COMPLEX OF CURVES- OVERVIEW

Main object of study: MCG(Sg) = To Homeo+(Sg) "mapping class"

= Homeo+(Sg)/homotopy class "group"

Motivation: ① MCG(Sg) = Out Th(Sg) Dehn-Nielsen-Baer than

MCG(Sg) is analog of GLn72 = Out 72

② MCG(Sg) = Thorb (Mg) Mg = moduli space of hyp. surfs

③ MCG(Sg) classifies Sg-bundles

Sg-bundles over B ←> ThB→ MCG(Sg)

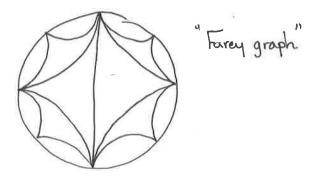
(already interesting for B=S¹).

Main tool: Complex of curves

C(Sg) vertices: homotopy classes of simple closed curves in Sg edges: disjoint representatives.

We'll see C(Sq) is ① connected
② 00-diam
③ hyperbolic
but... ④ locally infinite.

For g=1 we modify the definition: disjoint ~ minimal



HYPERBOLICITY

A metric space is hyperbolic if for any geodesic Δ , at the δ -nod of any two sides contains the third.



Facts. (1) En is not o-hyp

- @ IH" is In(1+12)-hyp
- 3 Trees are O-hyp.

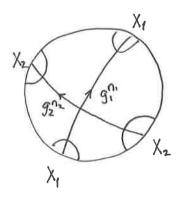
Will show C(Sg) is 17-hup (indep. of g!)

can import ideas from hyp manifolds to MCG,
for instance:

Prop. M = closed hyp n-man, $g_1, g_2 \in \pi$. MThen $\exists n_1, n_2 \text{ s.t. } g_1^{n_1}, g_2^{n_2}$ either commute or generate F_2 .

Ping Pong Lemma. X = set, $G \circlearrowleft X$, $g_1, g_2 \in G$ $X_1, X_2 \neq \emptyset$, $X_1 \cap X_2 = \emptyset$ $g_1^k(X_2) \subseteq X_1$, $g_2^k(X_1) \subseteq X_2$ $\forall k \neq 0$. Then $\langle g_1, g_2 \rangle \cong F_2$

If $\omega = \text{freely red word in } g_1, g_2$ $\text{Say } \omega = g_1^7 g_2^5 g_1^3 g_2 g_1$ Let $\chi \in g_2$. Note $\omega(\chi) \in \chi$, $\Rightarrow \omega(\chi) \neq \chi \Rightarrow \omega \neq id$. Pf of Prop. Apply PPL to:



W

This entire approach will generalize to MCG(Sg) Co C(Sg).

CURVES IN SURFACES

Q. How can we tell if two vertices of C(Sg) have disjoint reps?

simple closed curve

図

Prop (Bigon Criterion) Two transverse scc in Sq are in minimal position iff they do not form a bigon:



(minimal posn means smallest intersection number in homotopy classes).

Note: >> is easy: > ~

Lemma. If two scc do not form a bigon then a pair of lifts to IH2 can intersect in at most one pt.

Pf. If not, an (innermost) bigon in 1H2 projects to a bigon in 5g

Pf of Bigon Criterion (Sketch).

Assume $\alpha, \beta \subseteq Sg$ form no bigons Lemma \Rightarrow lifts can only intersect in 1 pt. Can argue these lifts must have distinct endpts

 \tilde{z} \tilde{z}

But isotopies So do not move pts at co So no isotopy can reduce intersection.

Geodesics

Prop. Every sec in Sy (9>2) is homotopic to a unique geodesic Prop. Geodesics in Sg are in minimal pos.

Change of Coordinates Principle

Configurations of curves can often be put into a Standard picture via homeo of Sg.

examples @ If x \subsection Sq is a nonsep sccin Sq, I he Homeo (Sq)

s.t h(x) = x.

② If $\alpha_1 \beta \leq Sg$ have $c(\alpha_1 \beta) = 1$ (geometric int num) then $\exists h \in Homeo(Sg)$ s.t. $h(\alpha_1 \beta) = (\alpha_0, \beta_0)$



Proofs use classification of surfaces.

CONNECTIVITY

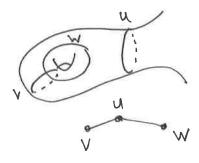
Thm C(Sg) is connected, 9 72.

Pf. Induction on i(v, w).

For i(v, w) = 0, nothing to do.

For i (v, w)=1, use change of coords:

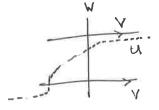
Now assume i(v,w) > 2. Orient the curves v,w. and assume minimal pos.



Look at two consecutive intersections along w.

Orientations can agree or disagree.

If they agree:



Note u is essential since i(u,v)=1. By induction u connected to v and w.

If they don't agree:

u is essential because otherwise & V, W not in min pos: By induction u conn. to V, W.

HYPERBOLLOTY

Thm (Masur-Minsky). C(Sg) is 5-hyp.

We'll show of can be taken indep of g (Hensel-Przytycki-Webb and others)
Proof from Sisto's blog.

Guessing geodesics lemma (Masur-Schleimer) X = metric graph. $X = \text{met$

O d(x,y) $\leq 1 \Rightarrow$ diam A(x,y) $\leq D$.

2 A(x,y) = ND (A(x,z) U A(x,y)) Y x,y,z.

Note. \Rightarrow easy: A(x,y) is any geodesic. $D=\max(\delta,1)$.

We will replace C(Sg) with C'(Sg). The latter has extra edges, namely, add edges between vertices a,b with i(a,b)=1.

To check: ① C'(Sg) is quasi-isometric to C(Sg).

(and constants do not depend on g)
② If X is δ -hyp, Y q i to X then Y is δ' -hyp

(δ' depends only on δ & gi constants).

Note: We need the guessing geodesics lemma precisely because we don't know how to find geodesics. And so it is hard to check f-hyp'ity directly.

Thmo. C'(Sg) is &-hyp.

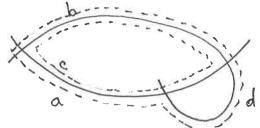
Pf. First: A(a,b) = {vertices of C'(Sg) formed from one arc of a, one arc of arcs should have distinct endpts

Claim. A(a,b) connected

Pf. Define a partial order c < d if b-arc of d contains the b-arc of c (so d is closer to being b)

Want for all $c \in A(a_1b)$ a $d \in A(a_1b)$ s.t. d > c and $c \in d$ life d find d, prolong one side of the b-arc of c and d if d again, shorten the d-arc of d :

this isn't quite a partial order as stated since two curves can have same b-arc but opposite a-arcs



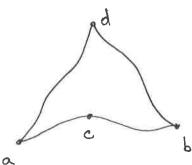
By defin, d > c. To see $i(c,d) \le 1$ note the worst that can happen is the prolonged arc ends up on the wrong side of c.

Notice the A(a,b) satisfy 1 since A(a,b) = {a,b} when a b

Claim. The A(a,b) form thin triangles as in 3

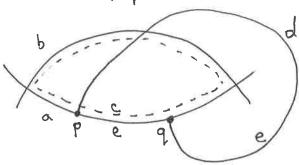
Pf. Fix a,b and (C = A(a,b) and d.

Need e = A(a,d) U A(d,b) close to C.



To find e: consider 3 consec. intersections of d with c (if fewer than 3, d is already close to c, so e=d). Say 2 of these intersections are on the a-arc.

call them pig:



Form e from the arc of d shown and the arc of cea shown.

Note $i(c,e) \le 2 \implies d(c,e) \le 2$.

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GUESSING GEODESICS

see Bowditch "Uniform hup" Prop 3.1 for a proof of the Stronger one.

We'll prove something a little weaker than the lemma used above.

JD s.t.

Lemma. (Hamenstadt) X = metric space. Suppose Y XM (X there is a path (X, Y) connecting them and so:

① diam p(x,y) ≤ D if d(x,y) ≤1

A x,y and x',y ∈ p(x,y), dHaus (p(x',y'), Subpeth of p(x,y) from x' to y') ≤ D

3 p(x,y) = No (p(x,z) · U p(z,y)) \ X,y,Z.

Then X is S-hyp.

So to prove the theorem, need to either prove the stronger lemma (i.e. eliminate @ above) or check @ for C'(Sg).

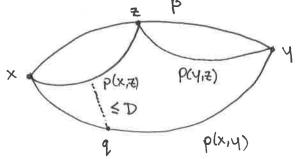
Idea: show the p(x,y) are (close to) geodesics

Pf. Two steps.

> Step 1. If B is any path $X \longrightarrow Y$ then $p(x,y) \subseteq N_R(B)$ where R~ log (length B).

> > recall: in IH" if a path leaves the R nood of a geodesic its length is ~ eR

let $q \in p(x, y)$ and To prove this, split B in half, draw the p paths. Note q is close to one; using condition 3).



Induct. Base case given by condition D.

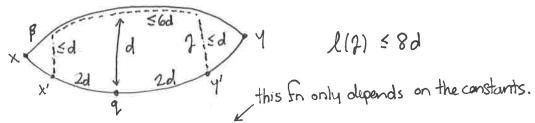
Step 2. Improve this when B is geodesic: p(x,y) is close to B.

Let q = furthest pt on p(x,y) from B. say d(q, B) = d.

actual dist, not dist along p(x,y)

Pick x', y' & p(x,y) before lafter q at distance 2d

Have:



→ d ≤ d(q, f) ≤ O(loyd) → d bounded above. I look at pic. by Step 1 and (2) applied to X', y'.

Step 3. B close to p(x,y) (similar)

VIII