

Assignment 1 - Daniel Lovegrove - 7763168

1. After running **ps xf** remotely on the lab machine, I see that the shell I'm using is TCSH, with process id 11080. Following back to pid 1:

PID	NAME
11080	tcsh
11079	sshd
11076	sshd
1053	sshd
1	systemd

It appears that at the base, systemd, the first process, has an SSH daemon (pid 1053) listening for connections. From the man page for SSHD, the SSH daemon is forked for each incoming connection, creating processes at 11076 and 11079. Then, a TCSH shell is opened with pid 11080 for communicating with the user.

Systemd has a parent process id 0, but there is no actual process with pid 0. This is because systemd is forked from the swapper, which is a direct part of the kernel and so is not a user process, hence why there is no `proc/0`.

2. Making an `exec` call within a process overwrites the entire address space of the process and runs the code that is to be executed. It does not make sense to make an `exec` call in a multi-threaded application because when one thread makes an `exec` call the whole program with all of its threads is stopped and overwritten to run the new program. If two threads make a call to `exec` there will be a race condition, since the same address space will be trying to be overwritten for two different programs, the outcome of which will be undefined and likely crash the code.

3. From the man pages for `strtok`, `strtok` uses a static buffer while parsing. So, if there are multiple threads running, each using `strtok`, they will all be using the same buffer and a race condition is created since we don't know which thread is accessing which part of the single static buffer at any time. `Strtok_r` overcomes this by keeping separate pointers that may point to the same string. Since `strtok_r` just uses pointers to a string, there is no issue of threads interfering with each other in regards to which part of the string they are processing.

4, 5, 6, 7 are coding questions. See the README.md document to see how to run the code.