

CSCI 2200 — Foundations of Computer Science (FoCS)

Homework 2 (document version 1.0)

Overview

- This homework is due by 11:59PM on Thursday, September 28
- You may work on this homework in a group of up to four students; unlike recitation problem sets, **your teammates may be in any section**
- You may use at most **two** late days on this assignment
- Please start this homework early and ask questions during office hours; also ask questions on the Discussion Forum
- Please be concise in your written answers; even if your solution is correct, if it is not well-presented, you may still lose points
- You can type or hand-write (or both) your solutions to the required graded problems below; **all work must be organized in one PDF that lists all teammate names**
- You are strongly encouraged to use LaTeX, in particular for mathematical symbols; see the corresponding `hw1.tex` file as a starting point and example

Grading

- For each assigned problem, a grade of 0, 1, or 2 is assigned as follows: 0 indicates no credit; 1 indicates half credit; and 2 indicates full credit
- No credit is assigned if a problem is not attempted or minimal work/progress is shown
- Half credit is assigned if a strong attempt was made toward a solution and/or only part of the problem was attempted or solved
- Full credit is assigned for a perfect or nearly perfect solution, i.e., only one or two minor typos/mistakes at most

Warm-up exercises

The problems below are good practice problems to work on. Do not submit these as part of your homework submission. **These are ungraded problems.**

- Problem 3.32.
- Problem 3.49.
- Problem 4.11.
- Problem 4.14.
- Problem 4.16.
- Problem 4.47.
- Problem 4.48(c). (See Problem 4.47.)
- Problem 5.10.
- Problem 6.3(a).
- Problem 6.15.
- Problem 6.32.
- Problem 7.3.
- Problem 7.4(a-c).
(In each subproblem, remove the recursion in your formula for A_n .)

Graded problems

The problems below are required and will be graded.

- *Problem 3.59 (Closure).
- *Problem 4.10(k-l).
- *Problem 5.12(d).
- *Problem 5.20.
- *Problem 5.39.
- *Problem 6.8.
- *Problem 6.43.

As you might not have the required textbook yet, all of the above problems (both graded and ungraded) are transcribed in the pages that follow.

Graded problems are noted with an asterisk (*).

If any typos exist below, please use the textbook description.

- **Problem 3.32.** Use truth tables to verify the rules for derivations in Figure 3.1 on page 29. Now use the rules in Figure 3.1 to show logical equivalence

$$\neg((p \wedge q) \vee r) \stackrel{\text{eqv}}{\equiv} (\neg p \wedge \neg r) \vee (\neg q \wedge \neg r).$$

- **Problem 3.49.** What is the difference between

$$\forall x : (\neg \exists y : P(x) \rightarrow Q(y)) \quad \text{and} \quad \neg \exists y : (\forall x : P(x) \rightarrow Q(y))?$$

- ***Problem 3.59 (Closure).** A set \mathcal{S} is closed under an operation if performing that operation on elements of \mathcal{S} returns an element in \mathcal{S} . Here are five examples of closure.

\mathcal{S} is closed under addition $\rightarrow \forall (x, y) \in \mathcal{S}^2 : x + y \in \mathcal{S}$.

\mathcal{S} is closed under subtraction $\rightarrow \forall (x, y) \in \mathcal{S}^2 : x - y \in \mathcal{S}$.

\mathcal{S} is closed under multiplication $\rightarrow \forall (x, y) \in \mathcal{S}^2 : xy \in \mathcal{S}$.

\mathcal{S} is closed under division $\rightarrow \forall (x, y \neq 0) \in \mathcal{S}^2 : x/y \in \mathcal{S}$.

\mathcal{S} is closed under exponentiation $\rightarrow \forall (x, y) \in \mathcal{S}^2 : x^y \in \mathcal{S}$.

Which of the five operations are the following sets closed under? (a) \mathbb{N} . (b) \mathbb{Z} . (c) \mathbb{Q} . (d) \mathbb{R} .

(a) Natural Numbers

- Closed Under Addition - Sum of two natural numbers will always be another natural number
- Not Closed Under Subtraction - Difference of two natural numbers will not always be another natural number (i.e., $9 - 12 = -3$)
- Closed Under Multiplication - Product of two natural numbers will always be another natural number
- Not Closed Under Division - Quotient of two natural numbers will not always be another natural number (i.e., $9/2 = 4.5$)
- Closed Under Exponentiation - A natural number to the power of another natural number will always result in a positive integer, making it natural.

(b) Integers

- Closed Under Addition - Sum of two integers will always be another integer
- Closed Under Subtraction - Difference of two natural numbers will always be another integer
- Closed Under Multiplication - Product of two integers will always be another integer
- Not Closed Under Division - Quotient of two integers will not always be another integer (i.e., $9/2 = 4.5$)
- Not Closed Under Exponentiation - A integer to the power of another integer will not always result in another integer (i.e., $2^{-3} = 0.125$)

(c) Rational Numbers

- Closed Under Addition - Sum of two rational numbers will always be another rational number
- Closed Under Subtraction - Difference of two rational numbers will always be another rational number
- Closed Under Multiplication - Product of two rational numbers will always be another rational number
- Closed Under Division - Quotient of two rational numbers will always be another rational number
- Not Closed Under Exponentiation - A rational number to the power of another rational number will not always result in another rational number (i.e., $(2)^{\frac{1}{2}} = \sqrt{2}$)

(d) Real Numbers

- Closed Under Addition - Sum of two real numbers will always be another real number
- Closed Under Subtraction - Difference of two real numbers will always be another real number
- Closed Under Multiplication - Product of two real numbers will always be another real number
- Closed Under Division - Quotient of two real numbers will always be another real number
- Closed Under Exponentiation - A real number to the power of another real number will always result in another real number

- ***Problem 4.10(k-l).** You may assume n is an integer. Prove by contraposition (explicitly state the contrapositive).

(k) $3 \text{ divides } n - 2 \rightarrow n \text{ is not a perfect square.}$

(l) If $p > 2$ is prime, then $p^2 + 1$ is composite.

(k) P(x) : (p)3 divides $n - 2 \rightarrow n$ (q)is not a perfect square.

Claim : If 3 divides $n - 2$, then n is not a perfect square

Contrapositive: Assume that n is a perfect square, then 3 does not divide $n - 2$

If n is a perfect square, then $n = k^2$ for some $k \in \mathbb{Z}$.

Now consider $n - 2$ in $(3 \text{ divides } n - 2)$, plug in k^2 for n

$$n - 2 = k^2 - 2$$

Now we want to show that 3 does not divide $k^2 - 2$. In order for k^2 to be divisible by 3, $(k^2 - 2) \% 3$ has to equal 0. Consider 3 cases, $k=0$, $k=1$, and $k=2$.

$k = 0$, $k^2 = 0$, $k^2 - 2 = -2$, $-2 \% 3 = 1$, remainder of 1 when divided by 3, not divisible by 3

$k = 1$, $k^2 = 1$, $k^2 - 2 = -1$, $-1 \% 3 = 2$, remainder of 2 when divided by 3, not divisibly by 3

$k = 2$, $k^2 = 4$, $k^2 - 2 = 2$, $2 \% 3 = 2$, remainder of 2 when divided by 3, not divisible by 3

In all cases, we see that $(k^2 - 2) \% 3$ does not equal to 0, none of these cases are divisible by 3
 We have proven that when n is a perfect square, then 3 does not divide $n - 2$, making 3 divides $n - 2$ false

Thus we have shown that p is false and $P(x)$ is true.

(l) $P(x)$: (p) If $p > 2$ is prime, (q) then $p^2 + 1$ is composite

Contrapositive: Assume that $p^2 + 1$ is not composite, then $p > 2$ is not prime or $p \leq 2$

When $p^2 + 1$ is not composite, it is either prime or a negative number, let's consider both cases

- **Problem 4.11.** For $x, y \in \mathbb{N}$, which statements below are contradictions (cannot possibly be true). Explain.

- (a) $x^2 < y$.
- (b) $x^2 = y/2$.
- (c) $x^2 - y^2 \leq 1$.
- (d) $x^2 + y^2 \leq 1$.
- (e) $2x + 1 = y^2 + 5y$.
- (f) $x^2 - y^2/2 = 1$.
- (g) $x^2 - y^2 = 1$.

- **Problem 4.14.** Prove: If $a, b, c \in \mathbb{Z}$ are odd, then for all $x \in \mathbb{Q}$, $ax^2 + bx + c \neq 0$. (Contradiction in a direct proof.)

- **Problem 4.47 (Without Loss of Generality (wlog)).** Consider the following claim.

If x and y have opposite parity (one is odd and one is even), then $x + y$ is odd.

Explain why, in a direct proof, we may assume that x is odd and y is even? Prove the claim. (Such a proof starts “Without loss of generality, assume x is odd and y is even. Then, ...”)

- **Problem 4.48(c).** Use the concept of “without loss of generality” to prove these claims.

(c) For any non-zero real number x , $x^2 + 1/x^2 \geq 2$.

- ***Problem 5.12(d).** For $n \geq 1$, prove by induction:

(d) $3^n > n^2$.

(d) **P(n) :** $3^n > n^2$

We use induction prove $\forall(n \geq 1) : P(n)$

[Base Case] $P(1) = 3^1 > 1^2 = 3 > 1$ which is True

[induction Step] Prove $P(n) \rightarrow P(n+1)$ for all $n \geq 1$

Specifically, if $3^n > n^2$, then $3^{n+1} > (n+1)^2$

Now we want to prove $3^{n+1} > (n+1)^2$

By Induction Assumption, we know that $3^n > n^2$

Which allows us to multiply both sides by 3, giving us $3^{n+1} > 3n^2$, looking at the LHS, we can substitute it in our next inequality

Simplify $(n+1)^2$ gives us $n^2 + 2n + 1$

Substitute everything into this next inequality so we can compare, we get

$$3^{n+1} > 3n^2 > n^2 + 2n + 1$$

We know that $3k^2$ is greater than the highest term in the most RHS's quadratic n^2 and that $3^{n+1} > 3n^2$ due to the induction hypothesis.

We can rewrite the inequality as $3^{n+1} > n^2 + 2n + 1$ which is equal/the same to what we wanted to prove, $3^{n+1} > (n+1)^2$, up in the 4th line

Thus, by induction, we have proved $P(n+1)$ to be true, making $P(n)$ to be true for all cases $n \geq 1$

- ***Problem 5.20.** Prove, by induction, that every $n \geq 1$ is a sum of distinct powers of 2.
- ***Problem 5.39.** Prove you can make any postage greater than 12¢ using only 4¢ and 5¢ stamps. (The USPS can set any postage above 12¢ and you don't have to buy any new stamps.)
- **Problem 6.3(a).** Strengthen the claim and prove by induction for $n \geq 1$:
 - (a) The sum of the first n odd numbers is a square. [Hint: Strengthen to a specific square.]

- ***Problem 6.8.** Prove $n^7 < 2^n$ for $n \geq 37$. (a) Use induction. (b) Use leaping induction.
- **Problem 6.15.** Prove that there are $2^{\lceil n/2 \rceil}$ distinct n -bit binary palindromes (strings that equal their reversal).
- **Problem 6.32.** We are back in L -tile land.
 - (a) This time, the potted plant needs more room than just one square. For $n \geq 1$, a $2^n \times 2^n$ grid-patio is missing a (large) 2×2 square in a corner as shown in the figure. Prove that the remainder of the patio can be L -tiled, for $n \geq 1$.
 - (b) We are no longer sure what the size of the potted plant is. The size may be $2^k \times 2^k$, and so a $2^k \times 2^k$ square will be missing from the corner of the $2^n \times 2^n$ grid-patio. Prove that the remainder of the patio can always be L -tiled, for $k \geq 1$ and $n \geq k$. [Hint: Tinker: try $k = 2; n = 3$ and $k = 2; n = 4$ to figure out what is going on.]

- ***Problem 6.43.** A sliding puzzle is a grid of 9 squares with 8 tiles. The goal is to get the 8 tiles into order (the target configuration). A move slides a tile into an empty square. Below, we show first a row move, then a column move.

Prove that no sequence of moves produces the target configuration. [Hint: The tiles form a sequence going left to right, top to bottom. An inversion is a pair that is out of order. Prove by induction that the number of inversions stays odd.]

- **Problem 7.3.** Give a recursive definition of the function $f(n) = n! \times 2^n$, where $n \geq 1$.

- **Problem 7.4(a-c).** Guess a formula for A_n and prove it by induction.

- $A_0 = 0$ and $A_n = A_{n-1} + 1$ for $n \geq 1$.
- $A_1 = 1$, $A_2 = 2$, and $A_n = A_{n-1} + 2A_{n-2}$ for $n \geq 2$.
- $A_0 = 1$; $A_1 = 2$; $A_n = 2A_{n-1} - A_{n-2} + 2$ for $n \geq 2$. [Hint: Method of differences.]