

# CSI 402 – Lecture 9

## (Linkers and Loaders – Continued)

# External Symbols in C

- Discussion of Handout 9.1.
- Compiler must also handle the `static` attribute for functions and variables.

## SIC/XE – A Representation for EDT and ERT:

- EDT and ERT as part of object code itself.
- **Definition Record** (D-record) for EDT: Each D-record has a symbol and its relative address.
- **Reference Record** (R-record) for ERT:
  - R-record contains all the external symbols referenced by the module.
  - Addresses will appear in modifier records.
  - R-records are strictly unnecessary; they are useful in making the linking process more efficient.

# Changes to SIC/XE Modifier Records

- Current format: Each modifier record contains starting byte address and number of bytes.
- Adequate for relocation but not for linking.
- New format: Each modifier record contains starting byte address, number of bytes, a flag and a local or external symbol.
- Flag can be '+' or '-'.
- New format can handle both relocation and linking.

**Example:** To be presented in class. (The example is based on Handout 8.3.)

# Local and External Symbols

**Example:** Numbers shown below are LC values.

	FUNCT	CSECT	
		EXTDEF	BUFFER
		EXTREF	SYMBTAB
0		LDA	#15
3		+ADD	BUFFER
7		+STA	SYMBTAB
		.	
		.	
	BUFFER	RESW	10
		END	

**M-Records for the above module:**

M	4	3	+	FUNCT
M	8	3	+	SYMBTAB

# Local and External Symbols (continued)

**Example (continued):** LC value is shown for the relevant statement.

	READ	CSECT	
		EXTREF	BUFFER
		.	
		.	
75		+STA	BUFFER
		.	
		.	
		END	

**An M-Record for the above module:**

M    76    3    +    BUFFER

# Local and External Symbols (continued)

## Notes:

- When the symbol BUFFER is referenced in the module where it is defined, the corresponding M-record uses the name of the module.
- When the symbol BUFFER is referenced in a module where it is an external symbol, the corresponding M-record uses the symbol BUFFER itself.

## Need for the '-' Flag:

- Operand field may have an expression containing external symbols.
- These expressions are generally simple. (They contain only '+' and '-' operators).

# Need for the '-' Flag (continued)

## Example:

FUNCT	CSECT	
	EXTREF	FIRST, LAST
	.	
	.	
SIZE	WORD	LAST-FIRST+1

- Suppose the LC value of SIZE is 70.
- The assembler initializes bytes 70, 71 and 72 to 0, 0 and 1 respectively.
- It writes the following M-records to the object file.

M	70	3	+	LAST
M	70	3	-	FIRST

# Algorithm for a Linking Loader

- **Input:** Object code for each module organized as Header record, D-record, R-record, Text records, M-records and End record.
- **Assumptions:**
  - Only one module specifies a starting address in its End record.
  - Modules loaded successively with no gaps.
- Algorithm uses two passes.
- This version of algorithm ignores R-records.
- **Pass I:** Builds the External Symbol Table by combining the EDTs of the modules.
- **Pass II:** Resolves the external symbols, performs relocation and loading and starts execution.
- **Algorithm Outline:** Handout 9.2.
- **Reading assignment:** Figures 3.11(a) and 3.11(b) of [Beck].



# Usefulness of R-records

- With the revised format for modifier records, R-records are strictly unnecessary.
- R-records can be used to improve the efficiency of the linking process.

## Example:

- Suppose the External Symbol Table (EST) has 100 symbols.
- **Fact:** If we have a sorted EST with  $n$  entries, we can use binary search to look for a symbol; in the worst case, this uses  $\lceil \log_2(n+1) \rceil$  probes. (Each probe involves a call to a function such as `strcmp`.)
- For the above example, the number of probes needed for each search is  $\lceil \log_2 100 \rceil = 7$ .

# Usefulness of R-records (continued)

## Example (continued):

- Suppose a particular module references only 3 of the symbols, say A, B and C, of the EST. (We would know this from the R-record for the module.)
- Let there be 5 M-records each for the symbols A, B and C. (Total: 15 M-records.)

## Method I:

- Ignore the R-record and directly search the (EST) for each of the 15 M-records.
- Total number of probes in Method I =  $15 \times 7 = 105$ .

# Usefulness of R-records (continued)

## Method II:

- The R-record of the module is:

R      A, B, C

- Search the EST and find the addresses of A, B and C. The number of probes used in this step =  $3 \times 7 = 21$ .
- Create a smaller EST containing just these three symbols and their addresses.
- For each of the 15 M-records, search the smaller EST. Each search uses at most 3 probes (even when sequential search is used). So, the number of probes used in this step =  $15 \times 3 = 45$ .
- Total number of probes used in Method II =  $21 + 45 = 66$ .

**Conclusion:** Method II uses about 40% fewer probes than Method I.

# Linking C programs with Libraries

- Functions from `<stdlib.h>` (e.g. `malloc`, `free`, `exit`) are linked automatically.
- Functions from `<stdio.h>` (e.g. `fprintf`, `fscanf`) are linked if `<stdio.h>` is included as a header.
- If functions from other libraries (e.g. `<math.h>`) are used, then the corresponding header must be included and the library must be specified as part of the `gcc` command.

**Example:** Suppose a program (say `prog.c`) uses the function `sqrt` from the `math` library.

- The program must include `<math.h>`.
- The Command for producing load module is:

```
% gcc    prog.c    -lm
```

# Linking C programs with Libraries (continued)

## Notes:

- The linker searches the libraries only after obtaining information about all the functions defined in the various files.
- This allows a user to redefine a library function. (If this is really necessary, it must be done with great care.)

## Dynamic Linking:

- **Static Linking:** Linking done before execution.
- **Dynamic Linking:** Linking done during execution.
- Dynamic linking is useful when a program calls only a few of a large collection of routines from a library (e.g. a library containing error handling routines).
- Advantage: The size of the load module is smaller.
- Disadvantage: Additional runtime overhead.

# Dynamic Linking (continued)

## Steps used in Dynamic Linking:

- 1 A user program calls a routine (say `sqrt`) which is not part of the load module.
- 2 The runtime system makes a service request to the OS to load the requested function. The request is handled by a part of the OS called the “Dynamic Loader” (DL).
- 3 DL checks if the requested routine is already in memory; if not, loads the routine at an appropriate part of memory. DL also starts the execution of the routine.
- 4 When the routine completes, it returns control to DL. DL decides whether or not to leave the routine in memory or reclaim the memory.
- 5 DL passes the control back to the user program.

## Additional Remarks:

- **Binding:** Association of an attribute value with a name.
- Binding can happen at compile time, at load time or at execution time.
- With static linking, the binding between a function name and its address happens at load time (before execution begins). For a given execution, this binding does not change with time (static binding).
- With dynamic linking, the binding between a function name and its address may happen at execution time; this binding may also change with time (dynamic binding).

**Issue:** How to load the very first program into memory.

- A small part of memory is implemented as Read Only Memory (ROM) which contains a program. (This program is an absolute loader.)
- When power is turned on (or when the system is rebooted), the ROM program begins to execute. The program reads a block of bytes from a specific device (e.g. hard disk).
- The block (called the **boot block**) read in by the ROM program contains another program which can load a larger program.
- Thus, the program in the boot block loads a larger program, ..., and so on, until the OS itself is loaded.
- This sequence of loading larger and larger programs until the OS is loaded is called **bootstrapping**.
- Systems also perform “Power on self-tests” during booting.



# Bootstrap Loader (continued)

## Remarks:

- Some people refer to the ROM program as the bootstrap loader.
- Others refer to the program in the boot block as the bootstrap loader.

Reading Assignment: Figure 3.3 of [Beck].

## Notes on the program in Figure 3.3 of [Beck]:

- It is a bootstrap loader (i.e., a simple absolute loader) for SIC/XE (that can be stored in ROM).
- It reads bytes from device F1 (hex) and loads the bytes from address 80 (hex).
- The device returns 04 (hex) when EOF occurs.
- At that time, the program causes a jump to location 80 (hex) to start the program that was just loaded.