

JOHNSON I

Thm $I(S_g)$ is fin. gen. by Dehn twists for $g \geq 3$.

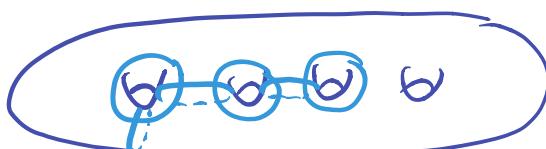
Basic Strategy

1. List prospective generators $\{g_i\}$.
s.t. g_i is a BP map of genus 1.
2. Show $\langle g_i \rangle \trianglelefteq \text{Mod}(S_g)$.

This suffices since $\langle\langle g_i \rangle\rangle_{\text{Mod}(S_g)} = I(S_g)$.

Chains and BP maps.

A chain:



→ BP map.

Given a chain, can resolve intersections to get another chain. Can also take subchains.

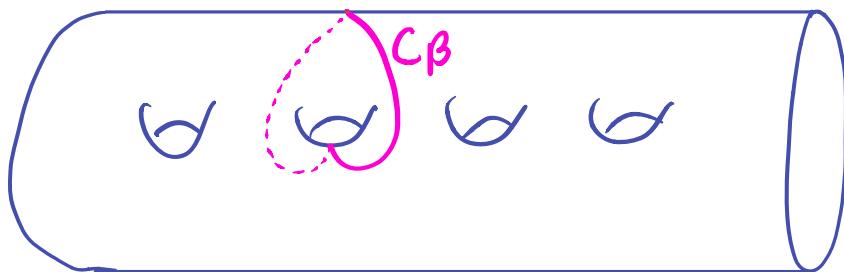
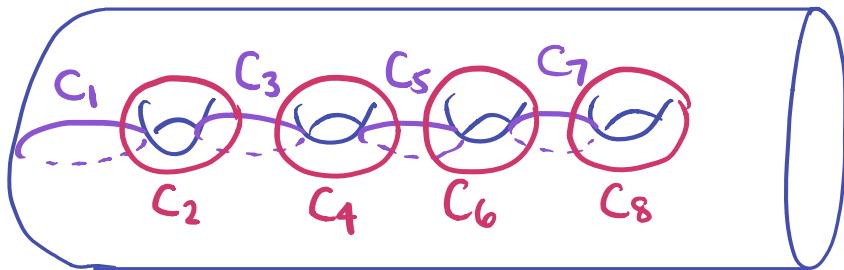
Given a chain $ch(c_1, \dots, c_n)$

$ch(i_1 i_2 \dots i_{k+1})$

denotes the chain you get by combining c_1, \dots, c_{i_2-1}
 $c_{i_2}, \dots, c_{i_3-1}$ etc. dropping c_{k+1}, \dots, c_n . "subchain"

Denote the BP-map $[i_1 i_2 \dots i_{k+1}]$

LISTING THE GENERATORS



Consider the chains:

(c_1, \dots, c_{2g}) straight chain

$(c_\beta, c_5, \dots, c_{2g})$ β -chain

Use same notation for subchains of β -chain:

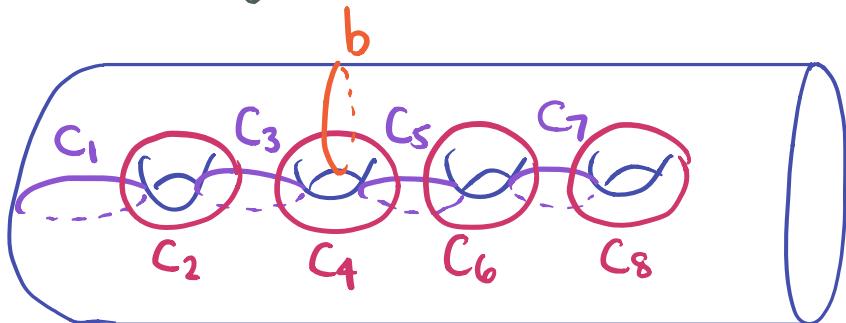
$(\beta_{i,}) =$ surger $c_\beta, c_5, \dots, c_{i-1}$

Theorem. For $g \geq 3$ the odd subchain maps of straight chain & β -chain generate $I(S'_g)$

Since $I(S'_g) \rightarrow I(S_g)$ this gives closed case as well.

SETUP.

Humphries generators :



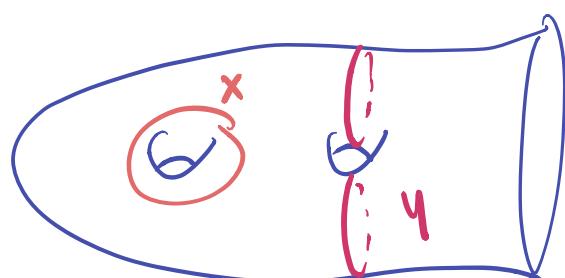
Let $\mathbb{J}(S_g')$ & $\mathbb{J}(S_g)$ denote groups gen by Johnson's generators.

As above, need to show

$$T_x * y = T_x y \overline{T_x}^{-1} \in \mathbb{J}(S_g)$$

$\forall x \in$ Humphries set
 $y \in$ Johnson set

In many cases $T_x * y$ equals y or is another Johnson gen:



CONJUGATING BY A POSITIVE TWIST IN THE CHAIN

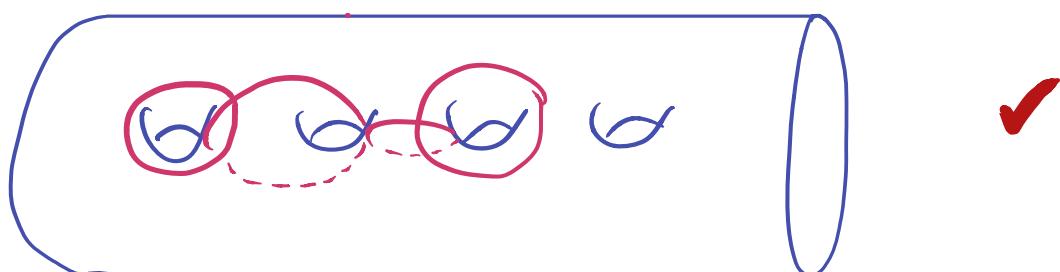
Prop.

| $j \text{ in } \{i_1, \dots\}$ | $j+1 \text{ in } \{i_1, \dots\}$ | $T_{Cj} * [i_1 i_2 \dots]$ |
|--------------------------------|----------------------------------|-------------------------------------|
| ✓ | ✓ | $[i_1 i_2 \dots]$ |
| ✗ | ✗ | $[i_1 i_2 \dots]$ |
| i_m | ✗ | $[i_1 \dots i_{m-1} i_{m+1} \dots]$ |
| ✗ | i_m | $[i_1 \dots i_{m-1} i_m \dots]$ |

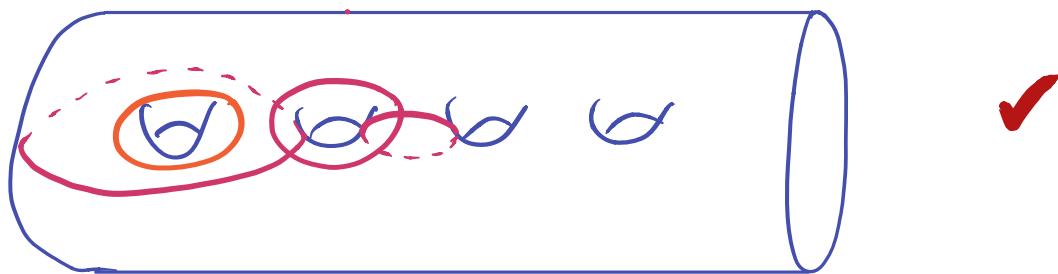
The lemma completely characterizes commuting among straight Johnson & Humphries gens.

PF (by examples)

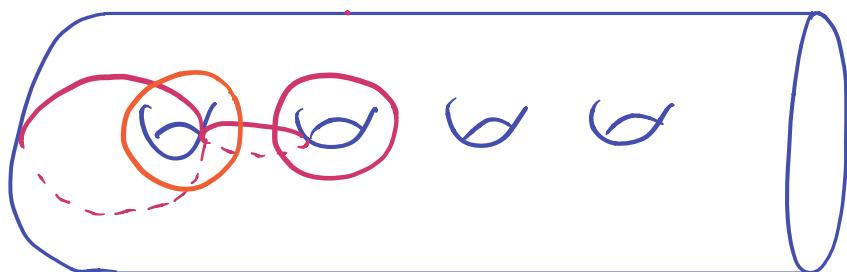
Example $j=2$ $[2 3 5 6]$



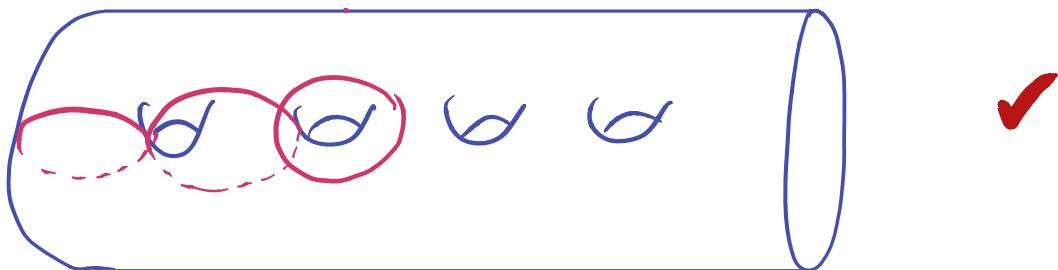
Example $j=2$ $[1 \ 4 \ 5 \ 6]$



Example $j=2$ $[1 \ 3 \ 4 \ 5]$



$[1 \ 2 \ 4 \ 5]$



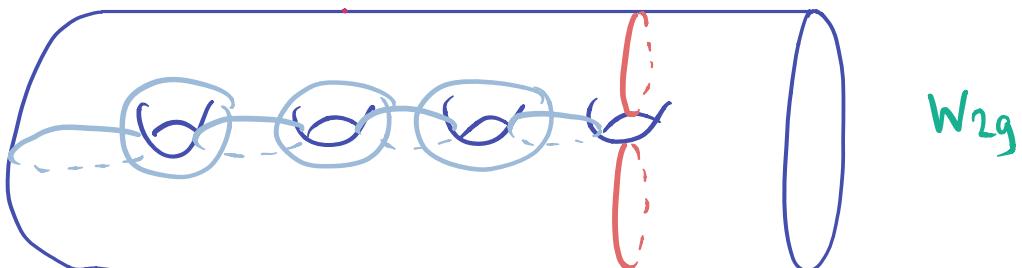
Third case similar.

Upshot of Lemma: If we conj a chain map by a positive twist about a curve in the chain we get a subchain map. What about negative twists?

GENERATING THE KERNEL OF $I(S'_g) \rightarrow I(S_g)$

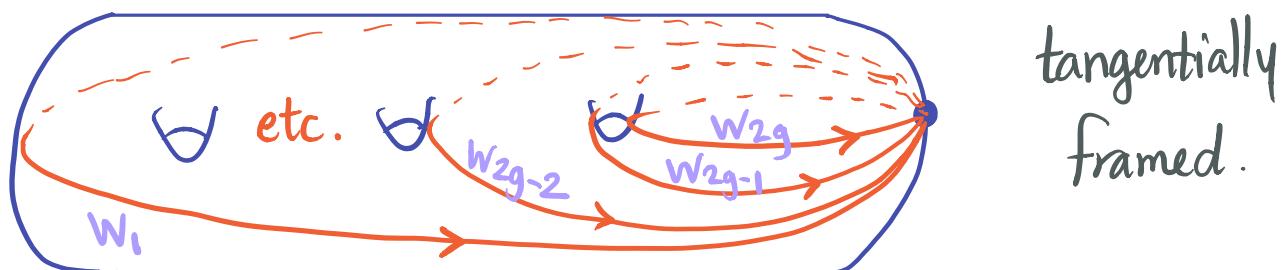
Define $W_i = [1 \ 2 \ 3 \ \dots \ \hat{i} \ \dots \ 2g+1]$
 for $i \in \{1, \dots, 2g+1\}$. "maximal odd
 subchain maps"

Each lies in $\pi_1 \text{UT}(S_g)$



Lemma. $\pi_1 \text{UT}(S_g)$ is gen by $W_1, \dots, W_{2g}, T_b * W_1$

Pf. The corresponding push maps are



These clearly generate $\tilde{\pi}_1$. Need to get
 the fiber of $\tilde{\pi}_1 \text{UT} \rightarrow \tilde{\pi}_1$.

Claim. $T_b * W_1 = W_4 W_3^{-1} W_2 T_b$ □

Cor. $\pi_1 \text{UT}(S'_g) \leq \mathcal{I}(S'_g)$

Pf. Use lantern to get $T_b * W_1$ (Cor 1 in paper).

CONJUGATING BY A NEGATIVE TWIST IN THE CHAIN

Prop. Fix a chain (c_i) . The subgroup of $I(S_g')$
 $\langle \text{odd subchain maps of length } 2k-1 \rangle$
is normalized by $\langle T_{c_i} \rangle$.

Lemma. $\langle w_1, \dots, w_{2g+1} \rangle$ is normalized by
 $\langle T_{c_1}, \dots, T_{c_{2g}} \rangle$

Pf. Need to check $T_{c_i}^{\pm 1}(w_j) \in \langle w_1, \dots, w_{2g+1} \rangle$
in $\Pi_1 \cup T(S_g')$.

Pf of Prop. Let $f = [i_1 \ i_2 \ \dots \ i_{2k}]$
Fix a c_j in original chain.
As above $T_{c_j} * f = f$ unless
exactly one of $j, j+1$ is an i_\square .
Enlarge $(i_1, i_2, \dots, i_{2k})$ to include both j & $j+1$.
 \rightsquigarrow subchain of length $2k$ (c'_i) w/ c_j basic

Lemma $\Rightarrow T_{c_j}^{\pm 1} * f$ is a product of maximal
subchain maps of (c'_i) . But these
are all subchains of original chain \square

Two More Tools.

① INDUCTION

If our Johnson & Humphries gens both lie in S_{g-1}' , can apply induction. Works because both sets of gens agree with $S_{g-1}' \hookrightarrow S_g'$ (this is one specific inclusion).

N.B. Base case is subgp of $I(S_2')$ gen.
by BP maps i.e. $\pi_1, UT(S_2)$
i.e. $\langle w_1, \dots, w_5 \rangle$.

② CONTROLLED CHANGE OF COORDS

Suppose $T_{ci} \in$ Humphries

$f \in \text{Mod}(S_g')$ already known to
normalize $J(S_g')$.

$h \in$ Johnson

Then checking $T_{ci} * h \in J(S_g')$ is equiv. to
checking $f * (T_{ci} * h)$
 $= T_{ci} * (f * h) \in J(S_g')$

Proof of THEOREM

Need to conjugate all Johnson gens by
 $T_{c_i}^{\pm 1}, T_b^{\pm 1} \in$ Humphries.

Step 1. $T_{c_1}, T_{c_2}, T_{c_5}, \dots, T_{c_{2g}}$ normalize $J(S_g')$.

T_{c_1}, T_{c_2} in the straight chain, disjoint from β -chain.

$T_{c_5}, \dots, T_{c_{2g}}$ in both chains.

Apply Prop.

Step 2. T_{c_3} normalizes $J(S_g')$.

T_{c_3} normalizes straight chains by Prop.

By controlled change of coords, can restrict to consecutive β -chain maps.

$[\beta \ 5 \ 6 \ \dots]$.

If chain not maximal, use induction.

If it is, need a relation (Cor 2 in paper).

But this is one check.

Step 3. T_{c_4} normalizes $J(S_g')$ similar to T_{c_3}

Step 4. T_b normalizes $S(S_g')$.

β -chains: controlled change of coords
reduces consecutive β -chains &
induction reduces to

$$T_b^{\pm 1} * [\beta \ 5 \ \dots \ 2g+1] \quad \text{single check}$$

Straight chains $[i_1 \ \dots \ i_{2k}]$

Can assume $i_1 \leq 4$ $i_{2k} \geq 5$
(otherwise T_b commutes)

Change of coords reduces to: consecutive
straight chain starting with 1, 2, 3, or 4.
and going up to $2g$ or $2g+1$ (induction).

If it starts with 1 or 2 it is maximal,
i.e. in $\Pi_1 \cup \Pi_2$.

Only two remaining cases:

$$[4 \ 5 \ \dots \ 2g+1] \ \& \ [3 \ 4 \ \dots \ 2g].$$

