

# Parametricity and Free Theorems for Nested Types

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Abstract goes here

## 1 INTRODUCTION

Suppose we wanted to prove some property of programs over an algebraic data type (ADT) such as that of lists, coded in Agda as

```
data List (A : Set) : Set where
  nil : List A
  Cons : A → List A → List A
```

A natural approach to the problem uses structural induction on the input data structure in question. This requires knowing not just the definition of the ADT of which the input data structure is an instance, but also the program text for the functions involved in the properties to be proved. For example, to prove by induction that mapping a polymorphic function over a list and then reversing the resulting list is the same as reversing the original list and then mapping the function over the result, we unwind the (recursive) definitions of the reverse and map functions over lists to according to the inductive structure of the input list. Such data-driven induction proofs over ADTs are so routine that they are often included in, say, undergraduate functional programming courses.

An alternative technique for proving results like the above map-reverse property for lists is to use parametricity, a formalization of extensional type-uniformity in polymorphic languages. Parametricity captures the intuition that a polymorphic program must act uniformly on all of its possible type instantiations; it is formalized as the requirement that every polymorphic program preserves all relations between any pair of types that it is instantiated with. Parametricity was originally put forth by Reynolds [Reynolds 1983] for System F [Girard et al. 1989], the formal calculus at the core of all polymorphic functional languages. It was later popularized as Wadler’s “theorems for free” [Walder 1989] because it allows the deduction of many properties of programs in such languages solely from their types, i.e., with no knowledge whatsoever of the text of the programs involved. To get interesting free theorems, Wadler’s calculus included, implicitly, built-in list types; indeed, most of the free theorems in [Walder 1989] are consequences of naturality for polymorphic list-processing functions. However, parametricity can also be used to prove naturality properties for non-list ADTs, as well as results, like correctness of program optimizations like *short cut fusion* [Gill et al. 1993; Johann 2002, 2003], that go beyond simple naturality.

This paper is about parametricity and free theorems for a polymorphic calculus with explicit syntax not just for ADTs, but for nested types as well. An ADT defines a *family of inductive data types*, one for each input type. For example, the List data type definition above defines a collection of data types List A, List B, List (A × B), List (List A), etc., each independent of all the others. By contrast, a nested type [Bird and Meertens 1998] is an *inductive family of data types* that is defined over, or is defined mutually recursively with, (other) such data types. Since the structures of the data type at one type can depend on those at other types, the entire family of types must be defined at once. Examples of nested types include, trivially, ordinary ADTs, such as list and tree types; simple nested types, such as the data type

```
data PTree (A : Set) : Set where
  pleaf : A → PTree A
  pnode : PTree (A × A) → PTree A
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reversePTree : ∀{A : Set} → PTree A → PTree A
reversePTree {A} = foldPTree {A} {PTree}
  pleaf
  (λp → pnode (mapPTree swap p))

foldPTree : ∀{A : Set} → {F : Set → Set} →
  ({B : Set} → B → FB) →
  ({B : Set} → F(B × B) → FB) →
  PTree A → F A
foldPTree n c (pleaf x) = n x
foldPTree n c (pnode p) = c (foldPTree n c p)

mapPTree : ∀{AB : Set} → (A → B) → PLeaves A → PLeaves B
mapPTree f (pleaf x) = pleaf (f x)
mapPTree f (pnode p) = pnode (mapPTree f (λp → (f(π1 p), f(π2 p))) p)

swap : ∀{A : Set} → (A × A) → (A × A)
swap (x, y) = (y, x)

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Fig. 1. reversePTree and auxiliary functions in Agda

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reverseBush : ∀{A : Set} → Bush A → Bush A
reverseBush {A} = foldBush {A} {Bush} bnil balg

foldBush : ∀{A : Set} → {F : Set → Set} →
  ({B : Set} → FB) →
  ({B : Set} → B → F (F B) → F B) →
  Bush A → F A
foldBush bn bc bnil = bn
foldBush bn bc (bcons x bb) =
  bc x (foldBush bn bc (mapBush (foldBush bn bc) bb))

mapBush : ∀{AB : Set} → (A → B) → (Bush A) → (Bush B)
mapBush _ bnil = bnil
mapBush f (bcons x bb) = bcons (f x) (mapBush (mapBush f) bb)

balg : ∀{B : Set} → B → Bush (Bush B) → Bush B
balg x bnil = bcons x bnil
balg x (bcons bnll bbbx) = bcons x (bcons bnll bbbx)
balg x (bcons (bcons y bx) bbbx) =
  bcons y (bcons (bcons x bx) bbbx)

```

Fig. 2. reverseBush and auxiliary functions in Agda

of perfect trees, whose recursive occurrences never appear below other type constructors; “deep” nested types [Johann and Polonsky 2020], such as the data type

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data Forest (A : Set) : Set where
  fempty : Forest A
  fnode : A → PTree (Forest A) → Forest A

```

of perfect forests, whose recursive occurrences appear below type constructors for other nested types; and truly nested types<sup>1</sup>, such as the data type

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data Bush (A : Set) : Set where
  bnil : Bush A
  bcons : A → Bush (Bush A) → Bush A

```

of bushes (also called *bootstrapped heaps* in [Okasaki 1999]), whose recursive occurrences appear below their own type constructors.

Suppose we now want to prove properties of functions over nested types. We might, for example, want to prove a map-reverse property for the functions on perfect trees in Figure 1, or for those on bushes<sup>2</sup> in Figure 2. A few well-chosen examples quickly convince us that such a property should indeed hold for perfect trees, and, drawing inspiration from the situation for ADTs, we easily construct a proof by induction on the input perfect tree. To formally establish this result, we could even prove it in Coq or Agda: each of these provers actually generates an induction rule for perfect trees and the generated rule gives the expected result because proving properties of perfect trees requires only that we induct over the top-level perfect tree in the recursive position, leaving any data internal to the input tree untouched.

Unfortunately, it is nowhere near as clear that analogous intuitive or formal inductive arguments can be made for the map-reverse property for bushes. Indeed, a proof by induction on the input bush must recursively induct over the bushes that are internal to the top-level bush in the recursive position. This is sufficiently delicate that no induction rule for bushes or other truly nested types was known until very recently, when *deep induction* [Johann and Polonsky 2020] was developed as a way to induct over *all* of the structured data present in an input. Deep induction thus not only gave the first principled and practically useful structural induction rules for bushes and other truly

<sup>1</sup>Nested types that are defined over themselves are known as *truly nested types*.

<sup>2</sup>To define the foldBush and mapBush functions in Figure 2 it is necessary to turn off Agda’s termination checker.

nested types, and has also opened the way for incorporating automatic generation of such rules for (truly) nested data types — and, eventually, even GADTs — into modern proof assistants.

Of course it is great to know that we *can*, at last, prove properties of programs over (truly) nested types by induction. But recalling that inductive proofs over ADTs can sometimes be circumvented in the presence of parametricity, we might naturally ask:

*Can we derive properties of functions over (truly) nested types from parametricity?*

This paper answers the above question in the affirmative by constructing a parametric model for a polymorphic calculus that provides primitives for constructing nested types.

We introduce our calculus in Section 2. At the type level, it is the level-2-truncation of the higher-kinded calculus from [Johann and Polonsky 2019], augmented with a primitive type of natural transformations. To represent nested types, it constructs type expressions not just from standard type variables, but also from type constructor variables of various arities. In addition, it includes an explicit  $\mu$ -construct to represent type-level recursion with respect to these variables. The class of nested types thus represented is very robust and includes all (truly) nested types known from the literature. In Section 3 we construct set and relational interpretations for the types from Section 2. As is usual when modeling parametricity, types are interpreted as functors from environments interpreting their type variable contexts to set or relations, as appropriate. But in order to ensure that these functors satisfy the cocontinuity properties needed to know that the fixpoints interpreting  $\mu$ -types exist, set environments must map each  $k$ -ary type constructor variable to an appropriately cocontinuous  $k$ -ary functor on sets and relation environments must map each  $k$ -ary type constructor variable to an appropriately cocontinuous  $k$ -ary relation transformer, and these cocontinuity conditions must be threaded throughout our type interpretations in such a way that the resulting model satisfies an appropriate Identity Extension Lemma (Theorem 23). Properly propagating cocontinuity turns out to be both subtle and challenging, and Section 4, where it is done, is where the bulk of the work in our model construction lies. At the term level, our calculus includes constructs representing the actions on morphisms of functors interpreting types, initial algebras for fixpoints of these functors, and structured recursion over elements of these initial algebras (i.e., map, in, and fold constructs, respectively). While our calculus does not support general recursion at the term level, it is strongly normalizing, so does perhaps edge us toward the kind of provably total practical programming language proposed at the end of [Walder 1989]. In Section 5, we construct set and relational interpretations for the terms of our calculus. As usual in parametric models, terms are interpreted as natural transformations from interpretations of the term contexts in which they are formed to the interpretations of their types, and these must cohere in what is essentially a fibrational way [Ghani et al. 2015]. Immediately from the definition of our interpretation we prove in Section 5.4 a scheme deriving free theorems that are consequences of naturality of functions that are polymorphic over nested types. This scheme is very general, is parameterized over both the data type and the polymorphic function at hand, and it has all of the above map-reverse results as instances. The relationship between naturality and parametricity has long been of interest, and our inclusion of a primitive type of natural transformations makes it possible for us clearly delineate those free theorems that are consequences of naturality, and thus would hold even in non-parametric models of our calculus, from those that use the full power of parametricity to go beyond naturality. In Section 5.5 we prove that our model satisfies an Abstraction Theorem (Theorem 29), from which we derive several of this latter kind of free theorem in Section 6. Specifically, we state and prove (non-)inhabitation results in Sections 6.1 and 6.2, a free theorem for the type of a filter function on generalized rose trees in Section 6.4, and the correctness of short cut fusion for nested types in Section 6.7.

We are not the first to consider parametricity at higher kinds. Atkey [Atkey 2012] constructs a parametric model for full System F $\omega$ , but within the impredicative Calculus of Inductive Constructions (iCIC) rather than in a semantic category. Although Atkey’s construction is similar to ours, he does not provide primitives for constructing data types, so he is not concerned that his type constructors represent functors or that fixpoints for these functors exist. As a result, his relation transformers include no cocontinuity conditions, and no such conditions are propagated throughout his model construction. Atkey does, however, verify the existence of initial algebras in iCIC for (what are intended to be) syntactic reflections of functors, provided each is *given*, together with an identity- and composition-preserving *fmap* function. (Unfortunately, the *fmap* function given with *F* is not otherwise required to behave like the functorial action of *F*.) Atkey’s “functors” are therefore postulated rather than constructed. He also does not in any way indicate which type constructors can be endowed with the required *fmap* functions, and we suspect that doing so would essentially result in a higher-kinded extension of the results presented here for a calculus giving an explicit syntax for the higher-kinded types it guarantees are properly functorial by construction.

Like us, Pitts [Pitts 1998, 2000] extends parametricity from pure System F to accommodate direct representations of data types. Only list types are added in [Pitts 2000], although other ADTs are easily accommodated as in [Pitts 1998]. But even there, Pitts considers only polynomial data types, all of which can all be modeled as fixpoints of *first-order* functors. Part of Pitts’ motivation is to show that ADTs and their Church encodings have the same operational behavior in his system, which he accomplishes via an operational analogue of short cut fusion for ADTs. (We cannot even ask this question about our system, since it does not support Church encodings of nested types, or even of ADTs such as list types, or even of pair or sum types.) Particularly interesting would be the operational analogues of functoriality and cocontinuity required to extend Pitts’ parametricity results to a calculus supplying primitives for constructing nested types and other data types modeled as fixpoints of *higher-order* functors.

Finally, there is a long line of work on categorical models of parametricity for System F; see, for example, [Bainbridge et al. 1990; Birkedal and Møgelberg 2005; Dunphy and Reddy 2004; Ghani et al. 2015; Hasegawa 1994; Jacobs 1999; Ma and Reynolds 1992; Robinson and Rosolini 1994]. Much of this work does ultimately extend parametricity to data types that are modeled as fixpoints of first-order functors, but to our knowledge this has never been done by building the data types directly into the calculus whose parametricity is to be modeled categorically. Indeed, all models we know of are constructed using the categorical equivalent of Atkey’s approach, i.e., by verifying in the just-constructed parametric model the existence of initial algebras for type constructors obtained by reflecting back into syntax the functors induced by the data types’ Church encodings and supporting self-related universal strengths. The present paper draws on this rich tradition of categorical models of parametricity for System F, but takes care to ensure that the functoriality needed to support primitive nested types interacts nicely with System F’s full impredicative polymorphism.

## 2 THE CALCULUS

### 2.1 Types

For each  $k \geq 0$ , we assume countable sets  $\mathbb{T}^k$  of *type constructor variables of arity  $k$*  and  $\mathbb{F}^k$  of *functorial variables of arity  $k$* , all mutually disjoint. The sets of all type constructor variables and functorial variables are  $\mathbb{T} = \bigcup_{k \geq 0} \mathbb{T}^k$  and  $\mathbb{F} = \bigcup_{k \geq 0} \mathbb{F}^k$ , respectively, and a *type variable* is any element of  $\mathbb{T} \cup \mathbb{F}$ . We use lower case Greek letters for type variables, writing  $\phi^k$  to indicate that  $\phi \in \mathbb{T}^k \cup \mathbb{F}^k$ , and omitting the arity indicator  $k$  when convenient, unimportant, or clear from context. We reserve letters from the beginning of the alphabet to denote type variables of arity 0, i.e., elements of  $\mathbb{T}^0 \cup \mathbb{F}^0$ . We write  $\bar{\zeta}$  for either a set  $\{\zeta_1, \dots, \zeta_n\}$  of type constructor variables or a

set of functorial variables when the cardinality  $n$  of the set is unimportant or clear from context. If  $P$  is a set of type variables we write  $P, \bar{\phi}$  for  $P \cup \bar{\phi}$  when  $P \cap \bar{\phi} = \emptyset$ . We omit the vector notation for a singleton set, thus writing  $\phi$ , instead of  $\bar{\phi}$ , for  $\{\phi\}$ .

DEFINITION 1. Let  $V$  be a finite subset of  $\mathbb{T}$ , let  $P$  be a finite subset of  $\mathbb{F}$ , let  $\bar{\alpha}$  be a finite subset of  $\mathbb{F}^0$  disjoint from  $P$ , and let  $\phi^k \in \mathbb{F}^k \setminus P$ . The sets  $\mathcal{T}(V)$  of type constructor expressions over  $V$  and  $\mathcal{F}^P(V)$  of functorial expressions over  $P$  and  $V$  are given by

$$\mathcal{T}(V) ::= V \mid \text{Nat}^{\bar{\alpha}} \mathcal{F}^{\bar{\alpha}}(V) \mid \overline{V\mathcal{T}(V)}$$

and

$$\begin{aligned} \mathcal{F}^P(V) ::= & \mathcal{T}(V) \mid 0 \mid 1 \mid \overline{P\mathcal{F}^P(V)} \mid \overline{V\mathcal{F}^P(V)} \mid \mathcal{F}^P(V) + \mathcal{F}^P(V) \mid \mathcal{F}^P(V) \times \mathcal{F}^P(V) \\ & \mid \left( \mu\phi^k . \lambda\alpha_1 \dots \alpha_k . \mathcal{F}^{P, \alpha_1, \dots, \alpha_k, \phi}(V) \right) \overline{\mathcal{F}^P(V)} \end{aligned}$$

A type over  $P$  and  $V$  is any element of  $\mathcal{T}(V) \cup \mathcal{F}^P(V)$ .

The notation for types entails that an application  $\tau\tau_1 \dots \tau_k$  is allowed only when  $\tau$  is a type variable of arity  $k$ , or  $\tau$  is a subexpression of the form  $\mu\phi^k . \lambda\alpha_1 \dots \alpha_k . \tau'$ . Moreover, if  $\tau$  has arity  $k$  then  $\tau$  must be applied to exactly  $k$  arguments. Accordingly, an overbar indicates a sequence of subexpressions whose length matches the arity of the type applied to it. The fact that types are always in  $\eta$ -long normal form avoids having to consider  $\beta$ -conversion at the level of types. In a subexpression  $\text{Nat}^{\bar{\alpha}} \sigma \tau$ , the  $\text{Nat}$  operator binds all occurrences of the variables in  $\bar{\alpha}$  in  $\sigma$  and  $\tau$ . Similarly, in a subexpression  $\mu\phi^k . \lambda\bar{\alpha} . \tau$ , the  $\mu$  operator binds all occurrences of the variable  $\phi$ , and the  $\lambda$  operator binds all occurrences of the variables in  $\bar{\alpha}$ , in the body  $\tau$ .

A type constructor context is a finite set  $\Gamma$  of type constructor variables, and a functorial context is a finite set  $\Phi$  of functorial variables. In Definition 2, a judgment of the form  $\Gamma; \Phi \vdash \tau : \mathcal{T}$  or  $\Gamma; \Phi \vdash \tau : \mathcal{F}$  indicates that the type  $\tau$  is intended to be functorial in the variables in  $\Phi$  but not necessarily in the variables in  $\Gamma$ .

DEFINITION 2. The formation rules for the set  $\mathcal{T} \subseteq \bigcup_{V \subseteq \mathbb{T}} \mathcal{T}(V)$  of well-formed type constructor expressions are

$$\begin{array}{c} \overline{\Gamma, \alpha^0; \emptyset \vdash \alpha^0 : \mathcal{T}} \\ \hline \Gamma; \bar{\alpha}^0 \vdash \sigma : \mathcal{F} \quad \Gamma; \bar{\alpha}^0 \vdash \tau : \mathcal{F} \\ \hline \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}^0} \sigma \tau : \mathcal{T} \end{array}$$

The formation rules for the set  $\mathcal{F} \subseteq \bigcup_{V \subseteq \mathbb{T}, P \subseteq \mathbb{F}} \mathcal{F}^P(V)$  of well-formed functorial expressions are

$$\begin{array}{c} \Gamma; \emptyset \vdash \tau : \mathcal{T} \\ \hline \Gamma; \emptyset \vdash \tau : \mathcal{F} \quad \Gamma; \Phi, \alpha^0 \vdash \alpha^0 : \mathcal{F} \quad \Gamma; \Phi \vdash 0 : \mathcal{F} \quad \Gamma; \Phi \vdash 1 : \mathcal{F} \\ \hline \phi^k \in \Gamma \cup \Phi \quad \Gamma; \Phi \vdash \tau : \mathcal{F} \\ \hline \Gamma; \Phi \vdash \phi^k \bar{\tau} : \mathcal{F} \\ \hline \Gamma; \Phi, \bar{\alpha}, \phi^k \vdash \tau : \mathcal{F} \quad \Gamma; \Phi \vdash \tau : \mathcal{F} \\ \hline \Gamma; \Phi \vdash (\mu\phi^k . \lambda\bar{\alpha} . \tau) \bar{\tau} : \mathcal{F} \\ \hline \Gamma; \Phi \vdash \sigma : \mathcal{F} \quad \Gamma; \Phi \vdash \tau : \mathcal{F} \quad \Gamma; \Phi \vdash \sigma : \mathcal{F} \quad \Gamma; \Phi \vdash \tau : \mathcal{F} \\ \hline \Gamma; \Phi \vdash \sigma + \tau : \mathcal{F} \quad \Gamma; \Phi \vdash \sigma \times \tau : \mathcal{F} \end{array}$$

A type  $\tau$  is well-formed if it is either a well-formed type constructor expression or a well-formed functorial expression.

If  $\tau$  is a closed type we may write  $\vdash \tau$ , rather than  $\emptyset; \emptyset \vdash \tau$ , for the judgment that it is well-formed. Definition 2 ensures that the expected weakening rules for well-formed types hold — although weakening does not change the contexts in which Nat-types can be formed. If  $\Gamma; \emptyset \vdash \sigma : \mathcal{T}$  and  $\Gamma; \emptyset \vdash \tau : \mathcal{T}$ , then our rules allow formation of the type  $\Gamma; \emptyset \vdash \text{Nat}^0 \sigma \tau$ . Since a type  $\Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} \sigma \tau$  represents a natural transformation in  $\bar{\alpha}$  from  $\sigma$  to  $\tau$ , the type  $\Gamma; \emptyset \vdash \text{Nat}^0 \sigma \tau$  represents the standard arrow type  $\Gamma \vdash \sigma \rightarrow \tau$  in our calculus. We similarly represent a standard  $\forall$ -type  $\Gamma; \emptyset \vdash \forall \bar{\alpha}. \tau$  as  $\Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} \mathbb{1} \tau : \mathcal{F}$  in our calculus. However, if  $\bar{\alpha}$  is non-empty then  $\tau$  cannot be of the form  $\text{Nat}^{\bar{\beta}} H K$  since  $\Gamma; \bar{\alpha} \vdash \text{Nat}^{\bar{\beta}} H K$  is not a valid type judgment in our calculus (except by weakening).

Can have list of forall types or arrows. Definition 2 allows the formation of all of the (closed) nested types from the introduction:

$$\begin{aligned} \text{List } \alpha &= \mu\beta. \mathbb{1} + \alpha \times \beta = (\mu\phi. \lambda\beta. \mathbb{1} + \beta \times \phi\beta) \alpha \\ \text{PTree } \alpha &= (\mu\phi. \lambda\beta. \beta + \phi(\beta \times \beta)) \alpha \\ \text{Forest } \alpha &= (\mu\phi. \lambda\beta. \mathbb{1} + \beta \times \text{PTree}(\phi\beta)) \alpha \\ \text{Bush } \alpha &= (\mu\phi. \lambda\beta. \mathbb{1} + \beta \times \phi(\phi\beta)) \alpha \end{aligned}$$

Each of these types can either be natural in  $\alpha$  or not, according to whether  $\alpha \in \Gamma$  or  $\alpha \in \Phi$ . For example, if  $\emptyset; \alpha \vdash \text{List } \alpha$ , then the type  $\vdash \text{Nat}^{\alpha} \mathbb{1} (\text{List } \alpha) : \mathcal{T}$  is well-formed; If  $\alpha; \emptyset \vdash \text{List } \alpha$ , then it is not. Definition 2 also allows the derivation of, e.g., the type  $\alpha; \emptyset \vdash \text{Nat}^{\alpha} (\text{List } \alpha) (\text{Tree } \alpha \gamma)$  representing a natural transformation from lists to trees that is natural in  $\alpha$  but not necessarily in  $\gamma$ . We emphasize that types can be functorial in variables of arity greater than 0. For example, the type  $\text{GRose } \phi \alpha = \mu\beta. \alpha \times \phi\beta$  can be functorial in  $\phi$  if  $\phi \in \Phi$ . As usual, whether  $\phi \in \Gamma$  or  $\phi \in \Phi$  determines whether types such as  $\text{Nat}^{\alpha} (\text{GRose } \phi \alpha) (\text{List } \alpha)$  are well-formed. But even if  $\text{GRose}$  is functorial in  $\phi$ , it still cannot be the (co)domain of a Nat type representing a natural transformation in  $\phi$ . This is because our calculus does not allow naturality in variables of arity greater than 0.

Definition 2 explicitly considers types in  $\mathcal{T}$  to be types in  $\mathcal{F}$  that are functorial in no variables. It is not hard to see that this definition also supports the demotion of functorial variables in a well-formed type  $\tau$  to non-functorial status. The proof is by induction on the structure of  $\tau$ .

LEMMA 3. If  $\Gamma; \Phi, \phi^k \vdash \tau : \mathcal{F}$ , then  $\Gamma, \psi^k; \Phi \vdash \tau[\phi^k := \psi^k]$  is also derivable. Here,  $\tau[\phi := \psi]$  is the textual replacement of  $\phi$  in  $\tau$ , meaning that all occurrences of  $\phi\bar{\sigma}$  in  $\tau$  become  $\psi\bar{\sigma}$ .

In addition to textual replacement, we also have a proper substitution operation on types. If  $\tau$  is a type over  $P$  and  $V$ , if  $P$  and  $V$  contain only type variables of arity 0, and if  $k = 0$  for every occurrence of  $\phi^k$  bound by  $\mu$  in  $\tau$ , then we say that  $\tau$  is *first-order*; otherwise we say that  $\tau$  is *second-order*. Substitution for first-order types is the usual capture-avoiding textual substitution. We write  $\tau[\alpha := \sigma]$  for the result of substituting  $\sigma$  for  $\alpha$  in  $\tau$ , and  $\tau[\alpha_1 := \tau_1, \dots, \alpha_k := \tau_k]$ , or  $\tau[\bar{\alpha} := \bar{\tau}]$  when convenient, for  $\tau[\alpha_1 := \tau_1][\alpha_2 := \tau_2, \dots, \alpha_k := \tau_k]$ . Substitution for second-order types is defined below, where we adopt a similar notational convention for vectors of types.

DEFINITION 4. If  $\phi^k \in \Gamma \cup \Phi$  with  $k \geq 1$ , if  $\Gamma; \Phi \vdash F : \mathcal{F}$ , and if  $\Gamma, \bar{\beta}; \Phi, \bar{\alpha} \vdash H : \mathcal{F}$  with  $|\bar{\alpha}| + |\bar{\beta}| = k$ , then  $\Gamma \setminus \phi^k; \Phi \setminus \phi^k \vdash F[\phi := \bar{\beta}. \bar{\alpha} H] : \mathcal{F}$ , where the operation  $(\cdot)[\phi := H]$  of second-order

type substitution is defined by:

$$\begin{aligned}
(\text{Nat}^Y G K)[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \text{Nat}^Y (G[\phi :=_{\bar{\beta}, \bar{\alpha}} H]) (K[\phi :=_{\bar{\beta}, \bar{\alpha}} H]) \\
\mathbb{1}[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \mathbb{1} \\
\mathbb{0}[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \mathbb{0} \\
(\psi \overline{\sigma \tau})[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \begin{cases} \psi \overline{\tau[\phi :=_{\bar{\beta}, \bar{\alpha}} H]} & \text{if } \psi \neq \phi \\ \overline{H[\alpha := \tau[\phi :=_{\bar{\beta}, \bar{\alpha}} H]] [\beta := \sigma[\phi :=_{\bar{\beta}, \bar{\alpha}} H]]} & \text{if } \psi = \phi \end{cases} \\
(\sigma + \tau)[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \sigma[\phi :=_{\bar{\beta}, \bar{\alpha}} H] + \tau[\phi :=_{\bar{\beta}, \bar{\alpha}} H] \\
(\sigma \times \tau)[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= \sigma[\phi :=_{\bar{\beta}, \bar{\alpha}} H] \times \tau[\phi :=_{\bar{\beta}, \bar{\alpha}} H] \\
((\mu\psi.\lambda\bar{y}.G)\bar{\tau})[\phi :=_{\bar{\beta}, \bar{\alpha}} H] &= (\mu\psi.\lambda\bar{y}.G[\phi :=_{\bar{\beta}, \bar{\alpha}} H]) \overline{\tau[\phi :=_{\bar{\beta}, \bar{\alpha}} H]}
\end{aligned}$$

We omit the variable subscripts in second-order type constructor substitution when convenient.

## 2.2 Terms

We assume an infinite set  $\mathcal{V}$  of term variables disjoint from  $\mathbb{T}$  and  $\mathbb{F}$ . If  $\Gamma$  be a type constructor context and  $\Phi$  is a functorial context, then a *term context* for  $\Gamma$  and  $\Phi$  is a finite set of bindings of the form  $x : \tau$ , where  $x \in \mathcal{V}$  and  $\Gamma; \Phi \vdash \tau : \mathcal{F}$ . We adopt the same conventions for denoting disjoint unions and for vectors in term contexts as for type constructor contexts and functorial contexts.

DEFINITION 5. Let  $\Delta$  be a term context for  $\Gamma$  and  $\Phi$ . The formation rules for the set of well-formed terms over  $\Delta$  are

$$\begin{array}{c}
\frac{\Gamma; \emptyset \vdash \tau : \mathcal{F}}{\Gamma; \emptyset \mid \Delta, x : \tau \vdash x : \tau} \quad \frac{\Gamma; \Phi \vdash \tau : \mathcal{F}}{\Gamma; \Phi \mid \Delta, x : \tau \vdash x : \tau} \\
\\
\frac{}{\Gamma; \Phi \mid \Delta \vdash \top : \mathbb{1}} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : \mathbb{0} \quad \Gamma; \Phi \vdash \tau : \mathcal{F}}{\Gamma; \Phi \mid \Delta \vdash \perp_{\tau} t : \tau} \\
\\
\frac{\Gamma; \Phi \mid \Delta \vdash s : \sigma}{\Gamma; \Phi \mid \Delta \vdash \text{inl } s : \sigma + \tau} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : \tau}{\Gamma; \Phi \mid \Delta \vdash \text{inr } t : \sigma + \tau} \\
\\
\frac{\Gamma; \Phi \vdash \tau, \sigma : \mathcal{F} \quad \Gamma; \Phi \mid \Delta \vdash t : \sigma + \tau \quad \Gamma; \Phi \mid \Delta, x : \sigma \vdash l : \gamma \quad \Gamma; \Phi \mid \Delta, y : \tau \vdash r : \gamma}{\Gamma; \Phi \mid \Delta \vdash \text{case } t \text{ of } \{\text{inl } x \mapsto l; \text{inr } y \mapsto r\} : \gamma} \\
\\
\\
\frac{\Gamma; \Phi \mid \Delta \vdash s : \sigma \quad \Gamma; \Phi \mid \Delta \vdash t : \tau}{\Gamma; \Phi \mid \Delta \vdash (s, t) : \sigma \times \tau} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau}{\Gamma; \Phi \mid \Delta \vdash \pi_1 t : \sigma} \quad \frac{\Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau}{\Gamma; \Phi \mid \Delta \vdash \pi_2 t : \tau}
\end{array}$$



$$\begin{array}{c}
\frac{\Gamma; \bar{\alpha} \vdash F : \mathcal{F} \quad \Gamma; \bar{\alpha} \vdash G : \mathcal{F} \quad \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G}{\Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G} \\
\\
\frac{\Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \quad \overline{\Gamma; \Phi \vdash \tau : \mathcal{F}} \quad \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \tau]}{\Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \tau]} \\
\\
\frac{\Gamma; \bar{\phi}, \bar{\gamma} \vdash H : \mathcal{F} \quad \overline{\Gamma; \bar{\beta}, \bar{\gamma} \vdash F : \mathcal{F}} \quad \overline{\Gamma; \bar{\beta}, \bar{\gamma} \vdash G : \mathcal{F}}}{\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{F}, \bar{G}}^{\bar{F}, \bar{G}} : \text{Nat}^{\emptyset} (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} \bar{F}]) H[\bar{\phi} :=_{\bar{\beta}} \bar{G}])} \\
\\
\frac{\Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H : \mathcal{F}}{\Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta}} \\
\\
\frac{\Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H : \mathcal{F} \quad \Gamma; \bar{\beta}, \bar{\gamma} \vdash F : \mathcal{F}}{\Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^{\emptyset} (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}]) F (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta}) \bar{\beta}}
\end{array}$$

In the rule for  $L_{\bar{\alpha}} x. t$ , the  $L$  operator binds all occurrences of the type variables in  $\bar{\alpha}$  in the type of the term variable  $x$  and in the body  $t$ , as well as all occurrences of  $x$  in  $t$ . In the rule for  $t_{\bar{\tau}} s$  there is one functorial expression  $\tau$  for every functorial variable  $\alpha$ . In the rule for  $\text{map}_{\bar{F}, \bar{G}}^{\bar{F}, \bar{G}}$  there is one functorial expression  $F$  and one functorial expression  $G$  for each functorial variable in  $\bar{\phi}$ . Moreover, for each  $\phi^k \in \bar{\phi}$  the number of functorial variables  $\beta$  in the judgments for its corresponding functorial expressions  $F$  and  $G$  is  $k$ . In the rules for  $\text{in}_H$  and  $\text{fold}_H^F$ , the functorial variables in  $\bar{\beta}$  are fresh with respect to  $H$ , and there is one  $\beta$  for every  $\alpha$ . (Recall from above that, in order for the types of  $\text{in}_H$  and  $\text{fold}_H^F$  to be well-formed, the length of  $\alpha$  must equal the arity of  $\phi$ .) Substitution for terms is the obvious extension of the usual capture-avoiding textual substitution, and Definition 5 ensures that the expected weakening rules for well-formed terms hold.

Using Definition 5 we can represent the `reversePTree` function from Figure 1 in our calculus as

$$\vdash \text{fold}_{\beta + \phi(\beta \times \beta)}^{PTree \alpha} (\text{in}_{\beta + \phi(\beta \times \beta)} \circ s) : \text{Nat}^{\alpha} (PTree \alpha) (PTree \alpha)$$

where

$$\begin{array}{ll}
\vdash \text{fold}_{\beta + \phi(\beta \times \beta)}^{PTree \alpha} & : \text{Nat}^{\emptyset} (\text{Nat}^{\alpha} (\alpha + PTree(\alpha \times \alpha)) (PTree \alpha)) (\text{Nat}^{\alpha} (PTree \alpha) (PTree \alpha)) \\
\vdash \text{in}_{\beta + \phi(\beta \times \beta)} & : \text{Nat}^{\alpha} (\alpha + PTree(\alpha \times \alpha)) (PTree \alpha) \\
\vdash \text{map}_{PTree \alpha}^{\alpha \times \alpha, \alpha \times \alpha} & : \text{Nat}^{\emptyset} (\text{Nat}^{\alpha} (\alpha \times \alpha) (\alpha \times \alpha)) (\text{Nat}^{\alpha} (PTree(\alpha \times \alpha)) (PTree(\alpha \times \alpha)))
\end{array}$$

and  $\text{swap}$  and  $s$  are the terms

$$\vdash L_{\alpha} p. (\pi_2 p, \pi_1 p) : \text{Nat}^{\alpha} (\alpha \times \alpha) (\alpha \times \alpha)$$

and

$$\vdash L_{\alpha} t. \text{case } t \text{ of } \{b \mapsto \text{inl } b; t' \mapsto \text{inr} (\text{map}_{PTree \alpha}^{\alpha \times \alpha, \alpha \times \alpha} \text{swap } t')\} : \text{Nat}^{\alpha} (\alpha + PTree(\alpha \times \alpha)) (\alpha + PTree(\alpha \times \alpha))$$

respectively. We can similarly represent the `reverseBush` function from Figure 2 as

$$\vdash \text{fold}_{\mathbb{1} + \beta \times \phi(\phi \beta)}^{Bush \alpha} (\text{in}_{\mathbb{1} + \beta \times \phi(\phi \beta)} \circ (\mathbb{1} + t \circ i \circ i')) : \text{Nat}^{\alpha} (Bush \alpha) (Bush \alpha)$$

where

$$\begin{array}{ll}
\vdash \text{fold}_{\mathbb{1} + \beta \times \phi(\phi \beta)}^{Bush \alpha} & : \text{Nat}^{\emptyset} (\text{Nat}^{\alpha} (\mathbb{1} + \alpha \times Bush(Bush \alpha)) (Bush \alpha)) (\text{Nat}^{\alpha} (Bush \alpha) (Bush \alpha)) \\
\vdash \text{in}_{\mathbb{1} + \beta \times \phi(\phi \beta)} & : \text{Nat}^{\alpha} (\mathbb{1} + \alpha \times Bush(Bush \alpha)) (Bush \alpha)
\end{array}$$



and  $bnil$ ,  $bcons$ ,  $\text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)}^{-1}$ ,  $t$ ,  $i$ , and  $i'$  are the terms

$$\begin{aligned}
& \vdash \text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)} \circ (L_\alpha x. \text{inl } x) : \text{Nat}^\alpha \mathbb{1} (Bush \alpha) \\
& \vdash \text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)} \circ (L_\alpha x. \text{inr } x) : \text{Nat}^\alpha (\alpha \times Bush (Bush \alpha)) (Bush \alpha) \\
& \vdash \text{fold}_{\mathbb{1}+\beta \times \phi(\phi\beta)}^{(\mathbb{1}+\beta \times \phi(\phi\beta))[\phi := Bush \alpha]} (\text{map}_{\mathbb{1}+\beta \times \phi(\phi\beta)}^{(\mathbb{1}+\beta \times \phi(\phi\beta))[\phi := Bush \alpha][\beta := \alpha], Bush \alpha} \text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)}) \\
& \quad : \text{Nat}^\alpha (Bush \alpha) (\mathbb{1} + \alpha \times Bush (Bush \alpha)) \\
& \vdash L_\alpha (b, s). \text{case } s \{ \quad * \mapsto bcons_\alpha b (bnil_\alpha *); \\
& \quad \quad (s', u) \mapsto \text{case } s' \{ \quad * \mapsto bcons_\alpha b (bcons_{Bush \alpha} (bnil_\alpha *) u); \\
& \quad \quad \quad (b', u') \mapsto bcons_\alpha b' (bcons_{Bush \alpha} (bcons_\alpha b u) u') \} \} \\
& \quad : \text{Nat}^\alpha (\alpha \times (\mathbb{1} + (\mathbb{1} + \alpha \times Bush (Bush \alpha))) \times Bush (Bush (Bush \alpha))) (\alpha \times Bush (Bush \alpha)) \\
& \vdash \alpha \times (\mathbb{1} + \text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)}^{-1} \times Bush (Bush (Bush \alpha))) \\
& \quad : \text{Nat}^\alpha (\alpha \times (\mathbb{1} + Bush \alpha \times Bush (Bush (Bush \alpha)))) \\
& \quad (\alpha \times (\mathbb{1} + (\mathbb{1} + \alpha \times Bush (Bush \alpha)) \times Bush (Bush (Bush \alpha)))) \\
& \vdash \alpha \times (L_\alpha x. (\text{in}_{\mathbb{1}+\beta \times \phi(\phi\beta)}^{-1})_{Bush \alpha} x) \\
& \quad : \text{Nat}^\alpha (\alpha \times Bush (Bush \alpha)) (\alpha \times (\mathbb{1} + Bush (\alpha \times Bush (Bush (Bush \alpha)))))
\end{aligned}$$

respectively. Here,  $\Gamma; \emptyset \mid \Delta \vdash \sigma + \eta : \text{Nat}^{\bar{\alpha}}(\sigma + F) (\sigma + G)$  and  $\Gamma; \emptyset \mid \Delta \vdash \sigma \times \eta : \text{Nat}^{\bar{\alpha}}(\sigma \times F) (\sigma \times G)$  for  $\sigma + \eta := L_{\bar{\alpha}} x. \text{case } x \text{ of } \{s \mapsto \text{inl } s; t \mapsto \text{inr } (\eta_{\bar{\alpha}} t)\}$  and  $\sigma \times \eta := L_{\bar{\alpha}} x. (\pi_1 x, \eta_{\bar{\alpha}}(\pi_2 x))$  for  $\Gamma; \emptyset \mid \Delta \vdash \eta : \text{Nat}^{\bar{\alpha}} F G$  and  $\Gamma; \bar{\alpha} \vdash \sigma : \mathcal{F}$ .

Unfortunately, we cannot write functions, such as  $\text{concat} : PTree \alpha \rightarrow PTree \alpha \rightarrow PTree \alpha$ , that take as input more than one non-algebraic nested type. This is because  $\text{Nat}$ -types must be formed in empty functorial contexts, and this conflicts with the need to feed folds algebras. (Can't fold over pairs to get the right functions; can't get them by using continuation style because of the aforementioned typing conflict.) [Add Daniel's commentary about "real" involution reverse for bushes, too. Massage paragraph. Not a good restriction. Also, generalized folds don't help much.](#)

The presence of the "extra" functorial variables in  $\bar{\gamma}$  in the rules for  $\text{map}_{\bar{H}}^{\bar{F}, \bar{G}}$ ,  $\text{in}_H$ , and  $\text{fold}_H^F$  merit special mention. They allows us to map or fold polymorphic functions over nested types. Consider, for example, the function  $\text{flatten} : \text{Nat}^{\beta}(PTree \beta) (List \beta)$  that maps perfect trees to lists. Even in the absence of extra variables the instance of map required to map each non-functorial monomorphic instantiation of  $\text{flatten}$  over a list of perfect trees is well-typed:

$$\frac{\Gamma; \alpha \vdash List \alpha \quad \Gamma; \emptyset \vdash \sigma \quad \Gamma; \emptyset \vdash \tau \quad \Gamma; \emptyset \vdash PTree \sigma \quad \Gamma; \emptyset \vdash List \tau}{\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{List \alpha}^{PTree \sigma, List \tau} : \text{Nat}^{\emptyset} (\text{Nat}^{\emptyset} (PTree \sigma) (List \tau)) (\text{Nat}^{\emptyset} (List (PTree \sigma)) (List (List \tau)))}$$

But in the absence of  $\bar{\gamma}$ , the instance

$$\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{List \alpha}^{PTree \beta, List \beta} : \text{Nat}^{\emptyset} (\text{Nat}^{\beta} (PTree \beta) (List \beta)) (\text{Nat}^{\beta} (List (PTree \beta)) (List (List \beta)))$$

of map required to map the *polymorphic flatten* function over a list of perfect trees is not: in that setting the functorial contexts for  $F$  and  $G$  in the rule for  $\text{map}_H^{F, G}$  would have to be empty, but the fact that the polymorphic *flatten* function is functorial in some variable, say  $\delta$ , means that it cannot possibly have a type of the form  $\text{Nat}^{\emptyset} F G$  that would be required for it to be the function input to map. Since untypeability of this instance of map is unsatisfactory in a polymorphic calculus, where we naturally expect to be able to manipulate entire polymorphic functions rather than just their monomorphic instances, we use the "extra" variables in  $\bar{\gamma}$  to remedy the situation. Specifically, the rules from Definition 5 ensure that the instance of map needed to map the polymorphic *flatten* function is typeable as follows:

$$\frac{\Gamma; \alpha, \beta \vdash List \alpha \quad \Gamma; \beta \vdash PTree \beta \quad \Gamma; \beta \vdash List \beta}{\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{List}^{F, G} : \text{Nat}^{\emptyset} (\text{Nat}^{\beta} (PTree \beta) (List \beta)) (\text{Nat}^{\beta} (List (PTree \beta)) (List (List \beta)))}$$

Similar remarks explain the appearance of  $\bar{\gamma}$  in the typing rules for in and fold.

### 3 INTERPRETING TYPES

We denote the category of sets and functions by  $\text{Set}$ . The category  $\text{Rel}$  has as its objects triples  $(A, B, R)$  where  $R$  is a relation between the objects  $A$  and  $B$  in  $\text{Set}$ , i.e., a subset of  $A \times B$ , and has as its morphisms from  $(A, B, R)$  to  $(A', B', R')$  pairs  $(f : A \rightarrow A', g : B \rightarrow B')$  of morphisms in  $\text{Set}$  such that  $(fa, gb) \in R'$  whenever  $(a, b) \in R$ . We write  $R : \text{Rel}(A, B)$  in place of  $(A, B, R)$  when convenient. If  $R : \text{Rel}(A, B)$  we write  $\pi_1 R$  and  $\pi_2 R$  for the *domain*  $A$  of  $R$  and the *codomain*  $B$  of  $R$ , respectively. If  $A : \text{Set}$ , then we write  $\text{Eq}_A = (A, A, \{(x, x) \mid x \in A\})$  for the *equality relation* on  $A$ .

The key idea underlying Reynolds' parametricity is to give each type  $\tau(\alpha)$  with one free variable  $\alpha$  both an *object interpretation*  $\tau_0$  taking sets to sets and a *relational interpretation*  $\tau_1$  taking relations  $R : \text{Rel}(A, B)$  to relations  $\tau_1(R) : \text{Rel}(\tau_0(A), \tau_0(B))$ , and to interpret each term  $t(\alpha, x) : \tau(\alpha)$  with one free term variable  $x : \sigma(\alpha)$  as a map  $t_0$  associating to each set  $A$  a function  $t_0(A) : \sigma_0(A) \rightarrow \tau_0(A)$ . These interpretations are to be given inductively on the structures of  $\tau$  and  $t$  in such a way that they imply two fundamental theorems. The first is an *Identity Extension Lemma*, which states that  $\tau_1(\text{Eq}_A) = \text{Eq}_{\tau_0(A)}$ , **and is the essential property that makes a model relationally parametric rather than just induced by a logical relation**. The second is an *Abstraction Theorem*, which states that, for any  $R : \text{Rel}(A, B)$ ,  $(t_0(A), t_0(B))$  is a morphism in  $\text{Rel}$  from  $(\sigma_0(A), \sigma_0(B), \sigma_1(R))$  to  $(\tau_0(A), \tau_0(B), \tau_1(R))$ . The Identity Extension Lemma is similar to the Abstraction Theorem except that it holds for *all* elements of a type's interpretation, not just those that are interpretations of terms. Similar results are expected to hold for types and terms with any number of free variables.

The key to proving the Identity Extension Lemma (Theorem 23) in our setting is a familiar "cutting down" of the interpretations of universally quantified types, such as our  $\text{Nat}$ -types, to include only the "parametric" elements. This requires that set interpretations of types are defined simultaneously with their relational interpretations. We give set interpretations for our types in Section 3.1 and give their relational interpretations in Section 3.2. While the set interpretations are relatively straightforward, their relation interpretations are less so, mainly because of the cocontinuity conditions we must impose to ensure that they are well-defined. We take some effort to develop conditions in Section 3.2, which separates Definitions 7 and 16 in space, but otherwise has no impact on the fact that they are given by mutual induction.

#### 3.1 Interpreting Types as Sets

We will interpret the types in our calculus as  $\omega$ -cocontinuous functors on locally finitely presentable categories [Adámek and Rosický 1994]. Since functor categories of locally finitely presentable categories are again locally finitely presentable, this will ensure, in particular, that the fixed points interpreting  $\mu$ -types in  $\text{Set}$  and  $\text{Rel}$  exist, and thus that both the set and relational interpretations of all of the types in Definition 2 are well-defined [Johann and Polonsky 2019]. To bootstrap this process, we interpret type variables themselves as  $\omega$ -cocontinuous functors in Definitions 6 and 14. If  $\mathcal{C}$  and  $\mathcal{D}$  are locally finitely presentable categories, we write  $[\mathcal{C}, \mathcal{D}]$  for the set of  $\omega$ -cocontinuous functors from  $\mathcal{C}$  to  $\mathcal{D}$ .

**DEFINITION 6.** A set environment maps each type variable in  $\mathbb{T}^k \cup \mathbb{F}^k$  to an element of  $[\text{Set}^k, \text{Set}]$ . A morphism  $f : \rho \rightarrow \rho'$  for set environments  $\rho$  and  $\rho'$  with  $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$  maps each type constructor variable  $\psi^k \in \mathbb{T}$  to the identity natural transformation on  $\rho\psi^k = \rho'\psi^k$  and each functorial variable  $\phi^k \in \mathbb{F}$  to a natural transformation from the  $k$ -ary functor  $\rho\phi^k$  on  $\text{Set}$  to the  $k$ -ary functor  $\rho'\phi^k$  on  $\text{Set}$ . Composition of morphisms on set environments is given componentwise, with the identity morphism mapping each set environment to itself. This gives a category of set environments and morphisms between them, which we denote  $\text{SetEnv}$ .

When convenient we identify a functor  $F : [\text{Set}^0, \text{Set}]$  with the set that is its codomain and consider a set environment to map a type variable of arity 0 to a set. If  $\bar{\alpha} = \{\alpha_1, \dots, \alpha_k\}$  and  $\bar{A} = \{A_1, \dots, A_k\}$ , then we write  $\rho[\bar{\alpha} := \bar{A}]$  for the set environment  $\rho'$  such that  $\rho'\alpha_i = A_i$  for  $i = 1, \dots, k$  and  $\rho'\alpha = \rho\alpha$  if  $\alpha \notin \{\alpha_1, \dots, \alpha_k\}$ . If  $\rho$  is a set environment we write  $\text{Eq}_\rho$  for the relation environment (see Definition 14) such that  $\text{Eq}_\rho v = \text{Eq}_{\rho v}$  for every type variable  $v$ . The relational interpretations appearing in the second clause of Definition 7 are given in full in Definition 16.

DEFINITION 7. The set interpretation  $\llbracket \cdot \rrbracket^{\text{Set}} : \mathcal{F} \rightarrow [\text{SetEnv}, \text{Set}]$  is defined by

$$\begin{aligned}
& \llbracket \Gamma; \emptyset \vdash v \rrbracket^{\text{Set}} \rho = \rho v \text{ if } v \in \mathbb{T}^0 \\
& \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho = \{ \eta : \lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}] \Rightarrow \lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}] \\
& \quad | \forall \bar{A}, \bar{B} : \text{Set}. \forall R : \text{Rel}(\bar{A}, \bar{B}). \\
& \quad (\eta_{\bar{A}}, \eta_{\bar{B}}) : \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\alpha} := \bar{R}] \rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\alpha} := \bar{R}] \} \\
& \llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\text{Set}} \rho = 0 \\
& \llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Set}} \rho = 1 \\
& \llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Set}} \rho = (\rho \phi) \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \\
& \llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} \rho + \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \\
& \llbracket \Gamma; \Phi \vdash \sigma \times \tau \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} \rho \times \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \\
& \llbracket \Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau} \rrbracket^{\text{Set}} \rho = (\mu T_{H, \rho}^{\text{Set}}) \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \\
& \text{ where } T_{H, \rho}^{\text{Set}} F = \lambda \bar{A}. \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} \rho[\phi := F][\bar{\alpha} := \bar{A}] \\
& \text{ and } T_{H, \rho}^{\text{Set}} \eta = \lambda \bar{A}. \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} \text{id}_\rho[\phi := \eta][\bar{\alpha} := \text{id}_{\bar{A}}]
\end{aligned}$$

The interpretations in Definition 7 respect weakening, i.e., a type and its weakenings all have the same set interpretations. The same holds for the actions of these interpretations on morphisms in Definition 8 below. Moreover, the interpretation of Nat types ensures that  $\llbracket \Gamma \vdash \sigma \rightarrow \tau \rrbracket^{\text{Set}} \rho = \llbracket \Gamma \vdash \sigma \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma \vdash \tau \rrbracket^{\text{Set}} \rho$ , as expected. If  $\rho$  is a set environment and  $\vdash \tau : \mathcal{F}$  then we may write  $\llbracket \vdash \tau \rrbracket^{\text{Set}}$  instead of  $\llbracket \vdash \tau \rrbracket^{\text{Set}} \rho$  since the environment is immaterial. We note that the second clause of Definition 7 does indeed define a set: local finite presentability of Set and  $\omega$ -cocontinuity of  $\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho$  ensure that  $\{ \eta : \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho \Rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho \}$  (which contains  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$ ) is a subset of  $\{ (\llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{S}])^{(\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{S}])} \mid \bar{S} = (S_1, \dots, S_{|\bar{\alpha}|}) \text{, and } S_i \text{ is a finite set for } i = 1, \dots, |\bar{\alpha}| \}$ . There are countably many choices for tuples  $\bar{S}$ , and each of these gives rise to a morphism from  $\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{S}]$  to  $\llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{S}]$ . But there are only Set-many choices of morphisms between these (or any) two objects because Set is locally small.

In order to make sense of the last clause in Definition 7, we need to know that, for each  $\rho \in \text{SetEnv}$ ,  $T_{H, \rho}^{\text{Set}}$  is an  $\omega$ -cocontinuous endofunctor on  $[\text{Set}^k, \text{Set}]$ , and thus admits a fixed point. Since  $T_{H, \rho}^{\text{Set}}$  is defined in terms of  $\llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}}$ , this means that interpretations of types must be such functors, which in turn means that the actions of set interpretations of types on objects and on morphisms in SetEnv are intertwined. Fortunately, we know from [Johann and Polonsky 2019] that, for every  $\Gamma; \bar{\alpha} \vdash \tau : \mathcal{F}$ ,  $\llbracket \Gamma; \bar{\alpha} \vdash \tau \rrbracket^{\text{Set}}$  is actually in  $[\text{Set}^k, \text{Set}]$  where  $k = |\bar{\alpha}|$ . This means that for each  $\llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}}$ , the corresponding operator  $T_H^{\text{Set}}$  can be extended to a functor from SetEnv to  $[[\text{Set}^k, \text{Set}], [\text{Set}^k, \text{Set}]]$ . The action of  $T_H^{\text{Set}}$  on an object  $\rho \in \text{SetEnv}$  is given by the higher-order functor  $T_{H, \rho}^{\text{Set}}$ , whose actions on objects (functors in  $[\text{Set}^k, \text{Set}]$ ) and morphisms (natural transformations between such functors) are given in Definition 7. Its action on a morphism

$f : \rho \rightarrow \rho'$  is the higher-order natural transformation  $T_{H,f}^{\text{Set}} : T_{H,\rho}^{\text{Set}} \rightarrow T_{H,\rho'}^{\text{Set}}$ , whose action on  $F : [\text{Set}^k, \text{Set}]$  is the natural transformation  $T_{H,\rho}^{\text{Set}} F : T_{H,\rho}^{\text{Set}} F \rightarrow T_{H,\rho'}^{\text{Set}} F$  whose component at  $\bar{A}$  is  $(T_{H,f}^{\text{Set}} F)_{\bar{A}} = \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} f[\phi := id_F][\bar{\alpha} := id_A]$ . The next definition uses the functor  $T_H^{\text{Set}}$  to define the actions of functors interpreting types on morphisms between set environments.

**DEFINITION 8.** Let  $f : \rho \rightarrow \rho'$  for set environments  $\rho$  and  $\rho'$  (so that  $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$ ). The action  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f$  of  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}$  on the morphism  $f$  is given as follows:

- If  $\Gamma, v; \emptyset \vdash v$  then  $\llbracket \Gamma, v; \emptyset \vdash v \rrbracket^{\text{Set}} f = id_{\rho v}$
- If  $\Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G$  then  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} f = id_{\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho}$
- If  $\Gamma; \Phi \vdash 0$  then  $\llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\text{Set}} f = id_0$
- If  $\Gamma; \Phi \vdash 1$  then  $\llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Set}} f = id_1$
- If  $\Gamma; \Phi \vdash \phi \bar{\tau}$  then  $\llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Set}} f : \llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Set}} \rho' = (\rho \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \rightarrow (\rho' \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}$  is defined by  $\llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Set}} f = (f \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \circ (\rho \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} = (\rho' \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \circ (f \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}$ . The latter equality holds because  $\rho \phi$  and  $\rho' \phi$  are functors and  $f \phi : \rho \phi \rightarrow \rho' \phi$  is a natural transformation, so the following naturality square commutes:

$$\begin{array}{ccc} (\rho \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} & \xrightarrow{(f \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}} & (\rho' \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \\ \downarrow (\rho \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f} & & \downarrow (\rho' \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f} \\ (\rho \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} & \xrightarrow{(f \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}} & (\rho' \phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \end{array} \quad (1)$$

- If  $\Gamma; \Phi \vdash \sigma + \tau$  then  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Set}} f$  is defined by  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Set}} f(\text{inl } x) = \text{inl } (\llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} f x)$  and  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Set}} f(\text{inr } y) = \text{inr } (\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f y)$ .
- If  $\Gamma; \Phi \vdash \sigma \times \tau$  then  $\llbracket \Gamma; \Phi \vdash \sigma \times \tau \rrbracket^{\text{Set}} f = \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} f \times \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f$ .
- If  $\Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau}$  then  $\llbracket \Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau} \rrbracket^{\text{Set}} f : \llbracket \Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau} \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau} \rrbracket^{\text{Set}} \rho' = (\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \rightarrow (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}$  is defined by  $(\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \rightarrow (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}$  is defined by  $(\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} = (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \circ (\mu T_{H,f}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}$ . The latter equality holds because  $\mu T_{H,\rho}^{\text{Set}}$  and  $\mu T_{H,\rho'}^{\text{Set}}$  are functors and  $\mu T_{H,f}^{\text{Set}} : \mu T_{H,\rho}^{\text{Set}} \rightarrow \mu T_{H,\rho'}^{\text{Set}}$  is a natural transformation, so the following naturality square commutes:

$$\begin{array}{ccc} (\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} & \xrightarrow{(\mu T_{H,f}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}} & (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \\ \downarrow (\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f} & & \downarrow (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f} \\ (\mu T_{H,\rho}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} & \xrightarrow{(\mu T_{H,f}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}} & (\mu T_{H,\rho'}^{\text{Set}}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \end{array} \quad (2)$$

### 3.2 Interpreting Types as Relations

**DEFINITION 9.** A  $k$ -ary relation transformer  $F$  is a triple  $(F^1, F^2, F^*)$ , where  $F^1, F^2 : [\text{Set}^k, \text{Set}]$  are functors,  $F^* : [\text{Rel}^k, \text{Rel}]$  is a functor, if  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , then  $F^* \bar{R} : \text{Rel}(F^1 \bar{A}, F^2 \bar{B})$ , and if  $(\alpha_1, \beta_1) \in \text{Hom}_{\text{Rel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\text{Rel}}(R_k, S_k)$  then  $F^*(\alpha, \beta) = (F^1 \bar{\alpha}, F^2 \bar{\beta})$ . We define  $F \bar{R}$  to be  $F^* \bar{R}$  and  $F(\alpha, \beta)$  to be  $F^*(\alpha, \beta)$ .

The last clause of Definition 9 expands to: if  $(a, b) \in R$  implies  $(\alpha a, \beta b) \in S$  then  $(c, d) \in F^* \bar{R}$  implies  $(F^1 \bar{\alpha} c, F^2 \bar{\beta} d) \in F^* \bar{S}$ . When convenient we identify a 0-ary relation transformer  $(A, B, R)$

with  $R : \text{Rel}(A, B)$ . We may also write  $\pi_1 F$  for  $F^1$  and  $\pi_2 F$  for  $F^2$ . We extend these conventions to relation environments, introduced in Definition 14 below, in the obvious way.

DEFINITION 10. The category  $RT_k$  of  $k$ -ary relation transformers is given by the following data:

- An object of  $RT_k$  is a relation transformer.
- A morphism  $\delta : (G^1, G^2, G^*) \rightarrow (H^1, H^2, H^*)$  in  $RT_k$  is a pair of natural transformations  $(\delta^1, \delta^2)$  where  $\delta^1 : G^1 \rightarrow H^1$ ,  $\delta^2 : G^2 \rightarrow H^2$  such that, for all  $R : \text{Rel}(A, B)$ , if  $(x, y) \in G^* \bar{R}$  then  $(\delta^1_A x, \delta^2_B y) \in H^* \bar{R}$ .
- Identity morphisms and composition are inherited from the category of functors on  $\text{Set}$ .

DEFINITION 11. An endofunctor  $H$  on  $RT_k$  is a triple  $H = (H^1, H^2, H^*)$ , where

- $H^1$  and  $H^2$  are functors from  $[\text{Set}^k, \text{Set}]$  to  $[\text{Set}^k, \text{Set}]$
- $H^*$  is a functor from  $RT_k$  to  $[\text{Rel}^k, \text{Rel}]$
- for all  $R : \text{Rel}(A, B)$ ,  $\pi_1((H^*(\delta^1, \delta^2))_{\bar{R}}) = (H^1 \delta^1)_{\bar{A}}$  and  $\pi_2((H^*(\delta^1, \delta^2))_{\bar{R}}) = (H^2 \delta^2)_{\bar{B}}$
- The action of  $H$  on objects is given by  $H(F^1, F^2, F^*) = (H^1 F^1, H^2 F^2, H^*(F^1, F^2, F^*))$
- The action of  $H$  on morphisms is given by  $H(\delta^1, \delta^2) = (H^1 \delta^1, H^2 \delta^2)$  for  $(\delta^1, \delta^2) : (F^1, F^2, F^*) \rightarrow (G^1, G^2, G^*)$

Since the results of applying an endofunctor  $H$  to  $k$ -ary relation transformers and morphisms between them must again be  $k$ -ary relation transformers and morphisms between them, respectively, Definition 11 implicitly requires that the following three conditions hold: i) if  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , then  $H^*(F^1, F^2, F^*)_{\bar{R}} : \text{Rel}(H^1 F^1 \bar{A}, H^2 F^2 \bar{B})$ ; ii) if  $(\alpha_1, \beta_1) \in \text{Hom}_{\text{Rel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\text{Rel}}(R_k, S_k)$ , then  $H^*(F^1, F^2, F^*)_{\bar{R}}(\alpha, \beta) = (H^1 F^1 \bar{\alpha}, H^2 F^2 \bar{\beta})$ ; and  $(\delta^1, \delta^2) : (F^1, F^2, F^*) \rightarrow (G^1, G^2, G^*)$  and  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , then  $((H^1 \delta^1)_{\bar{A}} x, (H^2 \delta^2)_{\bar{B}} y) \in H^*(G^1, G^2, G^*)_{\bar{R}}$  whenever  $(x, y) \in H^*(F^1, F^2, F^*)_{\bar{R}}$ . Note, however, that this last condition is automatically satisfied because it is implied by the third bullet point of Definition 11.

DEFINITION 12. If  $H$  and  $K$  are endofunctors on  $RT_k$ , then a natural transformation  $\sigma : H \rightarrow K$  is a pair  $\sigma = (\sigma^1, \sigma^2)$ , where  $\sigma^1 : H^1 \rightarrow K^1$  and  $\sigma^2 : H^2 \rightarrow K^2$  are natural transformations between endofunctors on  $[\text{Set}^k, \text{Set}]$  and the component of  $\sigma$  at  $F \in RT_k$  is given by  $\sigma_F = (\sigma^1_{F^1}, \sigma^2_{F^2})$ .

Definition 12 entails that  $\sigma^i_{F^i}$  must be natural in  $F^i : [\text{Set}^k, \text{Set}]$ , and, for every  $F$ , both  $(\sigma^1_{F^1})_{\bar{A}}$  and  $(\sigma^2_{F^2})_{\bar{B}}$  must be natural in  $\bar{A}$ . Moreover, since the results of applying  $\sigma$  to  $k$ -ary relation transformers must be morphisms of  $k$ -ary relation transformers, Definition 12 implicitly requires that  $(\sigma_F)_{\bar{R}} = ((\sigma^1_{F^1})_{\bar{A}}, (\sigma^2_{F^2})_{\bar{B}})$  is a morphism in  $\text{Rel}$  for any  $k$ -tuple of relations  $\bar{R} : \text{Rel}(A, B)$ , i.e., that if  $(x, y) \in H^* \bar{R}$ , then  $((\sigma^1_{F^1})_{\bar{A}} x, (\sigma^2_{F^2})_{\bar{B}} y) \in K^* \bar{R}$ .

**More context? Weave cocontinuity requirement for fixed points through text better.** Critically, we can compute  $\omega$ -directed colimits in  $RT_k$ : indeed, if  $\mathcal{D}$  is an  $\omega$ -directed set, then  $\varinjlim_{d \in \mathcal{D}} (F^1_d, F^2_d, F^*_d) = (\varinjlim_{d \in \mathcal{D}} F^1_d, \varinjlim_{d \in \mathcal{D}} F^2_d, \varinjlim_{d \in \mathcal{D}} F^*_d)$ . We then define an endofunctor  $T = (T^1, T^2, T^*)$  on  $RT_k$  to be  $\omega$ -cocontinuous if  $T^1$  and  $T^2$  are  $\omega$ -cocontinuous endofunctors on  $[\text{Set}^k, \text{Set}]$  and  $T^*$  is an  $\omega$ -cocontinuous functor from  $RT_k$  to  $[\text{Rel}^k, \text{Rel}]$ , i.e., is in  $[RT_k, [\text{Rel}^k, \text{Rel}]]$ . Now, for any  $k$  and  $R : \text{Rel}(A, B)$ , let  $K^{\text{Rel}}_R$  be the constantly  $R$ -valued functor from  $\text{Rel}^k$  to  $\text{Rel}$ , and for any  $k$  and set  $A$ , let  $K^{\text{Set}}_A$  be the constantly  $A$ -valued functor from  $\text{Set}^k$  to  $\text{Set}$ , and let  $0$  denote either the initial object of  $\text{Set}$  or the initial object of  $\text{Rel}$ , as appropriate. Observing that, for every  $k$ ,  $K^{\text{Set}}_0$  is initial in  $[\text{Set}^k, \text{Set}]$ , and  $K^{\text{Rel}}_0$  is initial in  $[\text{Rel}^k, \text{Rel}]$ , we have that, for each  $k$ ,  $K_0 = (K^{\text{Set}}_0, K^{\text{Set}}_0, K^{\text{Rel}}_0)$  is initial in  $RT_k$ . Thus, if  $T = (T^1, T^2, T^*) : RT_k \rightarrow RT_k$  is an endofunctor on  $RT_k$  then we can define

the relation transformer  $\mu T$  to be  $\lim_{\rightarrow n \in \mathbb{N}} T^n K_0$ . It is not hard to see that  $\mu T$  is given explicitly as

$$\mu T = (\mu T^1, \mu T^2, \lim_{\rightarrow n \in \mathbb{N}} (T^n K_0)^*) \quad (3)$$

and that, as our notation suggests, it really is a fixpoint for  $T$  if  $T$  is  $\omega$ -cocontinuous:

LEMMA 13. *For any  $T : [RT_k, RT_k]$ ,  $\mu T \cong T(\mu T)$ .*

The isomorphism is given by the morphisms  $(in_1, in_2) : T(\mu T) \rightarrow \mu T$  and  $(in_1^{-1}, in_2^{-1}) : \mu T \rightarrow T(\mu T)$  in  $RT_k$ . The latter is always a morphism in  $RT_k$ , but the former need not be if  $T$  is not  $\omega$ -cocontinuous.

**CHECK!** It is worth noting that the third component in Equation (3) is the colimit in  $[\text{Rel}^k, \text{Rel}]$  of third components of relation transformers, rather than a fixpoint of an endofunctor on  $[\text{Rel}^k, \text{Rel}]$ . That there is an asymmetry between the first two components of  $\mu T$  and its third is an important conceptual observation, and reflects the fact that the third component of an endofunctor on  $RT_k$  need not be a functor on all of  $[\text{Rel}^k, \text{Rel}]$ . For example, although we can define  $T_{H,\rho} F$  for a relation transformer  $F$  in Definition 16 below, it is not clear how we could define it for  $F : [\text{Rel}^k, \text{Rel}]$ .

DEFINITION 14. *A relation environment maps each type variable in  $\mathbb{T}^k \cup \mathbb{F}^k$  to a  $k$ -ary relation transformer. A morphism  $f : \rho \rightarrow \rho'$  for relation environments  $\rho$  and  $\rho'$  with  $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$  maps each type constructor variable  $\psi^k \in \mathbb{T}$  to the identity morphism on  $\rho\psi^k = \rho'\psi^k$  and each functorial variable  $\phi^k \in \mathbb{F}$  to a morphism from the  $k$ -ary relation transformer  $\rho\phi$  to the  $k$ -ary relation transformer  $\rho'\phi$ . Composition of morphisms on relation environments is given componentwise, with the identity morphism mapping each relation environment to itself. This gives a category of relation environments and morphisms between them, which we denote  $\text{RelEnv}$ .*

When convenient we identify a 0-ary relation transformer with the relation (transformer) that is its codomain and consider a relation environment to map a type variable of arity 0 to a relation. We write  $\rho[\alpha := \bar{R}]$  for the relation environment  $\rho'$  such that  $\rho'\alpha_i = R_i$  for  $i = 1, \dots, k$  and  $\rho'\alpha = \rho\alpha$  if  $\alpha \notin \{\alpha_1, \dots, \alpha_k\}$ . If  $\rho$  is a relation environment, we write  $\pi_1\rho$  and  $\pi_2\rho$  for the set environments mapping each type variable  $\phi$  to the functors  $(\rho\phi)^1$  and  $(\rho\phi)^2$ , respectively.

We define, for each  $k$ , the notion of an  $\omega$ -cocontinuous functor from  $\text{RelEnv}$  to  $RT_k$ :

DEFINITION 15. *A functor  $H : [\text{RelEnv}, RT_k]$  is a triple  $H = (H^1, H^2, H^*)$ , where*

- $H^1$  and  $H^2$  are objects in  $[\text{SetEnv}, [\text{Set}^k, \text{Set}]]$
- $H^*$  is an object in  $[\text{RelEnv}, [\text{Rel}^k, \text{Rel}]]$
- for all  $\bar{R} : \text{Rel}(A, B)$  and morphisms  $f$  in  $\text{RelEnv}$ ,  $\pi_1(H^* f \bar{R}) = H^1(\pi_1 f) \bar{A}$  and  $\pi_2(H^* f \bar{R}) = H^2(\pi_2 f) \bar{B}$
- The action of  $H$  on  $\rho$  in  $\text{RelEnv}$  is given by  $H\rho = (H^1(\pi_1\rho), H^2(\pi_2\rho), H^*\rho)$
- The action of  $H$  on morphisms  $f : \rho \rightarrow \rho'$  in  $\text{RelEnv}$  is given by  $Hf = (H^1(\pi_1 f), H^2(\pi_2 f))$

Spelling out the last two bullet points above gives the following analogues of the three conditions immediately following Definition 11: i) if  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , then  $H^*\rho \bar{R} : \text{Rel}(H^1(\pi_1\rho) \bar{A}, H^2(\pi_2\rho) \bar{B})$ ; ii) if  $(\alpha_1, \beta_1) \in \text{Hom}_{\text{Rel}}(R_1, S_1), \dots, (\alpha_k, \beta_k) \in \text{Hom}_{\text{Rel}}(R_k, S_k)$ , then  $H^*\rho(\bar{\alpha}, \bar{\beta}) = (H^1(\pi_1\rho) \bar{\alpha}, H^2(\pi_2\rho) \bar{\beta})$ ; and iii) if  $f : \rho \rightarrow \rho'$  and  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , then  $(H^1(\pi_1 f) \bar{A} x, H^2(\pi_2 f) \bar{B} y) \in H^*\rho' \bar{R}$  whenever  $(x, y) \in H^*\rho \bar{R}$ . As before, the last condition is automatically satisfied because it is implied by the third bullet point of Definition 15.

Considering  $\text{RelEnv}$  as a product  $\prod_{\phi^k \in \mathbb{T} \cup \mathbb{F}} RT_k$ , we extend the computation of  $\omega$ -directed colimits in  $RT_k$  to compute colimits in  $\text{RelEnv}$  componentwise. We similarly extend the notion of an  $\omega$ -cocontinuous endofunctor on  $RT_k$  componentwise to give a notion of  $\omega$ -cocontinuity for functors from  $\text{RelEnv}$  to  $RT_k$ . Recalling from the start of this subsection that Definition 16 is given mutually inductively with Definition 7 we can, at last, define:



DEFINITION 16. The relational interpretation  $\llbracket \cdot \rrbracket^{\text{Rel}} : \mathcal{F} \rightarrow [\text{RelEnv}, \text{Rel}]$  is defined by

$$\begin{aligned}
 \llbracket \Gamma; \emptyset \vdash v \rrbracket^{\text{Rel}} \rho &= \rho v \text{ if } v \in \mathbb{T}^0 \\
 \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Rel}} \rho &= \{ \eta : \lambda \bar{R}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Rel}} \rho[\bar{\alpha} := \bar{R}] \Rightarrow \lambda \bar{R}. \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Rel}} \rho[\bar{\alpha} := \bar{R}] \} \\
 &= \{ (t, t') \in \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}(\pi_1 \rho) \times \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}(\pi_2 \rho) \mid \\
 &\quad \forall R_1 : \text{Rel}(A_1, B_1) \dots R_k : \text{Rel}(A_k, B_k). \\
 &\quad (t_{\bar{A}}, t'_{\bar{B}}) \in (\llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Rel}} \rho[\bar{\alpha} := \bar{R}])^{\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Rel}} \rho[\bar{\alpha} := \bar{R}]} \} \\
 \llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\text{Rel}} \rho &= 0 \\
 \llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Rel}} \rho &= 1 \\
 \llbracket \Gamma; \Phi \vdash \phi \bar{\tau} \rrbracket^{\text{Rel}} \rho &= (\rho \phi) \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \\
 \llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Rel}} \rho &= \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} \rho + \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \\
 \llbracket \Gamma; \Phi \vdash \sigma \times \tau \rrbracket^{\text{Rel}} \rho &= \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} \rho \times \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \\
 \llbracket \Gamma; \Phi \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\tau} \rrbracket^{\text{Rel}} \rho &= (\mu T_{H, \rho}) \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \\
 &\text{where } T_{H, \rho} = (T_{H, \pi_1 \rho}^{\text{Set}}, T_{H, \pi_2 \rho}^{\text{Set}}, T_{H, \rho}^{\text{Rel}}) \\
 &\text{and } T_{H, \rho}^{\text{Rel}} F = \lambda \bar{R}. \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}} \rho[\phi := F][\bar{\alpha} := \bar{R}] \\
 &\text{and } T_{H, \rho}^{\text{Rel}} \delta = \lambda \bar{R}. \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}} id_{\rho}[\phi := \delta][\bar{\alpha} := id_{\bar{R}}]
 \end{aligned}$$

The interpretations in Definition 16, as well as in Definition 17 below, respect weakening. Definition 16 also ensures that  $\llbracket \Gamma \vdash \sigma \rightarrow \tau \rrbracket^{\text{Rel}} \rho = \llbracket \Gamma \vdash \sigma \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \Gamma \vdash \tau \rrbracket^{\text{Rel}} \rho$ . If  $\rho$  is a relational environment and  $\vdash \tau : \mathcal{F}$ , then we write  $\llbracket \vdash \tau \rrbracket^{\text{Rel}}$  instead of  $\llbracket \vdash \tau \rrbracket^{\text{Rel}} \rho$  as for set interpretations. For the last clause in Definition 16 to be well-defined, we need to know that  $T_{\rho}$  is an  $\omega$ -cocontinuous endofunctor on  $RT$  so that, by Lemma 13, it admits a fixed point. Since  $T_{\rho}$  is defined in terms of  $\llbracket \Gamma; \Phi, \phi^k, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}}$ , this means that relational interpretations of types must be  $\omega$ -cocontinuous functors from  $\text{RelEnv}$  to  $RT_0$ , which in turn entails that the actions of relational interpretations of types on objects and on morphisms in  $\text{RelEnv}$  are intertwined. As for set interpretations, we know from [Johann and Polonsky 2019] that, for every  $\Gamma; \bar{\alpha} \vdash \tau : \mathcal{F}$ ,  $\llbracket \Gamma; \bar{\alpha} \vdash \tau \rrbracket^{\text{Set}}$  is actually in  $[\text{Rel}^k, \text{Rel}]$  where  $k = |\bar{\alpha}|$ . We first define the actions of each of these functors on morphisms between environments in Definition 17, and then argue that the functors given by Definitions 16 and 17 are well-defined and have the required properties. To do this, we extend  $T_H$  to a functor from  $\text{RelEnv}$  to  $[[\text{Rel}^k, \text{Rel}], [\text{Rel}^k, \text{Rel}]]$ . Its action on an object  $\rho \in \text{RelEnv}$  is given by the higher-order functor  $T_{H, \rho}^{\text{Rel}}$  whose actions on objects and morphisms are given in Definition 17. Its action on a morphism  $f : \rho \rightarrow \rho'$  is the higher-order natural transformation  $T_{H, f} : T_{H, \rho} \rightarrow T_{H, \rho'}$  whose action on any  $F : [\text{Rel}^k, \text{Rel}]$  is the natural transformation  $T_{H, f} F : T_{H, \rho} F \rightarrow T_{H, \rho'} F$  whose component at  $\bar{R}$  is  $(T_{H, f} F)_{\bar{R}} = \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}} f[\phi := id_F][\bar{\alpha} := id_{\bar{R}}]$ . The next definition uses the functor  $T_H$  to define the actions of functors interpreting types on morphisms between relation environments.

DEFINITION 17. Let  $f : \rho \rightarrow \rho'$  for relation environments  $\rho$  and  $\rho'$  (so that  $\rho|_{\mathbb{T}} = \rho'|_{\mathbb{T}}$ ). The action  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f$  of  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}}$  on the morphism  $f$  is given as follows:

- If  $\Gamma, v; \emptyset \vdash v$  then  $\llbracket \Gamma, v; \emptyset \vdash v \rrbracket^{\text{Rel}} f = id_{\rho v}$
- If  $\Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G$ , then  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Rel}} f = id_{\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Rel}} \rho}$
- If  $\Gamma; \Phi \vdash 0$  then  $\llbracket \Gamma; \Phi \vdash 0 \rrbracket^{\text{Rel}} f = id_0$
- If  $\Gamma; \Phi \vdash 1$  then  $\llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Rel}} f = id_1$



- If  $\Gamma; \Phi \vdash \phi\bar{\tau}$ , then  $\llbracket \Gamma; \Phi \vdash \phi\bar{\tau} \rrbracket^{\text{Rel}} f : \llbracket \Gamma; \Phi \vdash \phi\bar{\tau} \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash \phi\bar{\tau} \rrbracket^{\text{Rel}} \rho' = (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho} \rightarrow (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho'}$  is defined by  $\llbracket \Gamma; \Phi \vdash \phi\tau A \rrbracket^{\text{Rel}} f = (f\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho'} \circ (\rho\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f} = (\rho'\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f} \circ (f\phi) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho}$
- If  $\Gamma; \Phi \vdash \sigma + \tau$  then  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Rel}} f$  is defined by  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Rel}} f(\text{inl } x) = \text{inl } (\llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} f x)$  and  $\llbracket \Gamma; \Phi \vdash \sigma + \tau \rrbracket^{\text{Rel}} f(\text{inr } y) = \text{inr } (\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f y)$
- If  $\Gamma; \Phi \vdash \sigma \times \tau$  then  $\llbracket \Gamma; \Phi \vdash \sigma \times \tau \rrbracket^{\text{Rel}} f = \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} f \times \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f$
- If  $\Gamma; \Phi \vdash (\mu\phi^k. \lambda\bar{\alpha}. H)\bar{\tau}$  then  $\llbracket \Gamma; \Phi \vdash (\mu\phi. \lambda\bar{\alpha}. H)\bar{\tau} \rrbracket^{\text{Rel}} f = (\mu T_{H,f}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho' \circ (\mu T_{H,\rho}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f}} = (\mu T_{H,\rho'}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} f} \circ (\mu T_{H,f}) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho}$

To see that the functors given by Definitions 16 and 17 are well-defined we must show that, for every  $H$ ,  $T_{H,\rho} F$  is a relation transformer for any relation transformer  $F$ , and that  $T_{H,f} F : T_{H,\rho} F \rightarrow T_{H,\rho'} F$  is a morphism of relation transformers for every relation transformer  $F$  and every morphism  $f : \rho \rightarrow \rho'$  in  $\text{RelEnv}$ . This is an immediate consequence of

LEMMA 18. For every  $\Gamma; \Phi \vdash \tau$ ,  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket = (\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}, \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}, \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}}) \in [\text{RelEnv}, RT_0]$ .

The proof is a straightforward induction on the structure of  $\tau$ , using an appropriate result from [Johann and Polonsky 2019] to deduce  $\omega$ -cocontinuity of  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket$  in each case.

We can also prove by induction that our interpretations of types interact well with demotion of functorial variables and substitution. Indeed, we have

LEMMA 19. Let  $\rho, \rho' : \text{SetEnv}$  be such that  $\rho\phi = \rho'\psi = \rho'\phi = \rho'\psi$ , and let  $f : \rho \rightarrow \rho'$  be a morphism of set environments such that  $f\phi = f\psi = \text{id}_{\rho\phi}$ . If  $\Gamma; \Phi, \phi^k \vdash F : \mathcal{F}$ ,  $\Gamma; \Phi, \bar{\alpha} \vdash G : \mathcal{G}$ ,  $\Phi, \alpha_1 \dots \alpha_k \vdash H$ , and  $\Gamma; \Phi \vdash \tau$ , then

$$\llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \psi; \Phi \vdash F[\phi := \psi] \rrbracket^{\text{Set}} \rho \quad (4)$$

$$\llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\text{Set}} f = \llbracket \Gamma; \psi; \Phi \vdash F[\phi := \psi] \rrbracket^{\text{Set}} f \quad (5)$$

$$\llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \Phi, \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}] \quad (6)$$

$$\llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f = \llbracket \Gamma; \Phi, \bar{\alpha} \vdash G \rrbracket^{\text{Set}} f[\bar{\alpha} := \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f}] \quad (7)$$

$$\llbracket \Gamma; \Phi \vdash F[\phi := H] \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\text{Set}} \rho[\phi := \lambda\bar{A}. \llbracket \Gamma; \Phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}]] \quad (8)$$

$$\llbracket \Gamma; \Phi \vdash F[\phi := H] \rrbracket^{\text{Set}} f = \llbracket \Gamma; \Phi, \phi \vdash F \rrbracket^{\text{Set}} f[\phi := \lambda\bar{A}. \llbracket \Gamma; \Phi, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} f[\bar{\alpha} := \overline{\text{id}_{\bar{A}}}}] \quad (9)$$

Analogous identities hold for relational interpretations.

#### 4 THE IDENTITY EXTENSION LEMMA

Compare against Sec 3.3, 3.4 of Bob. R&R refl graphs are our elts of  $RT_0$ . Bob defines equality relations starting from refl graphs by induction on kind structure. We do it as a special case of graphs.

The standard definition of the *graph* of a morphism  $f : A \rightarrow B$  in  $\text{Set}$  is the relation  $\langle f \rangle : \text{Rel}(A, B)$  defined by  $(x, y) \in \langle f \rangle$  iff  $fx = y$ . This definition naturally generalizes to associate to each natural transformation between  $k$ -ary functors on  $\text{Set}$  a  $k$ -ary relation transformer as follows:

DEFINITION 20. If  $F, G : \text{Set}^k \rightarrow \text{Set}$  and  $\alpha : F \rightarrow G$  is a natural transformation, then the functor  $\langle \alpha \rangle^* : \text{Rel}^k \rightarrow \text{Rel}$  is defined as follows. Given  $R_1 : \text{Rel}(A_1, B_1), \dots, R_k : \text{Rel}(A_k, B_k)$ , let  $\iota_{R_i} : R_i \hookrightarrow A_i \times B_i$ , for  $i = 1, \dots, k$ , be the inclusion of  $R_i$  as a subset of  $A_i \times B_i$ , let  $h_{A \times B}$  be the unique

morphism making the diagram

$$\begin{array}{ccccc}
 F\bar{A} & \xleftarrow{F\pi_1} & F(\bar{A} \times \bar{B}) & \xrightarrow{F\pi_2} & F\bar{B} & \xrightarrow{\alpha_{\bar{B}}} & G\bar{B} \\
 & \searrow \pi_1 & \downarrow h_{\bar{A} \times \bar{B}} & \nearrow \pi_2 & & & \\
 & & F\bar{A} \times G\bar{B} & & & & 
 \end{array}$$

commute, and let  $h_{\bar{R}} : F\bar{R} \rightarrow F\bar{A} \times G\bar{B}$  be  $h_{\bar{A} \times \bar{B}} \circ F\bar{i}_{\bar{R}}$ . Further, let  $\alpha^{\wedge} \bar{R}$  be the subobject through which  $h_{\bar{R}}$  is factorized by the mono-epi factorization system in  $\mathbf{Set}$ , as shown in the following diagram:

$$\begin{array}{ccc}
 F\bar{R} & \xrightarrow{h_{\bar{R}}} & F\bar{A} \times G\bar{B} \\
 \searrow q_{\alpha^{\wedge} \bar{R}} & & \nearrow \iota_{\alpha^{\wedge} \bar{R}} \\
 & \alpha^{\wedge} \bar{R} & 
 \end{array}$$

Then  $\alpha^{\wedge} \bar{R} : \mathbf{Rel}(F\bar{A}, G\bar{B})$  by construction, so the action of  $\langle \alpha \rangle^*$  on objects can be given by  $\langle \alpha \rangle^*(A, B, R) = (F\bar{A}, G\bar{B}, \iota_{\alpha^{\wedge} \bar{R}} \alpha^{\wedge} \bar{R})$ . Its action on morphisms is given by  $\langle \alpha \rangle^*(\beta, \beta') = (F\bar{\beta}, G\bar{\beta}')$ .

The data in Definition 20 yield a relation transformer  $\langle \alpha \rangle = (F, G, \langle \alpha \rangle^*)$ , called the *graph relation transformer* for  $\alpha$ .

LEMMA 21. If  $F, G : [\mathbf{Set}^k, \mathbf{Set}]$  and  $\alpha : F \rightarrow G$  is a natural transformation, then  $\langle \alpha \rangle$  is in  $\mathbf{RT}_k$ .

PROOF. Clearly,  $\langle \alpha \rangle^*$  is  $\omega$ -cocontinuous, so  $\langle \alpha \rangle^* : [\mathbf{Rel}^k, \mathbf{Rel}]$ . Now, suppose  $\bar{R} : \mathbf{Rel}(A, B), \bar{S} : \mathbf{Rel}(C, D)$ , and  $(\bar{\beta}, \bar{\beta}') : \bar{R} \rightarrow \bar{S}$ . We want to show that there exists a morphism  $\epsilon : \alpha^{\wedge} \bar{R} \rightarrow \alpha^{\wedge} \bar{S}$  such that

$$\begin{array}{ccc}
 \alpha^{\wedge} \bar{R} & \xrightarrow{\iota_{\alpha^{\wedge} \bar{R}}} & F\bar{A} \times G\bar{B} \\
 \epsilon \downarrow & & \downarrow F\bar{\beta} \times G\bar{\beta}' \\
 \alpha^{\wedge} \bar{S} & \xrightarrow{\iota_{\alpha^{\wedge} \bar{S}}} & F\bar{C} \times G\bar{D}
 \end{array}$$

commutes. Since  $(\bar{\beta}, \bar{\beta}') : \bar{R} \rightarrow \bar{S}$ , there exist  $\bar{\gamma} : \bar{R} \rightarrow \bar{S}$  such that each diagram

$$\begin{array}{ccc}
 R_i & \xrightarrow{\iota_{R_i}} & A_i \times B_i \\
 \gamma_i \downarrow & & \downarrow \beta_i \times \beta'_i \\
 S_i & \xrightarrow{\iota_{S_i}} & C_i \times D_i
 \end{array}$$

commutes. Moreover, since both  $h_{\bar{C} \times \bar{D}} \circ F(\bar{\beta} \times \bar{\beta}')$  and  $(F\bar{\beta} \times G\bar{\beta}') \circ h_{\bar{A} \times \bar{B}}$  make

$$\begin{array}{ccccc}
 F\bar{C} & \xleftarrow{\pi_1} & F\bar{C} \times F\bar{D} & \xrightarrow{\pi_2} & F\bar{D} & \xrightarrow{\alpha_{\bar{D}}} & G\bar{D} \\
 & \searrow F\pi_1 \circ F(\bar{\beta} \times \bar{\beta}') & \uparrow \exists! & \nearrow \alpha_{\bar{D}} \circ F\pi_2 \circ F(\bar{\beta} \times \bar{\beta}') & & & \\
 & & F(\bar{A} \times \bar{B}) & & & & 
 \end{array}$$

commute, they must be equal. We therefore get that the right-hand square below commutes, and thus that the entire following diagram does as well:

$$\begin{array}{ccccc}
 & & h_{\overline{R}} & & \\
 & \nearrow & & \searrow & \\
 F\overline{R} & \xrightarrow{F\overline{I}_R} & F(\overline{A \times B}) & \xrightarrow{h_{\overline{A \times B}}} & F\overline{A} \times G\overline{B} \\
 F\overline{Y} \downarrow & & \downarrow F(\overline{\beta \times \beta'}) & & \downarrow F\overline{\beta} \times F\overline{\beta'} \\
 F\overline{S} & \xrightarrow{F\overline{I}_S} & F(\overline{C \times D}) & \xrightarrow{h_{\overline{C \times D}}} & F\overline{C} \times G\overline{D} \\
 & \nwarrow & & \nearrow & \\
 & & h_{\overline{S}} & & 
 \end{array}$$

Finally, by the left-lifting property of  $q_{F \wedge \overline{R}}$  with respect to  $\iota_{F \wedge \overline{S}}$  given by the epi-mono factorization system, there exists an  $\epsilon$  such that the following diagram commutes:

$$\begin{array}{ccccc}
 F\overline{R} & \xrightarrow{q_{\alpha \wedge \overline{R}}} & \alpha \wedge \overline{R} & \xrightarrow{\iota_{\alpha \wedge \overline{R}}} & F\overline{A} \times G\overline{B} \\
 F\overline{Y} \downarrow & & \downarrow \epsilon & & \downarrow F\overline{\beta} \times G\overline{\beta'} \\
 F\overline{S} & \xrightarrow{q_{\alpha \wedge \overline{S}}} & \alpha \wedge \overline{S} & \xrightarrow{\iota_{\alpha \wedge \overline{S}}} & F\overline{C} \times G\overline{D}
 \end{array}$$

□

It is not hard to see that if  $f : A \rightarrow B$  is a morphism in  $\mathbf{Set}$  then the standard definition of its graph  $\langle f \rangle$  in  $\mathbf{Rel}$  coincides with its definition as the graph relation transformer for  $f$  as a natural transformation between the 0-ary functors  $A$  and  $B$ . Graph relation transformers are thus a reasonable generalization of graphs in  $\mathbf{Rel}$ .

The action of a graph relation transformer on a graph relation is easily computed:

LEMMA 22. *If  $\alpha : F \rightarrow G$  is a morphism in  $[\mathbf{Set}^k, \mathbf{Set}]$  and  $f_1 : A_1 \rightarrow B_1, \dots, f_k : A_k \rightarrow B_k$ , then  $\langle \alpha \rangle^* \langle \overline{f} \rangle = \langle G\overline{f} \circ \alpha_{\overline{A}} \rangle = \langle \alpha_{\overline{B}} \circ F\overline{f} \rangle$ .*

PROOF. Since  $h_{\overline{A \times B}}$  is the unique morphism making the bottom triangle of this diagram commute

$$\begin{array}{ccccc}
 & & F\overline{A} & & \\
 & \nearrow & \downarrow F\langle \overline{id_A}, f \rangle & \searrow & \\
 F\overline{A} & \xleftarrow{F\overline{\pi_1}} & F(\overline{A \times B}) & \xrightarrow{F\overline{\pi_2}} & F\overline{B} \xrightarrow{\alpha_{\overline{B}}} G\overline{B} \\
 & \nwarrow \pi_1 & \downarrow h_{\overline{A \times B}} & \nearrow \pi_2 & \\
 & & F\overline{A} \times G\overline{B} & & 
 \end{array}$$

and since  $h_{\langle \overline{f} \rangle} = h_{\overline{A \times B}} \circ F\overline{\iota_{\langle f \rangle}} = h_{\overline{A \times B}} \circ F\overline{\iota_{\langle id_A, f \rangle}}$ , the universal property of the product

$$\begin{array}{ccccc}
 F\overline{A} & \xleftarrow{\pi_1} & F\overline{A} \times G\overline{B} & \xrightarrow{\pi_2} & G\overline{B} \\
 & \nwarrow & \uparrow \exists! & & \uparrow \alpha_{\overline{B}} \\
 & & F\overline{A} & \xrightarrow{F\overline{f}} & F\overline{B}
 \end{array}$$

gives  $h_{\langle \bar{f} \rangle} = \langle id_{F\bar{A}}, \alpha_{\bar{B}} \circ F\bar{f} \rangle : F\bar{A} \rightarrow F\bar{A} \times G\bar{B}$ . Moreover,  $\langle id_{F\bar{A}}, \alpha_{\bar{B}} \circ F\bar{f} \rangle$  is a monomorphism in  $\text{Set}$  because  $id_{F\bar{A}}$  is, so its epi-mono factorization gives  $\iota_{\alpha^\wedge \langle \bar{f} \rangle} = \langle id_{F\bar{A}}, \alpha_{\bar{B}} \circ F\bar{f} \rangle$ , and thus  $\alpha^\wedge \langle \bar{f} \rangle = F\bar{A}$ . Then  $\iota_{\alpha^\wedge \langle \bar{f} \rangle} \alpha^\wedge \langle \bar{f} \rangle = \langle id_{F\bar{A}}, \alpha_{\bar{B}} \circ F\bar{f} \rangle (F\bar{A}) = \langle \alpha_{\bar{B}} \circ F\bar{f} \rangle^*$ , so that  $\langle \alpha \rangle^* \langle \bar{f} \rangle = (F\bar{A}, G\bar{B}, \iota_{\alpha^\wedge \langle \bar{f} \rangle} \alpha^\wedge \langle \bar{f} \rangle) = (F\bar{A}, G\bar{B}, \langle \alpha_{\bar{B}} \circ F\bar{f} \rangle^*) = \langle \alpha_{\bar{B}} \circ F\bar{f} \rangle$ . Finally,  $\alpha_{\bar{B}} \circ F\bar{f} = G\bar{f} \circ \alpha_{\bar{A}}$  by naturality of  $\alpha$ .  $\square$

The *equality relation transformer* on  $F : [\text{Set}^k, \text{Set}]$  is defined to be  $\text{Eq}_F = \langle id_F \rangle$ . Specifically,  $\text{Eq}_F = (F, F, \text{Eq}_F^*)$  with  $\text{Eq}_F^* = \langle id_F \rangle^*$ , and Lemma 22 ensures that  $\text{Eq}_F^* \text{Eq}_A = \text{Eq}_{F\bar{A}}$  for all  $A : \text{Set}$ . Graph relation transformers in general, and equality relation transformers in particular, extend to relation environments in the obvious ways. Indeed, if  $\rho, \rho' : \text{SetEnv}$  and  $f : \rho \rightarrow \rho'$ , then the *graph relation environment*  $\langle f \rangle$  is defined pointwise by  $\langle f \rangle \phi = \langle f \phi \rangle$  for every  $\phi$ . This entails that  $\pi_1 \langle f \rangle = \rho$  and  $\pi_2 \langle f \rangle = \rho'$ . In particular, the *equality relation environment*  $\text{Eq}_\rho$  is defined to be  $\langle id_\rho \rangle$ . This entails that  $\text{Eq}_\rho \phi = \text{Eq}_{\rho \phi}$  for every  $\phi$ . With these definitions in hand, we can state and prove both an Identity Extension Lemma and a Graph Lemma for our calculus.

**THEOREM 23 (IEL).** *If  $\rho : \text{SetEnv}$  and  $\Gamma; \Phi \vdash \tau : \mathcal{F}$  then  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \text{Eq}_\rho = \text{Eq}_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}$ .*

The proof is by induction on the structure of  $\tau$ . Only the application and fixpoint cases are non-routine. Both use Lemma 22. The latter also uses the observation that, for every  $n \in \mathbb{N}$ , the following intermediate results can be proved by simultaneous induction with Theorem 23:

$$\begin{aligned} T_{\text{Eq}_\rho}^n K_0 \overline{\text{Eq}_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}} &= (\text{Eq}_{(T_\rho^{\text{Set}})^n K_0})^* \overline{\text{Eq}_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}} \text{ and } \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}} \text{Eq}_\rho [\phi := T_{\text{Eq}_\rho}^n K_0] [\bar{\alpha} := \text{Eq}_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}] = \\ \llbracket \Gamma; \Phi, \phi, \bar{\alpha} \vdash H \rrbracket^{\text{Rel}} \text{Eq}_\rho [\phi := \text{Eq}_{(T_\rho^{\text{Set}})^n K_0}] [\bar{\alpha} := \text{Eq}_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}]. \end{aligned}$$

**LEMMA 24 (GRAPH LEMMA).** *If  $\rho, \rho' : \text{SetEnv}$ ,  $f : \rho \rightarrow \rho'$ , and  $\Gamma; \Phi \vdash F : \mathcal{F}$ , then  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ .*

**PROOF.** Applying Lemma 18 to the morphisms  $(f, id_{\rho'}) : \langle f \rangle \rightarrow \text{Eq}_{\rho'}$  and  $(id_\rho, f) : \text{Eq}_\rho \rightarrow \langle f \rangle$  of relation environments gives that  $(\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_{\rho'}) = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} (f, id_{\rho'}) : \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle \rightarrow \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho'}$  and  $(\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_\rho, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f) = \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} (id_\rho, f) : \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \text{Eq}_\rho \rightarrow \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ . Expanding the first equation gives that if  $(x, y) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$  then  $(\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_{\rho'} y) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho'}$ . Then  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_{\rho'} y = id_{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho'} y = y$  and  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho'} = \text{Eq}_{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho'}$ , so if  $(x, y) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$  then  $(\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x, y) \in \text{Eq}_{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho'}$ , i.e.,  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x = y$ , i.e.,  $(x, y) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f$ . So, we have that  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle \subseteq \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f$ . Expanding the second equation gives that if  $x \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho$  then  $(\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_\rho x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ . Then  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} id_\rho x = id_{\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho} x = x$ , so for any  $x \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho$  we have that  $(x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ . Moreover,  $x \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho$  if and only if  $(x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f$  and, if  $x \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} \rho$  then  $(x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ , so if  $(x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f$  then  $(x, \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f x) \in \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ , i.e.,  $\llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Set}} f \subseteq \llbracket \Gamma; \Phi \vdash F \rrbracket^{\text{Rel}} \langle f \rangle$ .  $\square$

## 5 INTERPRETING TERMS

Here, we are using angle bracket notation for both the graph relation of a function and for the pairing of functions with the same domain. This is justified by the relationship between the two notions observed immediately after Lemma 21.

If  $\Delta = x_1 : \tau_1, \dots, x_n : \tau_n$  is a term context for  $\Gamma$  and  $\Phi$ , then the interpretations  $\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}}$  and  $\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}}$  are defined by

$$\begin{aligned} \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}} &= \llbracket \Gamma; \Phi \vdash \tau_1 \rrbracket^{\text{Set}} \times \dots \times \llbracket \Gamma; \Phi \vdash \tau_n \rrbracket^{\text{Set}} \\ \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} &= \llbracket \Gamma; \Phi \vdash \tau_1 \rrbracket^{\text{Rel}} \times \dots \times \llbracket \Gamma; \Phi \vdash \tau_n \rrbracket^{\text{Rel}} \end{aligned}$$

Every well-formed term  $\Gamma; \Phi \mid \Delta \vdash t : \tau$  then has, for every  $\rho \in \text{SetEnv}$ , set interpretations  $\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho$  as natural transformations from  $\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}} \rho$  to  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho$ , and, for every  $\rho \in \text{RelEnv}$ , relational interpretations  $\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho$  as natural transformations from  $\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho$  to  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho$ . These are given in the next two definitions.

**DEFINITION 25.** *If  $\rho$  is a set environment and  $\Gamma; \Phi \mid \Delta \vdash t : \tau$  then  $\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho$  is defined as follows:*

$$\begin{aligned} \llbracket \Gamma; \emptyset \mid \Delta, x : \tau \vdash x : \tau \rrbracket^{\text{Set}} \rho &= \pi_{|\Delta|+1} \\ \llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho &= \text{curry}(\llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_]) \\ \llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho &= \text{eval} \circ \langle \lambda d. (\llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho d)_{\llbracket \Gamma; \Phi \vdash \bar{\tau} \rrbracket^{\text{Set}} \rho} \rangle \\ &\quad \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho \rangle \\ \llbracket \Gamma; \Phi \mid \Delta, x : \tau \vdash x : \tau \rrbracket^{\text{Set}} \rho &= \pi_{|\Delta|+1} \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \perp_{\tau} t : \tau \rrbracket^{\text{Set}} \rho &= !_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}^0 \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : 0 \rrbracket^{\text{Set}} \rho, \text{ where} \\ &\quad !_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}^0 \text{ is the unique morphism from } 0 \\ &\quad \text{to } \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \top : 1 \rrbracket^{\text{Set}} \rho &= !_1^{\llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Set}} \rho}, \text{ where } !_1^{\llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Set}} \rho} \\ &\quad \text{is the unique morphism from } \llbracket \Gamma; \Phi \vdash 1 \rrbracket^{\text{Set}} \rho \text{ to } 1 \\ \llbracket \Gamma; \Phi \mid \Delta \vdash (s, t) : \sigma \times \tau \rrbracket^{\text{Set}} \rho &= \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Set}} \rho \times \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \pi_1 t : \sigma \rrbracket^{\text{Set}} \rho &= \pi_1 \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \pi_2 t : \sigma \rrbracket^{\text{Set}} \rho &= \pi_2 \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \text{case } t \text{ of } \{x \mapsto l; y \mapsto r\} : \gamma \rrbracket^{\text{Set}} \rho &= \text{eval} \circ \langle \text{curry}(\llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash l : \gamma \rrbracket^{\text{Set}} \rho, \\ &\quad \llbracket \Gamma; \Phi \mid \Delta, y : \tau \vdash r : \gamma \rrbracket^{\text{Set}} \rho), \\ &\quad \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma + \tau \rrbracket^{\text{Set}} \rho \rangle \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \text{inl } s : \sigma + \tau \rrbracket^{\text{Set}} \rho &= \text{inl} \circ \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \Phi \mid \Delta \vdash \text{inr } t : \sigma + \tau \rrbracket^{\text{Set}} \rho &= \text{inr} \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{F}, \bar{G}}^{\bar{F}, \bar{G}} : \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) &= \lambda d \bar{\eta} \bar{B}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := \bar{B}]}[\phi := \lambda \bar{A}. \eta_{\bar{A} \bar{B}}] \\ &\quad (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] &= \lambda d \bar{B} \bar{C}. (\text{in}_{T_{\rho[\bar{\gamma} := \bar{C}]}}^{\text{Set}})_{\bar{B}} \\ &\quad (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho \\ \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^{\bar{F}} : \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] [\bar{\alpha} := \bar{\beta}] F) &= \lambda d \bar{\eta} \bar{B} \bar{C}. (\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}}))_{\bar{B}} \\ &\quad (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} \rho \end{aligned}$$

Add return type for fold in last clause? Should be  $\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}]$ .

This interpretation gives that  $\llbracket \Gamma; \emptyset \mid \Delta \vdash \lambda x. t : \sigma \rightarrow \tau \rrbracket^{\text{Set}} \rho = \text{curry}(\llbracket \Gamma; \emptyset \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Set}} \rho)$  and  $\llbracket \Gamma; \emptyset \mid \Delta \vdash st : \tau \rrbracket^{\text{Set}} \rho = \text{eval} \circ \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash s : \sigma \rightarrow \tau \rrbracket^{\text{Set}} \rho, \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho \rangle$ , as expected.

DEFINITION 26. If  $\rho$  is a relation environment and  $\Gamma; \Phi \mid \Delta \vdash t : \tau$  then  $\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho$  is defined as follows:

$$\begin{aligned}
& \llbracket \Gamma; \emptyset \mid \Delta, x : \tau \vdash x : \tau \rrbracket^{\text{Rel}} \rho &= \pi_{|\Delta|+1} \\
& \llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Rel}} \rho &= \text{curry}(\llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \rho [\bar{\alpha} := \_]) \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Rel}} \rho &= \text{eval} \circ \langle \lambda e. (\llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Rel}} \rho e) \overline{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho} \rangle \\
& & \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Rel}} \rho \rangle \\
& \llbracket \Gamma; \Phi \mid \Delta, x : \tau \vdash x : \tau \rrbracket^{\text{Rel}} \rho &= \pi_{|\Delta|+1} \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \perp_{\tau} t : \tau \rrbracket^{\text{Rel}} \rho &= \text{!}_1^0 \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : 0 \rrbracket^{\text{Rel}} \rho, \text{ where} \\
& & \text{!}_1^0 \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \text{ is the unique morphism from } 0 \\
& & \text{to } \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \top : 1 \rrbracket^{\text{Rel}} \rho &= \text{!}_1^{\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho}, \text{ where } \text{!}_1^{\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho} \\
& & \text{is the unique morphism from } \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho \text{ to } 1 \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash (s, t) : \sigma \times \tau \rrbracket^{\text{Rel}} \rho &= \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Rel}} \rho \times \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \pi_1 t : \sigma \rrbracket^{\text{Rel}} \rho &= \pi_1 \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \pi_2 t : \sigma \rrbracket^{\text{Rel}} \rho &= \pi_2 \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma \times \tau \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \text{case } t \text{ of } \{x \mapsto l; y \mapsto r\} : \gamma \rrbracket^{\text{Rel}} \rho &= \text{eval} \circ \langle \text{curry} [\llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash l : \gamma \rrbracket^{\text{Rel}} \rho, \\
& & \llbracket \Gamma; \Phi \mid \Delta, y : \tau \vdash r : \gamma \rrbracket^{\text{Rel}} \rho], \\
& & \llbracket \Gamma; \Phi \mid \Delta \vdash t : \sigma + \tau \rrbracket^{\text{Rel}} \rho \rangle \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \text{inl } s : \sigma + \tau \rrbracket^{\text{Rel}} \rho &= \text{inl} \circ \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \Phi \mid \Delta \vdash \text{inr } t : \sigma + \tau \rrbracket^{\text{Rel}} \rho &= \text{inr} \circ \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho \\
& \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{F}, \bar{G}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) &= \lambda d \bar{\eta} \bar{R}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Rel}} \text{id}_{\rho[\bar{\gamma} := \bar{R}]} [\phi := \lambda \bar{S}. \eta \bar{S} \bar{R}] \\
& \quad (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} := \bar{\beta} \bar{F}] H[\bar{\phi} := \bar{\beta} \bar{G}]) \rrbracket^{\text{Rel}} \rho & \\
& \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] &= \lambda d \bar{R} \bar{S}. (\text{in}_{T_{\rho[\bar{\gamma} := \bar{S}]}})_{\bar{R}} \\
& \quad (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Rel}} \rho & \\
& \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := \bar{\beta} \bar{F}] [\bar{\alpha} := \bar{\beta}] F) &= \lambda d \bar{\eta} \bar{R} \bar{S}. (\text{fold}_{T_{\rho[\bar{\gamma} := \bar{S}]}} (\lambda \bar{Z}. \eta \bar{Z} \bar{S}))_{\bar{R}} \\
& \quad (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \bar{F}) \rrbracket^{\text{Rel}} \rho &
\end{aligned}$$

Add return type for fold in last clause? Should be  $\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Rel}} \rho [\bar{\gamma} := \bar{C}]$ .

If  $t$  is closed, i.e., if  $\emptyset; \emptyset \mid \emptyset \vdash t : \tau$ , then we write  $\llbracket \vdash t : \tau \rrbracket^{\text{Set}}$  instead of  $\llbracket \emptyset; \emptyset \mid \emptyset \vdash t : \tau \rrbracket^{\text{Set}}$ , and similarly for  $\llbracket \emptyset; \emptyset \mid \emptyset \vdash t : \tau \rrbracket^{\text{Rel}}$ .

## 5.1 Basic Properties of Term Interpretations

This interpretation gives that  $\llbracket \Gamma; \emptyset \mid \Delta \vdash \lambda x. t : \sigma \rightarrow \tau \rrbracket^{\text{Rel}} \rho = \text{curry}(\llbracket \Gamma; \emptyset \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Rel}} \rho)$  and  $\llbracket \Gamma; \emptyset \mid \Delta \vdash st : \tau \rrbracket^{\text{Rel}} \rho = \text{eval} \circ \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash s : \sigma \rightarrow \tau \rrbracket^{\text{Rel}} \rho, \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \sigma \rrbracket^{\text{Rel}} \rho \rangle$ , as expected.

The interpretations in Definitions 25 and 26 respect weakening, i.e., a term and its weakenings all have the same set and relational interpretations. In particular, for any  $\rho \in \text{SetEnv}$ ,

$$\llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Set}} \rho = (\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho) \circ \pi_{\Delta}$$

where  $\pi_{\Delta}$  is the projection  $\llbracket \Gamma; \Phi \vdash \Delta, x : \sigma \rrbracket^{\text{Set}} \rightarrow \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}}$ , and for any  $\rho \in \text{RelEnv}$ ,

$$\llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Rel}} \rho = (\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho) \circ \pi_{\Delta}$$

where  $\pi_{\Delta}$  is the projection  $\llbracket \Gamma; \Phi \vdash \Delta, x : \sigma \rrbracket^{\text{Rel}} \rightarrow \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}}$ . Moreover, if  $\Gamma, \alpha; \Phi \mid \Delta \vdash t : \tau$  and  $\Gamma; \Phi, \alpha \mid \Delta \vdash t' : \tau$  and  $\Gamma; \Phi \vdash \sigma : \mathcal{F}$  then

$$\bullet \llbracket \Gamma; \Phi \mid \Delta[\alpha := \sigma] \vdash t[\alpha := \sigma] : \tau[\alpha := \sigma] \rrbracket^{\text{Set}} \rho = \llbracket \Gamma, \alpha; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} \rho]$$

- $\llbracket \Gamma; \Phi \mid \Delta[\alpha := \sigma] \vdash t[\alpha := \sigma] : \tau[\alpha := \sigma] \rrbracket^{\text{Rel}} \rho = \llbracket \Gamma, \alpha; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} \rho]$
- $\llbracket \Gamma; \Phi \mid \Delta[\alpha := \sigma] \vdash t'[\alpha := \sigma] : \tau[\alpha := \sigma] \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \Phi, \alpha \mid \Delta \vdash t' : \tau \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Set}} \rho]$
- $\llbracket \Gamma; \Phi \mid \Delta[\alpha := \sigma] \vdash t'[\alpha := \sigma] : \tau[\alpha := \sigma] \rrbracket^{\text{Rel}} \rho = \llbracket \Gamma; \Phi, \alpha \mid \Delta \vdash t' : \tau \rrbracket^{\text{Rel}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \sigma \rrbracket^{\text{Rel}} \rho]$

and if  $\Gamma; \Phi \mid \Delta, x : \sigma \vdash t : \tau$  and  $\Gamma; \Phi \mid \Delta \vdash s : \sigma$  then

- $\lambda A. \llbracket \Gamma; \Phi \mid \Delta \vdash t[x := s] : \tau \rrbracket^{\text{Set}} \rho A = \lambda A. \llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Set}} \rho(A, \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Set}} \rho A)$
- $\lambda R. \llbracket \Gamma; \Phi \mid \Delta \vdash t[x := s] : \tau \rrbracket^{\text{Rel}} \rho R = \lambda R. \llbracket \Gamma; \Phi \mid \Delta, x : \sigma \vdash t : \tau \rrbracket^{\text{Rel}} \rho(R, \llbracket \Gamma; \Phi \mid \Delta \vdash s : \sigma \rrbracket^{\text{Rel}} \rho R)$

Direct calculation reveals that the set interpretations of terms also satisfy

- $\llbracket \Gamma; \Phi \mid \Delta \vdash (L_{\bar{\alpha}x}.t)_{\bar{\tau}} s \rrbracket^{\text{Set}} = \llbracket \Gamma; \Phi \mid \Delta \vdash t[\bar{\alpha} := \bar{\tau}][x := s] \rrbracket^{\text{Set}}$

Standard type extensionality  $\llbracket \Gamma; \Phi \vdash (L_{\alpha}x.t)_{\alpha} t \rrbracket^{\text{Set}} = \llbracket \Gamma; \Phi \vdash t \rrbracket^{\text{Set}}$  and  $\llbracket \Gamma; \Phi \vdash (L_{\alpha}x.t)_{\alpha} t \rrbracket^{\text{Rel}} = \llbracket \Gamma; \Phi \vdash t \rrbracket^{\text{Rel}}$ , as well as term extensionality  $\llbracket \Gamma; \Phi \vdash (L_{\alpha}x.t)_{\alpha} \top \rrbracket^{\text{Set}} = \llbracket \Gamma; \Phi \vdash t \rrbracket^{\text{Set}}$  and  $\llbracket \Gamma; \Phi \vdash (L_{\alpha}x.t)_{\alpha} \top \rrbracket^{\text{Rel}} = \llbracket \Gamma; \Phi \vdash t \rrbracket^{\text{Rel}}$ , for terms are immediate consequences.

## 5.2 Properties of Terms of Nat-Type

If we define, for  $\Gamma; \bar{\alpha} \vdash F$ , the term  $id_F$  to be  $\Gamma; \emptyset \mid \emptyset \vdash L_{\bar{\alpha}}x.x : \text{Nat}^{\bar{\alpha}} F F$  and, for terms  $\Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G$  and  $\Gamma; \emptyset \mid \Delta \vdash s : \text{Nat}^{\bar{\alpha}} G H$ , the *composition*  $s \circ t$  of  $t$  and  $s$  to be  $\Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}}x.s_{\bar{\alpha}}(t_{\bar{\alpha}}x) : \text{Nat}^{\bar{\alpha}} F H$ , then

- $\llbracket \Gamma; \emptyset \mid \emptyset \vdash id_F : \text{Nat}^{\bar{\alpha}} F F \rrbracket^{\text{Set}} \rho * = id_{\lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}]}$  for any set environment  $\rho$
- $\llbracket \Gamma; \emptyset \mid \Delta \vdash s \circ t : \text{Nat}^{\bar{\alpha}} F H \rrbracket^{\text{Set}} = \llbracket \Gamma; \emptyset \mid \Delta \vdash s : \text{Nat}^{\bar{\alpha}} G H \rrbracket^{\text{Set}} \circ \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}$

Moreover, terms of Nat type behave as natural transformations with respect to their source and target functorial types.

- $\llbracket \Gamma; \emptyset \mid x : \text{Nat}^{\bar{\alpha}, \bar{\gamma}} F G, y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash ((\text{map}_{\bar{G}}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) \circ (L_{\bar{\gamma}}z.x_{\bar{\sigma}, \bar{\gamma}}z) : \text{Nat}^{\bar{\gamma}} F[\bar{\alpha} := \bar{\sigma}] G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} = \llbracket \Gamma; \emptyset \mid x : \text{Nat}^{\bar{\alpha}, \bar{\gamma}} F G, y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash (L_{\bar{\gamma}}z.x_{\bar{\tau}, \bar{\gamma}}z) \circ ((\text{map}_{\bar{F}}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) : \text{Nat}^{\bar{\gamma}} F[\bar{\alpha} := \bar{\sigma}] G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$

As the special case of the previous equality when  $x = in_H$  we have

$$\begin{aligned} \text{THEOREM 27.} \quad & \bullet \quad \llbracket \Gamma; \emptyset \mid y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash ((\text{map}_{(\mu\phi. \lambda \bar{\alpha}. H)_{\bar{\beta}}}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) \circ (L_{\bar{\gamma}}z.(in_H)_{\bar{\tau}, \bar{\gamma}}z) : \xi \rrbracket^{\text{Set}} \\ & = \llbracket \Gamma; \emptyset \mid y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash (L_{\bar{\gamma}}z.(in_H)_{\bar{\tau}, \bar{\gamma}}z) \circ ((\text{map}_{H[\phi := (\mu\phi. \lambda \bar{\alpha}. H)_{\bar{\beta}}]}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) : \xi \rrbracket^{\text{Set}} \\ & \text{at type } \xi = \text{Nat}^{\bar{\gamma}} H[\phi := (\mu\phi. \lambda \bar{\alpha}. H)_{\bar{\beta}}][\bar{\alpha} := \bar{\sigma}](\mu\phi. \lambda \bar{\alpha}. H)_{\bar{\tau}} \end{aligned}$$

Analogous results hold for relational interpretations of terms and relational environments.

As we observe in Section 5.4, Theorem 27 gives a family of results that we normally derive as free theorems but actually are consequences of naturality. Most of Wadler's fall into this family, for example, but not the free theorem for filter (even for lists) or short cut fusion.

## 5.3 Properties of Initial Algebraic Constructs

We first observe that map-terms are interpreted as semantic *maps*: **Contrast with Bob's  $fmaps_F$ s, which are not required to be the functorial action of  $F$ , only to preserve identities and compositions.**

Let  $\Gamma; \bar{\phi}, \bar{\gamma} \vdash H : \mathcal{F}$ ,  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash F : \mathcal{F}$  and  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash G : \mathcal{F}$ . By definition of the semantic interpretation of map terms, we have

$$\begin{aligned} \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_H^{\bar{F}, \bar{G}} : \text{Nat}^{\emptyset} (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} := \bar{\beta} F] H[\bar{\phi} := \bar{\beta} G]) \rrbracket^{\text{Set}} \rho \\ = \lambda d \bar{\eta} \bar{B}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := \bar{B}]}[\bar{\phi} := \lambda \bar{A}. \bar{\eta}_{\bar{A} \bar{B}}] \quad (10) \end{aligned}$$



Then let  $\Gamma; \bar{\alpha} \vdash F : \mathcal{F}$ ,  $\Gamma; \emptyset \vdash \sigma : \mathcal{F}$ ,  $\Gamma; \emptyset \vdash \tau : \mathcal{F}$  and  $*$  be the unique element of  $\llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$ . As a special case of the above definition, we have

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{F}}^{\bar{\sigma}, \bar{\tau}} : \text{Nat}^0 (\text{Nat}^0 \sigma \tau) (\text{Nat}^0 F[\bar{\alpha} := \sigma] F[\bar{\alpha} := \tau]) \rrbracket^{\text{Set}} \rho * \\ &= \lambda f : \llbracket \Gamma; \emptyset \vdash \sigma \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \emptyset \vdash \tau \rrbracket^{\text{Set}} \rho. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} id_{\rho}[\bar{\alpha} := f] \\ &= \lambda f : \llbracket \Gamma; \emptyset \vdash \sigma \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \emptyset \vdash \tau \rrbracket^{\text{Set}} \rho. \text{map}_{\lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}]} \bar{f} \\ &= \text{map}_{\lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}]} \end{aligned}$$

where the first equality is by Equation 10, the second equality is obtained by noting that  $\lambda \bar{A}. \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \bar{A}]$  is a functor in  $\alpha$ , and  $\text{map}_G$  denotes the action of the functor  $G$  on morphisms.

We also have the expected relationships between interpretations of terms involving map, in, and fold:

- If  $\Gamma; \bar{\psi}, \bar{\gamma} \vdash H$ ,  $\Gamma; \bar{\alpha}, \bar{\gamma}, \bar{\phi} \vdash K$ ,  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash F$ , and  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash G$ , then

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{H[\bar{\psi} := K]}^{\bar{F}, \bar{G}} : \xi \rrbracket^{\text{Set}} = \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_H^{K[\bar{\phi} := F], K[\bar{\phi} := G]} \circ \text{map}_K^{\bar{F}, \bar{G}} : \xi \rrbracket^{\text{Set}}$$

at type  $\xi = \text{Nat}^0 (\text{Nat}^0 \bar{\beta}, \bar{\gamma} F G) (\text{Nat}^0 H[\bar{\psi} := K][\bar{\phi} := F] H[\bar{\psi} := K][\bar{\phi} := G])$

- If  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash H$ ,  $\Gamma; \bar{\beta}, \bar{\gamma} \vdash K$ ,  $\Gamma; \bar{\alpha}, \bar{\gamma} \vdash F$ ,  $\Gamma; \bar{\alpha}, \bar{\gamma} \vdash G$ ,  $\Gamma; \bar{\phi}, \bar{\psi}, \bar{\gamma} \vdash \tau$ ,  $\bar{I}$  is the sequence  $\bar{F}, H$  and  $\bar{J}$  is the sequence  $\bar{G}, K$  then

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash L_{\emptyset}(x, \bar{y}). L_{\bar{\gamma}} z. x \frac{\tau[\bar{\psi} := F][\bar{\phi} := H], \tau[\bar{\psi} := G][\bar{\phi} := K]}{\tau[\bar{\psi} := G][\bar{\phi} := K], \bar{\gamma}} \left( ((\text{map}_H^{\tau[\bar{\psi} := F][\bar{\phi} := H], \tau[\bar{\psi} := G][\bar{\phi} := K]})_{\emptyset} ((\text{map}_{\bar{\tau}}^{\bar{I}, \bar{J}})_{\emptyset}(x, \bar{y})) \right)_{\bar{\gamma}} z \rrbracket^{\text{Set}}$$

$$= \llbracket \Gamma; \emptyset \mid \emptyset \vdash L_{\emptyset}(x, \bar{y}). L_{\bar{\gamma}} z. ((\text{map}_K^{\tau[\bar{\psi} := F][\bar{\phi} := H], \tau[\bar{\psi} := G][\bar{\phi} := K]})_{\emptyset} ((\text{map}_{\bar{\tau}}^{\bar{I}, \bar{J}})_{\emptyset}(x, \bar{y})) \rrbracket^{\text{Set}} \left( x \frac{\tau[\bar{\psi} := F][\bar{\phi} := H], \tau[\bar{\psi} := G][\bar{\phi} := K]}{\tau[\bar{\psi} := F][\bar{\phi} := H], \bar{\gamma}} z \right) : \xi \rrbracket^{\text{Set}}$$

at type  $\xi = \text{Nat}^0 (\text{Nat}^0 \bar{\beta}, \bar{\gamma} H K \times \text{Nat}^0 \bar{\alpha}, \bar{\gamma} F G) (\text{Nat}^0 H[\bar{\beta} := \tau][\bar{\psi} := F][\bar{\phi} := H] K[\bar{\beta} := \tau][\bar{\psi} := G][\bar{\phi} := K])$ .

- $\llbracket \Gamma; \emptyset \mid x : \text{Nat}^0 \bar{\beta}, \bar{\gamma} H[\bar{\phi} := F][\bar{\alpha} := \bar{\beta}] F \vdash ((\text{fold}_{H, F})_{\emptyset} x) \circ \text{in}_H : \xi \rrbracket^{\text{Set}}$   
 $= \llbracket \Gamma; \emptyset \mid x : \text{Nat}^0 \bar{\beta}, \bar{\gamma} H[\bar{\phi} := F][\bar{\alpha} := \bar{\beta}] F \vdash x \circ ((\text{map}_{H[\bar{\alpha} := \bar{\beta}]}^{(\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}, F})_{\emptyset} ((\text{fold}_{H, F})_{\emptyset} x)) : \xi \rrbracket^{\text{Set}}$

at type  $\xi = \text{Nat}^0 \bar{\beta}, \bar{\gamma} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}][\bar{\alpha} := \bar{\beta}] F$

- $\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H \circ (\text{fold}_{H, H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}]}^{\bar{\beta}})_{\emptyset} ((\text{map}_H^{H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}][\bar{\alpha} := \bar{\beta}], (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}})_{\emptyset} \text{in}_H) : \xi \rrbracket^{\text{Set}}$   
 $= \llbracket \Gamma; \emptyset \mid \emptyset \vdash Id_{(\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}} : \xi \rrbracket^{\text{Set}}$

at type  $\xi = \text{Nat}^0 \bar{\beta}, \bar{\gamma} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}$

- $\llbracket \Gamma; \emptyset \mid \emptyset \vdash (\text{fold}_{H, H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}]}^{\bar{\beta}})_{\emptyset} ((\text{map}_H^{H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}][\bar{\alpha} := \bar{\beta}], (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}})_{\emptyset} \text{in}_H) \circ \text{in}_H : \xi \rrbracket^{\text{Set}}$   
 $= \llbracket \Gamma; \emptyset \mid \emptyset \vdash Id_{H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}]} : \xi \rrbracket^{\text{Set}}$

at type  $\xi = \text{Nat}^0 \bar{\beta}, \bar{\gamma} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}]$ .

Analogous results hold for relational interpretations of terms and relational environments. The set and relational interpretations of terms therefore respect the congruence closed equational theory obtained by adding these judgments to those generating the usual congruence closed equational theory induced by the other term formers.

## 5.4 Free Theorems Derived from Naturality

Foralls in Nat-types are at the object level, whereas the foralls in contexts are at the meta-level. So par results in subst theorem internalize parametricity in the calculus, whereas those parametricity results that do not follow from the interpretation of Nat-types are externalized at the meta-level.

**Make this not about *subst*** Note that the free theorem for a type is always independent of the particular term of that type, so the proof below is independent of the choice of function *subst*. In addition, it is independent of the particular data type — in this case,  $\text{Lam}$  — over which *subst* acts. Also independent of the functor arguments — in this case  $+1$  and *id* — to the data type. Indeed, the following result is just a consequence of naturality.

We already know from Theorem 27 that

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid x : \text{Nat}^{\bar{\alpha}, \bar{\gamma}} F G, y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash ((\text{map}_{\bar{G}}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) \circ (L_{\bar{\gamma} z}.x_{\bar{\sigma}, \bar{\gamma}} z) : \text{Nat}^{\bar{\gamma}} F[\bar{\alpha} := \bar{\sigma}] G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \\ &= \llbracket \Gamma; \emptyset \mid x : \text{Nat}^{\bar{\alpha}, \bar{\gamma}} F G, y : \text{Nat}^{\bar{\gamma}} \sigma \tau \vdash (L_{\bar{\gamma} z}.x_{\bar{\tau}, \bar{\gamma}} z) \circ ((\text{map}_{\bar{F}}^{\bar{\sigma}, \bar{\tau}})_{\emptyset} \bar{y}) : \text{Nat}^{\bar{\gamma}} F[\bar{\alpha} := \bar{\sigma}] G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \end{aligned} \quad (11)$$

In particular, if we instantiate  $x$  with any term *subst* of type  $\vdash \text{Nat}^{\alpha}(\text{Lam}(\alpha + 1) \times \text{Lam} \alpha) \text{Lam} \alpha$  (and thus there is a single  $\alpha$  and no  $\gamma$ 's) we have

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash ((\text{map}_{\text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y) \circ (L_{\emptyset z}.\text{subst}_{\sigma} z) : \text{Nat}^0(\text{Lam}(\sigma + 1) \times \text{Lam} \sigma) \text{Lam} \tau \rrbracket^{\text{Set}} \\ &= \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash (L_{\emptyset z}.\text{subst}_{\tau} z) \circ ((\text{map}_{\text{Lam}(\alpha + 1) \times \text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y) : \text{Nat}^0(\text{Lam}(\sigma + 1) \times \text{Lam} \sigma) \text{Lam} \tau \rrbracket^{\text{Set}} \end{aligned} \quad (12)$$

Thus, for any set environment  $\rho$  and any function  $f : \llbracket \Gamma; \emptyset \vdash \text{Nat}^0 \sigma \tau \rrbracket^{\text{Set}} \rho$ , we have that

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash ((\text{map}_{\text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y) \circ (L_{\emptyset z}.\text{subst}_{\sigma} z) : \text{Nat}^0(\text{Lam}(\sigma + 1) \times \text{Lam} \sigma) \text{Lam} \tau \rrbracket^{\text{Set}} \rho f \\ &= \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash ((\text{map}_{\text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y) \rrbracket^{\text{Set}} \rho f \circ \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash L_{\emptyset z}.\text{subst}_{\sigma} z \rrbracket^{\text{Set}} \rho f \\ &= \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\text{Lam} \alpha}^{\sigma, \tau} \rrbracket^{\text{Set}} \rho f \circ \llbracket \Gamma; \emptyset \mid \emptyset \vdash L_{\emptyset z}.\text{subst}_{\sigma} z \rrbracket^{\text{Set}} \rho \\ &= \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} f \circ (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_{\llbracket \Gamma; \emptyset \vdash \sigma \rrbracket^{\text{Set}} \rho} \end{aligned} \quad (13)$$

and

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash (L_{\emptyset z}.\text{subst}_{\tau} z) \circ ((\text{map}_{\text{Lam}(\alpha + 1) \times \text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y) : \text{Nat}^0(\text{Lam}(\sigma + 1) \times \text{Lam} \sigma) \text{Lam} \tau \rrbracket^{\text{Set}} \rho f \\ &= \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash L_{\emptyset z}.\text{subst}_{\tau} z \rrbracket^{\text{Set}} \rho f \circ \llbracket \Gamma; \emptyset \mid y : \text{Nat}^0 \sigma \tau \vdash (\text{map}_{\text{Lam}(\alpha + 1) \times \text{Lam} \alpha}^{\sigma, \tau})_{\emptyset} y \rrbracket^{\text{Set}} \rho f \\ &= \llbracket \Gamma; \emptyset \mid \emptyset \vdash L_{\emptyset z}.\text{subst}_{\tau} z \rrbracket^{\text{Set}} \rho \circ \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\text{Lam}(\alpha + 1) \times \text{Lam} \alpha}^{\sigma, \tau} \rrbracket^{\text{Set}} \rho f \\ &= (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_{\llbracket \Gamma; \emptyset \vdash \tau \rrbracket^{\text{Set}} \rho} \circ \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam}(\alpha + 1) \times \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} f \\ &= (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_{\llbracket \Gamma; \emptyset \vdash \tau \rrbracket^{\text{Set}} \rho} \circ (\text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} (f + 1) \times \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} f) \end{aligned} \quad (14)$$

So, we can conclude that

$$\begin{aligned} & \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} f \circ (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_{\llbracket \Gamma; \emptyset \vdash \sigma \rrbracket^{\text{Set}} \rho} \\ &= (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_{\llbracket \Gamma; \emptyset \vdash \tau \rrbracket^{\text{Set}} \rho} \circ (\text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} (f + 1) \times \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam} \alpha \rrbracket^{\text{Set}} [\alpha := \_]} f) \end{aligned} \quad (15)$$

Moreover, for any  $A, B : \text{Set}$ , we can choose  $\sigma = v$  and  $\tau = w$  to be variables such that  $\rho v = A$  and  $\rho w = B$ . Then for any function  $f : A \rightarrow B$  we have that

$$\begin{aligned} & \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam } \alpha \rrbracket^{\text{Set}}[\alpha := \_]} f \circ (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_A \\ &= (\llbracket \vdash \text{subst} \rrbracket^{\text{Set}})_B \circ (\text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam } \alpha \rrbracket^{\text{Set}}[\alpha := \_]} (f + \mathbb{1}) \times \text{map}_{\llbracket \emptyset; \alpha \vdash \text{Lam } \alpha \rrbracket^{\text{Set}}[\alpha := \_]} f) \quad (16) \end{aligned}$$

## 5.5 The Abstraction Theorem

To go beyond naturality and get *all* consequences of parametricity, we prove an Abstraction Theorem for our calculus. In fact, we actually prove a more general result in Theorem 28 about possibly open terms. We then recover the Abstraction Theorem as the special case of Theorem 28 for closed terms of closed type.

**THEOREM 28.** *Every well-formed term  $\Gamma; \Phi \mid \Delta \vdash t : \tau$  induces a natural transformation from  $\llbracket \Gamma; \Phi \vdash \Delta \rrbracket$  to  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket$ , i.e., a triple of natural transformations*

$$(\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}}, \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}}, \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}})$$

where

$$\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} : \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}} \rightarrow \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}$$

has as its component at  $\rho : \text{SetEnv}$  a morphism

$$\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} \rho : \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho$$

in  $\text{Set}$ , and

$$\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} : \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rightarrow \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}}$$

has as its component at  $\rho : \text{RelEnv}$  a morphism

$$\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho : \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho$$

in  $\text{Rel}$ , and for all  $\rho : \text{RelEnv}$ ,

$$\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Rel}} \rho = (\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}}(\pi_1 \rho), \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}}(\pi_2 \rho))$$

**PROOF.** We proceed by structural induction, showing only the interesting cases.

- We first consider  $\Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G$ .
  - To see that  $\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}$  is a natural transformation from  $\llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Set}}$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}$ , since the functorial part  $\Phi$  of the context is empty, we need only show that, for every  $\rho : \text{SetEnv}$ ,  $\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$  is a morphism in  $\text{Set}$  from  $\llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Set}} \rho$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$ . For this, recall that

$$\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\bar{\alpha}} x. t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho = \text{curry}(\llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_])$$

By the induction hypothesis,  $\llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket \rho[\bar{\alpha} := \_]$  induces a natural transformation

$$\begin{aligned} & \llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \\ & : \llbracket \Gamma; \bar{\alpha} \vdash \Delta, x : F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \\ & = \llbracket \Gamma; \bar{\alpha} \vdash \Delta \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \times \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \end{aligned}$$

and thus a family of morphisms

$$\begin{aligned} & \text{curry}(\llbracket \Gamma; \bar{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket \rho[\bar{\alpha} := \_]) \\ & : \llbracket \Gamma; \bar{\alpha} \vdash \Delta \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \rightarrow (\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_] \rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_]) \end{aligned}$$

That is, for each  $\overline{A} : \text{Set}$  and each  $d : \llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Set}} \rho = \llbracket \Gamma; \overline{\alpha} \vdash \Delta \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}]$  by weakening, we have

$$\begin{aligned} & (\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Set}} \rho d)_{\overline{A}} \\ &= \text{curry} (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}]) d \\ &: \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}] \rightarrow \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}] \end{aligned}$$

Moreover, these maps actually form a natural transformation  $\eta : \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \_] \rightarrow \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \_]$  because each

$$\eta_{\overline{A}} = \text{curry} (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}]) d$$

is the component at  $\overline{A}$  of the partial specialization to  $d$  of the natural transformation  $\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \_]$ .

To see that the components of  $\eta$  also satisfy the additional condition necessary for  $\eta$  to be in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Set}} \rho$ , let  $\overline{R} : \text{Rel}(A, B)$  and

$$(u, v) \in \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] = (\llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}], \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{B}])$$

Then the induction hypothesis on the term  $t$  ensures that

$$\begin{aligned} & \llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] \\ &: \llbracket \Gamma; \overline{\alpha} \vdash \Delta, x : F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] \rightarrow \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] \end{aligned}$$

and

$$\begin{aligned} & \llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] \quad (*) \\ &= (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}], \llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{B}]) \end{aligned}$$

Since  $(d, d) \in \llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Rel}} \text{Eq}_{\rho} = \llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}]$  we therefore have that

$$\begin{aligned} & (\eta_{\overline{A}} u, \eta_{\overline{B}} v) \\ &= (\text{curry} (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{A}]) d u, \text{curry} (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \overline{B}]) d v) \\ &= \text{curry} (\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}]) (d, d) (u, v) \\ &: \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\overline{\alpha} := \overline{R}] \end{aligned}$$

Here, the second equality is by  $(*)$ .

- The proofs that  $\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}}$  is a natural transformation from  $\llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Rel}}$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}}$  and that, for all  $\rho : \text{RelEnv}$  and  $d : \llbracket \Gamma; \emptyset \vdash \Delta \rrbracket^{\text{Rel}}$ ,

$$\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho d$$

is a natural transformation from  $\llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Rel}} \rho[\overline{\alpha} := \_]$  to  $\llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Rel}} \rho[\overline{\alpha} := \_]$ , are analogous.

- Finally, to see that  $\pi_i(\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho) = \llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Set}} (\pi_i \rho)$  we observe that  $\pi_1$  and  $\pi_2$  are surjective and compute

$$\begin{aligned} & \pi_i(\llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho) \\ &= \pi_i(\text{curry}(\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \rho[\overline{\alpha} := \_])) \\ &= \text{curry}(\pi_i(\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Rel}} \rho[\overline{\alpha} := \_])) \\ &= \text{curry}(\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} (\pi_i(\rho[\overline{\alpha} := \_]))) \\ &= \text{curry}(\llbracket \Gamma; \overline{\alpha} \mid \Delta, x : F \vdash t : G \rrbracket^{\text{Set}} (\pi_i \rho)[\overline{\alpha} := \_])) \\ &= \llbracket \Gamma; \emptyset \mid \Delta \vdash L_{\overline{\alpha}} x.t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Set}} (\pi_i \rho) \end{aligned}$$

- We now consider  $\Gamma; \Phi \mid \Delta \vdash t_{\overline{\tau}} s : G[\overline{\alpha} := \overline{\tau}]$ .

– To see that  $\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$  is a natural transformation from  $\llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Set}}$  to  $\llbracket \Gamma; \Phi \mid G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$  we need to show that, for every  $\rho : \text{SetEnv}$ ,  $\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$  is a morphism from  $\llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Set}} \rho$  to  $\llbracket \Gamma; \Phi \mid G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho$ , and that this family of morphisms is natural in  $\rho$ . Let  $d : \llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Set}} \rho$ . Then

$$\begin{aligned} & \llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho d \\ = & (\text{eval} \circ \langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho \_ \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho \rangle) d \\ = & \text{eval}(\langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho \_ \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} d, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho d) \\ = & \text{eval}(\langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho d \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho d) \end{aligned}$$

By the induction hypothesis,  $(\llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho d)_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}$  has type

$$\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] \rightarrow \llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho]$$

and  $\llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho d$  has type

$$\begin{aligned} & \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho \\ = & \llbracket \Gamma; \Phi, \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] \\ = & \llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] \end{aligned}$$

by Equation 6, and by weakening in the last step, since the type  $\Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G$  is only well-formed if  $\Gamma; \bar{\alpha} \vdash F : \mathcal{F}$  and  $\Gamma; \bar{\alpha} \vdash G : \mathcal{F}$ . Thus,  $\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho d$  has type  $\llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\alpha := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] = \llbracket \Gamma; \Phi \mid G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho$ , as desired.

To see that the family of maps comprising  $\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\bar{\tau}} s : G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$  form a natural transformation, i.e., are natural in their set environment argument, we need to show that the following diagram commutes:

$$\begin{array}{ccc} \llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Set}} \rho & \xrightarrow{\llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Set}} f} & \llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Set}} \rho' \\ \downarrow \langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho \rangle & \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} f \times \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f & \downarrow \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho' \rangle \\ \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho \times \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho & \xrightarrow{\quad} & \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho' \times \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho' \\ \downarrow \text{eval} \circ ((-)_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \times id) & & \downarrow \text{eval} \circ ((-)_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \times id) \\ \llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho & \xrightarrow{\llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f} & \llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho' \end{array}$$

The top diagram commutes because the induction hypothesis ensures  $\llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}$  and  $\llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}}$  are natural in  $\rho$ . To see that the bottom diagram commutes, we first note that since  $\rho|_{\Gamma} = \rho'|_{\Gamma}$ ,  $\Gamma; \bar{\alpha} \vdash F : \mathcal{F}$ , and  $\Gamma; \bar{\alpha} \vdash G : \mathcal{F}$  we can replace the instance of  $f$  in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} f$  with  $id$ . Then, using the fact that  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}}$  is a functor, we have that  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} id = id$ . To see that the bottom diagram commutes we must therefore prove that, for every  $\eta \in \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$  and  $x \in \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} \rho$ , we have

$$\llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f(\eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} x) = \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}(\llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f x)$$

i.e., that for every  $\eta \in \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$ ,

$$\llbracket \Gamma; \Phi \vdash G[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f \circ \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} = \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \circ \llbracket \Gamma; \Phi \vdash F[\bar{\alpha} := \bar{\tau}] \rrbracket^{\text{Set}} f$$

But this follows from the naturality of  $\eta$ . Indeed,  $\eta \in \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\alpha}} F G \rrbracket^{\text{Set}} \rho$  implies that  $\eta$  is a natural transformation from  $\llbracket \Gamma; \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_]$  to  $\llbracket \Gamma; \bar{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\bar{\alpha} := \_]$ . For each

$\tau$ , consider the morphism  $\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f : \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho \rightarrow \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'$ . The following diagram commutes by naturality of  $\eta$ :

$$\begin{array}{ccc}
 \llbracket \Gamma; \Phi \vdash F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}} \rho & & \llbracket \Gamma; \Phi \vdash G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}} \rho \\
 \parallel & & \parallel \\
 \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] & \xrightarrow{\eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho}} & \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho] \\
 \downarrow \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} id_{\rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho]} & & \downarrow \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} id_{\rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho]} \\
 \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'] & \xrightarrow{\eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'}} & \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho']
 \end{array}$$

That is,

$$\begin{aligned}
 & \llbracket \Gamma; \overline{\alpha} \vdash G \rrbracket^{\text{Set}} id_{\rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f]} \circ \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} \\
 = & \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \circ \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} id_{\rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f]}
 \end{aligned}$$

But since the only variables in the functorial contexts for  $F$  and  $G$  are  $\overline{\alpha}$ , we have that

$$\begin{aligned}
 & \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} id_{\rho[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f]} \\
 = & \llbracket \Gamma; \overline{\alpha} \vdash F \rrbracket^{\text{Set}} f[\overline{\alpha} := \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} f]} \\
 = & \llbracket \Gamma; \Phi \vdash F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}} f
 \end{aligned}$$

and similarly for  $G$ . Commutativity of this last diagram thus gives that  $\llbracket \Gamma; \Phi \vdash G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}} f \circ \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho} = \eta_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho'} \circ \llbracket \Gamma; \Phi \vdash F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}} f$ , as desired.

- The proof that  $\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\tau} s : G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}}$  is a natural transformation from  $\llbracket \Gamma; \Phi \mid \Delta \rrbracket^{\text{Rel}}$  to  $\llbracket \Gamma; \Phi \vdash G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}}$  is analogous.
- Finally, to see that  $\pi_i(\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\tau} s : G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}} \rho) = \llbracket \Gamma; \Phi \mid \Delta \vdash t_{\tau} s : G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}}(\pi_i \rho)$  we compute

$$\begin{aligned}
 & \pi_i(\llbracket \Gamma; \Phi \mid \Delta \vdash t_{\tau} s : G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}} \rho) \\
 = & \pi_i(\text{eval} \circ \langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho \_ \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho}, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}} \rho \rangle) \\
 = & \text{eval} \circ \langle \pi_i(\langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho \_ \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho}), \pi_i(\llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}} \rho) \rangle \\
 = & \text{eval} \circ \langle \pi_i(\llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Rel}} \rho \_ \rangle_{\pi_i(\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho)}, \pi_i(\llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Rel}} \rho) \rangle \\
 = & \text{eval} \circ \langle \langle \llbracket \Gamma; \emptyset \mid \Delta \vdash t : \text{Nat}^{\overline{\alpha}} F G \rrbracket^{\text{Set}}(\pi_i \rho) \_ \rangle_{\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}(\pi_i \rho)}, \llbracket \Gamma; \Phi \mid \Delta \vdash s : F[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}}(\pi_i \rho) \rangle \\
 = & \llbracket \Gamma; \Phi \mid \Delta \vdash t_{\tau} s : G[\overline{\alpha} := \overline{\tau}] \rrbracket^{\text{Set}}(\pi_i \rho)
 \end{aligned}$$

- We now consider  $\Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{F}, \overline{G}}^{\overline{F}, \overline{G}} : \text{Nat}^{\overline{\alpha}}(\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} := \overline{\beta} F] H[\overline{\phi} := \overline{\beta} G])$ .

- To see that  $\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{F}, \overline{G}}^{\overline{F}, \overline{G}} : \text{Nat}^{\overline{\alpha}}(\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} := \overline{\beta} F] H[\overline{\phi} := \overline{\beta} G]) \rrbracket^{\text{Set}}$  is a natural transformation from  $\llbracket \Gamma; \emptyset \mid \emptyset \rrbracket^{\text{Set}}$  to  $\llbracket \text{Nat}^{\overline{\alpha}}(\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} := \overline{\beta} F] H[\overline{\phi} := \overline{\beta} G]) \rrbracket^{\text{Set}}$ , since the functorial part  $\Phi$  of the context is empty, we need only show that, for every  $\rho : \text{SetEnv}$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{F}, \overline{G}}^{\overline{F}, \overline{G}} : \text{Nat}^{\overline{\alpha}}(\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} := \overline{\beta} F] H[\overline{\phi} := \overline{\beta} G]) \rrbracket^{\text{Set}} \rho$$

is a morphism in  $\text{Set}$  from  $\llbracket \Gamma; \emptyset \mid \emptyset \rrbracket^{\text{Set}} \rho$  to

$$\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\alpha}}(\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} := \overline{\beta} F] H[\overline{\phi} := \overline{\beta} G]) \rrbracket^{\text{Set}} \rho$$

i.e., that, for the unique  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{H}}^{\overline{F}, \overline{G}} : \text{Nat}^0 (\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G]) \rrbracket^{\text{Set}} \rho d$$

is a morphism from  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\beta}, \overline{\gamma}} F G \rrbracket^{\text{Set}} \rho$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G] \rrbracket^{\text{Set}} \rho$ .

For this we show that for all  $\eta : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\beta}, \overline{\gamma}} F G \rrbracket^{\text{Set}} \rho$  we have

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{H}}^{\overline{F}, \overline{G}} : \text{Nat}^0 (\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G]) \rrbracket^{\text{Set}} \rho d \overline{\eta} \\ & : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G] \rrbracket^{\text{Set}} \rho \end{aligned}$$

To this end, we note that, for any  $\overline{B}$ ,

$$\begin{aligned} & (\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{H}}^{\overline{F}, \overline{G}} : \text{Nat}^0 (\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G]) \rrbracket^{\text{Set}} \rho d \overline{\eta})_{\overline{B}} \\ = & \llbracket \Gamma; \overline{\phi}, \overline{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\overline{\gamma} := \overline{B}]} [\overline{\phi} := \lambda \overline{A}. \eta_{\overline{A} \overline{B}}] \end{aligned}$$

is indeed a morphism from

$$\begin{aligned} & \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := F] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] \\ = & \llbracket \Gamma; \overline{\gamma}, \overline{\phi} \vdash H \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\beta} := A]] \end{aligned}$$

to

$$\begin{aligned} & \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := G] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] \\ = & \llbracket \Gamma; \overline{\gamma}, \overline{\phi} \vdash H \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\beta} := A]] \end{aligned}$$

since  $\llbracket \Gamma; \overline{\phi}, \overline{\gamma} \vdash H \rrbracket^{\text{Set}}$  is a functor from  $\text{SetEnv}$  to  $\text{Set}$  and  $id_{\rho[\overline{\gamma} := \overline{B}]} [\overline{\phi} := \lambda \overline{A}. \eta_{\overline{A} \overline{B}}]$  is a morphism in  $\text{SetEnv}$  from

$$\rho[\overline{\gamma} := \overline{B}] [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\beta} := A]]$$

to

$$\rho[\overline{\gamma} := \overline{B}] [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash G \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\beta} := A]]$$

To see that this family of morphisms is natural in  $\overline{B}$  we first observe that if  $f : \overline{B} \rightarrow \overline{B'}$  then, writing  $t$  for  $\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\overline{H}}^{\overline{F}, \overline{G}} : \text{Nat}^0 (\text{Nat}^{\overline{\beta}, \overline{\gamma}} F G) (\text{Nat}^{\overline{\gamma}} H[\overline{\phi} :=_{\overline{\beta}} F] H[\overline{\phi} :=_{\overline{\beta}} G]) \rrbracket^{\text{Set}} \rho d \overline{\eta}$ , we have

$$\begin{array}{ccc} \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := F] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] & \xrightarrow{t_{\overline{B}}} & \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := G] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] \\ \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := F] \rrbracket^{\text{Set}} id_{\rho[\overline{\gamma} := f]} \downarrow & & \downarrow \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := G] \rrbracket^{\text{Set}} id_{\rho[\overline{\gamma} := f]} \\ \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := F] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B'}] & \xrightarrow{t_{\overline{B'}}} & \llbracket \Gamma; \overline{\gamma} \vdash H[\overline{\phi} := G] \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B'}] \end{array}$$

This diagram commutes because  $\llbracket \Gamma; \overline{\phi}, \overline{\gamma} \vdash H \rrbracket^{\text{Set}}$  is a functor from  $\text{SetEnv}$  to  $\text{Set}$  and because, letting

$$E_{F, B} = \rho[\overline{\gamma} := \overline{B}] [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\gamma} := \overline{B}] [\overline{\beta} := A]]$$

and

$$e_{F, f} = id_{\rho[\overline{\gamma} := f]} [\overline{\phi} := \lambda \overline{A}. \llbracket \Gamma; \overline{\gamma}, \overline{\beta} \vdash F \rrbracket^{\text{Set}} \rho[\overline{\gamma} := f] [\overline{\beta} := id_A]]$$



for all  $F$  and  $B$  and  $\bar{f} : B \rightarrow B'$ , the following diagram commutes by the fact that composition of environments is componentwise together with the naturality of  $\eta$ :

$$\begin{array}{ccc}
 E_{F,B} & \xrightarrow{id_\rho[\gamma := id_B][\phi := \lambda \bar{A}. \eta_{\bar{A} \bar{B}}]} & E_{G,B} \\
 \downarrow e_{F,f} & & \downarrow e_{G,f} \\
 E_{F,B'} & \xrightarrow{id_\rho[\gamma := id_{B'}][\phi := \lambda \bar{A}. \eta_{\bar{A} \bar{B}'}]} & E_{G,B'}
 \end{array}$$

We therefore have that

$$\lambda \bar{B}. \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{H}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Set}} \rho d \bar{\eta} \bar{B}$$

is natural in  $\bar{B}$  as desired.

- To see that, for every  $\rho : \text{SetEnv}$  and  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$ , and all  $\eta : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} F G \rrbracket^{\text{Set}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{H}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Set}} \rho d \bar{\eta}$$

satisfies the additional condition necessary for it to be in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}] \rrbracket^{\text{Set}} \rho$ ,

let  $\bar{R} : \text{Rel}(B, B')$  and  $\bar{S} : \text{Rel}(C, C')$ . Since each map in  $\bar{\eta}$  satisfies the extra condition necessary for it to be in its corresponding  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} F G \rrbracket^{\text{Set}} \rho$  – i.e., since

$$(\eta_{\bar{B} \bar{C}}, \eta_{\bar{B}' \bar{C}'} ) \in \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}] \rightarrow \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash G \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}]$$

– we have that

$$\begin{aligned}
 & ((\lambda e \nu \bar{Z}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := \bar{Z}]}[\phi := \lambda \bar{A}. \nu_{\bar{A} \bar{Z}}]) d \bar{\eta} \bar{B}, \\
 & (\lambda e \nu \bar{Z}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := \bar{Z}]}[\phi := \lambda \bar{A}. \nu_{\bar{A} \bar{Z}}]) d \bar{\eta} \bar{B}') \\
 = & (\llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := B]}[\phi := \lambda \bar{A}. \eta_{\bar{A} \bar{B}}]), \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{\rho[\bar{\gamma} := B']}[\phi := \lambda \bar{A}. \eta_{\bar{A} \bar{B}'}])
 \end{aligned}$$

has type

$$\begin{aligned}
 & (\llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{F}] \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{B}] \rightarrow \llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{G}] \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{B}], \\
 & \llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{F}] \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{B}'] \rightarrow \llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{G}] \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{B}']) \\
 = & \llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{F}] \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\gamma} := \bar{R}] \rightarrow \llbracket \Gamma; \bar{\gamma} \vdash H[\phi := \bar{G}] \rrbracket^{\text{Rel}} \text{Eq}_\rho[\bar{\gamma} := \bar{R}]
 \end{aligned}$$

as desired.

- The proofs that

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{H}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Rel}}$$

is a natural transformation from  $\llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}}$  to

$$\llbracket \Gamma; \emptyset \vdash \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Rel}}$$

and that, for every  $\rho : \text{RelEnv}$  and the unique  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{\bar{H}}^{\bar{F}, \bar{G}} : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G) (\text{Nat}^{\bar{\gamma}} H[\phi :=_{\bar{\beta}} \bar{F}] H[\phi :=_{\bar{\beta}} \bar{G}]) \rrbracket^{\text{Rel}} \rho d$$

is a morphism from  $\overline{\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} F G \rrbracket^{\text{Rel}} \rho}$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F] H[\bar{\phi} :=_{\bar{\beta}} G] \rrbracket^{\text{Rel}} \rho$ ,  
are analogous.

– Finally, to see that

$$\begin{aligned} & \pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{H^{\bar{F} \bar{G}}}^{\bar{F} \bar{G}} : \text{Nat}^{\emptyset}(\overline{\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G}) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F] H[\bar{\phi} :=_{\bar{\beta}} G]) \rrbracket^{\text{Rel}} \rho) \\ = & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{H^{\bar{F} \bar{G}}}^{\bar{F} \bar{G}} : \text{Nat}^{\emptyset}(\overline{\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G}) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F] H[\bar{\phi} :=_{\bar{\beta}} G]) \rrbracket^{\text{Set}} (\pi_i \rho) \end{aligned}$$

we compute

$$\begin{aligned} & \pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{H^{\bar{F} \bar{G}}}^{\bar{F} \bar{G}} : \text{Nat}^{\emptyset}(\overline{\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G}) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F] H[\bar{\phi} :=_{\bar{\beta}} G]) \rrbracket^{\text{Rel}} \rho) \\ = & \pi_i(\lambda e \bar{v} \bar{R}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Rel}} id_{\rho[\bar{\gamma} := \bar{R}]}[\bar{\phi} := \lambda \bar{S}. v_{\bar{S}} \bar{R}]) \\ = & \lambda e \bar{v} \bar{R}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{(\pi_i \rho)[\bar{\gamma} := \pi_i \bar{R}]}[\bar{\phi} := \lambda \bar{S}. (\pi_i v)_{\pi_i \bar{S} \pi_i \bar{R}}] \\ = & \lambda d \bar{\eta} \bar{B}. \llbracket \Gamma; \bar{\phi}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} id_{(\pi_i \rho)[\bar{\gamma} := \bar{B}]}[\bar{\phi} := \lambda \bar{A}. \eta_{\bar{A} \bar{B}}] \\ = & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{map}_{H^{\bar{F} \bar{G}}}^{\bar{F} \bar{G}} : \text{Nat}^{\emptyset}(\overline{\text{Nat}^{\bar{\beta}, \bar{\gamma}} F G}) (\text{Nat}^{\bar{\gamma}} H[\bar{\phi} :=_{\bar{\beta}} F] H[\bar{\phi} :=_{\bar{\beta}} G]) \rrbracket^{\text{Set}} (\pi_i \rho) \end{aligned}$$

- We now consider  $\Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}$ .  
– To see that if  $d : \llbracket \Gamma; \emptyset \mid \emptyset \rrbracket^{\text{Set}} \rho$  then

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho d$$

is in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho$ , we first note that, for all  $\bar{B}$  and  $\bar{C}$ ,

$$\begin{aligned} & (\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho d)_{\bar{B} \bar{C}} \\ = & (in_{T_{\rho[\bar{\gamma} := \bar{C}]}})_{\bar{B}} \end{aligned}$$

does indeed map

$$\begin{aligned} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}] [\bar{\gamma} := \bar{C}] \\ = & \llbracket \Gamma; \bar{\beta}, \bar{\gamma}, \bar{\alpha} \vdash H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}] [\bar{\gamma} := \bar{C}] [\bar{\alpha} := \bar{B}] \\ = & \llbracket \Gamma; \bar{\phi}, \bar{\beta}, \bar{\gamma}, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}] [\bar{\gamma} := \bar{C}] [\bar{\alpha} := \bar{B}] \\ & [\bar{\phi} := \lambda \bar{D}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma}, \bar{\alpha} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}] [\bar{\gamma} := \bar{C}] [\bar{\alpha} := \bar{B}] [\bar{\beta} := \bar{D}]] \\ = & \llbracket \Gamma; \bar{\phi}, \bar{\gamma}, \bar{\alpha} \vdash H \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}] [\bar{\alpha} := \bar{B}] \\ & [\bar{\phi} := \lambda \bar{D}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{D}] [\bar{\gamma} := \bar{C}]] \\ = & T_{\rho[\bar{\gamma} := \bar{C}]}(\lambda \bar{D}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{D}] [\bar{\gamma} := \bar{C}]) \bar{B} \\ = & T_{\rho[\bar{\gamma} := \bar{C}]}(\mu T_{\rho[\bar{\gamma} := \bar{C}]}) \bar{B} \end{aligned}$$

to

$$\begin{aligned} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}] [\bar{\gamma} := \bar{C}] \\ = & (\lambda \bar{D}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{D}] [\bar{\gamma} := \bar{C}]) \bar{B} \\ = & (\mu T_{\rho[\bar{\gamma} := \bar{C}]}) \bar{B} \end{aligned}$$

To see that

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\bar{\phi} := (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta}] [\bar{\alpha} := \bar{\beta}] (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho d \\ = & \lambda \bar{B} \bar{C}. (in_{T_{\rho[\bar{\gamma} := \bar{C}]}})_{\bar{B}} \end{aligned}$$

is natural in  $\bar{B}$  and  $\bar{C}$ , we observe that the following diagram commutes for all  $\bar{f} : B \rightarrow B'$  and  $\bar{g} : C \rightarrow C'$ :

$$\begin{array}{ccc}
 T_{\rho[\bar{Y}:=C]} (\mu T_{\rho[\bar{Y}:=C]}) \bar{B} & \xrightarrow{(in_{T_{\rho[\bar{Y}:=C]}})_{\bar{B}}} & (\mu T_{\rho[\bar{Y}:=C]}) \bar{B} \\
 \downarrow \sigma_{id_{\rho}[\bar{Y}:=g]} (\mu \sigma_{id_{\rho}[\bar{Y}:=g]}) \bar{B} & & \downarrow (\mu \sigma_{id_{\rho}[\bar{Y}:=g]}) \bar{B} \\
 T_{\rho[\bar{Y}:=C']} (\mu T_{\rho[\bar{Y}:=C']}) \bar{B} & \xrightarrow{(in_{T_{\rho[\bar{Y}:=C']}})_{\bar{B}}} & (\mu T_{\rho[\bar{Y}:=C']}) \bar{B} \\
 \downarrow T_{\rho[\bar{Y}:=C']} (\mu T_{\rho[\bar{Y}:=C']}) \bar{f} & & \downarrow (in_{T_{\rho[\bar{Y}:=C']}})_{\bar{B}} \\
 T_{\rho[\bar{Y}:=C']} (\mu T_{\rho[\bar{Y}:=C']}) \bar{B}' & \xrightarrow{(in_{T_{\rho[\bar{Y}:=C']}})_{\bar{B}'}} & (\mu T_{\rho[\bar{Y}:=C']}) \bar{B}'
 \end{array}$$

Indeed, naturality of  $in$  with respect to its functor argument ensures that the top diagram commutes, and naturality of  $in_{T_{\rho[\bar{Y}:=C]}}$  ensures that the bottom one commutes.

- To see that  $\llbracket \Gamma; \emptyset \mid \emptyset \vdash in_H : Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho d$  satisfies the additional property necessary for it to be in

$$\llbracket \Gamma; \emptyset \vdash Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho$$

let  $\bar{R} : \text{Rel}(\bar{B}, \bar{B}')$  and  $\bar{S} : \text{Rel}(\bar{C}, \bar{C}')$ . Then

$$\begin{aligned}
 & ((\llbracket \Gamma; \emptyset \mid \emptyset \vdash in_H : Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho d)_{\bar{B}, \bar{C}}, \\
 & \quad (\llbracket \Gamma; \emptyset \mid \emptyset \vdash in_H : Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho d)_{\bar{B}', \bar{C}'} \\
 = & ((in_{T_{\rho[\bar{Y}:=C]}})_{\bar{B}}, (in_{T_{\rho[\bar{Y}:=C']}})_{\bar{B}'})
 \end{aligned}$$

has type

$$\begin{aligned}
 & ((\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}] \rightarrow \\
 & \quad \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}], \\
 & \quad \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}'][\bar{\gamma} := \bar{C}'] \rightarrow \\
 & \quad \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}'][\bar{\gamma} := \bar{C}'])) \\
 = & (T_{\rho[\bar{Y}:=C]} (\mu T_{\rho[\bar{Y}:=C]}) \bar{B} \rightarrow (\mu T_{\rho[\bar{Y}:=C]}) \bar{B}, T_{\rho[\bar{Y}:=C']} (\mu T_{\rho[\bar{Y}:=C']}) \bar{B}' \rightarrow (\mu T_{\rho[\bar{Y}:=C']}) \bar{B}') \\
 = & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Rel}} \text{Eq}_{\rho} [\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}] \rightarrow \\
 & \quad \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}} \text{Eq}_{\rho} [\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}]
 \end{aligned}$$

- The proofs that  $\llbracket \Gamma; \emptyset \mid \emptyset \vdash in_H : Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}}$  is a natural transformation from  $\llbracket \Gamma; \emptyset \mid \emptyset \rrbracket^{\text{Rel}}$  to  $\llbracket \Gamma; \emptyset \vdash Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}}$  and that, for all  $\rho : \text{RelEnv}$  and  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}}$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash in_H : Nat^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}} \rho d$$

is a natural transformation from  $\lambda \bar{R} \bar{S}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Rel}} \rho[\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}]$  to  $\lambda \bar{R} \bar{S}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}} \rho[\bar{\beta} := \bar{R}][\bar{\gamma} := \bar{S}]$ , are analogous.

- Finally, to see that  $\pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}} \rho d) = \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} (\pi_i \rho) (\pi_i d)$  we first note that  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho$  and  $\pi_i d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} (\pi_i \rho)$  are uniquely determined. Further, the definition of natural transformations in Rel ensures that, for any  $\bar{R}$  and  $\bar{S}$ ,

$$\begin{aligned} & (in_{T_{\rho[\bar{Y} := \bar{S}]}})_{\bar{R}} \\ &= ((in_{\pi_1(T_{\rho[\bar{Y} := \bar{S}]})})_{\pi_1 \bar{R}}, (in_{\pi_2(T_{\rho[\bar{Y} := \bar{S}]})})_{\pi_2 \bar{R}}) \\ &= ((in_{T_{(\pi_1 \rho)[\bar{Y} := \pi_1 \bar{S}]}}^{\text{Set}})_{\pi_1 \bar{R}}, (in_{T_{(\pi_2 \rho)[\bar{Y} := \pi_2 \bar{S}]}}^{\text{Set}})_{\pi_2 \bar{R}}) \end{aligned}$$

Observing that  $\pi_1$  and  $\pi_2$  are surjective, we therefore have that

$$\begin{aligned} & \pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Rel}} \rho d) \\ &= \pi_i(\lambda \bar{R} \bar{S}. (in_{T_{\rho[\bar{Y} := \bar{S}]}})_{\bar{R}}) \\ &= \lambda \bar{B} \bar{C}. (in_{T_{(\pi_i \rho)[\bar{Y} := \bar{C}]}}^{\text{Set}})_{\bar{B}} \\ &= \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{in}_H : \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta}][\bar{\alpha} := \bar{\beta}] (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} (\pi_i \rho) (\pi_i d) \end{aligned}$$

- We now consider  $\Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F)$ .  
– To see that  $\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}}$  is a natural transformation from  $\llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}}$  to

$$\llbracket \Gamma; \emptyset \vdash \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}}$$

since the functorial part  $\Phi$  of the context is empty, we need only show that, for every  $\rho : \text{SetEnv}$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}} \rho$$

is a morphism in Set from  $\llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$  to

$$\llbracket \Gamma; \emptyset \vdash \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}} \rho$$

i.e., that, for the unique  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}} \rho d$$

is a morphism from  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F \rrbracket^{\text{Set}} \rho$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F \rrbracket^{\text{Set}} \rho$ .

For this we show that for every  $\eta : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F \rrbracket^{\text{Set}} \rho$  we have

$$\begin{aligned} & \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta \\ & : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F \rrbracket^{\text{Set}} \rho \end{aligned}$$

To this end we show that, for any  $\bar{B}$  and  $\bar{C}$ ,

$$(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta)_{\bar{B} \bar{C}}$$

is a morphism from

$$\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu\phi.\lambda\bar{\alpha}.H)\bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}] = (\mu T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}})_{\bar{B}}$$

to

$$\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$$

To see this, we use Equations 6 and 8 for the first and second equalities below, together with weakening, to see that  $\eta$  is itself a natural transformation from

$$\begin{aligned}
 & \lambda \bar{B} \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}] \\
 = & \lambda \bar{B} \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma}, \bar{\alpha} \vdash H[\phi := F] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}][\bar{\alpha} := \bar{B}] \\
 = & \lambda \bar{B} \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma}, \bar{\alpha}, \phi \vdash H \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}][\bar{\alpha} := \bar{B}] \\
 & [\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma}, \bar{\alpha} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}][\bar{\alpha} := \bar{B}][\bar{\beta} := \bar{A}]] \\
 = & \lambda \bar{B} \bar{C}. \llbracket \Gamma; \bar{\gamma}, \bar{\alpha}, \phi \vdash H \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}][\bar{\alpha} := \bar{B}][\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}][\bar{\beta} := \bar{A}]] \\
 = & \lambda \bar{B} \bar{C}. T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}]) \bar{B}
 \end{aligned}$$

to

$$\lambda \bar{B} \bar{C}. (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}]) \bar{B} = \lambda \bar{B} \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$$

Thus, if  $x : \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}] = (\mu T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}) \bar{B}$ , then

$$\begin{aligned}
 & (\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta)_{\bar{B} \bar{C}} x \\
 = & (\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}}))_{\bar{B}} x \\
 : & (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}]) \bar{B}
 \end{aligned}$$

i.e., for each  $\bar{B}$  and  $\bar{C}$

$$(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu \phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta)_{\bar{B} \bar{C}}$$

is a morphism from  $(\mu T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}) \bar{B}$  to  $\llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$ .

To see that this family of morphisms is natural in  $\bar{B}$  and  $\bar{C}$ , we observe that the following diagram commutes for all  $\bar{f} : \bar{B} \rightarrow \bar{B}'$  and  $\bar{g} : \bar{C} \rightarrow \bar{C}'$ :

$$\begin{array}{ccc}
 (\mu T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}) \bar{B} & \xrightarrow{(\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}}))_{\bar{B}}} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}][\bar{\beta} := \bar{B}] \\
 \downarrow (\mu \sigma_{\text{id} \rho[\bar{\gamma} := \bar{g}]}^{\text{Set}})_{\bar{B}} & & \downarrow \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\gamma} := \bar{g}][\bar{\beta} := \text{id}_{\bar{B}}]} \\
 (\mu T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}) \bar{B} & \xrightarrow{(\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}'}))_{\bar{B}}} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}'][\bar{\beta} := \bar{B}] \\
 \downarrow (\mu T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}) \bar{f} & & \downarrow \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\gamma} := \text{id}_{\bar{C}'}][\bar{\beta} := \bar{f}]} \\
 (\mu T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}) \bar{B}' & \xrightarrow{(\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}'}))_{\bar{B}'}} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\gamma} := \bar{C}'][\bar{\beta} := \bar{B}']
 \end{array}$$

Indeed, naturality of  $\text{fold}_{T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}'})$  ensures that the bottom diagram commutes. To see that the top one commutes is considerably more delicate.

To see that the top diagram commutes we first observe that, given a natural transformation  $\Theta : H \rightarrow K : [\text{Set}^k, \text{Set}] \rightarrow [\text{Set}^k, \text{Set}]$ , the fixpoint natural transformation  $\mu \Theta : \mu H \rightarrow$

$\mu K : \text{Set}^k \rightarrow \text{Set}$  is defined to be  $\text{fold}_H(\Theta(\mu K) \circ \text{in}_K)$ , i.e., the unique morphism making the following square commute:

$$\begin{array}{ccc} H(\mu H) & \xrightarrow{H(\mu\Theta)} & H(\mu K) \\ \text{in}_H \downarrow & & \downarrow \Theta(\mu K) \\ \mu H & \xrightarrow{\mu\Theta} & \mu K \end{array}$$

Taking  $\Theta = \sigma_f^{\text{Set}} : T_\rho^{\text{Set}} \rightarrow T_{\rho'}^{\text{Set}}$  gives that the following diagram commutes for any morphism of set environments  $f : \rho \rightarrow \rho'$ :

$$\begin{array}{ccc} T_\rho^{\text{Set}}(\mu T_\rho^{\text{Set}}) & \xrightarrow{T_\rho^{\text{Set}}(\mu\sigma_f^{\text{Set}})} & T_{\rho'}^{\text{Set}}(\mu T_{\rho'}^{\text{Set}}) \\ \text{in}_{T_\rho^{\text{Set}}} \downarrow & & \downarrow \sigma_f^{\text{Set}}(\mu T_{\rho'}^{\text{Set}}) \\ \mu T_\rho^{\text{Set}} & \xrightarrow{\mu\sigma_f^{\text{Set}}} & \mu T_{\rho'}^{\text{Set}} \end{array} \quad (17)$$

We next observe that the action of the functor

$$\lambda \bar{B}. \lambda \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$$

on the morphisms  $\bar{f} : \bar{B} \rightarrow \bar{B}', \bar{g} : \bar{C} \rightarrow \bar{C}'$  is given by

$$\begin{aligned} & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash H[\phi := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \text{id}_\rho[\bar{\beta} := \bar{f}][\bar{\gamma} := \bar{g}] \\ = & \llbracket \Gamma; \bar{\alpha}, \bar{\gamma} \vdash H[\phi := F] \rrbracket^{\text{Set}} \text{id}_\rho[\bar{\alpha} := \bar{f}][\bar{\gamma} := \bar{g}] \\ = & \llbracket \Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} \text{id}_\rho[\bar{\alpha} := \bar{f}][\bar{\gamma} := \bar{g}][\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\beta} := \bar{A}]}[\bar{\gamma} := \bar{g}]] \\ = & \llbracket \Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} (\text{id}_{\rho[\bar{\gamma} := \bar{C}']}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']][\bar{\alpha} := \bar{f}] \\ & \quad \circ \text{id}_{\rho[\bar{\alpha} := \bar{B}]}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']][\bar{\gamma} := \bar{g}]] \\ & \quad \circ \text{id}_{\rho[\bar{\alpha} := \bar{B}][\bar{\gamma} := \bar{C}]}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\beta} := \bar{A}]}[\bar{\gamma} := \bar{g}]] \\ = & \llbracket \Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\gamma} := \bar{C}']}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']][\bar{\alpha} := \bar{f}] \\ & \quad \circ \llbracket \Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\alpha} := \bar{B}]}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']][\bar{\gamma} := \bar{g}] \\ & \quad \circ \llbracket \Gamma; \phi, \bar{\alpha}, \bar{\gamma} \vdash H \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\alpha} := \bar{B}][\bar{\gamma} := \bar{C}]}[\phi := \lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\beta} := \bar{A}]}[\bar{\gamma} := \bar{g}]] \\ = & T_{\rho[\bar{\gamma} := \bar{C}']}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']) \bar{f} \\ & \quad \circ (\sigma_{\text{id}_{\rho[\bar{\gamma} := \bar{g}]}}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{A}][\bar{\gamma} := \bar{C}']))_{\bar{B}} \\ & \quad \circ (T_{\rho[\bar{\gamma} := \bar{C}]}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \text{id}_{\rho[\bar{\beta} := \bar{A}]}[\bar{\gamma} := \bar{g}]))_{\bar{B}} \end{aligned}$$

So, if  $\eta$  is a natural transformation from

$$\lambda \bar{B}. \lambda \bar{C}. \llbracket \Gamma; \bar{\alpha}, \bar{\gamma} \vdash H[\phi := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$$

to

$$\lambda \bar{B}. \lambda \bar{C}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}]$$

then, by naturality,

$$\begin{aligned}
 & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := f][\bar{\gamma} := g]} \circ \eta_{\bar{B}, \bar{C}} \\
 &= \eta_{\bar{B}', \bar{C}'} \circ \llbracket \Gamma; \bar{\alpha}, \bar{\gamma} \vdash H[\bar{\alpha} := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := f][\bar{\gamma} := g]} \\
 &= \eta_{\bar{B}', \bar{C}'} \circ T_{\rho[\bar{\gamma} := C']}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := A][\bar{\gamma} := C']) \bar{f} \\
 &\quad \circ (\sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := A][\bar{\gamma} := C']))_{\bar{B}} \\
 &\quad \circ (T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := A][\bar{\gamma} := g]}))_{\bar{B}}
 \end{aligned}$$

As a special case when  $\bar{f} = id_{\bar{B}}$  we have

$$\begin{aligned}
 & \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := B][\bar{\gamma} := g]} \circ \eta_{\bar{B}, \bar{C}} \\
 &= \eta_{\bar{B}, \bar{C}'} \circ (\sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := A][\bar{\gamma} := C']))_{\bar{B}} \\
 &\quad \circ (T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := A][\bar{\gamma} := g]}))_{\bar{B}}
 \end{aligned}$$

i.e.,

$$\begin{aligned}
 & \lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := B][\bar{\gamma} := g]} \circ \lambda \bar{B}. \eta_{\bar{B}, \bar{C}} \\
 &= \lambda \bar{B}. \eta_{\bar{B}, \bar{C}'} \circ \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := A][\bar{\gamma} := C']) \quad (18) \\
 &\quad \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\lambda \bar{A}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} id_{\rho[\bar{\beta} := A][\bar{\gamma} := g]})
 \end{aligned}$$

Now, to see that the top diagram in the diagram on page 35 commutes we first note that the diagram

$$\begin{array}{ccc}
 T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\mu T_{\rho[\bar{\gamma} := C]}^{\text{Set}}) & \xrightarrow{T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ \mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}}} & T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']) \\
 \downarrow \text{in}_{T_{\rho[\bar{\gamma} := C]}^{\text{Set}}} & & \downarrow \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']) \\
 & & T_{\rho[\bar{\gamma} := C']}^{\text{Set}} (\lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']) \\
 & & \downarrow \lambda \bar{A}. \eta_{\bar{A}, \bar{C}'} \\
 \mu T_{\rho[\bar{\gamma} := C]}^{\text{Set}} & \xrightarrow{\mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}}} \mu T_{\rho[\bar{\gamma} := C']}^{\text{Set}} \xrightarrow{\text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'}))} & \lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']
 \end{array}$$

commutes because

$$\begin{aligned}
 & \lambda \bar{A}. \eta_{\bar{A}, \bar{C}'} \circ \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']) \\
 &\quad \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ \mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} \\
 &= \lambda \bar{A}. \eta_{\bar{A}, \bar{C}'} \circ \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\lambda \bar{B}. \llbracket \Gamma; \bar{\beta}, \bar{\gamma} \vdash F \rrbracket^{\text{Set}} \rho[\bar{\beta} := B][\bar{\gamma} := C']) \\
 &\quad \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}}) \\
 &= \lambda \bar{A}. \eta_{\bar{A}, \bar{C}'} \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\mu T_{\rho[\bar{\gamma} := C']}^{\text{Set}}) \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}}) \\
 &= \text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ \text{in}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} \circ \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} (\mu T_{\rho[\bar{\gamma} := C']}^{\text{Set}}) \circ T_{\rho[\bar{\gamma} := C]}^{\text{Set}} (\mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}}) \\
 &= \text{fold}_{T_{\rho[\bar{\gamma} := C']}^{\text{Set}}} (\lambda \bar{A}. \eta_{\bar{A}, \bar{C}'})) \circ \mu \sigma_{id_{\rho[\bar{\gamma} := g]}}^{\text{Set}} \circ \text{in}_{T_{\rho[\bar{\gamma} := C]}^{\text{Set}}}
 \end{aligned}$$



Here, the first equality is by functoriality of  $T_{\rho[\gamma:=C]}^{\text{Set}}$ , the second equality is by naturality of  $\sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}$ , the third equality by the universal property of  $\text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})$  and the last equality by Equation 17. That is, we have

$$\begin{aligned} & \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}) \circ \mu\sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}} \\ &= \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}} \circ \sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}'])) \end{aligned} \quad (19)$$

Next, we note that the diagram

$$\begin{array}{ccc} T_{\rho[\gamma:=C]}^{\text{Set}}(\mu T_{\rho[\gamma:=C]}^{\text{Set}}) & \xrightarrow{T_{\rho[\gamma:=C]}^{\text{Set}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}))} & T_{\rho[\gamma:=C]}^{\text{Set}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}']) \\ \downarrow \text{in}_{T_{\rho[\gamma:=C]}^{\text{Set}}} & & \downarrow \sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}']) \\ \mu T_{\rho[\gamma:=C]}^{\text{Set}} & \xrightarrow{\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}']} & \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}'] \\ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}) & \searrow & \downarrow \lambda\bar{A}.\eta_{\bar{A},\bar{C}} \\ & & \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \end{array}$$

commutes because

$$\begin{aligned} & \lambda\bar{A}.\eta_{\bar{A},\bar{C}} \circ \sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}']) \\ & \quad \circ T_{\rho[\gamma:=C]}^{\text{Set}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})) \\ &= \lambda\bar{A}.\eta_{\bar{A},\bar{C}} \circ \sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}']) \\ & \quad \circ T_{\rho[\gamma:=C]}^{\text{Set}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}]) \circ T_{\rho[\gamma:=C]}^{\text{Set}}(\text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})) \\ &= \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \lambda\bar{A}.\eta_{\bar{A},\bar{C}} \circ T_{\rho[\gamma:=C]}^{\text{Set}}(\text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})) \\ &= \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}) \circ \text{in}_{T_{\rho[\gamma:=C]}^{\text{Set}}} \end{aligned}$$

Here, the first equality is by functoriality of  $T_{\rho[\gamma:=C]}^{\text{Set}}$ , the second equality is by Equation 18, and the last equality is by the universal property of  $\text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})$ . That is, we have

$$\begin{aligned} & \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}) \\ &= \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}} \circ \sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}}(\lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} \rho[\bar{\beta} := \bar{B}][\bar{\gamma} := \bar{C}'])) \end{aligned} \quad (20)$$

Combining Equations 19 and 20 we get that

$$\text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}}) \circ \mu\sigma_{id_{\rho[\gamma:=g]}^{\text{Set}}} = \lambda\bar{B}.\llbracket\Gamma;\bar{\beta},\bar{\gamma} \vdash F\rrbracket^{\text{Set}} id_{\rho[\bar{\beta}:=\bar{B}]}[\bar{\gamma}:=\bar{g}] \circ \text{fold}_{T_{\rho[\gamma:=C]}^{\text{Set}}}(\lambda\bar{A}.\eta_{\bar{A},\bar{C}})$$

i.e., that the top diagram in the diagram on page 35 commutes. We therefore have that

$$(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta)_{\bar{B}\bar{C}}$$

is natural in  $\bar{B}$  and  $\bar{C}$  as desired.

– To see that, for every  $\rho : \text{SetEnv}$ ,  $d \in \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Set}} \rho$ , and  $\eta : \llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F \rrbracket^{\text{Set}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} \rho d \eta$$

satisfies the additional condition necessary for it to be in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F \rrbracket^{\text{Set}} \rho$ , let  $\bar{R} : \text{Rel}(\bar{B}, \bar{B}')$  and  $\bar{S} : \text{Rel}(\bar{C}, \bar{C}')$ . Since  $\eta$  satisfies the additional condition necessary for it to be in  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} (H[\phi := F][\bar{\alpha} := \bar{\beta}]) F \rrbracket^{\text{Set}} \rho$  – i.e., since

$$\begin{aligned} (\eta_{\bar{B}\bar{C}}, \eta_{\bar{B}'\bar{C}'}) &\in \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash H[\phi := F][\bar{\alpha} := \bar{\beta}] \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \rightarrow \\ &\quad \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \\ &= T_{\text{Eq}_{\rho}[\bar{\gamma} := \bar{S}]}(\llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}]) \rightarrow \\ &\quad \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \end{aligned}$$

– we have that

$$((\text{fold}_{T^{\text{Set}}_{\rho[\bar{\gamma} := \bar{C}]}}(\lambda \bar{A}. \eta_{\bar{A}\bar{C}}))_{\bar{B}}, (\text{fold}_{T^{\text{Set}}_{\rho[\bar{\gamma} := \bar{C}']}}(\lambda \bar{A}. \eta_{\bar{A}\bar{C}'}))_{\bar{B}'})$$

has type

$$\begin{aligned} &(\mu T_{\text{Eq}_{\rho}[\bar{\gamma} := \bar{S}]} \bar{R}) \rightarrow \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \\ &= (\mu T_{\text{Eq}_{\rho}[\bar{\gamma} := \bar{S}]} \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash \beta \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}]) \rightarrow \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \\ &= \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \rightarrow \llbracket \Gamma; \bar{\gamma}, \bar{\beta} \vdash F \rrbracket^{\text{Rel}} \text{Eq}_{\rho}[\bar{\gamma} := \bar{S}][\bar{\beta} := \bar{R}] \end{aligned}$$

as desired.

– The proofs that

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}} \rho$$

is a natural transformation from  $\llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}}$  to

$$\llbracket \Gamma; \emptyset \vdash \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}} \rho$$

and that, for all  $\rho : \text{RelEnv}$  and the unique  $d : \llbracket \Gamma; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho$ ,

$$\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}} \rho d$$

is a morphism from  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F \rrbracket^{\text{Rel}} \rho$  to  $\llbracket \Gamma; \emptyset \vdash \text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F \rrbracket^{\text{Rel}} \rho$ , are analogous.

– Finally, to see that

$$\begin{aligned} &\pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}} \rho) \\ &= \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi :=_{\bar{\beta}} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda \bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Set}} (\pi_i \rho) \end{aligned}$$

we compute

$$\begin{aligned}
& \pi_i(\llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := \bar{\beta} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda\bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}}) \\
&= \pi_i(\lambda e \eta \bar{R} \bar{S}. (\text{fold}_{T_{\rho[\bar{\gamma} := \bar{S}]}} (\lambda \bar{Z}. \eta \bar{Z} \bar{S}))_{\bar{R}}) \\
&= \lambda e \eta \bar{R} \bar{S}. (\text{fold}_{T_{(\pi_i \rho)[\bar{\gamma} := \pi_i \bar{S}]}} (\lambda \bar{Z}. (\pi_i \eta)_{\pi_i \bar{Z} \pi_i \bar{S}}))_{\pi_i \bar{R}} \\
&= \lambda d \eta \bar{B} \bar{C}. (\text{fold}_{T_{(\pi_i \rho)[\bar{\gamma} := \bar{C}]}} (\lambda \bar{A}. \eta_{\bar{A} \bar{C}}))_{\bar{B}} \\
&= \llbracket \Gamma; \emptyset \mid \emptyset \vdash \text{fold}_H^F : \text{Nat}^0 (\text{Nat}^{\bar{\beta}, \bar{\gamma}} H[\phi := \bar{\beta} F][\bar{\alpha} := \bar{\beta}] F) (\text{Nat}^{\bar{\beta}, \bar{\gamma}} (\mu\phi. \lambda\bar{\alpha}. H) \bar{\beta} F) \rrbracket^{\text{Rel}}(\pi_i \rho)
\end{aligned}$$

Here, we are again using the fact that  $\pi_1$  and  $\pi_2$  are surjective.

□

The Abstraction Theorem is now the special case of Theorem 28 for closed terms of close type:

State more generally as: If  $(a, b) \in \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho$  then  $(\llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}(\pi_1 \rho)a, \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}}(\pi_2 \rho)b) \in \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Set}} \rho$ . Get the next theorem as a corollary for closed terms of closed type.

THEOREM 29. If  $\vdash \tau : \mathcal{F}$  and  $\vdash t : \tau$ , then  $(\llbracket \vdash t : \tau \rrbracket^{\text{Set}}, \llbracket \vdash t : \tau \rrbracket^{\text{Set}}) \in \llbracket \vdash \tau \rrbracket^{\text{Rel}}$ .

## 6 FREE THEOREMS FOR NESTED TYPES

### 6.1 Free Theorem for Type of Polymorphic Bottom

Suppose  $\vdash g : \text{Nat}^\alpha \perp \alpha$ , let  $G^{\text{Set}} = \llbracket \vdash g : \text{Nat}^\alpha \perp \alpha \rrbracket^{\text{Set}}$ , and let  $G^{\text{Rel}} = \llbracket \vdash g : \text{Nat}^\alpha \perp \alpha \rrbracket^{\text{Rel}}$ . By Theorem 29,  $(G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) = G^{\text{Rel}} \rho$ . Thus, for all  $\rho \in \text{RelEnv}$  and any  $(a, b) \in \llbracket \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$ , eliding the only possible instantiations of  $a$  and  $b$  gives that

$$\begin{aligned}
(G^{\text{Set}}, G^{\text{Set}}) &= (G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) \in \llbracket \vdash \text{Nat}^\alpha \perp \alpha \rrbracket^{\text{Rel}} \rho \\
&= \{\eta : K_1 \Rightarrow \text{id}\} \\
&= \{(\eta_1 : K_1 \Rightarrow \text{id}, \eta_2 : K_1 \Rightarrow \text{id})\}
\end{aligned}$$

That is,  $G^{\text{Set}}$  is a natural transformation from the constantly 1-valued functor to the identity functor in Set. In particular, for every  $S : \text{Set}$ ,  $G_S^{\text{Set}} : 1 \rightarrow S$ . Note, however, that if  $S = \emptyset$ , then there can be no such morphism, so no such natural transformation can exist in Set, and thus no term  $\vdash g : \text{Nat}^\alpha \perp \alpha$  can exist in our calculus. That is, our calculus does not admit any terms with the closed type  $\text{Nat}^\alpha \perp \alpha$  of the polymorphic bottom.

### 6.2 Free Theorem for Type of Polymorphic Identity

Suppose  $\vdash g : \text{Nat}^\alpha \alpha \alpha$ , let  $G^{\text{Set}} = \llbracket \vdash g : \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Set}}$ , and let  $G^{\text{Rel}} = \llbracket \vdash g : \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Rel}}$ . By Theorem 29,  $(G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) = G^{\text{Rel}} \rho$ . Thus, for all  $\rho \in \text{RelEnv}$  and any  $(a, b) \in \llbracket \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$ , eliding the only possible instantiations of  $a$  and  $b$  gives that

$$\begin{aligned}
(G^{\text{Set}}, G^{\text{Set}}) &= (G^{\text{Set}}(\pi_1 \rho), G^{\text{Set}}(\pi_2 \rho)) \in \llbracket \vdash \text{Nat}^\alpha \alpha \alpha \rrbracket^{\text{Rel}} \rho \\
&= \{\eta : \text{id} \Rightarrow \text{id}\} \\
&= \{(\eta_1 : \text{id} \Rightarrow \text{id}, \eta_2 : \text{id} \Rightarrow \text{id})\}
\end{aligned}$$

That is,  $G^{\text{Set}}$  is a natural transformation from the identity functor on Set to itself.

Now let  $S$  be any set. If  $S = \emptyset$ , then there is exactly one morphism  $\text{id}_S : S \rightarrow S$ , so  $G_S^{\text{Set}} : S \rightarrow S$  must be  $\text{id}_S$ . If  $S \neq \emptyset$ , then if  $a$  is any element of  $S$  and  $K_a : S \rightarrow S$  is the constantly  $a$ -valued morphism on  $S$ , then instantiating the naturality square implied by the above equality gives that  $G_S^{\text{Set}} \circ K_a = K_a \circ G_S^{\text{Set}}$ , i.e.,  $G_S^{\text{Set}} a = a$ , i.e.,  $G_S^{\text{Set}} = \text{id}_S$ . Putting these two cases together we have that for every  $S : \text{Set}$ ,  $G_S^{\text{Set}} = \text{id}_S$ , i.e.,  $G^{\text{Set}}$  is the identity natural transformation for the identity functor on Set. So every closed term  $g$  of closed type  $\text{Nat}^\alpha \alpha \alpha$  always denotes the identity natural

transformation for the identity functor on  $\text{Set}$ , i.e., every closed term  $g$  of type  $\text{Nat}^\alpha \alpha$  denotes the polymorphic identity function.

### 6.3 Free Theorem for Type of filter for Lists

Let  $\text{List } \alpha = (\mu\phi.\lambda\beta.\mathbb{1} + \beta \times \phi\beta)\alpha$ , and let  $\text{map} = \text{map}_{\lambda A. \llbracket \emptyset; \alpha \vdash \text{List } \alpha \rrbracket^{\text{Set}} \rho[\alpha := A]}$ .

LEMMA 30. *If  $g : A \rightarrow B$ ,  $\rho : \text{RelEnv}$ , and  $\rho\alpha = (A, B, \langle g \rangle)$ , then  $\llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho = \langle \text{map } g \rangle$*

PROOF.

$$\begin{aligned}
 & \llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho \\
 &= \mu T_\rho (\llbracket \alpha; \emptyset \vdash \alpha \rrbracket^{\text{Rel}} \rho) \\
 &= \mu T_\rho (A, B, \langle g \rangle) \\
 &= (\mu T_{\pi_1 \rho} A, \mu T_{\pi_2 \rho} B, \lim_{n \in \mathbb{N}} (T_\rho^n K_0)^* (A, B, \langle g \rangle)) \\
 &= (\text{List } A, \text{List } B, \lim_{n \in \mathbb{N}} \sum_{i=0}^n (A, B, \langle g \rangle)^i) \\
 &= (\text{List } A, \text{List } B, \text{List } (A, B, \langle g \rangle)) \\
 &= (\text{List } A, \text{List } B, \langle \text{map } g \rangle)
 \end{aligned}$$

The first equality is by Definition 16, the third equality is by Equation 3, and the fourth and sixth equalities are by Equations 21 and 22 below.

The following sequence of equalities shows

$$(T_\rho^n K_0)^* R = \sum_{i=0}^n R^i \quad (21)$$

by induction on  $n$ :

$$\begin{aligned}
 & (T_\rho^n K_0)^* R \\
 &= T_\rho^{\text{Rel}} (T_\rho^{n-1} K_0)^* R \\
 &= \llbracket \alpha; \phi, \beta \vdash \mathbb{1} + \beta \times \phi\beta \rrbracket^{\text{Set}} \rho[\phi := (T_\rho^{n-1} K_0)^*][\beta := R] \\
 &= \mathbb{1} + R \times (T_\rho^{n-1} K_0)^* R \\
 &= \mathbb{1} + R \times (\sum_{i=0}^{n-1} R^i) \\
 &= \sum_{i=0}^n R^i
 \end{aligned}$$

The following reasoning shows

$$\text{List } (A, B, \langle g \rangle) = \langle \text{map } g \rangle \quad (22)$$

By showing that  $(xs, xs') \in \text{List } (A, B, \langle g \rangle)$  if and only if  $(xs, xs') \in \langle \text{map } g \rangle$ :

$$\begin{aligned}
 & (xs, xs') \in \text{List } (A, B, \langle g \rangle) \\
 & \iff \forall i. (xs_i, xs'_i) \in \langle g \rangle \\
 & \iff \forall i. xs'_i = g(xs_i) \\
 & \iff xs' = \text{map } g \, xs \\
 & \iff (xs, xs') \in \langle \text{map } g \rangle
 \end{aligned}$$

□

THEOREM 31. *If  $\Gamma; \Phi \mid \Delta \vdash t : \tau$  and  $\rho \in \text{RelEnv}$ , and if  $(a, b) \in \llbracket \Gamma; \Phi \vdash \Delta \rrbracket^{\text{Rel}} \rho$ , then  $(\llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} (\pi_1 \rho) a, \llbracket \Gamma; \Phi \mid \Delta \vdash t : \tau \rrbracket^{\text{Set}} (\pi_2 \rho) b) \in \llbracket \Gamma; \Phi \vdash \tau \rrbracket^{\text{Rel}} \rho$*

PROOF. Immediate from Theorem 28 (at-gen).  $\square$

THEOREM 32. If  $g : A \rightarrow B$ ,  $\rho : \text{RelEnv}$ ,  $\rho\alpha = (A, B, \langle g \rangle)$ ,  $(a, b) \in \llbracket \alpha; \emptyset \vdash \Delta \rrbracket^{\text{Rel}} \rho$ ,  $(s \circ g, s) \in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 \alpha \text{ Bool} \rrbracket^{\text{Rel}} \rho$ , and, for some well-formed term filter,

$$t = \llbracket \alpha; \emptyset \mid \Delta \vdash \text{filter} : \text{Nat}^0 (\text{Nat}^0 \alpha \text{ Bool}) (\text{Nat}^0 (\text{List } \alpha) (\text{List } \alpha)) \rrbracket^{\text{Set}}, \text{ then}$$

$$\text{map } g \circ t(\pi_1 \rho) a (s \circ g) = t(\pi_2 \rho) b s \circ \text{map } g$$

PROOF. By Theorem 31,  $(t(\pi_1 \rho) a, t(\pi_2 \rho) b) \in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 (\text{Nat}^0 \alpha \text{ Bool}) (\text{Nat}^0 (\text{List } \alpha) (\text{List } \alpha)) \rrbracket^{\text{Rel}} \rho$ . Thus if  $(s, s') \in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 \alpha \text{ Bool} \rrbracket^{\text{Rel}} \rho = \rho\alpha \rightarrow \text{Eq}_{\text{Bool}}$ , then

$$\begin{aligned} (t(\pi_1 \rho) a s, t(\pi_2 \rho) b s') &\in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 (\text{List } \alpha) (\text{List } \alpha) \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho \end{aligned}$$

So if  $(xs, xs') \in \llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho$  then,

$$(t(\pi_1 \rho) a s xs, t(\pi_2 \rho) b s' xs') \in \llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho \quad (23)$$

Consider the case in which  $\rho\alpha = (A, B, \langle g \rangle)$ . Then  $\llbracket \alpha; \emptyset \vdash \text{List } \alpha \rrbracket^{\text{Rel}} \rho = \langle \text{map } g \rangle$ , by Lemma 30, and  $(xs, xs') \in \langle \text{map } g \rangle$  implies  $xs' = \text{map } g xs$ . We also have that  $(s, s') \in \langle g \rangle \rightarrow \text{Eq}_{\text{Bool}}$  implies  $\forall (x, gx) \in \langle g \rangle. sx = s'(gx)$  and thus  $s = s' \circ g$  due to the definition of morphisms between relations. With these instantiations, Equation 24 becomes

$$\begin{aligned} (t(\pi_1 \rho) a (s' \circ g) xs, t(\pi_2 \rho) b s' (\text{map } g xs)) &\in \langle \text{map } g \rangle, \\ \text{i.e.,} \\ \text{map } g (t(\pi_1 \rho) a (s' \circ g) xs) &= t(\pi_2 \rho) b s' (\text{map } g xs), \\ \text{i.e.,} \\ \text{map } g \circ t(\pi_1 \rho) a (s' \circ g) &= t(\pi_2 \rho) b s' \circ \text{map } g \end{aligned}$$

as desired.  $\square$

## 6.4 Free Theorem for Type of filter for GRose

THEOREM 33. Let  $g : A \rightarrow B$  be a function,  $\eta : F \rightarrow G$  a natural transformation of Set functors,  $\rho : \text{RelEnv}$ ,  $\rho\alpha = (A, B, \langle g \rangle)$ ,  $\rho\psi = (F, G, \langle \eta \rangle)$ ,  $(a, b) \in \llbracket \alpha, \psi; \emptyset \vdash \Delta \rrbracket^{\text{Rel}} \rho$ , and  $(s \circ g, s) \in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 \alpha \text{ Bool} \rrbracket^{\text{Rel}} \rho$ . Then, for any well-formed term filter, if we call

$$t = \llbracket \alpha, \psi; \emptyset \mid \Delta \vdash \text{filter} : \text{Nat}^0 (\text{Nat}^0 \alpha \text{ Bool}) (\text{Nat}^0 (\text{List } \alpha) (\text{List } \alpha)) \rrbracket^{\text{Set}}$$

we have that

$$\text{map } \eta (g + 1) \circ t(\pi_1 \rho) a (s \circ g) = t(\pi_2 \rho) b s \circ \text{map } \eta g$$

PROOF. By Theorem 31,

$$(t(\pi_1 \rho) a, t(\pi_2 \rho) b) \in \llbracket \alpha, \psi; \emptyset \vdash \text{Nat}^0 (\text{Nat}^0 \alpha \text{ Bool}) (\text{Nat}^0 (\text{List } \alpha) (\text{List } \alpha)) \rrbracket^{\text{Rel}} \rho$$

Thus if  $(s, s') \in \llbracket \alpha; \emptyset \vdash \text{Nat}^0 \alpha \text{ Bool} \rrbracket^{\text{Rel}} \rho = \rho\alpha \rightarrow \text{Eq}_{\text{Bool}}$ , then

$$\begin{aligned} (t(\pi_1 \rho) a s, t(\pi_2 \rho) b s') &\in \llbracket \alpha, \psi; \emptyset \vdash \text{Nat}^0 (\text{GRose } \psi \alpha) (\text{GRose } \psi (\alpha + \mathbb{1})) \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \alpha, \psi; \emptyset \vdash \text{GRose } \psi \alpha \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \alpha, \psi; \emptyset \vdash \text{GRose } \psi (\alpha + \mathbb{1}) \rrbracket^{\text{Rel}} \rho \end{aligned}$$

So if  $(xs, xs') \in \llbracket \alpha; \emptyset \vdash \text{GRose } \psi \alpha \rrbracket^{\text{Rel}} \rho$  then,

$$(t(\pi_1 \rho) a \ s \ xs, t(\pi_2 \rho) b \ s' \ xs') \in \llbracket \alpha, \psi; \emptyset \vdash \text{GRose } \psi (\alpha + \mathbb{1}) \rrbracket^{\text{Rel}} \rho \quad (24)$$

Since  $\rho \alpha = (A, B, \langle g \rangle)$  and  $\rho \psi = (F, G, \langle \psi \rangle)$ , then  $\llbracket \alpha, \psi; \emptyset \vdash \text{GRose } \psi \alpha \rrbracket^{\text{Rel}} \rho = \langle \text{map } \eta g \rangle$  and  $\llbracket \alpha, \psi; \emptyset \vdash \text{GRose } \psi (\alpha + \mathbb{1}) \rrbracket^{\text{Rel}} \rho = \langle \text{map } \eta (g + 1) \rangle$ , by Lemma 30. Moreover,  $(xs, xs') \in \langle \text{map } \eta g \rangle$  implies  $xs' = \text{map } \eta g \ xs$ . We also have that  $(s, s') \in \langle g \rangle \rightarrow \text{Eq}_{\text{Bool}}$  implies  $\forall (x, gx) \in \langle g \rangle. \ sx = s'(gx)$  and thus  $s = s' \circ g$  due to the definition of morphisms between relations. With these instantiations, Equation 24 becomes

$$(t(\pi_1 \rho) a \ (s' \circ g) \ xs, t(\pi_2 \rho) b \ s' \ (\text{map } \eta g \ xs)) \in \langle \text{map } \eta (g + 1) \rangle,$$

i.e.,

$$\text{map } \eta (g + 1) (t(\pi_1 \rho) a \ (s' \circ g) \ xs) = t(\pi_2 \rho) b \ s' \ (\text{map } \eta g \ xs),$$

i.e.,

$$\text{map } \eta (g + 1) \circ t(\pi_1 \rho) a \ (s' \circ g) = t(\pi_2 \rho) b \ s' \circ \text{map } \eta g$$

as desired. □

## 6.5 Short Cut Fusion for Lists

**THEOREM 34.** *Let  $\vdash \tau : \mathcal{F}, \vdash \tau' : \mathcal{F}$ , and  $\beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta$ . If*

$$G = \llbracket \beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta \rrbracket^{\text{Set}}$$

*then*

$$\text{fold}_{1+\tau \times \_} n \ c \ (G \ (\text{List } \llbracket \vdash \tau \rrbracket^{\text{Set}}) \ \text{nil} \ \text{cons}) = G \ \llbracket \vdash \tau' \rrbracket^{\text{Set}} n \ c$$

**PROOF.** Let  $\vdash \tau : \mathcal{F}$  and  $\vdash \tau' : \mathcal{F}$ , let

$$\beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta$$

and let

$$G = \llbracket \beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta \rrbracket^{\text{Set}}$$

Then Theorem 29 gives that, for any relation environment  $\rho$  and any  $(a, b) \in \llbracket \beta; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$ , then (eliding the only possible instantiations of  $a$  and  $b$ ) we have

$$(G(\pi_1 \rho), G(\pi_2 \rho)) \in \llbracket \beta; \emptyset \vdash \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta \rrbracket^{\text{Rel}} \rho$$

Since

$$\begin{aligned} & \llbracket \beta; \emptyset \vdash \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta) \beta \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \beta; \emptyset \vdash \text{Nat}^0(\mathbb{1} + \tau \times \beta) \beta \rrbracket^{\text{Rel}} \rho \rightarrow \rho \beta \\ &= (\llbracket \beta; \emptyset \vdash \mathbb{1} + \tau \times \beta \rrbracket^{\text{Rel}} \rho \rightarrow \rho \beta) \rightarrow \rho \beta \\ &= ((\mathbb{1} + \llbracket \vdash \tau \rrbracket^{\text{Rel}} \rho \times \rho \beta) \rightarrow \rho \beta) \rightarrow \rho \beta \\ &\cong (((\llbracket \vdash \tau \rrbracket^{\text{Rel}} \rho \times \rho \beta) \rightarrow \rho \beta) \times \rho \beta) \rightarrow \rho \beta \end{aligned}$$

we have that if  $(c', c) \in \llbracket \vdash \tau \rrbracket^{\text{Rel}} \rho \times \rho \beta \rightarrow \rho \beta$  and  $(n', n) \in \rho \beta$ , then

$$(G(\pi_1 \rho) n' \ c', G(\pi_2 \rho) n \ c) \in \rho \beta$$

Now note that

$$\llbracket \vdash \text{fold}_{1+\tau \times \beta}^{\tau'} : \text{Nat}^0(\text{Nat}^0(\mathbb{1} + \tau \times \tau') \tau') (\text{Nat}^0(\mu \alpha. \mathbb{1} + \tau \times \alpha) \tau') \rrbracket^{\text{Set}} = \text{fold}_{1+\tau \times \_}$$

and observe that if  $c \in \llbracket \vdash \tau \rrbracket^{\text{Set}} \times \llbracket \vdash \tau' \rrbracket^{\text{Set}} \rightarrow \llbracket \vdash \tau' \rrbracket^{\text{Set}}$  and  $n \in \llbracket \vdash \tau' \rrbracket^{\text{Set}}$ , then

$$(n, c) \in \llbracket \vdash \text{Nat}^0(\mathbb{1} + \tau \times \tau') \tau' \rrbracket^{\text{Set}}$$

Consider the instantiation:

$$\begin{aligned} \pi_1 \rho \beta &= \llbracket \vdash \mu \alpha. \mathbb{1} + \tau \times \alpha \rrbracket^{\text{Set}} = \text{List } \llbracket \vdash \tau \rrbracket^{\text{Set}} \\ \pi_2 \rho \beta &= \llbracket \vdash \tau' \rrbracket^{\text{Set}} \\ \rho \beta &= \langle \text{fold}_{1+\tau \times \_} n c \rangle : \text{Rel}(\pi_1 \rho \beta, \pi_2 \rho \beta) \\ c' &= \text{cons} \\ n' &= \text{nil} \end{aligned}$$

Clearly,  $(\text{nil}, n) \in \rho \beta = \langle \text{fold}_{1+\tau \times \_} n c \rangle$  because  $\text{fold}_{1+\tau \times \_} n c \text{ nil} = n$ . Moreover,  $(\text{cons}, c) \in \llbracket \vdash \tau \rrbracket^{\text{Rel}} \times \rho \beta \rightarrow \rho \beta$  since if  $(x, x') \in \llbracket \vdash \tau \rrbracket^{\text{Rel}}$ , i.e.,  $x = x'$ , and if  $(y, y') \in \rho \beta = \langle \text{fold}_{1+\tau \times \_} n c \rangle$ , i.e.,  $y' = \text{fold}_{1+\tau \times \_} n c y$ , then

$$(\text{cons } x y, c x (\text{fold}_{1+\tau \times \_} n c y)) \in \langle \text{fold}_{1+\tau \times \_} n c \rangle$$

i.e.,

$$c x (\text{fold}_{1+\tau \times \_} n c y) = \text{fold}_{1+\tau \times \_} n c (\text{cons } x y)$$

holds by definition of  $\text{fold}_{1+\tau \times \_}$ . We therefore conclude that

$$(G (\text{List } \llbracket \vdash \tau \rrbracket^{\text{Set}}) \text{ nil } \text{cons}, G \llbracket \vdash \tau' \rrbracket^{\text{Set}} n c) \in \langle \text{fold}_{1+\tau \times \_} n c \rangle$$

i.e., that

$$\text{fold}_{1+\tau \times \_} n c (G (\text{List } \llbracket \vdash \tau \rrbracket^{\text{Set}}) \text{ nil } \text{cons}) = G \llbracket \vdash \tau' \rrbracket^{\text{Set}} n c$$

□

## 6.6 Short Cut Fusion for Arbitrary ADTs

**THEOREM 35.** *Let  $\vdash \tau : \mathcal{F}$ , let  $\vdash \tau' : \mathcal{F}$ , let  $\overline{\alpha} : \mathcal{F}$ ,  $\beta \vdash F : \mathcal{F}$ , and let  $\beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau]) \beta$ . If we regard*

$$\begin{aligned} H &= \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}} \\ G &= \llbracket \beta; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau]) \beta \rrbracket^{\text{Set}} \end{aligned}$$

*as functors in  $\beta$ , then for every  $B \in H \llbracket \vdash \tau' \rrbracket^{\text{Set}} \rightarrow \llbracket \vdash \tau' \rrbracket^{\text{Set}}$  we have*

$$\text{fold}_H B (G \mu H \text{ in}_H) = G \llbracket \vdash \tau' \rrbracket^{\text{Set}} B$$

**PROOF.** We first note that the type of  $g$  is well-formed, since  $\emptyset; \beta \vdash F[\overline{\alpha} := \tau] : \mathcal{F}$  so our promotion theorem gives that  $\beta; \emptyset \vdash F[\overline{\alpha} := \tau] : \mathcal{F}$ , and  $\emptyset; \beta \vdash \beta : \mathcal{F}$  so that our promotion theorem gives  $\beta; \emptyset \vdash \beta : \mathcal{F}$ . From these facts we deduce that  $\beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta : \mathcal{T}$ , and thus that  $\beta; \emptyset \vdash \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau]) \beta : \mathcal{T}$ .

Theorem 29 gives that, for any relation environment  $\rho$  and any  $(a, b) \in \llbracket \beta; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$ , eliding the only possible instantiations of  $a$  and  $b$  gives that

$$(G (\pi_1 \rho), G (\pi_2 \rho)) \in \llbracket \beta; \emptyset \vdash \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau]) \beta \rrbracket^{\text{Rel}} \rho$$

Since

$$\begin{aligned} &\llbracket \beta; \emptyset \vdash \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau]) \beta \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Rel}} \rho \rightarrow \rho \beta \end{aligned}$$

we have that if  $(A, B) \in \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Rel}} \rho$  then

$$(G (\pi_1 \rho) A, G (\pi_2 \rho) B) \in \rho \beta$$



Now note that

$$\llbracket \vdash \text{fold}_{F[\overline{\alpha} := \tau]}^{\tau'} : \text{Nat}^0(\text{Nat}^0 F[\overline{\alpha} := \tau][\beta := \tau'] \tau')(\text{Nat}^0(\mu\beta.F[\overline{\alpha} := \tau] \tau'))^{\text{Set}} = \text{fold}_H$$

and consider the instantiation

$$\begin{aligned} A &= in_H : H(\mu H) \rightarrow \mu H \\ B &: H[\llbracket \vdash \tau' \rrbracket^{\text{Set}} \rightarrow \llbracket \vdash \tau' \rrbracket^{\text{Set}} \\ \rho\beta &= \langle \text{fold}_H B \rangle \end{aligned}$$

(Note that all the types here are well-formed.) This gives

$$\begin{aligned} \pi_1 \rho\beta &= \llbracket \vdash \mu\beta.F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}} = \mu H \\ \pi_2 \rho\beta &= \llbracket \vdash \tau' \rrbracket^{\text{Set}} \\ \rho\beta &: \text{Rel}(\pi_1 \rho\beta, \pi_2 \rho\beta) \\ A &: \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Set}}(\pi_1 \rho) \\ B &: \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Set}}(\pi_2 \rho) \end{aligned}$$

since

$$\begin{aligned} A = in_H &: H(\mu H) \rightarrow \mu H \\ &= \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}}(\mu \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}}) \rightarrow \mu \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}} \\ &= \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}}(\pi_1 \rho) \rightarrow \llbracket \emptyset; \beta \vdash \beta \rrbracket^{\text{Set}}(\pi_1 \rho) \\ &= \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}}(\pi_1 \rho) \rightarrow \llbracket \beta; \emptyset \vdash \beta \rrbracket^{\text{Set}}(\pi_1 \rho) \quad \text{Daniel's trick; now a theorem} \\ &= \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Set}}(\pi_1 \rho) \end{aligned}$$

where “Daniel’s trick” is the observation that a functor can be seen as non-functorial when we only care about its action on objects. This is now a theorem. We also have

$$\begin{aligned} (A, B) = (in_H, B) &\in \llbracket \beta; \emptyset \vdash \text{Nat}^0 F[\overline{\alpha} := \tau] \beta \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \beta; \emptyset \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Rel}} \rho[\beta := \langle \text{fold}_H B \rangle] \rightarrow \langle \text{fold}_H B \rangle \\ &= \llbracket \beta; \emptyset \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Rel}} \langle \text{fold}_H B \rangle \rightarrow \langle \text{fold}_H B \rangle \\ &= \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Rel}} \langle \text{fold}_H B \rangle \rightarrow \langle \text{fold}_H B \rangle \quad \text{Daniel's trick; now a theorem} \\ &= \langle \llbracket \emptyset; \beta \vdash F[\overline{\alpha} := \tau] \rrbracket^{\text{Set}}(\text{fold}_H B) \rangle \rightarrow \langle \text{fold}_H B \rangle \quad \text{by the graph lemma} \\ &= \langle \text{map}_H(\text{fold}_H B) \rangle \rightarrow \langle \text{fold}_H B \rangle \end{aligned}$$

since if  $(x, y) \in \langle \text{map}_H(\text{fold}_H B) \rangle$ , i.e., if  $\text{map}_H(\text{fold}_H B)x = y$ , then  $\text{fold}_H B(in_H x) = By = B(\text{map}_H(\text{fold}_H B)x)$  by the definition of  $\text{fold}_H B$  as a (indeed, the unique) morphism from  $in_H$  to  $B$ . Thus,

$$(G(\pi_1 \rho) A, G(\pi_2 \rho) B) \in \langle \text{fold}_H B \rangle$$

i.e.,

$$\text{fold}_H B(G(\pi_1 \rho) in_H) = G(\pi_2 \rho) B$$

Since  $\beta$  is the only free variable in  $G$ , this simplifies to

$$\text{fold}_H B(G \mu H in_H) = G \llbracket \vdash \tau' \rrbracket^{\text{Set}} B$$

□

## 6.7 Short Cut Fusion for Arbitrary Nested Types

fold/build rule shows that Church encodings for nested types are iso (have the same interp) as the nested types themselves. Just like Ex 2.8 in *parpoe*, but categorical rather than operational semantics. But built-in data types are more versatile in implementations. They can be inducted on, e.g. And they are stored in the heap, not the run-time stack, and are therefore more efficient.

“Most higher order type languages meant for human programmers eschew fully impredicative polymorphism” says Andy on p322.

Can take  $\emptyset; \alpha \vdash c$  with  $\llbracket \emptyset; \alpha \vdash c \rrbracket^{\text{Set}} \rho = C$  for all  $\rho$ , i.e., can take  $c$  to denote a constant  $C$ . We then get a free theorem whose conclusion is  $\text{fold}_H B \circ G \mu H \text{ in}_H = G \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} B$ .

Can do Hinze's bit-reversal protocol in our system with

$\text{cat} :: \alpha; \emptyset \vdash \text{Nat}^0(\text{Nat}^0(\text{List } \alpha)(\text{List } \alpha))(\text{List } \alpha)$   
 $\text{zip} :: \alpha; \emptyset \vdash \text{Nat}^0(\text{Nat}^0(\text{List } \alpha)(\text{List } \beta))(\text{List } (\alpha \times \beta))$   
 $?$

**THEOREM 36.** Let  $\emptyset; \phi, \alpha \vdash F : \mathcal{F}$ , let  $\emptyset; \alpha \vdash K : \mathcal{F}$ , and let  $\phi; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^\alpha F(\phi\alpha))(\text{Nat}^\alpha \mathbb{1}(\phi\alpha))$ . If we let  $H : [\text{Set}, \text{Set}] \rightarrow [\text{Set}, \text{Set}]$  be defined by

$$H f x = \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}} [\phi := f][\alpha := x]$$

and let

$$G = \llbracket \phi; \emptyset \mid \emptyset \vdash g : \text{Nat}^0(\text{Nat}^\alpha F(\phi\alpha))(\text{Nat}^\alpha \mathbb{1}(\phi\alpha)) \rrbracket^{\text{Set}}$$

then we have that, for every  $B \in H \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} \rightarrow \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}}$ ,

$$\text{fold}_H B (G \mu H \text{ in}_H) = G \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} B$$

**PROOF.** We first note that the type of  $g$  is well-formed since  $\emptyset; \phi, \alpha \vdash F : \mathcal{F}$  so our promotion theorem gives that  $\phi; \alpha \vdash F : \mathcal{F}$ , and  $\phi; \alpha \vdash \phi\alpha : \mathcal{F}$ , so that  $\phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) : \mathcal{T}$  and  $\phi; \emptyset \vdash \text{Nat}^\alpha \mathbb{1}(\phi\alpha) : \mathcal{T}$ . Then  $\phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) : \mathcal{F}$  and  $\phi; \emptyset \vdash \text{Nat}^\alpha \mathbb{1}(\phi\alpha) : \mathcal{F}$  also hold, and, finally,  $\phi; \emptyset \vdash \text{Nat}^0(\text{Nat}^\alpha F(\phi\alpha))(\text{Nat}^\alpha \mathbb{1}(\phi\alpha)) : \mathcal{T}$

Theorem 29 gives that, for any relation environment  $\rho$  and any  $(a, b) \in \llbracket \phi, \alpha; \emptyset \vdash \emptyset \rrbracket^{\text{Rel}} \rho = 1$ , eliding the only possible instantiations of  $a$  and  $b$  gives that

$$\begin{aligned} (G(\pi_1 \rho), G(\pi_2 \rho)) &\in \llbracket \phi; \emptyset \vdash \text{Nat}^0(\text{Nat}^\alpha F(\phi\alpha))(\text{Nat}^\alpha \mathbb{1}(\phi\alpha)) \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Rel}} \rho \rightarrow \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha \mathbb{1}(\phi\alpha) \rrbracket^{\text{Rel}} \rho \\ &= \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Rel}} \rho \rightarrow (\lambda A. 1 \Rightarrow \lambda A. (\rho\phi)A) \\ &= \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Rel}} \rho \rightarrow (1 \Rightarrow \rho\phi) \\ &= \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Rel}} \rho \rightarrow \rho\phi \end{aligned}$$

So if  $(A, B) \in \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Rel}} \rho$  then

$$(G(\pi_1 \rho) A, G(\pi_2 \rho) B) \in \rho\phi$$

Now note that

$$\llbracket \vdash \text{fold}_F^K : \text{Nat}^0(\text{Nat}^\alpha F[\phi := K] K) (\text{Nat}^\alpha ((\mu\phi. \lambda\alpha. F)\alpha) K) \rrbracket^{\text{Set}} = \text{fold}_H$$

and consider the instantiation

$$\begin{aligned} A &= \text{in}_H : H(\mu H) \Rightarrow \mu H \\ B &: H \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} \Rightarrow \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} \\ \rho\phi &= \langle \text{fold}_H B \rangle \quad \text{a graph of a natural transformation, defined in Enrico's notes} \end{aligned}$$

(Note that all the types here are well-formed.) This gives

$$\begin{aligned} \pi_1 \rho\phi &= \mu H \\ \pi_2 \rho\phi &= \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} \\ \rho\phi &: \text{Rel}(\pi_1 \rho\phi, \pi_2 \rho\phi) \\ A &: \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Set}}(\pi_1 \rho) \\ B &: \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi\alpha) \rrbracket^{\text{Set}}(\pi_2 \rho) \end{aligned}$$

since

$$\begin{aligned}
 A = in_H & : H(\mu H) \Rightarrow \mu H \\
 & = \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}}[\phi := \mu \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}}] \Rightarrow \mu \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}} \\
 & = \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}}(\pi_1 \rho) \Rightarrow \llbracket \emptyset; \phi, \alpha \vdash \phi \alpha \rrbracket^{\text{Set}}(\pi_1 \rho) \\
 & = \llbracket \phi; \alpha \vdash F \rrbracket^{\text{Set}}(\pi_1 \rho) \Rightarrow \llbracket \phi; \alpha \vdash \phi \alpha \rrbracket^{\text{Set}}(\pi_1 \rho) \quad \text{Daniel's trick; now a theorem} \\
 & = \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi \alpha) \rrbracket^{\text{Set}}(\pi_1 \rho)
 \end{aligned}$$

We also have

$$\begin{aligned}
 (A, B) = (in_H, B) & \in \llbracket \phi; \emptyset \vdash \text{Nat}^\alpha F(\phi \alpha) \rrbracket^{\text{Rel}} \rho \\
 & = \lambda A. \llbracket \phi; \alpha \vdash F \rrbracket^{\text{Rel}} \rho[\alpha := A] \Rightarrow \lambda A. (\rho \phi) A \\
 & = \lambda A. \llbracket \phi; \alpha \vdash F \rrbracket^{\text{Rel}}[\phi := \langle fold_H B \rangle][\alpha := A] \Rightarrow \langle fold_H B \rangle \\
 & = \lambda A. \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Rel}}[\phi := \langle fold_H B \rangle][\alpha := A] \Rightarrow \langle fold_H B \rangle \quad \text{Daniel's trick; now a theorem} \\
 & = \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Rel}} \langle fold_H B \rangle \Rightarrow \langle fold_H B \rangle \\
 & = \langle \llbracket \emptyset; \phi, \alpha \vdash F \rrbracket^{\text{Set}}(fold_H B) \rangle \Rightarrow \langle fold_H B \rangle \quad \text{Graph Lemma} \\
 & = \langle map_H(fold_H B) \rangle \Rightarrow \langle fold_H B \rangle
 \end{aligned}$$

since if  $(x, y) \in \langle map_H(fold_H B) \rangle$ , i.e., if  $map_H(fold_H B)x = y$ , then  $fold_H B(in_H x) = By = B(map_H(fold_H B)x)$  by the definition of  $fold_H B$  as a (indeed, the unique) morphism from  $in_H$  to  $B$ . Thus,

$$(G(\pi_1 \rho) A, G(\pi_2 \rho) B) \in \langle fold_H B \rangle$$

i.e.,

$$fold_H B(G(\pi_1 \rho) in_H) = G(\pi_2 \rho) B$$

Since  $\phi$  is the only free variable in  $G$ , this simplifies to

$$fold_H B(G \mu H in_H) = G \llbracket \emptyset; \alpha \vdash K \rrbracket^{\text{Set}} B$$

□

## 7 CONCLUSION AND DIRECTIONS FOR FUTURE WORK

We have forall-types.

Can do everything in abstract locally presentable cartesian closed category.

Give definitions for arb lpccc, but compute free theorems in Set/Rel.

Future Work (in progress): extend calculus to GADTs

Add more polymorphisms (all foralls), even though most free theorems only use one level (or maybe two, like short cut).

Couldn't do this before [Johann and Polonsky 2019] because we didn't know before that nested types (and then some) always have well-defined interpretations in locally finitely presentable categories like Set and Rel. In fact, could extend results here to “locally presentable fibrations”, where these are yet to be defined, but would at least have locally presentable base and total categories with the locally presentable structure preserved by the fibration and appropriate reflection of the total category in the base (as in Alex's effects paper?).

fixed points at term level ala Pitts

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