

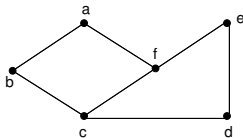
# CSI 445/660 – Part 1

## (Graph Theory Basics)

**Ref:** Chapter 2 of [Easley & Kleinberg].

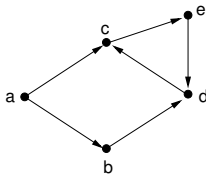
# Types of Graphs

## ■ Undirected and Directed.



### Undirected graph:

- **Example:** Friendship relation among people.
- A **symmetric** relationship.

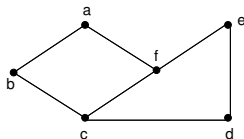


### Directed graph:

- **Example:** Follower relationship in Twitter.
- May not be symmetric.

# Undirected Graphs: Notation and Definitions

## Example:



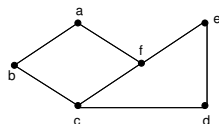
**Notation:**  $G(V, E)$

$$V = \{a, b, c, d, e, f\} \quad (\text{nodes or vertices})$$

$$E = \{ \{a,b\}, \{a,f\}, \{b,c\}, \{c,d\}, \{c,f\}, \{d,e\}, \{e,f\} \} \\ (\text{edges})$$

$$|V| = \text{No. of nodes} = 6 \qquad |E| = \text{No. of edges} = 7$$

# Notation and Definitions (continued)



**Definition:** The **degree** of a node  $v$  is the number of edges **incident on**  $v$ .

**Example:** Degree of  $a = 2$ , degree of  $f = 3$ .

## Some observations:

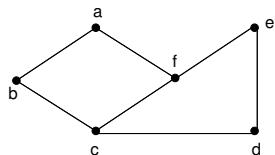
- Sum of the degrees of all the nodes
$$\begin{aligned} &= \text{Degree}(a) + \text{Degree}(b) + \dots + \text{Degree}(f) \\ &= 2 + 2 + 3 + 2 + 2 + 3 = 14 \text{ (even)} \\ &= 2 \times \text{No. of edges.} \end{aligned}$$
- Nodes with **odd** degree =  $\{c, f\}$ ; thus, the number of nodes of odd degree is **even**.

# Notation and Definitions (continued)

## Theorem: [First Theorem of Graph Theory]

In any undirected graph, **the sum of the degrees of all the nodes is equal to twice the number of edges.**

**Corollary:** In any undirected graph, the **number of nodes of odd degree is even.**



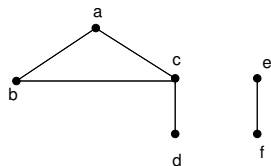
Examples of paths in graph  $G$ :

■  $a - f - e - d$

■  $a - b - c - f - e - d$

- There is a path between every pair of nodes.
- Graph  $G$  is **connected**.

# Notation and Definitions (continued)

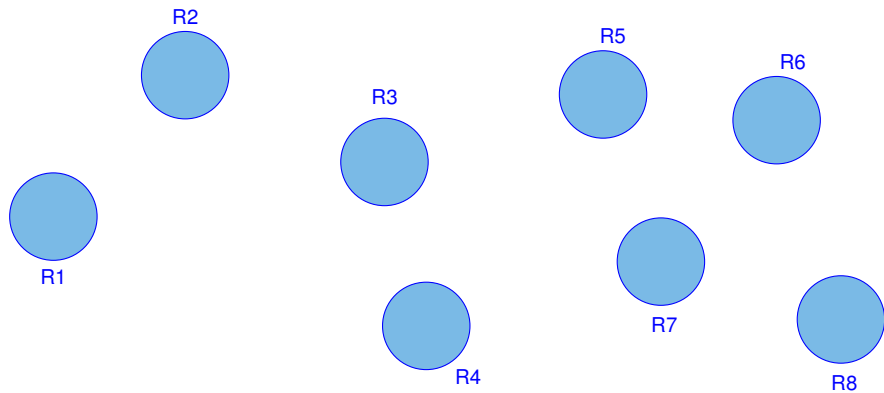


- **Disconnected** graph.
- Has two **connected components**.

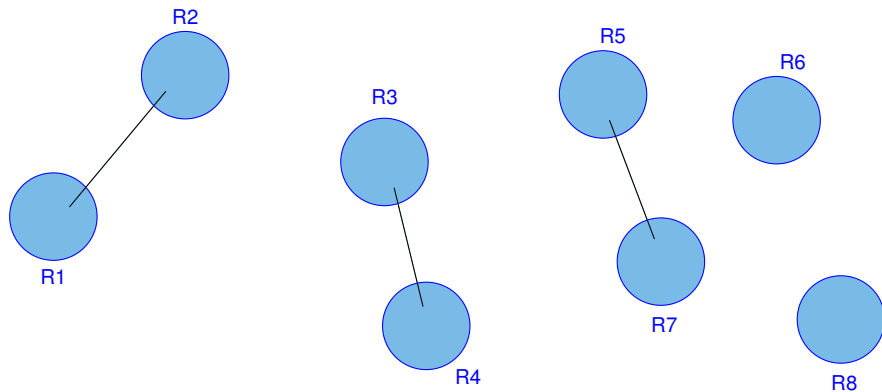
**Evolution of a large social network:** Imagine the following global friendship graph.

- One node per person in the world  
(No. of nodes  $\approx 7.3$  billion).
- An edge between each pair of friends.

# Friendship Network Evolution

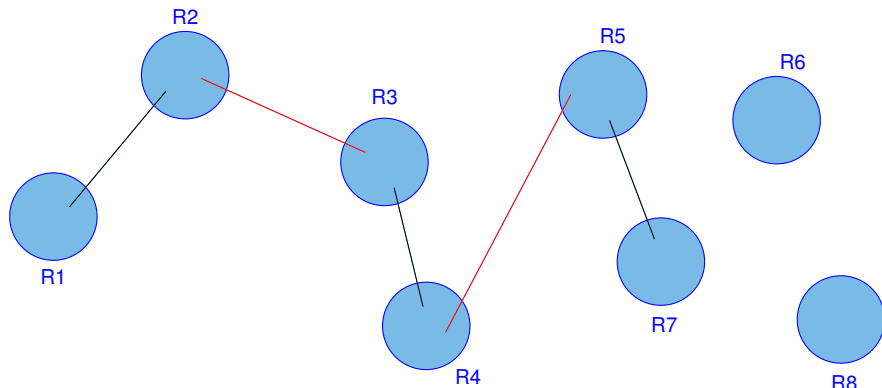


## Friendship Network Evolution (continued)





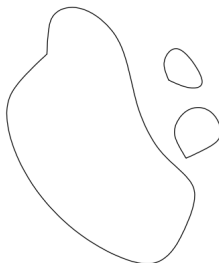
## Friendship Network Evolution (continued)



# Friendship Network Evolution (continued)

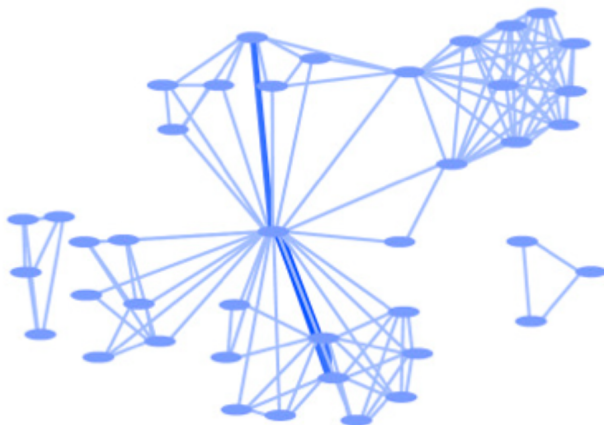
- Components get **merged** over time.
- The graph is likely to contain paths between people in remote parts of the world.
- A large subset of the nodes are in one component, called the **giant component**. (This is typical of many social networks arising in practice.)

**An Illustration by Prof. Alistair Sinclair (UC Berkeley):**

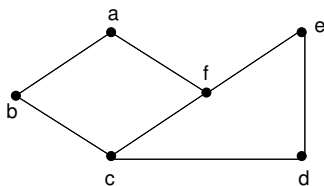


# Giant Component: Another Example

**Collaboration graph at a research center (from [EK]):**



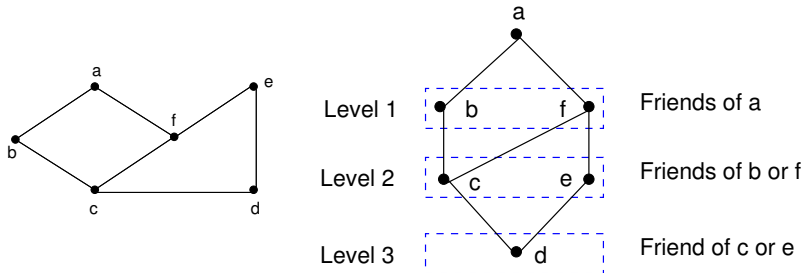
# Shortest Paths



**Paths between a and e:**

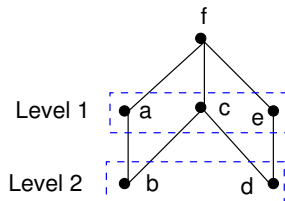
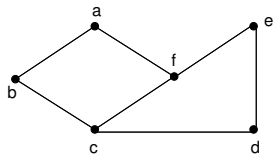
- **a – f – e** : Length = 2 (No. of edges)
- **a – b – c – f – e** : Length = 4
- There is no path between a and e with length  $< 2$ .
- So, **a – f – e** is a **shortest path** between a and e.
- Shortest paths can be found using a procedure called **breadth-first-search** (BFS).

# Breadth-First-Search: Example I



**Observation:** Each node is within a distance of 3 from node a.

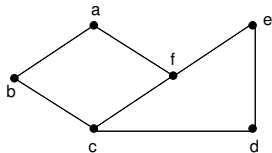
# Breadth-First-Search: Example II



**Observation:** Each node is within a distance of 2 from node f.

# Definition of Diameter

## Shortest Path Lengths: (Partial list)



Node pair	Shortest Distance
a, b	1
a, c	2
a, d	3
⋮	⋮
b, e	3
⋮	⋮
e, f	1

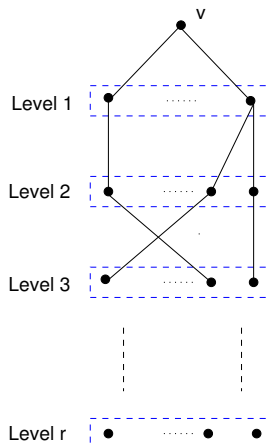
- **Diameter:** Maximum among the shortest path lengths.
- Diameter of the above graph = 3.

# Some Notes About Diameter

- Diameter is meaningful only for **connected** graphs. (Some references use  $\infty$  as the diameter of a disconnected graph.)
- If a graph is disconnected, one needs to consider the diameter each connected component.
- For a connected graph with  $n$  nodes, the diameter is at most  $n - 1$ .
- In communication networks, diameter gives an indication of the worst-case delay for message delivery.
- Typically, giant components of social networks have small diameters (**small world phenomenon**).



# BFS and Diameter



**Observation:** For any **connected** graph, if a BFS produces  $r$  levels, then the diameter of the graph is at most  $2r$ .

## Example: Erdős Collaboration Network



- Paul Erdős (1913 – 1996)
  - Hungarian Mathematician
- 
- Each node is a researcher and edge  $\{x, y\}$  means that researchers  $x$  and  $y$  co-authored at least one paper.
  - **Level 0:** Node corresponding to Erdős.
  - **Level 1:** Nodes corresponding to researchers who have published a paper with Erdős.

# Erdős Collaboration Network (continued)

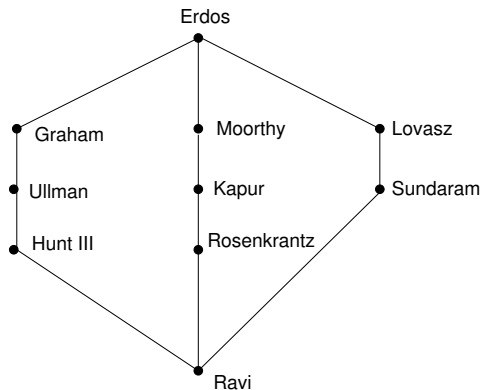
- **Level 2:** Nodes corresponding to researchers who have published a paper with some researcher in Level 1.

$\vdots$                        $\vdots$

- **Level  $j$ :** Nodes corresponding to researchers who have published a paper with some researcher in Level  $j - 1$ .
- **Erdős Number** of a researcher: The level number in the graph for the node corresponding to the researcher.

Largest known Erdős Number = 8.

# Example for Erdős Number



- Ravi's Erdős Number  $\leq 3$ .
- Erdős Numbers of Teri Harrison, Catherine Dumas and Dan Lamanna  $\leq 4$ .