

# Homework 3

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## Boolean Algebra

1

(a)  $A * (\bar{A} + B * B) + \overline{(B + A)} * (\bar{A} + B)$

$$\begin{aligned} & A * (\bar{A} + B * B) + (\bar{B} * \bar{A}) * (\bar{A} + B) \quad \text{DeMorgan's} \\ & A * (\bar{A} + B) + (\bar{B} * \bar{A}) * (\bar{A} + B) \quad \text{Idempotent Law } (B * B = B) \\ & (A + (\bar{B} * \bar{A}))(\bar{A} + B) \quad \text{Distributive Law (Inversed)} \\ & (A + (\bar{A} * \bar{B}))(\bar{A} + B) \quad \text{Commutative Law} \\ & (A + \bar{B})(\bar{A} + B) \quad \text{Redundancy Law } (A + (\bar{A} * \bar{B}) = A + \bar{B}) \\ & A(\bar{A} + B) + \bar{B}(\bar{A} + B) \quad \text{Distributive Law} \\ & (A * \bar{A} + A * B) + \bar{B}(\bar{A} + B) \quad \text{Distributive Law} \\ & (0 + A * B) + \bar{B}(\bar{A} + B) \quad \text{Inverse Law } (A * \bar{A} = 0) \\ & (A * B) + \bar{B}(\bar{A} + B) \quad \text{Identity Law } (0 + A * B = A * B) \\ & (A * B) + (\bar{B} * \bar{A} + \bar{B} * B) \quad \text{Distributive Law} \\ & (A * B) + (\bar{B} * \bar{A} + 0) \quad \text{Inverse Law } (\bar{B} * B = 0) \\ & (A * B) + (\bar{B} * \bar{A}) \quad \text{Identity Law } (\bar{B} * \bar{A} + 0 = \bar{B} * \bar{A}) \end{aligned}$$

$$\boxed{A * B + \bar{B} * \bar{A}}$$

(b)  $\overline{C * B} + (A * B * C) + \overline{A + C + \bar{B}}$

$$\begin{aligned} & \overline{C} + \bar{B} + (A * B * C) + \overline{A + C + \bar{B}} \quad \text{DeMorgan's} \\ & \overline{C} + (A * B * C) + \bar{B} + \overline{A + C + \bar{B}} \quad \text{Commutative Law} \\ & \overline{C} + (C * A * B) + \bar{B} + \overline{A + C + \bar{B}} \quad \text{Commutative Law} \\ & \overline{C} + (A * B) + \bar{B} + \overline{A + C + \bar{B}} \quad \text{Redundancy Law } (\overline{C} + (C * A * B) = \overline{C} + (A * B)) \\ & \overline{C} + \bar{B} + (A * B) + \overline{A + C + \bar{B}} \quad \text{Commutative Law} \\ & \overline{C} + \bar{B} + (B * A) + \overline{A + C + \bar{B}} \quad \text{Commutative Law} \\ & \overline{C} + (\bar{B} + A) + \overline{A + C + \bar{B}} \quad \text{Redundancy Law } (\bar{B} + (B * A) = \bar{B} + A) \\ & \overline{C} + (\bar{B} + A) + \bar{A} * \bar{C} * \bar{B} \quad \text{DeMorgan's} \\ & \overline{C} + (\bar{B} + A) + \bar{A} * \bar{C} * B \quad \text{Double Negation Law} \\ & \overline{C} + (\bar{A} * \bar{C} * B) + (\bar{B} + A) \quad \text{Commutative Law} \\ & \overline{C} + (\bar{C} * \bar{A} * B) + (\bar{B} + A) \quad \text{Commutative Law} \\ & \overline{C} + (\bar{B} + A) \quad \text{Absorption Law } (\overline{C} + (\bar{C} * \bar{A} * B) = \overline{C}) \end{aligned}$$

$$\boxed{\overline{C} + \bar{B} + A}$$

$$(c) (A + B) * (\overline{A} + C) * (\overline{C} + B)$$

$$\begin{aligned}
& (A(\overline{A} + C) + B(\overline{A} + C)) * (\overline{C} + B) \quad \text{Distributive Law} \\
& ((A * \overline{A} + A * C) + B(\overline{A} + C)) * (\overline{C} + B) \quad \text{Distributive Law} \\
& (A * \overline{A} + A * C + B * \overline{A} + B * C) * (\overline{C} + B) \quad \text{Distributive Law} \\
& ((A * \overline{A} * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A} * B) + (B * C * B)) \quad \text{Distributive Law} \\
& ((A * \overline{A} * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A} * B) + (B * B * C)) \quad \text{Commutative Law} \\
& ((A * \overline{A} * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A} * B) + (B * C)) \quad \text{Idempotent Law} \\
& ((A * \overline{A} * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * B * \overline{A}) + (B * C)) \quad \text{Commutative Law} \\
& ((A * \overline{A} * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Idempotent Law} \\
& ((0 * \overline{C}) + (A * C * \overline{C}) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Inverse Law} \\
& ((0 * \overline{C}) + (A * 0) + (B * \overline{A} * \overline{C}) + (B * C * \overline{C})) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Inverse Law} \\
& ((0 * \overline{C}) + (A * 0) + (B * \overline{A} * \overline{C}) + (B * 0)) + \\
& ((A * \overline{A} * B) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Inverse Law} \\
& ((0 * \overline{C}) + (A * 0) + (B * \overline{A} * \overline{C}) + (B * 0)) + \\
& ((0 * B) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Inverse Law} \\
& ((0) + (0) + (B * \overline{A} * \overline{C}) + (0)) + \\
& ((0) + (A * C * B) + (B * \overline{A}) + (B * C)) \quad \text{Law of Zeros (Too much to do them one by one ...)} \\
& (B * \overline{A} * \overline{C}) + (A * C * B) + (B * \overline{A}) + (B * C) \quad \text{Identity Law (Too much to do them one by one...)} \\
& (B * \overline{A}) + (B * \overline{A} * \overline{C}) + (A * C * B) + (B * C) \quad \text{Commutative Law} \\
& (B * \overline{A}) + (A * C * B) + (B * C) \quad \text{Absorption Law } ((B * \overline{A}) + (B * \overline{A} * \overline{C}) = (B * \overline{A})) \\
& (B * \overline{A}) + (B * C) + (A * C * B) \quad \text{Commutative Law} \\
& (B * \overline{A}) + (B * C) + (B * C * A) \quad \text{Commutative Law} \\
& (B * \overline{A}) + (B * C) \quad \text{Absorption Law } ((B * C) + (B * C * A) = (B * C)) \\
& B * (\overline{A} + C) \quad \text{Distributive Law (Inversed)}
\end{aligned}$$

$B * (\overline{A} + C)$

(a)  $(\bar{A} + C) * (\bar{B} + D + A) * (D + A * \bar{C}) * (\bar{D} + A) = 1$

$$\begin{aligned}
& (\bar{A} + C) * (\bar{B} + D + A) * (D + \overline{\bar{A} + \bar{C}}) * (\bar{D} + A) && \text{DeMorgan's} \\
& (\bar{A} + C) * (\bar{B} + D + A) * (D + \overline{\bar{A} + \bar{C}}) * (\bar{D} + A) && \text{Double Negation Law} \\
& (\bar{A} + C) * (D + \overline{\bar{A} + \bar{C}}) * (\bar{B} + D + A) * (\bar{D} + A) && \text{Commutative Law} \\
& (\bar{A} + C) * (\overline{\bar{A} + \bar{C}} + D) * (\bar{B} + D + A) * (\bar{D} + A) && \text{Commutative Law} \\
& ((\bar{A} + C) * D) * (\bar{B} + D + A) * (\bar{D} + A) && \text{Redundancy Law } ((\bar{A} + C) * (\overline{\bar{A} + \bar{C}} + D) = (\bar{A} + C) * D) \\
& (\bar{A} + C) * D * (D + A + \bar{B}) * (\bar{D} + A) && \text{Commutative Law} \\
& (\bar{A} + C) * D * (\bar{D} + A) && \text{Absorption Law } (D * (D + A + \bar{B}) = D) \\
& (\bar{A} + C) * (D * A) && \text{Redundancy Law } (D * (\bar{D} + A) = D * A) \\
& A * (\bar{A} + C) * D && \text{Commutative Law} \\
& A * C * D && \text{Redundancy Law } (A * (\bar{A} + C) = A * C) \\
& A * C * D = 1
\end{aligned}$$

$A = 1, C = 1, D = 1, B = 0 \text{ or } 1$
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(b)  $((\bar{K} * L * N) * (L + M)) + ((\bar{K} + L + N) * (K * \bar{L} * \bar{M})) * (\bar{K} + \bar{N}) = 1$

$$\begin{aligned}
& ((\bar{K} * L * (L + M) * N) + ((\bar{K} + L + N) * (K * \bar{L} * \bar{M}))) * (\bar{K} + \bar{N}) && \text{Commutative Law} \\
& ((\bar{K} * L * N) + ((\bar{K} + L + N) * (K * \bar{L} * \bar{M}))) * (\bar{K} + \bar{N}) && \text{Absorption Law } (L * (L + M) = L) \\
& ((\bar{K} * L * N) + (K * (\bar{K} + L + N) * \bar{L} * \bar{M})) * (\bar{K} + \bar{N}) && \text{Commutative Law} \\
& ((\bar{K} * L * N) + (K * (L + N) * \bar{L} * \bar{M})) * (\bar{K} + \bar{N}) && \text{Redundancy Law } (K * (\bar{K} + L + N) = K * (L + N)) \\
& ((\bar{K} * L * N) + (K * \bar{L} * (L + N) * \bar{M})) * (\bar{K} + \bar{N}) && \text{Commutative Law} \\
& ((\bar{K} * L * N) + (K * \bar{L} * N * \bar{M})) * (\bar{K} + \bar{N}) && \text{Redundancy Law } (\bar{L} * (L + N) = L * N) \\
& N * ((\bar{K} * L) + (K * \bar{L} * \bar{M})) * (\bar{K} + \bar{N}) && \text{Distributive Law (Inversed)} \\
& N * (\bar{K} + \bar{N}) * ((\bar{K} * L) + (K * \bar{L} * \bar{M})) && \text{Commutative Law} \\
& N * (\bar{N} + \bar{K}) * ((\bar{K} * L) + (K * \bar{L} * \bar{M})) && \text{Commutative Law} \\
& N * \bar{K} * ((\bar{K} * L) + (K * \bar{L} * \bar{M})) && \text{Redundancy Law } (N * (\bar{N} + \bar{K}) = N * \bar{K}) \\
& (\bar{K} * L * N * \bar{K}) + (K * \bar{L} * \bar{M} * N * \bar{K}) && \text{Distributive Law} \\
& (\bar{K} * \bar{K} * L * N) + (K * \bar{L} * \bar{M} * N * \bar{K}) && \text{Commutative Law} \\
& (\bar{K} * L * N) + (K * \bar{L} * \bar{M} * N * \bar{K}) && \text{Idempotent Law} \\
& (\bar{K} * L * N) + (K * \bar{K} * \bar{L} * \bar{M} * N) && \text{Commutative Law} \\
& (\bar{K} * L * N) + (0 * \bar{L} * \bar{M} * N) && \text{Inverse Law} \\
& (\bar{K} * L * N) + (0) && \text{Law of Zeros} \\
& (\bar{K} * L * N) && \text{Identity Law} \\
& \bar{K} * L * N = 1
\end{aligned}$$

$K = 0, L = 1, N = 1, M = 0 \text{ or } 1$
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### 3

#### Truth Table

$X$	$Y$	$Z$	$Q$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

#### K-Map

$Z$	$XY$			
	00	01	11	10
0	0	1	1	1
1	1	0	0	0

$Z$	$XY$			
	00	01	11	10
0	0	1	1	1
1	1	0	0	0

The merge shown by the yellow means "if  $Z$  is 0 and  $Y$  is 1, it doesn't matter what the value of  $X$  is."  
 In the function, it expresses –  $X * Y * \overline{Z} + \overline{X} * Y * \overline{Z} = Y * \overline{Z}$

$Z$	$XY$			
	00	01	11	10
0	0	1	1	1
1	1	0	0	0

The merge shown by the yellow means "if  $Z$  is 0 and  $X$  is 1, it doesn't matter what the value of  $Y$  is."  
 In the function, it expresses –  $X * Y * \overline{Z} + X * \overline{Y} * \overline{Z} = X * \overline{Z}$

After applying those two reductions, we get

Original :  $\overline{X} * \overline{Y} * Z + X * Y * \overline{Z} + \overline{X} * Y * \overline{Z} + X * \overline{Y} * \overline{Z}$

Reduced:

$$Y * \overline{Z} + X * \overline{Z} + \overline{X} * \overline{Y} * Z$$

# Logical Circuits

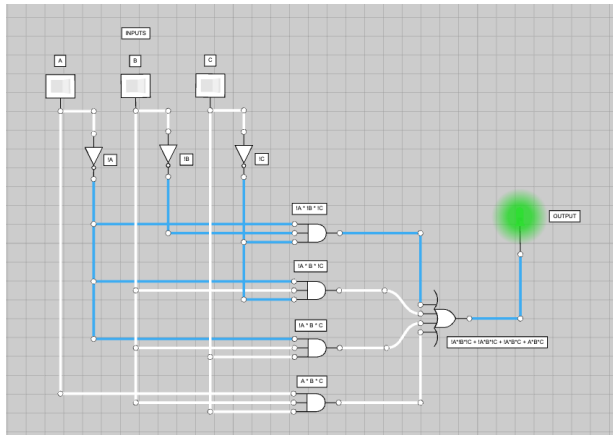
4

To obtain the sums of products representation, we look for the rows where the output is 1 and then form a product term for each row.

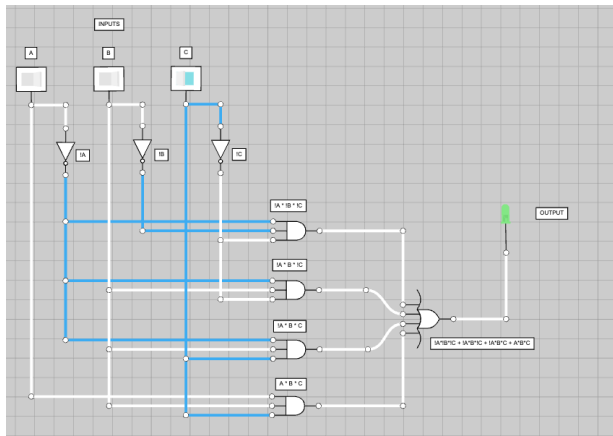
From the truth table, we get

$$(\overline{A} * \overline{B} * \overline{C}) + (\overline{A} * B * \overline{C}) + (\overline{A} * B * C) + (A * B * C)$$

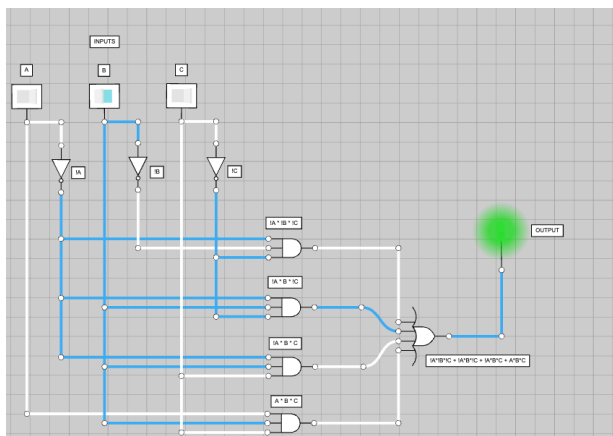
6



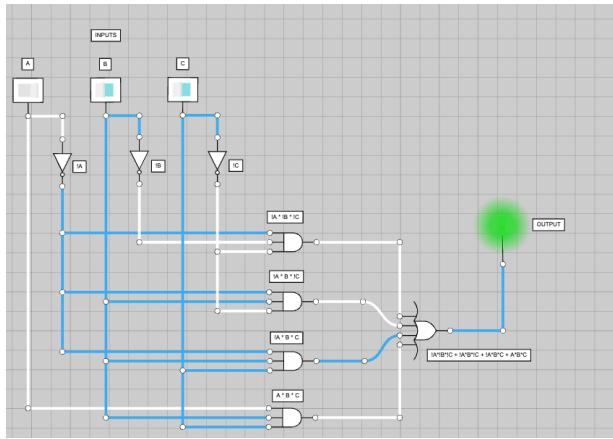
When input  $A = 0$ ,  $B = 0$ ,  $C = 0$ , the output is 1, like in the truth table (Also satisfies expression  $(\overline{A} * \overline{B} * \overline{C})$  as it should)



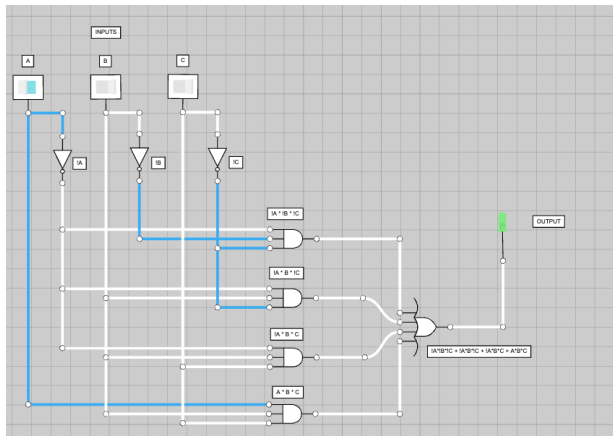
When input  $A = 0$ ,  $B = 0$ ,  $C = 1$ , the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



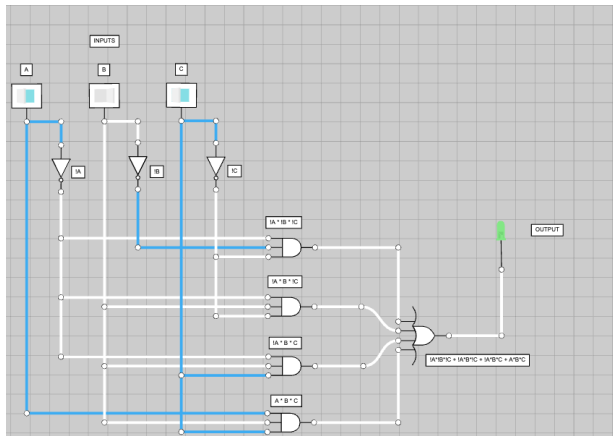
When input  $A = 0$ ,  $B = 1$ ,  $C = 0$ , the output is 1, like in the truth table (Also satisfies expression  $(\overline{A} * B * \overline{C})$ )



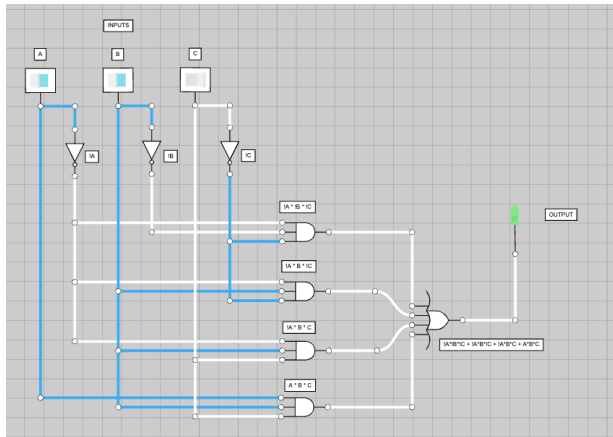
When input  $A = 0$ ,  $B = 1$ ,  $C = 1$ , the output is 1, like in the truth table (Also satisfies expression  $(\bar{A} * B * C)$ )



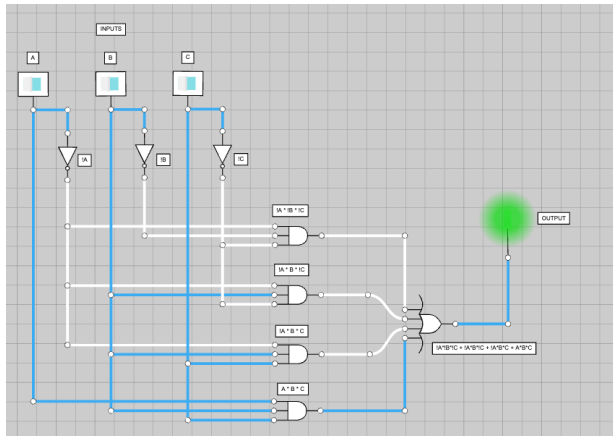
When input  $A = 1$ ,  $B = 0$ ,  $C = 0$ , the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input  $A = 1$ ,  $B = 0$ ,  $C = 1$ , the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input  $A = 1$ ,  $B = 1$ ,  $C = 0$ , the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input  $A = 1$ ,  $B = 1$ ,  $C = 1$ , the output is 1, like in the truth table (Also satisfies expression  $(A * B * C)$ )

After going through all possible combination of inputs, I find that the expected output with each different set of inputs is exactly like it is in the truth table. When the output is 1, the appropriate AND components also lights up, matching the appropriate expression in the boolean function (i.e.,  $(A * B * C)$  is an expression that lights up when  $A = 1$ ,  $B = 1$ ,  $C = 1$ ). Through these findings, I believe there is sufficient proof that my logical circuit does in fact implement the required Boolean function, and that the boolean function also correctly represents the truth table.

# Numerical Conversions and Arithmetic

8

**50.4375**

50.4375

$$50 = \underline{25} * 2 + 0$$

$$25 = \underline{12} * 2 + 1$$

$$12 = \underline{6} * 2 + 0$$

$$6 = \underline{3} * 2 + 0$$

$$3 = \underline{1} * 2 + 1$$

$$1 = 0 * 2 + 1$$

We get **110010**, which is 50 in binary representation

50.4375

$$.4375 * 2 = 0.875$$

$$.0875 * 2 = 1.75$$

$$.75 * 2 = 1.5$$

$$.5 * 2 = 1.0$$

We get **0111**, which is .4375 in binary representation

Combining, we get – **110010.0111**

Convert to scientific notation, we get – **1.100100111 \* 2<sup>5</sup>**

Where **100100111** is the mantisa

And **5** is the exponent

Convert the exponent to the bias(127 + 5 = **132**) in binary, we get **1000 0100**

We have everything we need, lets list them out

Sign = **0**, because positive

Exponent Bias Binary = **1000 0100**

Mantisa/Fraction = **100100111000000000000000**

Combine to get Single Precision IEEE 754 Floating Point Representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

<b>0100 0010 0100 1001 1100 0000 0000 0000</b>
--

Get the hexadecimal value by converting each set of 4 binaries to respective value

<b>32-bit Hexadecimal – 0x4249c000</b>
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**0.0**

Sign = **0**, because positive

Exponent = 0, Bias Binary = **0000 0000**

Mantisa/Fraction = **000000000000000000000000**

Combine to get Single Precision IEEE 754 Floating Point Representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

<b>0000 0000 0000 0000 0000 0000 0000 0000</b>
--

Get the hexadecimal value by converting each set of 4 binaries to respective value

<b>32-bit Hexadecimal – 0x00000000</b>
--



## -Infinity

Sign = **1**, because Negative

Exponent = 128, Bias(128 + 127 = **255**) Binary = **1111 1111**

Mantisa/Fraction = **000000000000000000000000**

Combine to get Single Precision IEEE 754 Floating Point Representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

<b>1111 1111 1000 0000 0000 0000 0000 0000</b>
--

Get the hexadecimal value by converting each set of 4 binaries to respective value

<b>32-bit Hexadecimal – 1xff800000</b>
--

## 1.0000001

1.0000001

$1 = 0 * 2 + 1$

We get **1**, which is 1 in binary representation

1.0000001

$.0000001 * 2 = 0.0000002$

$.0000002 * 2 = 0.0000004$

$.0000004 * 2 = 0.0000008$

$.0000008 * 2 = 0.0000016$

$.0000016 * 2 = 0.0000032$

$.0000032 * 2 = 0.0000064$

$.0000064 * 2 = 0.0000128$

$.0000128 * 2 = 0.0000256$

$.0000256 * 2 = 0.0000512$

$.0000512 * 2 = 0.0001024$

$.0001024 * 2 = 0.0002048$

$.0002048 * 2 = 0.0004096$

$.0004096 * 2 = 0.0008192$

$.0008192 * 2 = 0.0016384$

$.0016384 * 2 = 0.0032768$

$.0032768 * 2 = 0.0065536$

$.0065536 * 2 = 0.0131072$

$.0131072 * 2 = 0.0262144$

$.0262144 * 2 = 0.0524288$

$.0524288 * 2 = 0.1048576$

$.1048576 * 2 = 0.2097152$

$.2097152 * 2 = 0.4194304$

$.4194304 * 2 = 0.8388608$

$.8388608 * 2 = 1.6777216$

We get **000000000000000000000001**, which is 0.0000001 in binary representation

Combining, we get **- 1.000000000000000000000001**

Convert to scientific notation, we get **- 1.000000000000000000000001** \*  $2^0$

Where **000000000000000000000001** is the mantisa/fraction

And **0** is the exponent

Convert the exponent to the bias(127 + 0 = **127**) in binary, we get **0111 1111**

We have everything we need, lets list them out

Sign = **0**, because positive

Exponent Bias Binary = **0111 1111**

Mantisa/Fraction = **000000000000000000000001**

Combine to get Single Precision IEEE 754 Floating Point representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

**0011 1111 1000 0000 0000 0000 0000 0001**

Get the hexadecimal value by converting each set of 4 binaries to respective value

**32-bit Hexadecimal – 0x3f800001**

## 9

### 0xc349a000

IEEE 7354 Representation – **1100 0011 0100 1001 1010 0000 0000 0000**

Sign = 1

Exponent Bias Binary = **1000 0110**

Mantisa/Fraction = **100100110100000000000000**

Exponent Bias Decimal Value = 134

Exponent = 134 - 127 = **7**

Fraction Value From Mantisa =  $2^{-1} + 2^{-4} + 2^{-7} + 2^{-8} + 2^{-10} = \frac{589}{1024}$

Equation to plug into  $-1^{Sign} * (1 + (FractionValue)) * 2^{Exponent}$

Plug in, we get

$-1^1 * (1 + \frac{589}{1024}) * 2^7 = -(\frac{1613}{1024}) * 2^7 = -201.625$

**-201.625**

### 0xffe00001

IEEE 7354 Representation – **1111 1111 1110 0000 0000 0000 0000 0001**

Sign = 1

Exponent Bias Binary = **1111 1111**

Mantisa/Fraction = **110000000000000000000001**

Because the exponent field is all 1s and the fraction field is non-zero, we can conclude what we are trying to find is **NaN**

**NaN(Not a Number)**

### 0x80000000

IEEE 7354 Representation – **1000 0000 0000 0000 0000 0000 0000 0000**

Sign = 1

Exponent Bias Binary = **0000 0000**

Mantisa/Fraction = **000000000000000000000000**

Both Fraction Value and Exponent Value is 0, so Decimal Value = -0, negative because sign is 1

**-0**

**0x00400000**

IEEE 7354 Representation – **0000 0000 0100 0000 0000 0000 0000 0000**

Sign = 0

Exponent Bias Binary = **0000 0000**

Mantisa/Fraction = **100000000000000000000000**

Exponent Bias Decimal Value = 0

Fraction Value From Mantisa =  $2^{-1} = \frac{1}{2}$

Because Fraction Value is non-zero and the Exponent Value is 0, the number is **subnormal**

We can then use this formula  $-1^{Sign} * (FractionValue) * 2^{-126}$

Plug in, we get

$$-1^0 * (\frac{1}{2}) * 2^{-126} = \frac{1}{2} * 2^{-126} = \frac{1}{2^{127}}$$

$$\boxed{\frac{1}{2^{127}}}$$