## Homework 3

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# Boolean Algebra

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(a) 
$$A*(\overline{A}+B*B)+\overline{(B*A)}*(\overline{A}+B)$$
 DeMorgan's  $A*(\overline{A}+B*B)+(\overline{B}*\overline{A})*(\overline{A}+B)$  DeMorgan's  $A*(\overline{A}+B)+(\overline{B}*\overline{A})*(\overline{A}+B)$  Idempotent Law  $(B*B=B)$   $(A+(\overline{B}*\overline{A}))(\overline{A}+B)$  Distributive Law (Inversed)  $(A+(\overline{A}*\overline{B}))(\overline{A}+B)$  Commutative Law  $(A+\overline{B})(\overline{A}+B)$  Redundacy Law  $(A+(\overline{A}*\overline{B}))=A+\overline{B}$  A $(\overline{A}+B)+\overline{B}(\overline{A}+B)$  Distributive Law  $(A*\overline{A}+A*B)+\overline{B}(\overline{A}+B)$  Distributive Law  $(A*\overline{A}+A*B)+\overline{B}(\overline{A}+B)$  Inverse Law  $(A*\overline{A}=0)$   $(A*B)+\overline{B}(\overline{A}+B)$  Identity Law  $(0+A*B=A*B)$   $(A*B)+(\overline{B}*\overline{A}+\overline{B}*B)$  Distributive Law  $(A*B)+(\overline{B}*\overline{A}+\overline{B}*B)$  Distributive Law  $(A*B)+(\overline{B}*\overline{A}+\overline{B}*A)$  Identity Law  $(\overline{B}*\overline{A}+0)=\overline{B}*\overline{A}$  (b)  $\overline{C*B}+(A*B*C)+\overline{A+C+\overline{B}}$  DeMorgan's

$$\overline{C} * \overline{B} + (A * B * C) + \overline{A + C + B}$$

$$\overline{C} + \overline{B} + (A * B * C) + \overline{A + C + B}$$

$$\overline{C} + (A * B * C) + \overline{B} + \overline{A + C + B}$$

$$\overline{C} + (A * B * C) + \overline{B} + \overline{A + C + B}$$

$$\overline{C} + (C * A * B) + \overline{B} + \overline{A + C + B}$$

$$\overline{C} + (A * B) + \overline{B} + \overline{A + C + B}$$

$$\overline{C} + (A * B) + \overline{A + C + B}$$

$$\overline{C} + \overline{B} + (A * B) + \overline{A + C + B}$$

$$\overline{C} + \overline{B} + (A * B) + \overline{A + C + B}$$

$$\overline{C} + \overline{B} + (B * A) + \overline{A + C + B}$$

$$\overline{C} + (\overline{B} + A) + \overline{A} + \overline{C} + \overline{B}$$

$$\overline{C} + (\overline{B} + A) + \overline{A} + \overline{C} + \overline{B}$$

$$\overline{C} + (\overline{B} + A) + \overline{A} * \overline{C} * \overline{B}$$

$$\overline{C} + (\overline{B} + A) + \overline{A} * \overline{C} * \overline{B}$$

$$\overline{C} + (\overline{B} + A) + \overline{A} * \overline{C} * B$$

$$\overline{C} + (\overline{A} * \overline{C} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{A} * \overline{C} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{B} + A)$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C}$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C}$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C} * \overline{C}$$

$$\overline{C} + (\overline{C} * \overline{A} * B) + (\overline{C} * \overline{C} * \overline{C}$$

```
(c) (A+B)*(\overline{A}+C)*(\overline{C}+B)
    (A(\overline{A}+C)+B(\overline{A}+C))*(\overline{C}+B) Distributive Law
    ((A*\overline{A}+A*C)+B(\overline{A}+C))*(\overline{C}+B) Distributive Law
    (A*\overline{A} + A*C + B*\overline{A} + B*C)*(\overline{C} + B) Distributive Law
    ((A*\overline{A}*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}*B) + (B*C*B))
                                                                                 Distributive Law
    ((A*\overline{A}*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}*B) + (B*B*C))
                                                                                 Commutative Law
    ((A*\overline{A}*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B)+(A*C*B)+(B*\overline{A}*B)+(B*C)) Idempotent Law
    ((A*\overline{A}*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B)+(A*C*B)+(B*B*\overline{A})+(B*C)) Commutative Law
    ((A*\overline{A}*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}) + (B*C))
                                                                      Idempotent Law
    ((0*\overline{C}) + (A*C*\overline{C}) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}) + (B*C))
                                                                       Inverse Law
    ((0*\overline{C}) + (A*0) + (B*\overline{A}*\overline{C}) + (B*C*\overline{C})) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}) + (B*C))
                                                                       Inverse Law
    ((0*\overline{C}) + (A*0) + (B*\overline{A}*\overline{C}) + (B*0)) +
    ((A*\overline{A}*B) + (A*C*B) + (B*\overline{A}) + (B*C))
                                                                       Inverse Law
    ((0 * \overline{C}) + (A * 0) + (B * \overline{A} * \overline{C}) + (B * 0)) +
    ((0*B) + (A*C*B) + (B*\overline{A}) + (B*C))
                                                               Inverse Law
    ((0) + (0) + (B * \overline{A} * \overline{C}) + (0)) +
    ((0) + (A * C * B) + (B * \overline{A}) + (B * \overline{C})) Law of Zeros (Too much to do them one by one ...)
    (B*\overline{A}*\overline{C}) + (A*C*B) + (B*\overline{A}) + (B*C)
                                                                   Identity Law (Too much to do them one by one...)
    (B*\overline{A}) + (B*\overline{A}*\overline{C}) + (A*C*B) + (B*C) Commutative Law
    (B*\overline{A}) + (A*C*B) + (B*C) Absorption Law ((B*\overline{A}) + (B*\overline{A}*\overline{C}) = (B*\overline{A}))
    (B*\overline{A}) + (B*C) + (A*C*B) Commutative Law
    (B*\overline{A}) + (B*C) + (B*C*A)
                                                  Commutative Law
    (B*\overline{A}) + (B*C) Absorption Law ((B*C) + (B*C*A) = (B*C))
    B*(\overline{A}+C) Distributive Law (Inversed)
     B*(\overline{A}+C)
```

(a) 
$$(\overline{A}+C)*(\overline{B}+D+A)*(D+A*\overline{C})*(\overline{D}+A)=1$$

$$(\overline{A}+C)*(\overline{B}+D+A)*(D+\overline{A}+\overline{C})*(\overline{D}+A)=1$$

$$(\overline{A}+C)*(\overline{B}+D+A)*(D+\overline{A}+\overline{C})*(\overline{D}+A)=0$$
DeMorgan's
$$(\overline{A}+C)*(\overline{B}+D+A)*(D+\overline{A}+C)*(\overline{D}+A)=0$$
Double Negation Law
$$(\overline{A}+C)*(D+\overline{A}+C)*(\overline{B}+D+A)*(\overline{D}+A)=0$$
Commutative Law
$$(\overline{A}+C)*(\overline{A}+C+D)*(\overline{B}+D+A)*(\overline{D}+A)=0$$
Commutative Law
$$((\overline{A}+C)*D)*(\overline{B}+D+A)*(\overline{D}+A)=0$$
Commutative Law
$$(\overline{A}+C)*D*(D+A+\overline{B})*(\overline{D}+A)=0$$
Commutative Law
$$(\overline{A}+C)*D*(\overline{D}+A)=0$$
Absorption Law  $(D*(D+A+\overline{B})=D)=0$ 

$$(\overline{A}+C)*D*(D+A)=0$$
Redundacy Law  $(D*(D+A)=D*A)=0$ 

$$A*(\overline{A}+C)*D=0$$
Commutative Law
$$A*C*D=0$$
Redundacy Law  $(A*(\overline{A}+C)=A*C)=0$ 

$$A*C*D=1$$

$$(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)=0$$
Commutative Law
$$(A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D=0$$
Commutative Law
$$(A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D=0$$
Commutative Law
$$(A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D*(D+A)=0$$
Commutative Law
$$(A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D*(\overline{A}+C)=0$$

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$$A*(\overline{A}+C)*D*(D+A)=0$$

$$A*(\overline{A}+C)*D*(D+A)=$$

$$((\overline{K}*L*N) + (K*\overline{L}*N*\overline{M}))*(\overline{K}+\overline{N}) \quad \text{Redundacy Law } (\overline{L}*(L+N) = L*N) \\ N*((\overline{K}*L) + (K*\overline{L}*\overline{M}))*(\overline{K}+\overline{N}) \quad \text{Distributive Law (Inversed)} \\ N*(\overline{K}+\overline{N})*((\overline{K}*L) + (K*\overline{L}*\overline{M})) \quad \text{Commutative Law} \\ N*(\overline{N}+\overline{K})*((\overline{K}*L) + (K*\overline{L}*\overline{M})) \quad \text{Commutative Law} \\ N*\overline{K}*((\overline{K}*L) + (K*\overline{L}*\overline{M})) \quad \text{Redundacy Law } (N*(\overline{N}+\overline{K}) = N*\overline{K}) \\ (\overline{K}*L*N*\overline{K}) + (K*\overline{L}*\overline{M}*N*\overline{K}) \quad \text{Distributive Law} \\ (\overline{K}*\overline{K}*L*N) + (K*\overline{L}*\overline{M}*N*\overline{K}) \quad \text{Commutative Law} \\ (\overline{K}*L*N) + (K*\overline{L}*\overline{M}*N*\overline{K}) \quad \text{Idempotent Law} \\ (\overline{K}*L*N) + (K*\overline{K}*\overline{L}*\overline{M}*N) \quad \text{Commutative Law} \\ (\overline{K}*L*N) + (O*\overline{L}*\overline{M}*N) \quad \text{Inverse Law}$$

$$(\overline{K}*L*N) + (0)$$
 Law of Zeros

 $(\overline{K} * L * N)$  Identity Law

$$\overline{K} * L * N = 1$$

# K = 0, L = 1, N = 1, M = 0 or 1

Truth Table

X	Y	Z	Q
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

#### K-Map

$\mathbf{z}$	XY			
"	00	01	11	10
0	0	1	1	1
1	1	0	0	0

$\mathbf{z}$	XY			
	00	01	11	10
0	0	1	1	1
1	1	0	0	0

The merge shown by the yellow means "if Z is 0 and Y is 1, it doesn't matter what the value of X is." In the function, it expresses  $-X*Y*\overline{Z}+\overline{X}*Y*\overline{Z}=Y*\overline{Z}$ 

	$\mathbf{z}$	XY			
		00	01	11	10
	0	0	1	1	1
	1	1	0	0	0

The merge shown by the yellow means "if Z is 0 and X is 1, it doesn't matter what the value of Y is." In the function, it expresses  $-X*Y*\overline{Z}+X*\overline{Y}*\overline{Z}=X*\overline{Z}$ 

After applying those two reductions, we get

Original : 
$$\overline{X} * \overline{Y} * Z + X * Y * \overline{Z} + \overline{X} * Y * \overline{Z} + X * \overline{Y} * \overline{Z}$$

Reduced:

$$\boxed{Y*\overline{Z} + X*\overline{Z} + \overline{X}*\overline{Y}*Z}$$

# Logical Circuits

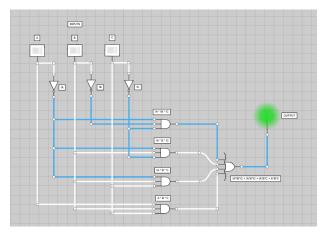
4

To obtain the sums of products representation, we look for the rows where the output is 1 and then form a product term for each row.

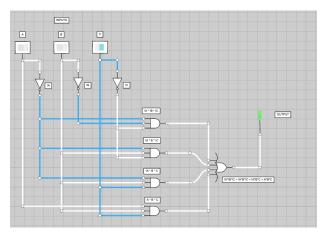
From the truth table, we get

$$(\overline{A}*\overline{B}*\overline{C}) + (\overline{A}*B*\overline{C}) + (\overline{A}*B*C) + (A*B*C)$$

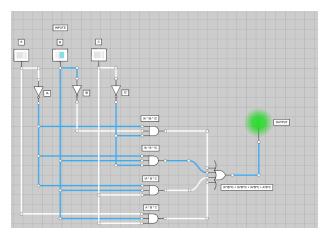
6



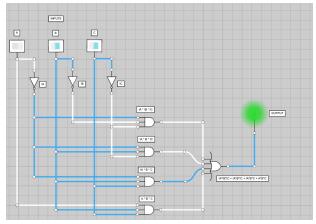
When input A = 0, B = 0, C = 0, the output is 1, like in the truth table (Also satisfies expression  $(\overline{A} * \overline{B} * \overline{C})$  as it should)



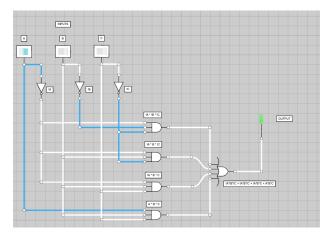
When input A = 0, B = 0, C = 1, the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



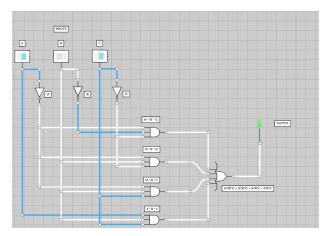
When input A = 0, B = 1, C = 0, the output is 1, like in the truth table (Also satisfies expression  $(\overline{A}*B*\overline{C})$ )



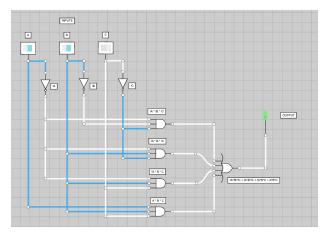
When input  $A=0,\,B=1,\,C=1,$  the output is 1, like in the truth table (Also satisfies expression  $(\overline{A}*B*C)$ )



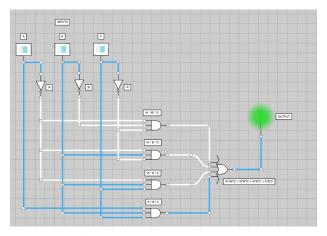
When input A = 1, B = 0, C = 0, the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input A = 1, B = 0, C = 1, the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input A = 1, B = 1, C = 0, the output is 0, like in the truth table (Doesn't satisfy none of the expressions)



When input A = 1, B = 1, C = 1, the output is 1, like in the truth table (Also satisfies expression (A \* B \* C))

After going through all possible combination of inputs, I find that the expected output with each different set of inputs is exactly like it is in the truth table. When the output is 1, the appropriate AND components also lights up, matching the appropriate expression in the boolean function (i.e., (A\*B\*C) is an expression that lights up when A=1, B=1, C=1). Through these findings, I believe there is sufficient proof that my logical circuit does in fact implement the required Boolean function, and that the boolean function also correctly represents the truth table.

#### **Numerical Conversions and Arithmetic**

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#### 50.4375

```
50.4375
50 = 25 * 2 + 0
25 = 12 * 2 + 1
12 = 6 * 2 + 0
6 = 3 * 2 + 0
3 = 1 * 2 + 1
1 = 0 * 2 + 1
```

We get 110010, which is 50 in binary representation

#### 50.4375

.4375 \* 2 = 0.875 .0875 \* 2 = 1.75 .75 \* 2 = 1.5.5 \* 2 = 1.0

We get **0111**, which is .4375 in binary representation

Combining, we get -110010.0111

Convert to scientific notation, we get  $-1.100100111 * 2^{5}$ 

Where 100100111 is the mantisa

And **5** is the exponent

Convert the exponent to the bias(127 + 5 = 132) in binary, we get 1000 0100

We have everything we need, lets list them out

Sign = 0, because positive

Exponent Bias Binary = 10000100

 ${\rm Mantisa/Fraction} = 10010011100000000000000$ 

Combine to get Single Precision IEEE 754 Floating Point Representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

#### 0100 0010 0100 1001 1100 0000 0000 0000

Get the hexidecimal value by converting each set of 4 binaries to respective value

#### 32-bit Hexadecimal -0x4249c000

#### 0.0

```
Sign = 0, because positive
```

Exponent = 0, Bias Binary = 00000000

Combine to get Single Precision IEEE 754 Floating Point Representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

### $0000\ 0000\ 0000\ 0000\ 0000\ 0000\ \overline{0000}$

Get the hexidecimal value by converting each set of 4 binaries to respective value

#### 32-bit Hexadecimal -0x00000000

#### -Infinity

```
Sign = 1, because Negative
Exponent = 128, Bias(128 + 127 = 255) Binary = 1111 1111
Combine to get Single Precision IEEE 754 Floating Point Representation
Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits
1111 1111 1000 0000 0000 0000 0000 0000
Get the hexidecimal value by converting each set of 4 binaries to respective value
```

# 32-bit Hexadecimal – 1xff800000

```
1.0000001
   1.0000001
   1 = 0 * 2 + 1
   We get 1, which is 1 in binary representation
   1.0000001
   .0000001 * 2 = \mathbf{0}.0000002
   .0000002 * 2 = 0.0000004
   .0000004 * 2 = 0.0000008
   .0000008 * 2 = \mathbf{0}.0000016
   .0000016 * 2 = \mathbf{0}.0000032
   .0000032 * 2 = 0.0000064
   .0000064 * 2 = \mathbf{0}.0000128
   .0000128 * 2 = 0.0000256
   .0000256 * 2 = \mathbf{0}.0000512
   .0000512 * 2 = 0.0001024
   .0001024 * 2 = \mathbf{0}.0002048
   .0002048 * 2 = \mathbf{0}.0004096
   .0004096 * 2 = \mathbf{0}.0008192
   .0008192 * 2 = 0.0016384
   .0016384 * 2 = \mathbf{0}.0032768
   .0032768 * 2 = \mathbf{0}.0065536
   .0065536 * 2 = \mathbf{0}.0131072
   .0131072 * 2 = 0.0262144
   .0262144 * 2 = \mathbf{0}.0524288
   .0524288 * 2 = 0.1048576
   .1048576 * 2 = \mathbf{0}.2097152
   .2097152 * 2 = 0.4194304
   .4194304 * 2 = \mathbf{0}.8388608
   .8388608 * 2 = 1.6777216
   And 0 is the exponent
```

Convert the exponent to the bias (127 + 0 = 127) in binary, we get **0111 1111** 

We have everything we need, lets list them out

Sign = 0, because positive

Exponent Bias Binary = 01111111

Combine to get Single Precision IEEE 754 Floating Point representation

Where Sign = 1 bit, Exponent = 8 bits, Mantisa/Fraction = 23 bits

#### $0011\ 1111\ 1000\ 0000\ 0000\ 0000\ 0000\ 0001$

Get the hexidecimal value by converting each set of 4 binaries to respective value

#### 32-bit Hexadecimal -0x3f800001

9

#### 0xc349a000

Sign = 1

Exponent Bias Binary = 10000110

Exponent Bias Decimal Value = 134

Exponent = 134 - 127 = 7

Fraction Value From Mantisa =  $2^{-1} + 2^{-4} + 2^{-7} + 2^{-8} + 2^{-10} = \frac{589}{1024}$ 

Equation to plug into  $-1^{Sign} * (1 + (FractionValue)) * 2^{Exponent}$ 

Plug in, we get

$$-1^{1} * (1 + \frac{589}{1024}) * 2^{7} = -(\frac{1613}{1024}) * 2^{7} = -201.625$$

-201.625

#### 0xffe00001

Sign = 1

Exponent Bias Binary = 11111111

Because the exponent field is all 1s and the fraction field is non-zero, we can conclude what we are trying to find is NaN

# $NaN(Not\ a\ Number)$

#### 0x80000000

IEEE 7354 Representation – 1000 0000 0000 0000 0000 0000 0000

Sign = 1

Exponent Bias Binary = 00000000

Both Fraction Value and Exponent Value is 0, so Decimal Value = -0, negative because sign is 1

-0

#### 0x00400000

Exponent Bias Binary = 00000000

 $\operatorname{Mantisa/Fraction} = 10000000000000000000000$ 

Exponent Bias Decimal Value = 0

Fraction Value From Mantisa =  $2^{-1} = \frac{1}{2}$ 

Because Fraction Value is non-zero and the Exponent Value is 0, the number is  $\mathbf{subnormal}$ 

We can then use this formula  $-1^{Sign} * (FractionValue) * 2^{-126}$ 

Plug in, we get

$$-1^{0} * (\frac{1}{2}) * 2^{-126} = \frac{1}{2} * 2^{-126} = \frac{1}{2^{127}}$$
 
$$\frac{1}{2^{127}}$$

$$\frac{1}{2^{127}}$$