma_{F^2}^2)$.

Definition 14 entails that $\sigma_{F^i}^i$ must be natural in $F^i : [\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$, and, for every F , $(\sigma_{F^1}^1)_{\bar{A}}$ and $(\sigma_{F^2}^2)_{\bar{B}}$ must be natural in \bar{A} and \bar{B} , respectively. Moreover, since the results of applying σ to k -ary ω CPO relation transformers must be morphisms of k -ary relation transformers, Definition 14 implicitly requires that $(\sigma_F)_{\bar{R}} = ((\sigma_{F^1}^1)_{\bar{A}}, (\sigma_{F^2}^2)_{\bar{B}})$ is a morphism in ωCPOrel for any k -tuple of relations $R : \text{Rel}(A, B)$, i.e., that if $(x, y) \in H^* \bar{F} \bar{R}$, then $((\sigma_{F^1}^1)_{\bar{A}} x, (\sigma_{F^2}^2)_{\bar{B}} y) \in K^* \bar{F} \bar{R}$.

Critically, we can compute ω_1 -directed colimits in ωCPORT_k : it is not hard to see that if \mathcal{D} is an ω_1 -directed set then $\lim_{d \in \mathcal{D}} (F_d^1, F_d^2, F_d^*) = (\lim_{d \in \mathcal{D}} F_d^1, \lim_{d \in \mathcal{D}} F_d^2, \lim_{d \in \mathcal{D}} F_d^*)$. We define an endofunctor $T = (T^1, T^2, T^*)$ on ωCPORT_k to be ω_1 -cocontinuous if T^1 and T^2 are ω_1 -cocontinuous endofunctors on $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$ and T^* is an ω_1 -cocontinuous functor from ωCPORT_k to $[\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$, i.e., is in $[\omega\text{CPORT}_k, [\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}]_{\omega_1}$.

Now, for any k , any $A : \omega\text{CPO}$, and any $R : \omega\text{CPOrel}(A, B)$, let $K_A^{\omega\text{CPO}}$ be the constantly A -valued functor from ωCPO^k to ωCPO and $K_R^{\omega\text{CPOrel}}$ be the constantly R -valued functor from ωCPOrel^k to ωCPOrel . Also let 0 denote either the initial object of ωCPO or the initial object of ωCPOrel , as appropriate. Observing that, for every k , $K_0^{\omega\text{CPO}}$ is initial in $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$, and $K_0^{\omega\text{CPOrel}}$

is initial in $[\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$, we have that, for each k , $K_0 = (K_0^{\omega\text{CPO}}, K_0^{\omega\text{CPO}}, K_0^{\omega\text{CPOrel}})$ is initial in ωCPORT_k . Thus, if $T = (T^1, T^2, T^*) : \omega\text{CPORT}_k \rightarrow \omega\text{CPORT}_k$ is an endofunctor on ωCPORT_k then we can define the ωCPO relation transformer μT to be $\lim_{\rightarrow n \in \mathbb{N}} T^n K_0$. It is not hard to see that μT is given explicitly as

$$\mu T = (\mu T^1, \mu T^2, \lim_{\rightarrow n \in \mathbb{N}} (T^n K_0)^*) \quad (3)$$

and that, as our notation suggests, it really is a fixpoint for T if T is ω_1 -cocontinuous:

LEMMA 15. *For any $T : [\omega\text{CPORT}_k, \omega\text{CPORT}_k]_{\omega_1}$, $\mu T \cong T(\mu T)$.*

The isomorphism is given by the morphisms $(in_1, in_2) : T(\mu T) \rightarrow \mu T$ and $(in_1^{-1}, in_2^{-1}) : \mu T \rightarrow T(\mu T)$ in ωCPORT_k . The latter is always a morphism in ωCPORT_k , but the former need not be if T is not ω_1 -cocontinuous.

It is worth noting that the third component in Equation (3) is the colimit in $[\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$ of third components of ωCPO relation transformers, rather than a fixpoint of an endofunctor on $[\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}$. That there is an asymmetry between the first two components of μT and its third reflects the important conceptual observation that the third component of an endofunctor on ωCPORT_k need not be a functor on all of $[\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$. In particular, although we can define $T_{H,\rho} F$ for an ωCPO relation transformer F in Definition 18 below, it is not clear how we could define it for an arbitrary $F : [\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}$.

DEFINITION 16. *An ωCPO relation environment maps each type variable in $\mathbb{T}^k \cup \mathbb{F}^k$ to a k -ary ωCPO relation transformer. A morphism $f : \rho \rightarrow \rho'$ between ωCPO relation environments ρ and ρ' with $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$ maps each type constructor variable $\psi^k \in \mathbb{T}$ to the identity morphism on $\rho\psi^k = \rho'\psi^k$ and each functorial variable $\phi^k \in \mathbb{F}$ to a morphism from the k -ary ωCPO relation transformer $\rho\phi$ to the k -ary ωCPO relation transformer $\rho'\phi$. Composition of morphisms on ωCPO relation environments is given componentwise, with the identity morphism mapping each ωCPO relation environment to itself. This gives a category of ωCPO relation environments and morphisms between them, which we denote $\omega\text{CPOrelEnv}$.*

When convenient we identify a 0-ary ωCPO relation transformer with the ωCPO relation (transformer) that is its codomain and consider an ωCPO relation environment to map a type variable of arity 0 to an ωCPO relation. We write $\rho[\alpha := \bar{R}]$ for the ωCPO relation environment ρ' such that $\rho'\alpha_i = R_i$ for $i = 1, \dots, k$ and $\rho'\alpha = \rho\alpha$ if $\alpha \notin \{\alpha_1, \dots, \alpha_k\}$. If ρ is an ωCPO relation environment, we write $\pi_1\rho$ and $\pi_2\rho$ for the ωCPO environments mapping each type variable ϕ to the functors $(\rho\phi)^1$ and $(\rho\phi)^2$, respectively.

We define, for each k , the notion of an ω_1 -cocontinuous functor from $\omega\text{CPOrelEnv}$ to ωCPORT_k :

DEFINITION 17. *A functor $H : [\omega\text{CPOrelEnv}, \omega\text{CPORT}_k]_{\omega_1}$ is a triple $H = (H^1, H^2, H^*)$, where*

- H^1 and H^2 are objects in $[\omega\text{CPOEnv}, [\omega\text{CPO}^k, \omega\text{CPO}]_{\omega_1}]_{\omega_1}$
- H^* is an object in $[\omega\text{CPOrelEnv}, [\omega\text{CPOrel}^k, \omega\text{CPOrel}]_{\omega_1}]_{\omega_1}$
- for all $\bar{R} : \omega\text{CPOrel}(A, B)$ and morphisms f in $\omega\text{CPOrelEnv}$, $\pi_1(H^* f \bar{R}) = H^1(\pi_1 f) \bar{A}$ and $\pi_2(H^* f \bar{R}) = H^2(\pi_2 f) \bar{B}$
- The action of H on ρ in $\omega\text{CPOrelEnv}$ is given by $H\rho = (H^1(\pi_1\rho), H^2(\pi_2\rho), H^*\rho)$
- The action of H on morphisms $f : \rho \rightarrow \rho'$ in $\omega\text{CPOrelEnv}$ is given by $Hf = (H^1(\pi_1 f), H^2(\pi_2 f))$

Spelling out the last two bullet points above gives the following analogues of the three conditions immediately following Definition 13: i) if $R_1 : \omega\text{CPOrel}(A_1, B_1), \dots, R_k : \omega\text{CPOrel}(A_k, B_k)$, then $H^*\rho \bar{R} : \omega\text{CPOrel}(H^1(\pi_1\rho) \bar{A}, H^2(\pi_2\rho) \bar{B})$; ii) if $(\alpha_1, \beta_1) \in \text{Hom}_{\omega\text{CPOrel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\omega\text{CPOrel}}(R_k, S_k)$, then $H^*\rho(\alpha, \beta) = (H^1(\pi_1\rho) \bar{\alpha}, H^2(\pi_2\rho) \bar{\beta})$; and iii) if $f : \rho \rightarrow \rho'$ and $R_1 :$

$\omega\text{CPORel}(A_1, B_1), \dots, R_k : \omega\text{CPORel}(A_k, B_k)$, then $(H^1(\pi_1 f) \bar{A}x, H^2(\pi_2 f) \bar{B}y) \in H^* \rho' \bar{R}$ whenever $(x, y) \in H^* \rho \bar{R}$. As before, the last condition is automatically satisfied because it is implied by the third bullet point of Definition 17.

Considering $\omega\text{CPORelEnv}$ as a product $\prod_{\phi^k \in \mathbb{T} \cup \mathbb{F}} \omega\text{CPORT}_k$, we extend the computation of ω_1 -directed colimits in ωCPORT_k to compute colimits in $\omega\text{CPORelEnv}$ componentwise. We similarly extend the notion of an ω_1 -cocontinuous endofunctor on ωCPORT_k componentwise to give a notion of ω_1 -cocontinuity for functors from $\omega\text{CPORelEnv}$ to ωCPORT_k . Recalling from the start of this subsection that Definition 18 is given mutually inductively with Definition 8 we can, at last, define:

DEFINITION 18. The relational interpretation $\llbracket \cdot \rrbracket^{\omega\text{CPORel}} : \mathcal{F} \rightarrow [\omega\text{CPORelEnv}, \omega\text{CPORel}]_{\omega_1}$ is defined by

$$\begin{aligned}
 & \llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\omega\text{CPORel}} \rho = 0 \\
 & \llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\omega\text{CPORel}} \rho = 1 \\
 & \llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPORel}} \rho = \{ \eta : \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}] \Rightarrow \lambda \bar{K}. \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}] \} \\
 & \quad = \{ (t, t') \in \llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPO}}(\pi_1 \rho) \times \llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPO}}(\pi_2 \rho) \mid \\
 & \quad \quad \forall K = (K^0, K^1, K^*) : \omega\text{CPORT}_k. \\
 & \quad \quad (t_{\bar{K}^0}, t'_{\bar{K}^1}) \in (\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}])^{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}]} \} \\
 & \llbracket \Gamma; \Phi \vdash \phi \bar{F} \rrbracket^{\omega\text{CPORel}} \rho = (\rho \phi) \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho \\
 & \llbracket \Gamma; \Phi \vdash F + G \rrbracket^{\omega\text{CPORel}} \rho = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho + \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho \\
 & \llbracket \Gamma; \Phi \vdash F \times G \rrbracket^{\omega\text{CPORel}} \rho = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho \times \llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho \\
 & \llbracket \Gamma; \Phi, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{G} \rrbracket^{\omega\text{CPORel}} \rho = (\mu_{T_{H, \rho}}) \llbracket \Gamma; \Phi, \bar{\gamma} \vdash G \rrbracket^{\omega\text{CPORel}} \rho \\
 & \quad \text{where } T_{H, \rho} = (T_{H, \pi_1 \rho}^{\text{Set}}, T_{H, \pi_2 \rho}^{\text{Set}}, T_{H, \rho}^{\text{Rel}}) \\
 & \text{and } T_{H, \rho}^{\omega\text{CPORel}} F = \lambda \bar{R}. \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPORel}} \rho [\phi := F] [\bar{\alpha} := \bar{R}] \\
 & \text{and } T_{H, \rho}^{\omega\text{CPORel}} \delta = \lambda \bar{R}. \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPORel}} \rho [\phi := \delta] [\bar{\alpha} := \bar{id}_{\bar{R}}] \\
 & \llbracket \Gamma; \Phi \vdash (\text{Lan}_{\bar{K}} \bar{\alpha} F) \bar{A} \rrbracket^{\omega\text{CPORel}} \rho = (\text{Lan}_{\llbracket \Gamma; \bar{\alpha} \vdash K \rrbracket^{\omega\text{CPORel}} \rho [\bar{\alpha} := _]} \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPORel}} \rho [\bar{\alpha} := _]) \llbracket \Gamma; \Phi \vdash A \rrbracket^{\omega\text{CPORel}} \rho
 \end{aligned}$$

The interpretations in Definitions 18 and 19 below respect weakening, and also ensure that $\llbracket \Gamma \vdash F \rightarrow G \rrbracket^{\omega\text{CPORel}}$ and $\llbracket \Gamma \vdash \forall \bar{\alpha}. F \rrbracket^{\omega\text{CPORel}}$ are as expected in any parametric model. Observe that $\llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPORel}} \rho$ is in ωCPORel . Same comments here as above. It is indeed a relation between ωCPOs , and it preserves sups of ω -chains of natural transformations because their sups are computed pointwise. Concretely, if $(t_i, t'_i) : \llbracket \Gamma; \emptyset \vdash \text{Nat}^\Phi F G \rrbracket^{\omega\text{CPORel}} \rho$ for all $i < \omega$, then for all $K = (K^0, K^1, K^*) : \omega\text{CPORT}_k$ we have that

$$((t_i)_{\bar{K}^0}, (t'_i)_{\bar{K}^1}) \in (\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}])^{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}]}$$

for every $i < \omega$. In addition, $((t_i)_{\bar{K}^0})_{i < \omega}$ is an ω -chain in $(\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}^0])^{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}^0]}$ and $((t'_i)_{\bar{K}^1})_{i < \omega}$ is an ω -chain in $(\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}^1])^{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPO}} \rho [\bar{\Phi} := \bar{K}^1]}$, so their sups are not only well-defined but must also be related in $(\llbracket \Gamma; \Phi \vdash G \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}])^{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} \rho [\bar{\Phi} := \bar{K}]}$.

This gives that

$$\left(\bigvee_{i < \omega} t_i \right)_{\overline{K^0}}, \left(\bigvee_{i < \omega} t'_i \right)_{\overline{K^1}} = \left(\bigvee_{i < \omega} (t_i)_{\overline{K^0}}, \bigvee_{i < \omega} (t'_i)_{\overline{K^1}} \right) \in ([\Gamma; \Phi \vdash G]^{\omega\text{CPORel}} \rho[\Phi := \overline{K}])^{\omega\text{CPORel}} [\Gamma; \Phi \vdash F]^{\omega\text{CPORel}} \rho[\Phi := \overline{K}]$$

Moreover, as for ωCPO interpretations, ω_1 -cocontinuity of each of the above interpretations follows from Corollary 12 of [Johann and Polonsky 2019] if we APPROPRIATELY RESTRICT THE SUBSCRIPTS of the Lans. For Nat-types, $[\Gamma; \emptyset \vdash \text{Nat}^\Phi F G]^{\omega\text{CPORel}}$ is an ω_1 -cocontinuous functor because it is constant on ω_1 -directed sets. If $\rho \in \omega\text{CPORelEnv}$ and $\vdash F$, then we write $\llbracket \vdash F \rrbracket^{\omega\text{CPORel}}$ instead of $\llbracket \vdash F \rrbracket^{\omega\text{CPORel}} \rho$. For the second-to-last clause in Definition 18 to be well-defined we need $T_{H,\rho}$ to be an ω_1 -cocontinuous endofunctor on ωCPORT so that, by Lemma 15, it admits a fixpoint. Since $T_{H,\rho}$ is defined in terms of $\llbracket \Gamma; \bar{\gamma}, \phi^k, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPORel}}$, this means that relational interpretations of types must be ω_1 -cocontinuous functors from $\omega\text{CPORelEnv}$ to ωCPORT_0 , which in turn entails that the actions of relational interpretations of types on objects and on morphisms in $\omega\text{CPORelEnv}$ are intertwined. As for ωCPO interpretations, we know from [Johann and Polonsky 2019] that, for every $\Gamma; \bar{\alpha} \vdash F$, $\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\omega\text{CPORel}}$ is actually in $[\omega\text{CPORel}^k, \omega\text{CPORel}]_{\omega_1}$, where $k = |\bar{\alpha}|$. We first define the actions of each of these functors on morphisms between ωCPO relation environments in Definition 19, and then argue that the functors given by Definitions 18 and 19 are well-defined and have the required properties. To do this, we extend T_H to a functor from $\omega\text{CPORelEnv}$ to $[[\omega\text{CPORel}^k, \omega\text{CPORel}]_{\omega_1}, [\omega\text{CPORel}^k, \omega\text{CPORel}]_{\omega_1}]_{\omega_1}$. Its action on an object $\rho \in \omega\text{CPORelEnv}$ is given by the higher-order functor $T_{H,\rho}$ whose actions on objects and morphisms are given in Definition 19. Its action on a morphism $f : \rho \rightarrow \rho'$ is the higher-order natural transformation $T_{H,f} : T_{H,\rho} \rightarrow T_{H,\rho'}$ whose action on any $F : [\omega\text{CPORel}^k, \omega\text{CPORel}]_{\omega_1}$ is the higher-order natural transformation $T_{H,f} F : T_{H,\rho} F \rightarrow T_{H,\rho'} F$ whose component at \bar{R} is $(T_{H,f} F)_{\bar{R}} = \llbracket \Gamma; \bar{\gamma}, \phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPORel}} f[\phi := id_F][\bar{\alpha} := id_R]$. The next definition uses T_H to define the actions of functors interpreting types on morphisms between ωCPO relation environments.

DEFINITION 19. Let $f : \rho \rightarrow \rho'$ for ωCPO relation environments ρ and ρ' (so that $\rho|_{\top} = \rho'|_{\top}$). The action $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}} f$ of $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\omega\text{CPORel}}$ on the morphism f is given exactly as in Definition 10, except that all interpretations are ωCPO relational interpretations and all occurrences of $T_{H,f}^{\omega\text{CPO}}$ are replaced by $T_{H,f}$.

To see that the functors given by Definitions 18 and 19 are well-defined we must show that, for every H , $T_{H,\rho} F$ is an ωCPO relation transformer for any ωCPO relation transformer F , and that $T_{H,f} F : T_{H,\rho} F \rightarrow T_{H,\rho'} F$ is a morphism of ωCPO relation transformers for every ωCPO relation transformer F and every morphism $f : \rho \rightarrow \rho'$ in $\omega\text{CPORelEnv}$. This is an immediate consequence of the following Lemma.

LEMMA 20. For every $\Gamma; \Phi \vdash F$, $\llbracket \Gamma; \Phi \vdash F \rrbracket = ([\Gamma; \Phi \vdash F]^{\omega\text{CPO}}, [\Gamma; \Phi \vdash F]^{\omega\text{CPO}}, [\Gamma; \Phi \vdash F]^{\omega\text{CPORel}}) \in [\omega\text{CPORelEnv}, \omega\text{CPORT}_0]_{\omega_1}$.

The proof is a straightforward induction on the structure of F , using an appropriate result from [Johann and Polonsky 2019] to deduce ω_1 -cocontinuity of $\llbracket \Gamma; \Phi \vdash F \rrbracket$ in each case, together with Lemma 15 and Equation 3 for μ -types. **Anything needed for Lan types??**

We can also prove by simultaneous induction that our interpretations of types interact well with demotion of functorial variables. Indeed, we have that, if $\rho, \rho' : \omega\text{CPOEnv}$, $f : \rho \rightarrow \rho'$, $\rho\phi = \rho'\phi = \rho'\psi$, $f\phi = f\psi = id_{\rho\phi}$, $\Gamma; \Phi, \phi^k \vdash F$, $\Gamma; \Phi, \bar{\alpha} \vdash G$, $\Gamma; \Phi, \alpha_1 \dots \alpha_k \vdash H$, and $\Gamma; \Phi \vdash K$,

then

$$\llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\omega\text{CPO}} \rho = \llbracket \Gamma, \psi; \Phi \vdash F[\phi := \psi] \rrbracket^{\omega\text{CPO}} \rho \quad (4)$$

$$\llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\omega\text{CPO}} f = \llbracket \Gamma, \psi; \Phi \vdash F[\phi := \psi] \rrbracket^{\omega\text{CPO}} f \quad (5)$$

$$\llbracket \Gamma; \Phi \vdash G[\overline{\alpha := \bar{K}}] \rrbracket^{\omega\text{CPO}} \rho = \llbracket \Gamma; \Phi, \bar{\alpha} \vdash G \rrbracket^{\omega\text{CPO}} \rho[\overline{\alpha := \llbracket \Gamma; \Phi \vdash K \rrbracket^{\omega\text{CPO}} \rho}] \quad (6)$$

$$\llbracket \Gamma; \Phi \vdash G[\overline{\alpha := \bar{K}}] \rrbracket^{\omega\text{CPO}} f = \llbracket \Gamma; \Phi, \bar{\alpha} \vdash G \rrbracket^{\omega\text{CPO}} f[\overline{\alpha := \llbracket \Gamma; \Phi \vdash K \rrbracket^{\omega\text{CPO}} f}] \quad (7)$$

$$\llbracket \Gamma; \Phi \vdash F[\phi := H] \rrbracket^{\omega\text{CPO}} \rho = \llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\omega\text{CPO}} \rho[\phi := \lambda \bar{A}. \llbracket \Gamma; \Phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}} \rho[\overline{\alpha := \bar{A}}]] \quad (8)$$

$$\llbracket \Gamma; \Phi \vdash F[\phi := H] \rrbracket^{\omega\text{CPO}} f = \llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\omega\text{CPO}} f[\phi := \lambda \bar{A}. \llbracket \Gamma; \Phi, \bar{\alpha} \vdash H \rrbracket^{\omega\text{CPO}} f[\overline{\alpha := id_{\bar{A}}}]] \quad (9)$$

Identities analogous to (4) through (9) hold for ωCPO relational interpretations as well.

Some questions/issues:

- Can we write `zipBush` and `appendBush` with ∂ and \int ? We could already represent the uncurried type of `appendBush` (although not its curried type), but couldn't recurse over both input bushes because folds take natural transformations as inputs.
- More generally, how do we compute with ∂ and \int ? Can we use the colimit formulation of `Lans` (see Lemma 6.3.7 of [Riehl 2016]) to get a handle on this?
- What is the connection between exponentials and natural transformations? (Should we assume only small objects are exponentiable?) Do we want the former or the latter for computational purposes? (I suspect the latter.)
[From `nlab`: In a functor category D^C , a natural transformation $\alpha : F \rightarrow G$ is exponentiable if (though probably not “only if”) it is cartesian and each component $\alpha_c : Fc \rightarrow Gc$ is exponentiable in D . Given $H \rightarrow F$ we define $(\Pi_\alpha H)c = \Pi_{\alpha_c}(Hc)$; then for $u : c \rightarrow c'$ to obtain a map $\Pi_{\alpha_c}(Hc) \rightarrow \Pi_{\alpha_{c'}}(Hc')$ we need a map $\alpha_{c'}^*(\Pi_{\alpha_c}(Hc)) \rightarrow Hc'$. But since α is cartesian, $\alpha_{c'}^*(\Pi_{\alpha_c}(Hc)) \cong \alpha_c^*(\Pi_{\alpha_c}(Hc))$, so we have the counit $\alpha_c^*(\Pi_{\alpha_c}(Hc)) \rightarrow Hc$ that we can compose with Hu .]
- After we understand what we can do with `Lans` and folds on GADTs we might want to try to extend calculus with term-level fixpoints. This would give a categorical analogue for GADTs of [Pitts 1998, 2000] for ADTs. Would it also more accurately reflect how GADTs are used in practice, or are functions over GADTs usually folds? Investigate applications in the literature and/or in implementations.
- ωCPO is a natural choice for modeling general recursion. We know $(\text{Lan}_C^\gamma \mathbb{1})D$ is $C \rightarrow D$ for any closed type C . (Also for select classes of open types?) So can model $\text{Nat} \rightarrow \gamma$. But the functor $NX = \text{Nat} \rightarrow X$ isn't ω -cocontinuous. It also doesn't preserve ω_1 -presentable objects, i.e., countable ωCPO s since $\text{Nat} \rightarrow \text{Nat}$ is not countable. So we cannot have a functor like N as the subscript to `Lan` and expect the resulting `Lan` to be ω_1 -cocontinuous.
- What functors can be subscripts to `Lan` and produce ω_1 -cocontinuous functors? We can use functors that preserve presentable objects by theorem in [Johann and Polonsky 2019], and possibly others as well. These include polynomial functors, ADTs and nested types seen as functors, certain (which?) GADTs seen as functors? How big can GADTs get?

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