Complexidade de Algoritmos

$$\begin{bmatrix} \exists X \exists \end{bmatrix} A \begin{cases} \Gamma(n) = \Gamma(n/2) + n \\ \Gamma(1) = 1 \end{cases}$$

$$T(n) = T(\frac{n}{2}) + n = [T(\frac{n}{2}) + \frac{n}{4}] + n$$

$$T(n) = [T(\frac{n}{2}) + \frac{n}{4}] + n = [T(\frac{n}{2}) + \frac{n}{2}] + n$$

$$T(n) = T(n/2k) + n/2k-1 + n/2k-2 + ... + n/2 + n(2k-1) = 1$$

flasumindo 
$$\frac{n}{2^k} = 1$$
, então temos  $n = 2^k e^{-k} = \log n$ 

$$T(n) = T(1) + n \left[ \frac{1}{2^{K-1}} + \frac{1}{2^{K-2}} + \dots + \frac{1}{2} + 1 \right] = 1 + n \left[ 1 + 1 \right]$$

$$\left[ T(n) = 1 + 2n \right] = O(n)$$

$$B[T(n) = aT(n-1) + n$$
  
 $T(1) = 1$ 

$$T(n) = a(aT(n-a) + n-1) + n = 4T(n-a) + a(n-1) + n$$
  
 $T(n) = 8T(n-3) + 4(n-a) + a(n-1) + n$ 

$$T(n) = 2^{n-1} + (1) + \sum_{i=0}^{n-2} 2^{i}(n-i) = 2 + \sum_{i=0}^{n-1} 2^{i}(n-i)$$

$$T(n) = 2^{n-1} + \left( n \sum_{i=0}^{n-2} 2^{i} - \sum_{i=0}^{n-2} i \cdot 2^{i} \right) = 2 + \left( n \cdot (2^{n-1} - 1) - \frac{n \cdot 2^{n} - 3 \cdot 2^{n} + 4}{2} \right)$$

$$T(n) = (n(2^{n-1} - 1) - (n \cdot 2^{n-1} - 3 \cdot 2^{n-1} + 2)) + 2^{n-1}$$

$$T(n) = (3 \cdot 2^{n-1} - n - 2) + 2^{n-1} = 2 \cdot 2^{n-1} - n - 2 = 2^n - n - 2$$

$$T(n) = 2^n - n - 2$$