

Internet Protocols EBU5403

The Data Link Layer Part 2

D2

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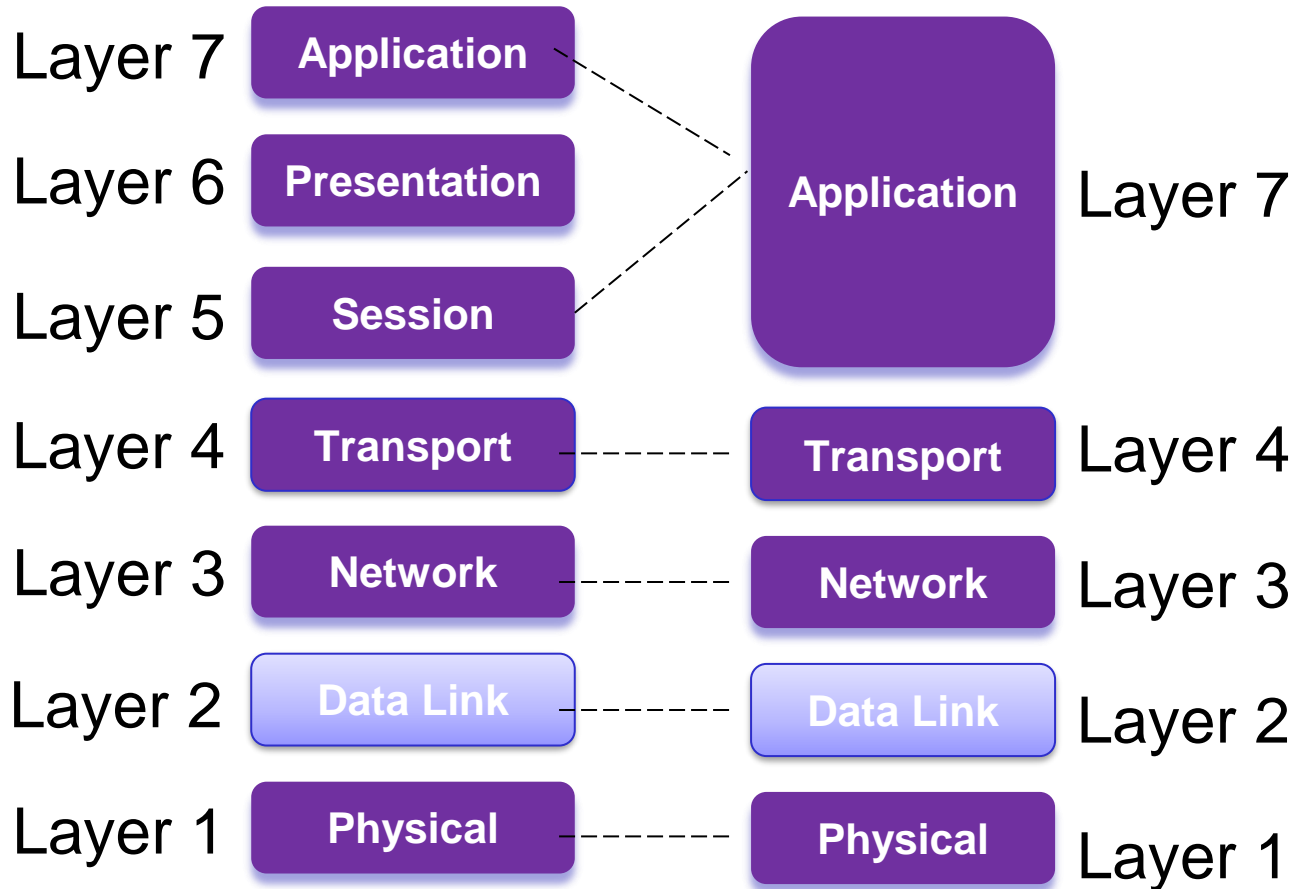
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	Part 1	Part 2	Part 3	Part 4
Ecommerce + Telecoms 1	Richard Clegg		Cunhua Pan	
Telecoms 2				

Structure of course

- Part A
 - Introduction to IP Networks
 - The Transport layer (part I)
- Part B
 - The Transport layer (part II)
 - The Network layer (part I)
 - Class test (open book exam in class)
- Part C
 - The Network layer (part II)
 - The Data link layer (part I)
 - Router lab tutorial (assessed labwork after this week)
- Part D
 - The Data link layer (part II)
 - Security and network management
 - Class test

Data Link Layer



Link layer, LANs: outline

6.1 introduction, services

6.2 error detection,
correction

6.3 multiple access
protocols

6.4 LANs

- addressing, ARP
- Ethernet
- WiFi
- switches
- WLANS

6.5 link virtualization:
MPLS

6.6 a day in the life of a
web request

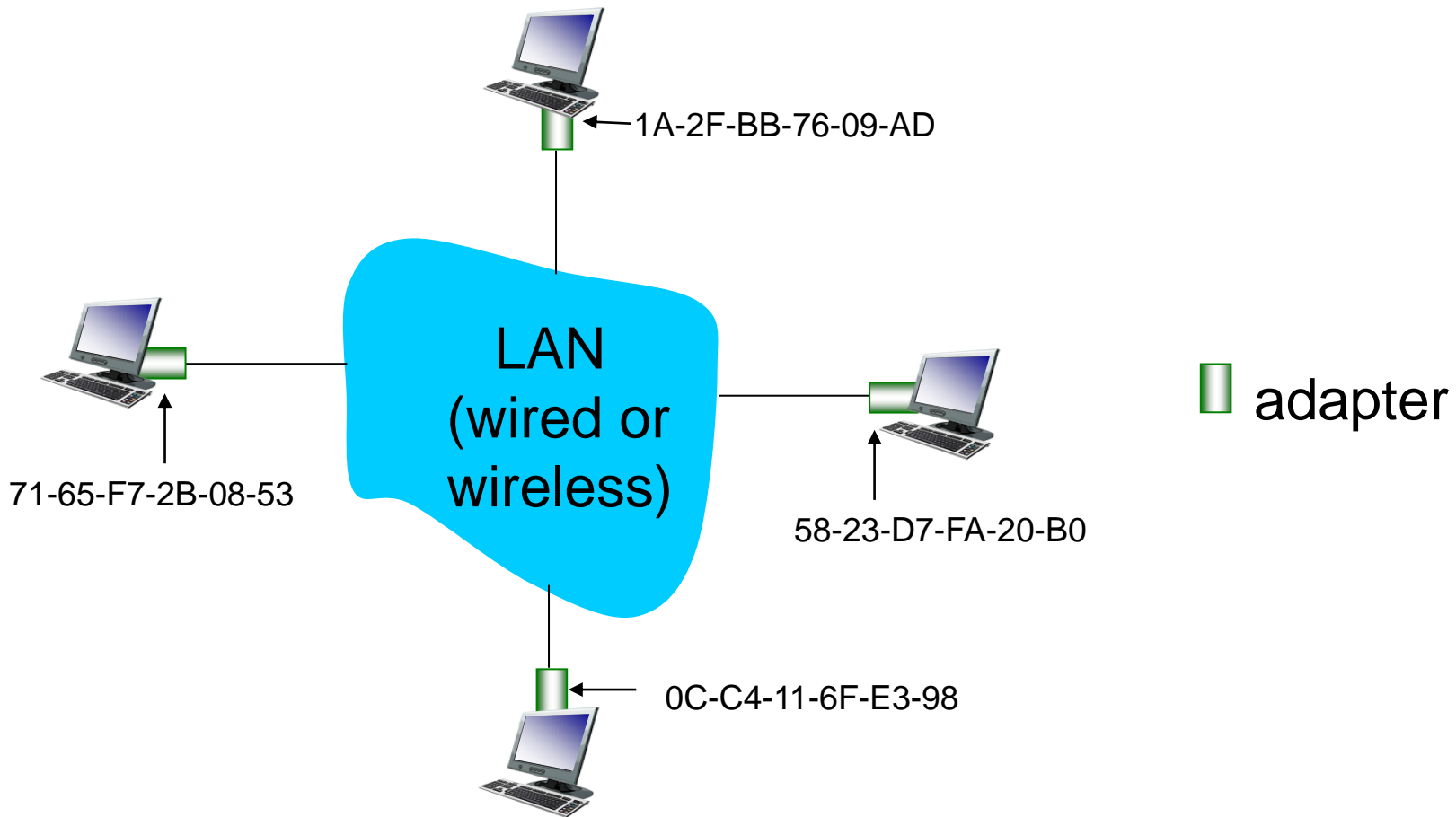
MAC addresses and ARP

- 32-bit IP address:
 - *network-layer* address for interface
 - used for layer 3 (network layer) forwarding
- MAC (or LAN or physical or Ethernet) address:
 - function: *used ‘locally’ to get frame from one interface to another physically-connected interface (same network, in IP-addressing sense)*
 - 48 bit MAC address (for most LANs) burned in NIC ROM, also sometimes software settable
 - e.g.: 1A-2F-BB-76-09-AD

hexadecimal (base 16) notation
(each “numeral” represents 4 bits)

LAN addresses and ARP

each adapter on LAN has unique **LAN** address

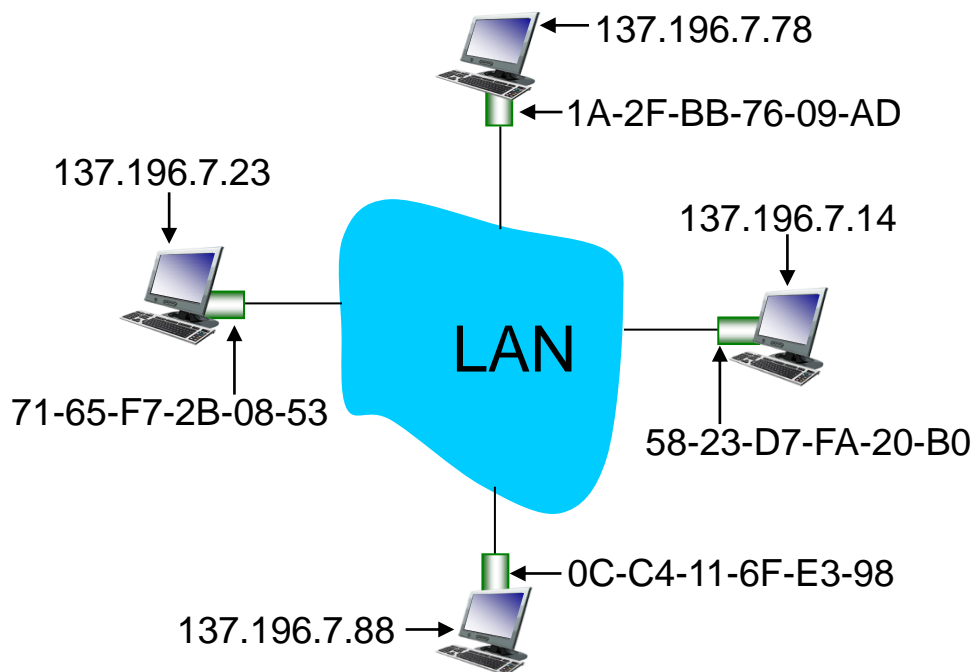


LAN addresses (more)

- MAC address allocation administered by IEEE
- manufacturer buys portion of MAC address space (to assure uniqueness)
- analogy:
 - MAC address: like Social Security Number
 - IP address: like postal address
- MAC flat address → portability
 - can move LAN card from one LAN to another
- IP hierarchical address *not* portable
 - address depends on IP subnet to which node is attached

ARP: address resolution protocol

Question: how to determine interface's MAC address, knowing its IP address?



ARP table: each IP node (host, router) on LAN has table

- IP/MAC address mappings for some LAN nodes:
< IP address; MAC address; TTL >
- TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min)

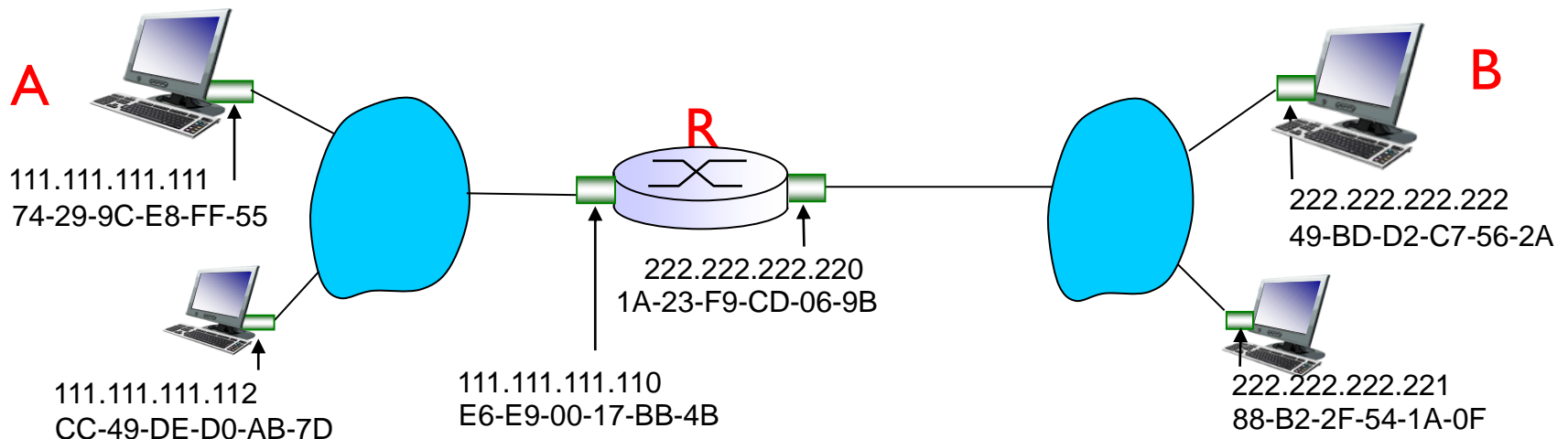
ARP protocol: same LAN

- A wants to send datagram to B
 - B's MAC address not in A's ARP table.
- A **broadcasts** ARP query packet, containing B's IP address
 - destination MAC address = FF-FF-FF-FF-FF-FF
 - all nodes on LAN receive ARP query
- B receives ARP packet, replies to A with its (B's) MAC address
 - frame sent to A's MAC address (unicast)
- A caches (saves) IP-to-MAC address pair in its ARP table until information becomes old (times out)
 - soft state: information that times out (goes away) unless refreshed
- ARP is “plug-and-play”:
 - nodes create their ARP tables *without intervention from net administrator*

Addressing: routing to another LAN

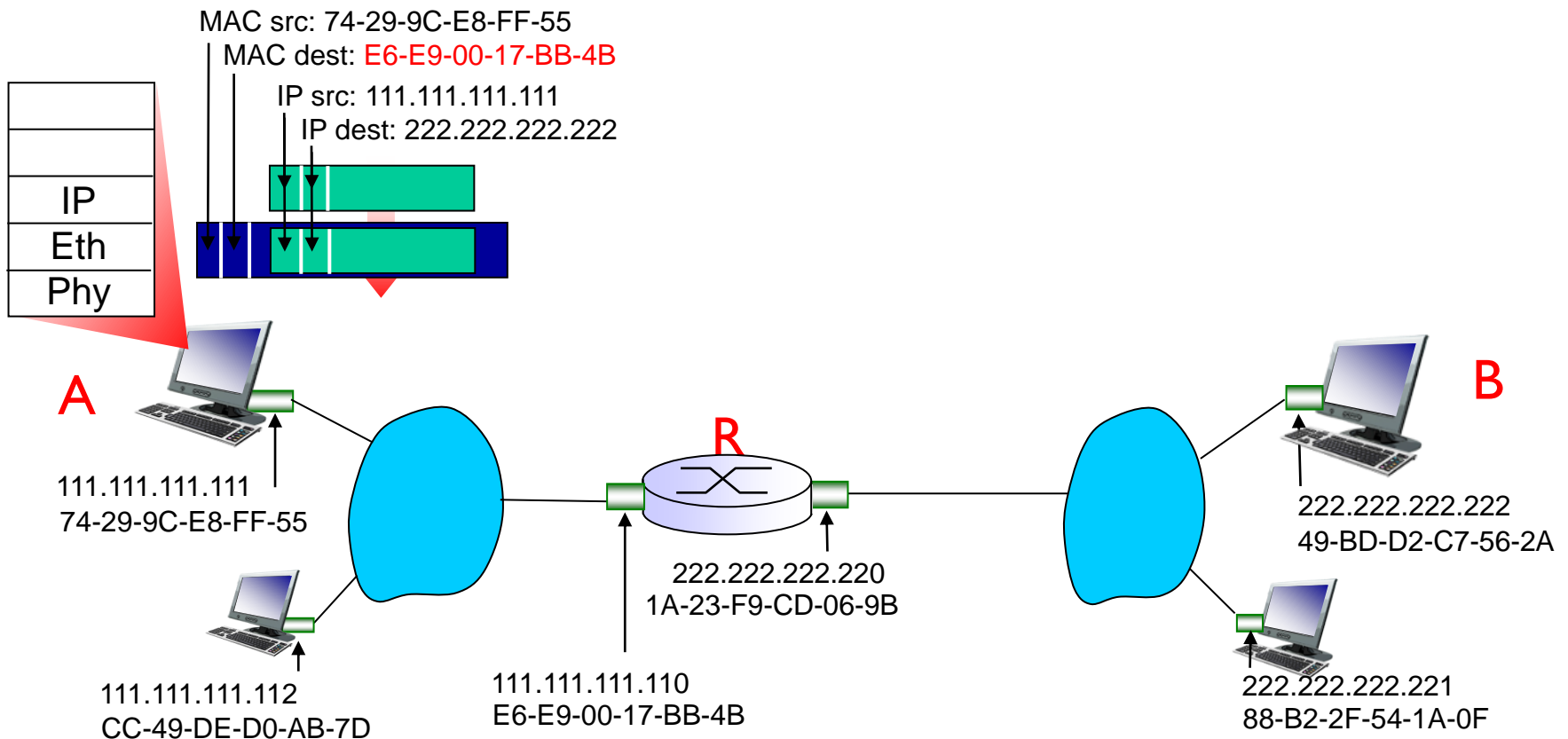
walkthrough: **send datagram from A to B via R**

- focus on addressing – at IP (datagram) and MAC layer (frame)
- assume A knows B's IP address
- assume A knows IP address of first hop router, R (how?)
- assume A knows R's MAC address (how?)



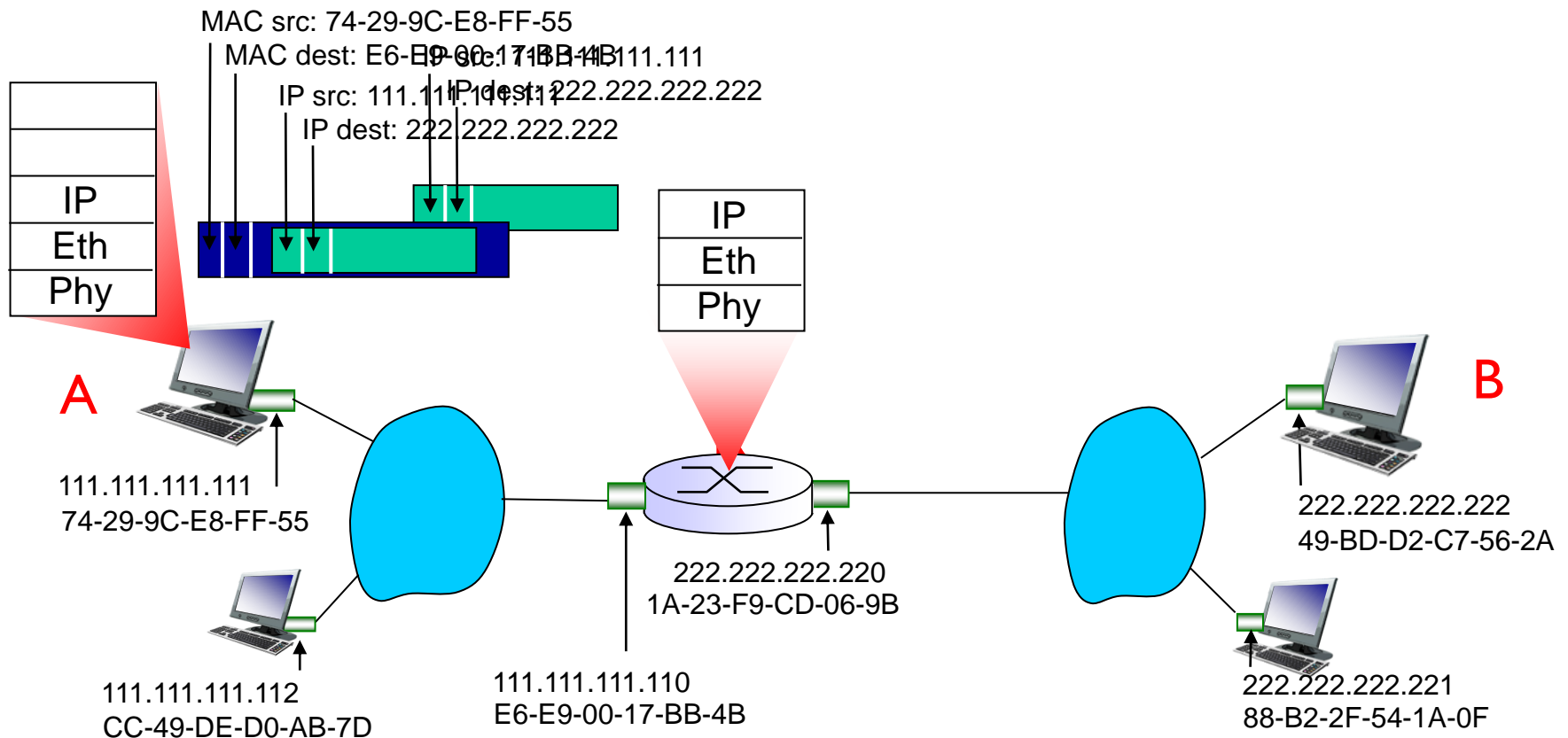
Addressing: routing to another LAN

- A creates IP datagram with IP source A, destination B
- A creates link-layer frame with R's MAC address as destination address, frame contains A-to-B IP datagram



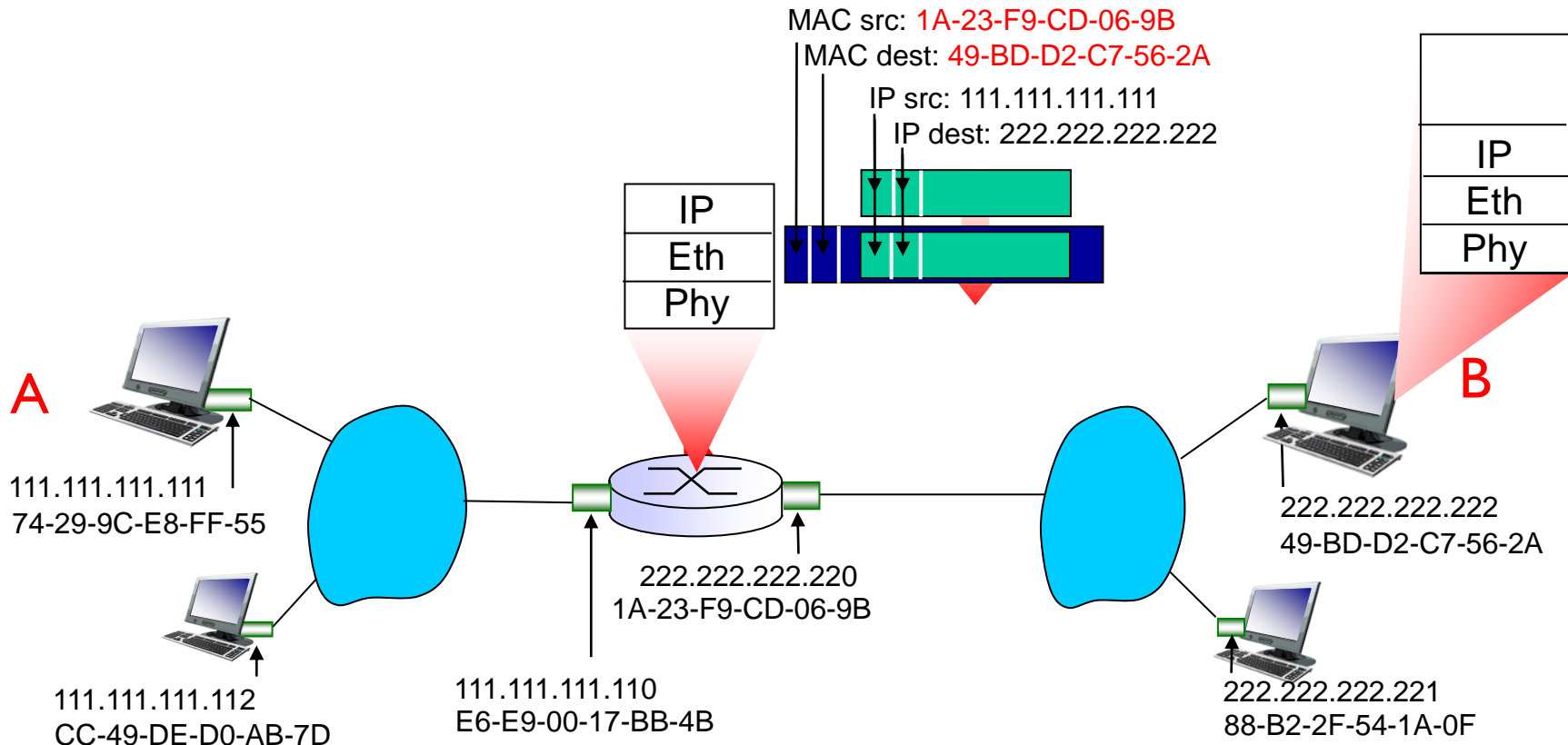
Addressing: routing to another LAN

- frame sent from A to R
- frame received at R, datagram removed, passed up to IP



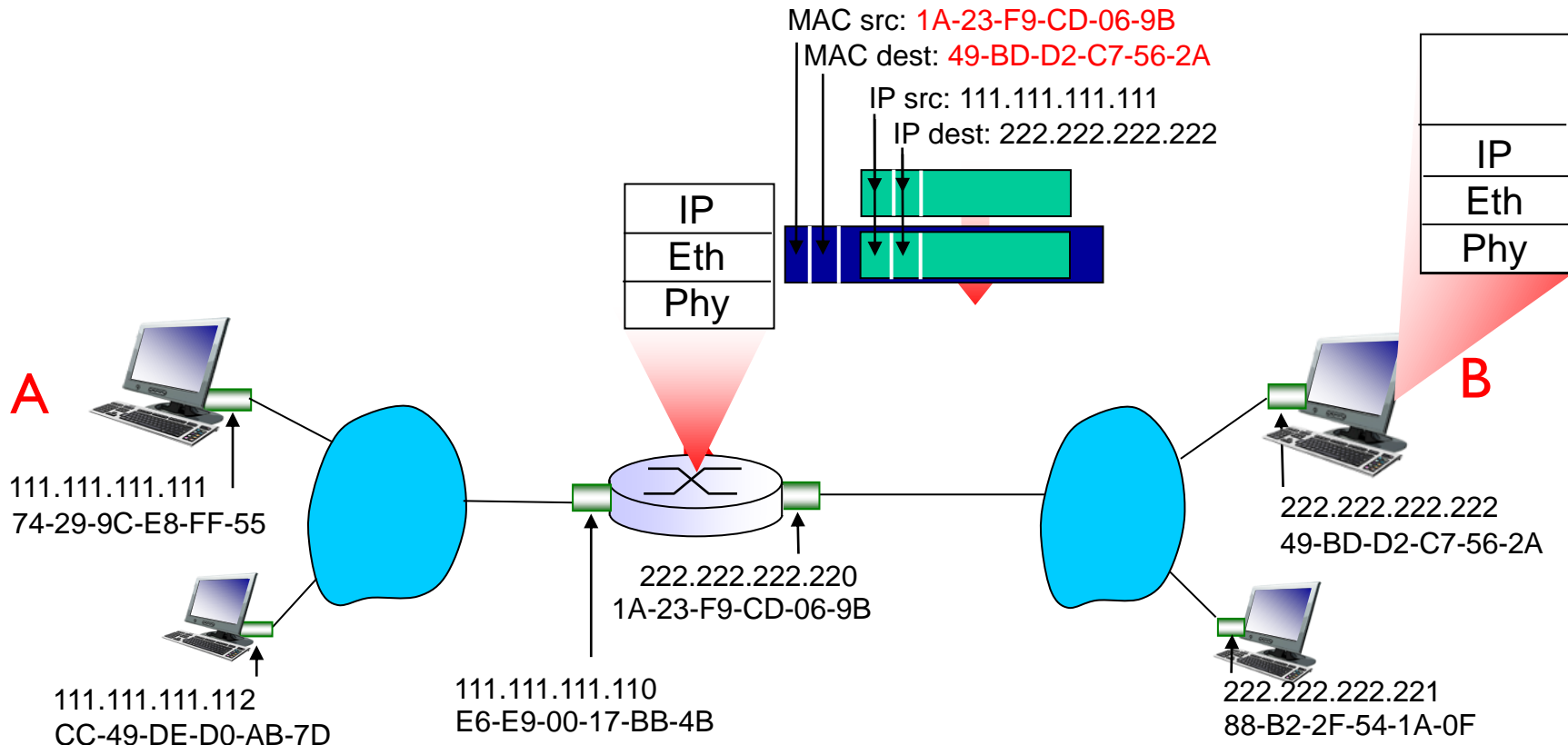
Addressing: routing to another LAN

- R forwards datagram with IP source A, destination B
- R creates link-layer frame with B's MAC address as destination address, frame contains A-to-B IP datagram



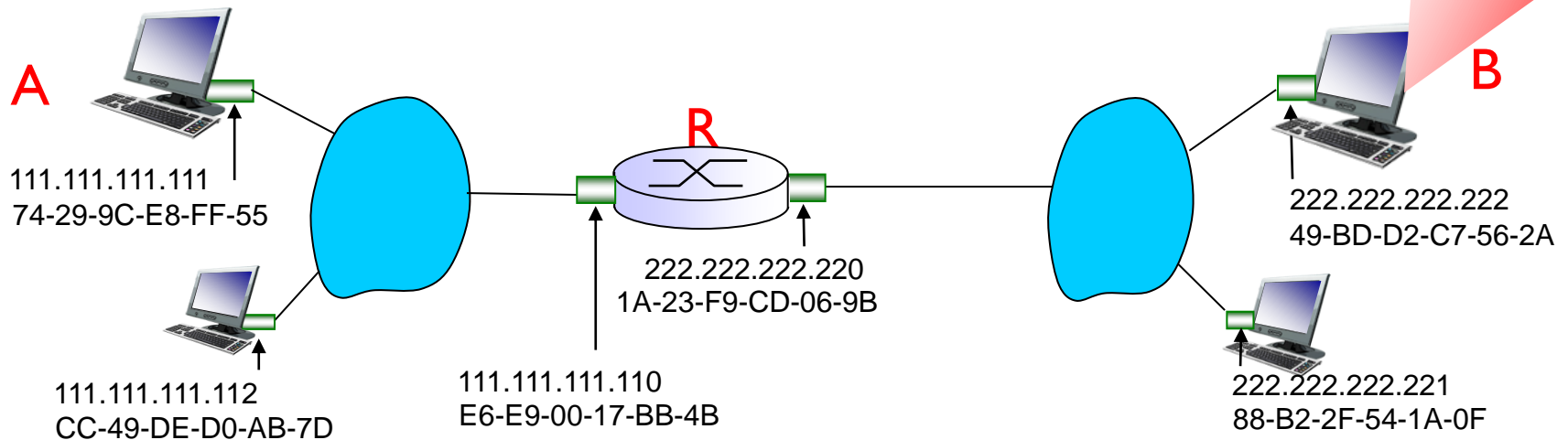
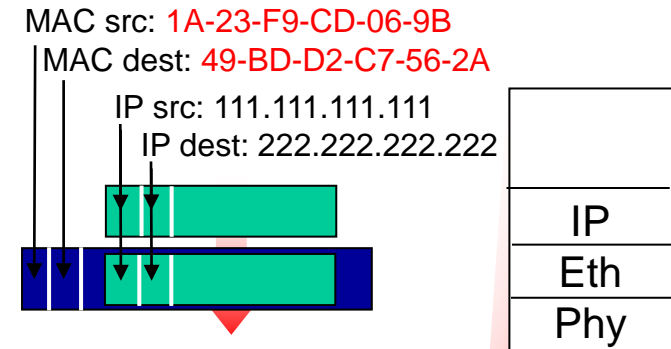
Addressing: routing to another LAN

- R forwards datagram with IP source A, destination B
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Addressing: routing to another LAN

- R forwards datagram with IP source A, destination B
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* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/

Test your understanding

- Address Resolution Protocol, ARP is used to find address when only address is known

Test your understanding

- Address Resolution Protocol, ARP is used to find ... the destination MAC ... address when only ... the destination IP address is known

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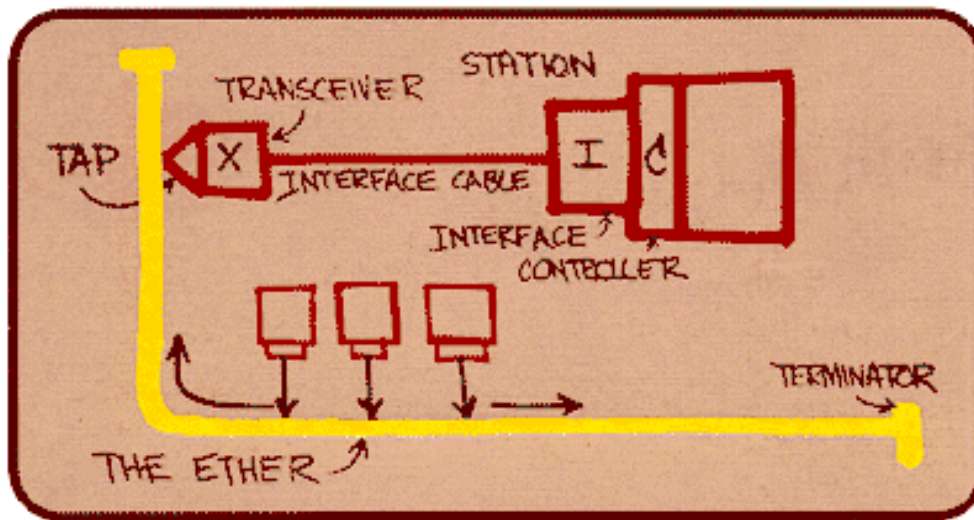
6.6 data center
networking

6.7 a day in the life of a
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Ethernet

“dominant” wired LAN technology:

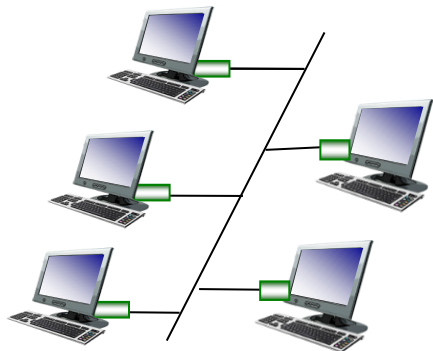
- single chip, multiple speeds (e.g., Broadcom BCM5761)
- first widely used LAN technology
- simpler, cheap
- kept up with speed race: 10 Mbps – 10 Gbps



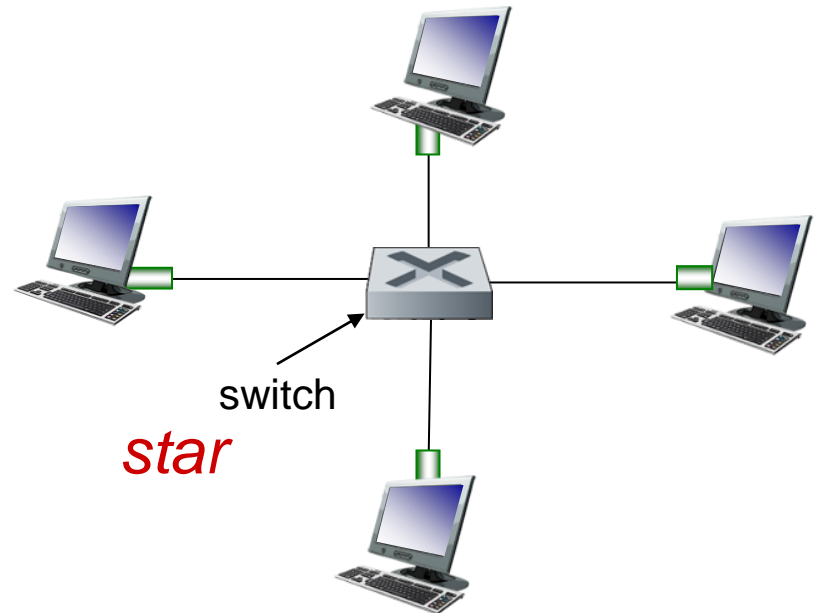
Metcalfe's Ethernet sketch

Ethernet: physical topology

- **bus:** popular through mid 90s
 - all nodes in same collision domain (can collide with each other)
- **star:** prevails today
 - active **switch** in center
 - each “spoke” runs a (separate) Ethernet protocol (nodes do not collide with each other)



bus: coaxial cable



Ethernet frame structure

sending adapter encapsulates IP datagram (or other network layer protocol packet) in **Ethernet frame**



preamble:

- 7 bytes with pattern 10101010 followed by one byte with pattern 10101011
- used to synchronize receiver, sender clock rates

Ethernet frame structure (more)

- **addresses:** 6 byte source, destination MAC addresses
 - if adapter receives frame with matching destination address, or with broadcast address (e.g. ARP packet), it passes data in frame to network layer protocol
 - otherwise, adapter discards frame
- **type:** indicates higher layer protocol (mostly IP but others possible, e.g., Novell IPX, AppleTalk)
- **CRC:** cyclic redundancy check at receiver
 - error detected: frame is dropped



Ethernet: unreliable, connectionless

- *connectionless*: no handshaking between sending and receiving NICs
- *unreliable*: receiving NIC doesn't send ACK or NACK to sending NIC
 - data in dropped frames recovered only if initial sender uses higher layer rdt (e.g., TCP), otherwise dropped data lost
- Ethernet's MAC protocol: unslotted *CSMA/CD with binary backoff*
- Various types of wired MAC protocol e.g.
 - Copper (100BASE-TX, 100BASE-T2, 100BASE-T4)
 - Optical (100BASE-SX, 100BASE-FX, 100BASE-BX)

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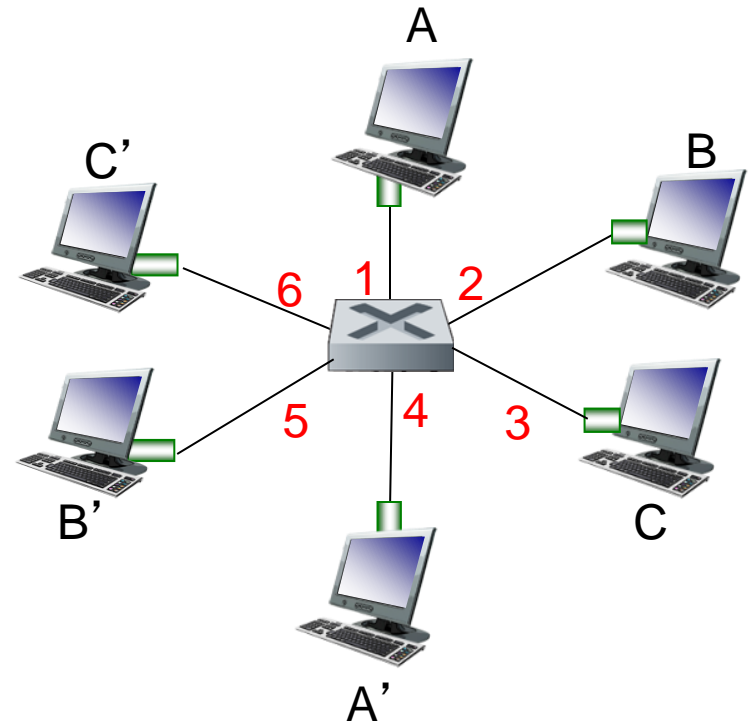
6.6 a day in the life of a
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Ethernet switch

- link-layer device: takes an *active* role
 - store, forward Ethernet frames
 - examine incoming frame's MAC address, *selectively* forward frame to one-or-more outgoing links when frame is to be forwarded on segment, uses CSMA/CD to access segment
- *transparent*
 - hosts are unaware of presence of switches
- *plug-and-play, self-learning*
 - switches do not need to be configured

Switch: *multiple* simultaneous transmissions

- hosts have dedicated, direct connection to switch
- switches buffer packets
- Ethernet protocol used on *each* incoming link, but no collisions; full duplex
 - each link is its own collision domain
- **switching:** A-to-A' and B-to-B' can transmit simultaneously, without collisions



*switch with six interfaces
(1,2,3,4,5,6)*

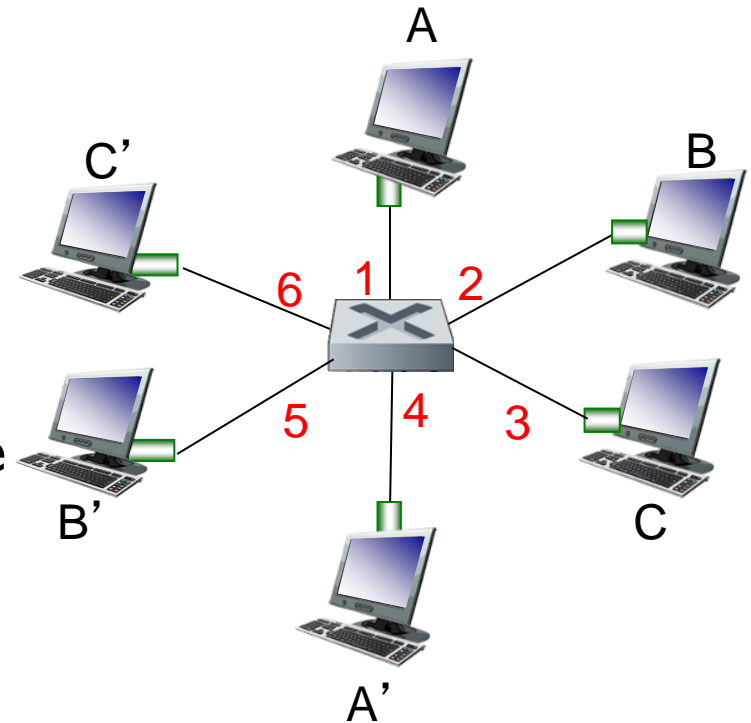
Switch forwarding table

Q: how does switch know A' reachable via interface 4, B' reachable via interface 5?

- A: each switch has a **switch table**, each entry:
 - (MAC address of host, interface to reach host, time stamp)
 - looks like a routing table!

Q: how are entries created, maintained in switch table?

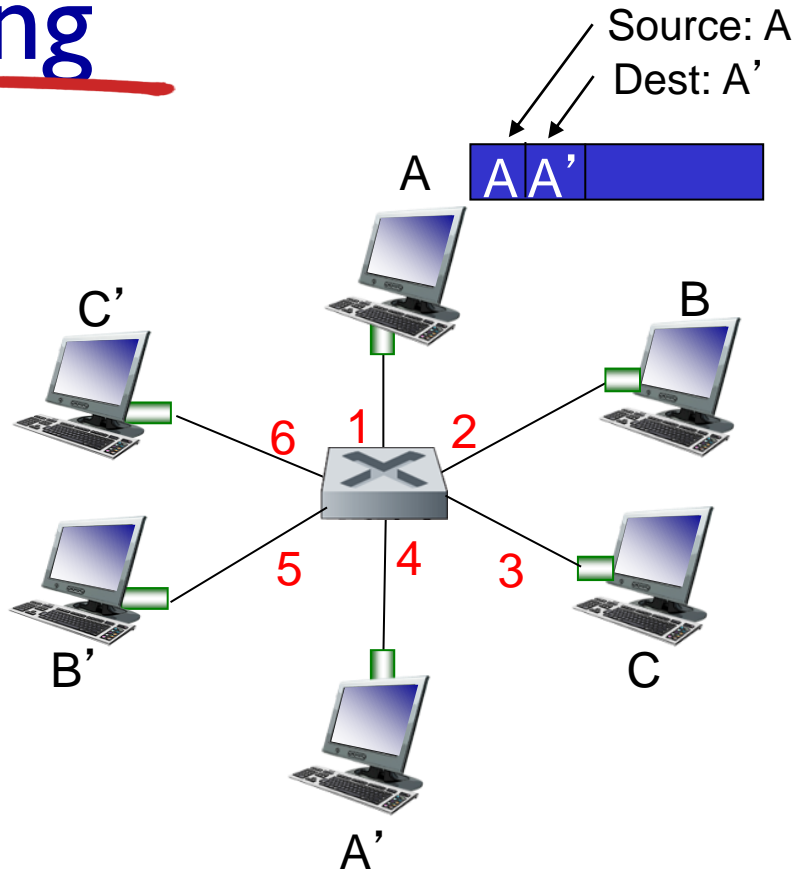
- something like a routing protocol?



*switch with six interfaces
(1,2,3,4,5,6)*

Switch: self-learning

- switch *learns* which hosts can be reached through which interfaces
 - when frame received, switch “learns” location of sender: incoming LAN segment
 - records sender/location pair in switch table



MAC addr	interface	TTL
A	1	60

*Switch table
(initially empty)*

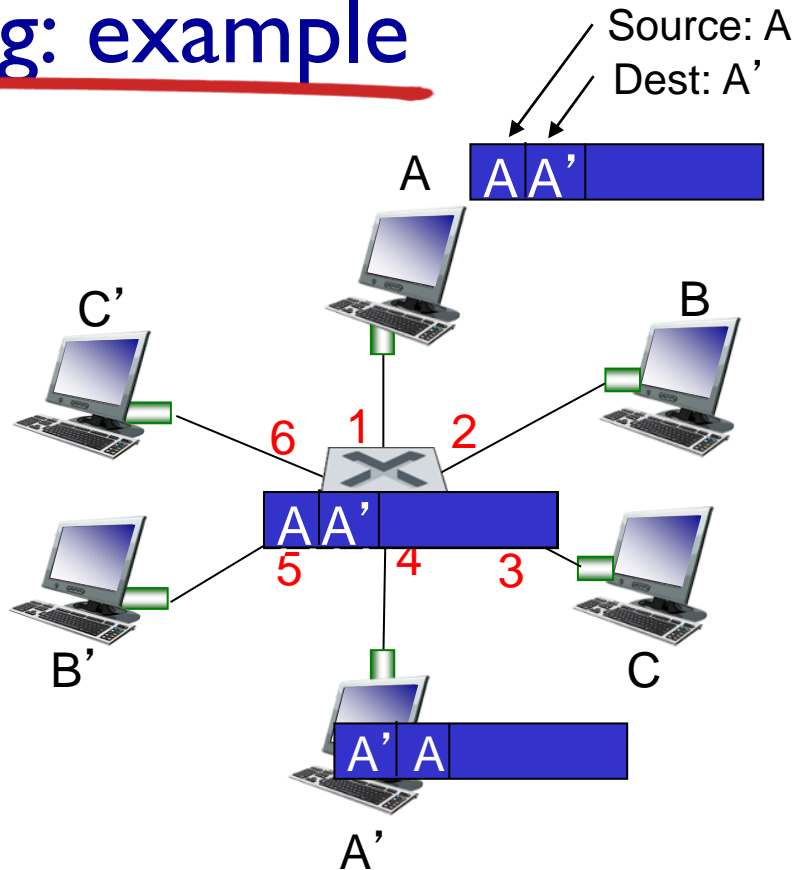
Switch: frame filtering/forwarding

when frame received at switch:

1. record incoming link, MAC address of sending host
2. index switch table using MAC destination address
3. **if** entry found for destination
 then {
 if destination on segment from which frame arrived
 then drop frame
 else forward frame on interface indicated by entry
 }
 else flood /* forward on all interfaces except arriving interface */

Self-learning, forwarding: example

- frame destination, A', location unknown: *flood*
- destination A location known: *selectively send on just one link*

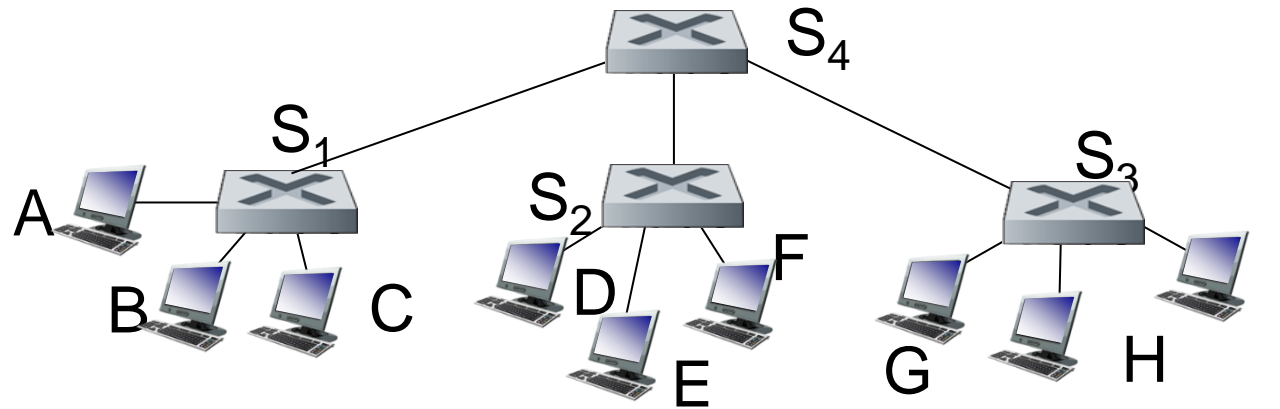


MAC addr	interface	TTL
A	1	60
A'	4	60

*switch table
(initially empty)*

Interconnecting switches

self-learning switches can be connected together:

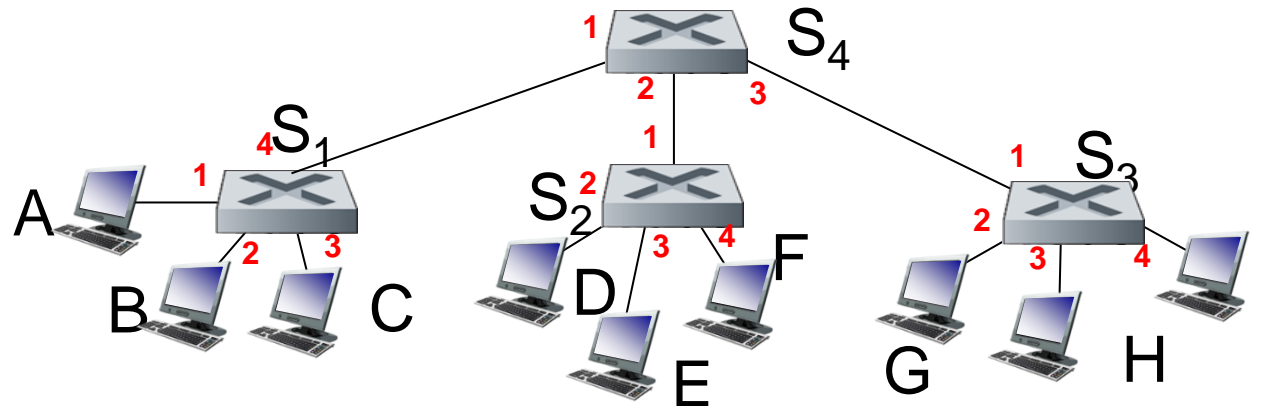


Q: sending from A to G - how does S₁ know to forward frame destined to G via S₄ and S₃?

- **A:** self learning! (works exactly the same as in single-switch case!)

Self-learning multi-switch example

Suppose C sends frame to I, I responds to C



- Q: show switch tables and packet forwarding in S₁, S₂, S₃, S₄

S1

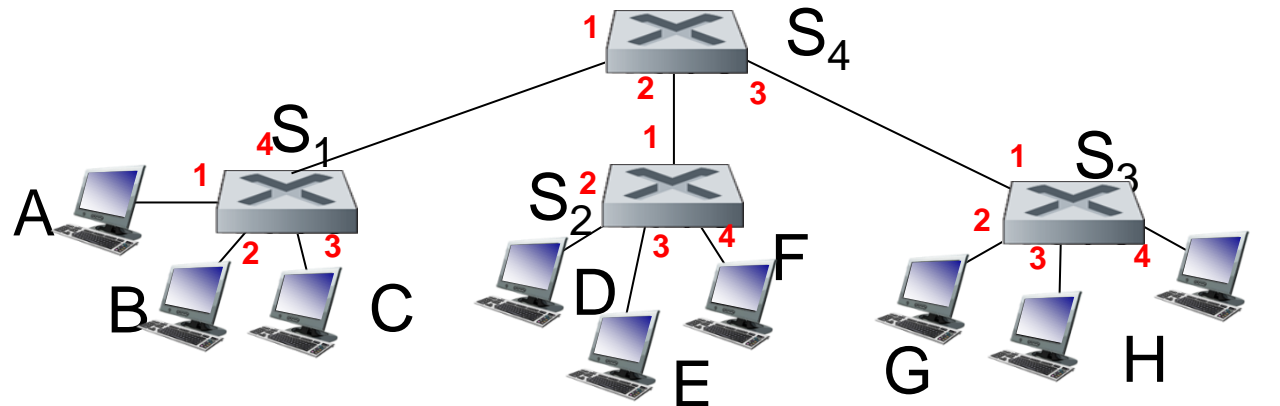
S2

S3

S4

Self-learning multi-switch example

Suppose C sends frame to I, I responds to C



- Q: show switch tables and packet forwarding in S_1, S_2, S_3, S_4

C	3
I	4

S1

C	I

S2

C	I
I	4

S3

C	I
I	3

S4

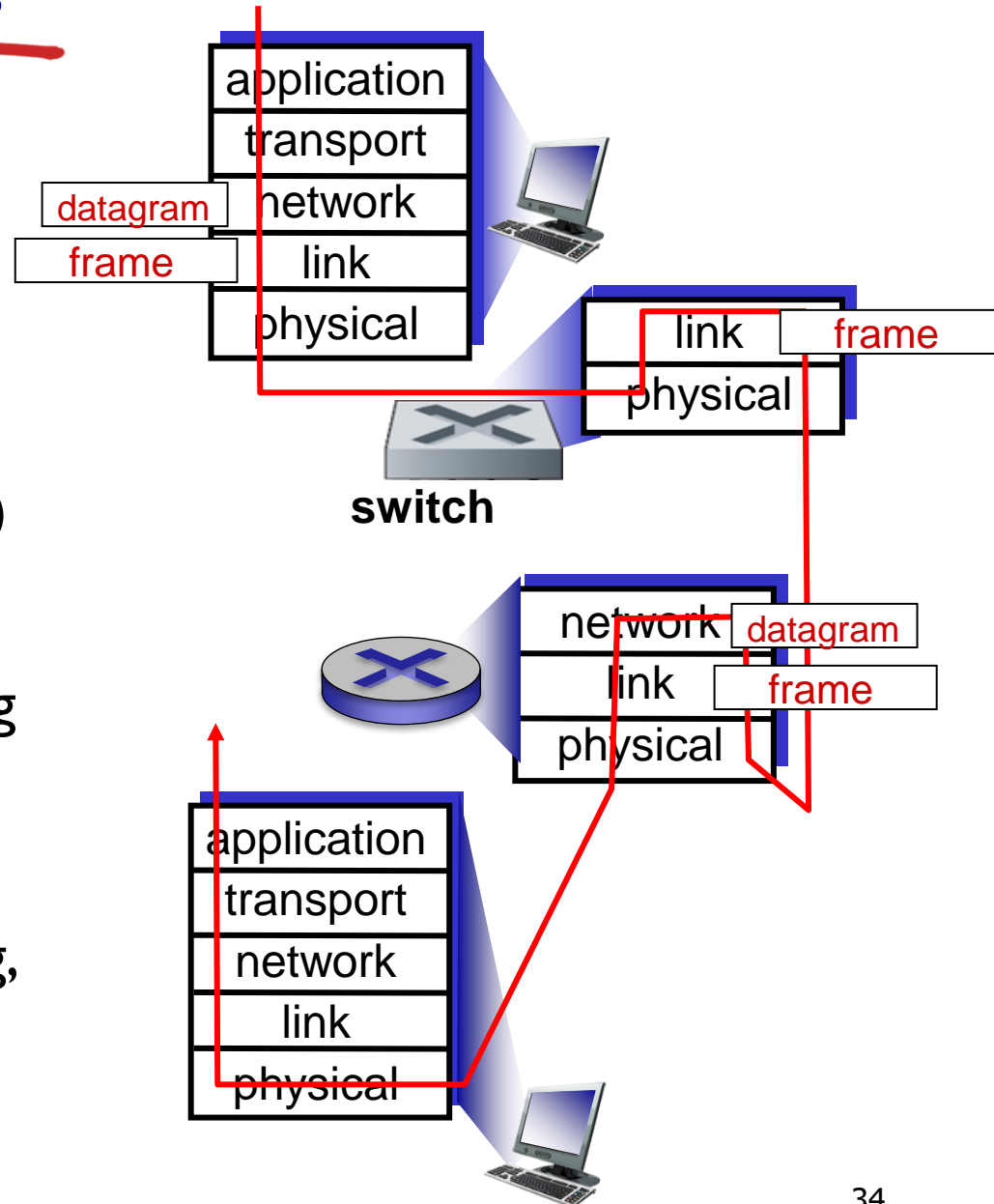
Switches vs. routers

both are store-and-forward:

- **routers:** network-layer devices (examine network-layer headers)
- **switches:** link-layer devices (examine link-layer headers)

both have forwarding tables:

- **routers:** compute tables using routing algorithms, IP addresses
- **switches:** learn forwarding table using flooding, learning, MAC addresses



What have we learned?

- *MAC address*
 - Format
- *ARP* : Determine MAC address, knowing its IP address
 - Addressing when routing to another networks
- *Ethernet Protocol*,
 - Physical topology
 - Frame structure
- *Ethernet Switches*
 - Forwarding table
 - Self learning
 - Interconnecting switches