

Operating Systems (INFR10079) 2023/2024 Semester 2

Structure (Operating System Structure)

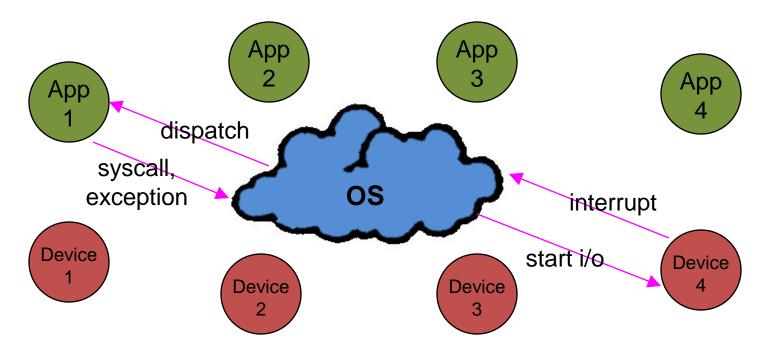
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Overview

- Architecture impact
- Application-Operating System interaction
- Operating System structure

OS Structure

- Application
 Operating System
 Hardware
- OS mediates access and abstracts away hardware diversity
 - Enables hardware sharing
- OS sits between applications and the hardware
 - Applications (App) request services
 - Explicitly via syscalls
 - Implicitly via exceptions
 - Devices (Device) request attention via interrupts



Operating System Design and Implementation

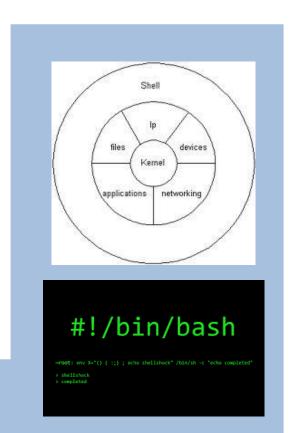
- Design and Implementation of OS not "solvable"
 - but some approaches have proven successful
- Internal structure of different Operating Systems can vary widely
- Start the design by defining goals and specifications
 - User goals: convenient to use, easy to learn, reliable, safe, and fast
 - System goals: easy to design, implement, and maintain, as well as flexible, reliable, error-free, and efficient
- Affected by choice of hardware, type of system

Operating System Design and Implementation

- Important principles to separate
- Policy: What will be done? (Algorithm)
- Mechanism: How to do it?
- Separation allows maximum flexibility
 - Policies are likely to change across places or over time
 - A general mechanism can support a wide range of policies
- Microkernel OSes (see later) are based on such principle
 - A core kernel implements the mechanisms
 - Policies are implemented outside the core kernel
 - Easily modifiable

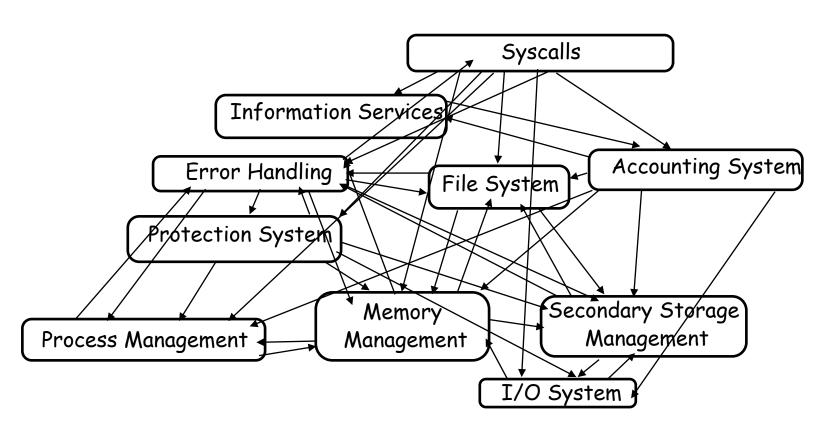
Major OS Services

- processes
- memory
- I/O
- secondary storage
- file systems
- protection
- networking
- shells (i.e., command interpreter)
- GUI
- etc. Systems programs outside the kernel



OS Structure #1

It's not always clear how to stitch OS services together

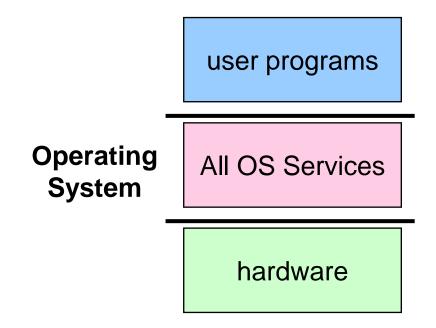


OS Structure #2

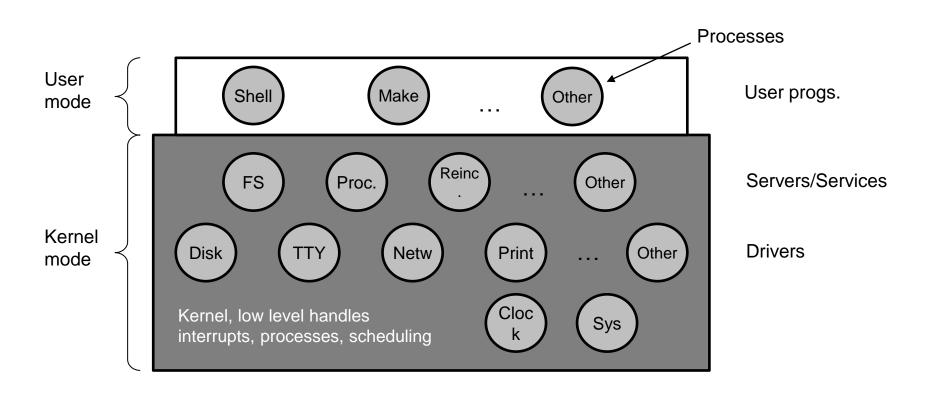
- Major issues
 - how do we organize all these?
 - what are all of the code modules, and where do they exist?
 - how do they cooperate?
- Massive software engineering and design problem
 - design a large, complex program that
 - performs well
 - is reliable
 - is extensible
 - is backwards compatible
 - etc.

Monolithic OS Design #1

- Likely the earliest OS organization
- UNIX was built as monolithic
 - Linux is built as monolithic



Monolithic Example: Linux



Monolithic OS Design #2

- Major advantage
 - cost of subsystems interactions is low (procedure call)
- Disadvantages
 - hard to understand
 - hard to modify
 - unreliable (no isolation between system modules)
 - hard to maintain
- What is the alternative?
 - find a way to organize OS subsystems to simplify its design and implementation

Layered OS Design

- Layering
 - implement OS as a set of layers
 - each layer presents an enhanced 'virtual machine' to the layer above
- The first description of this approach was Dijkstra's THE system
 - Layer 5: Job Managers
 - Execute users' programs
 - Layer 4: Device Managers
 - Handle devices and provide buffering
 - Layer 3: Console Manager
 - Implements virtual consoles
 - Layer 2: Page Manager
 - Implements virtual memories for each process
 - Layer 1: Kernel
 - Implements a virtual processor for each process
 - Layer 0: Hardware
- Each layer can be tested and verified independently



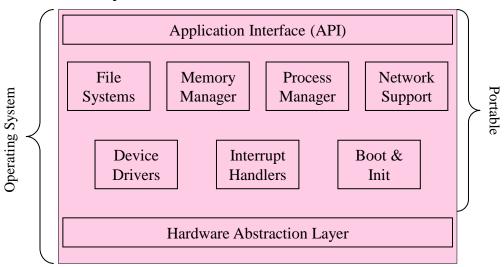
Problems with layering

- Imposes hierarchical structure
 - but real systems are more complex
 - File system requires virtual memory services
 - Virtual memory would like to use files for its backing store
 - strict layering isn't flexible enough
- Poor performance
 - each layer crossing has overhead associated with it
- Disjunction between model and reality
 - systems modeled as layers, but not really built that way



Hardware Abstraction Layer

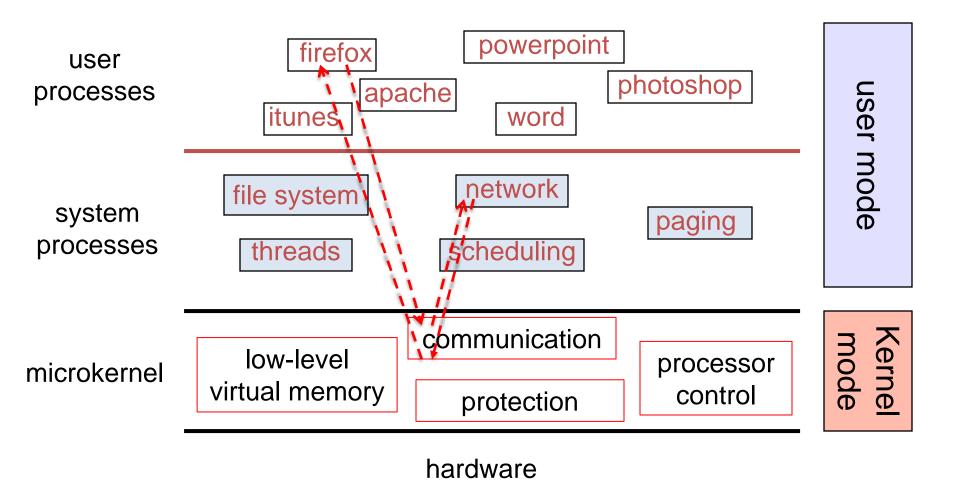
- An example of layering in modern operating systems
 - Windows, etc.
- Goal: separates hardware-specific routines from the core kernel of the OS
 - Provides portability
 - Improves readability



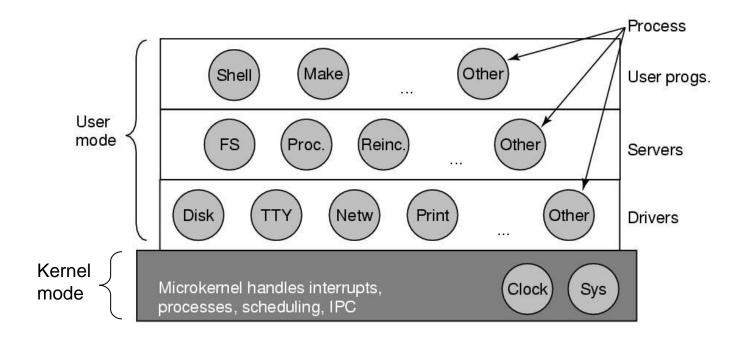
Microkernel OS Design

- Popular in the late 80's, early 90's
 - recent resurgence of popularity
- Goal
 - minimize what goes in kernel
 - organize rest of OS as user-level processes (services)
- This results in
 - better reliability (isolation between components)
 - ease of extension and customization
 - poor performance (user/kernel boundary crossings)
- First microkernel system was Hydra (CMU, 1970)
 - Follow-ons: Mach (CMU), Chorus (French UNIX-like OS), MINIX (UNIX-like OS from Amsterdam)

Microkernel Structure Illustrated

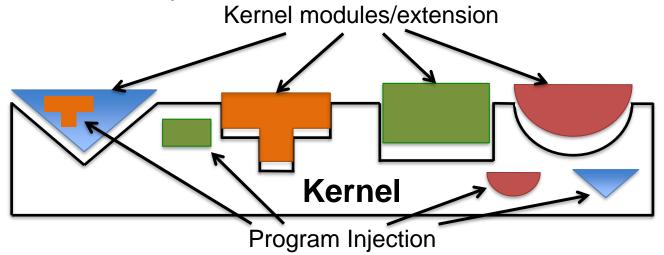


Microkernel Example: MINIX

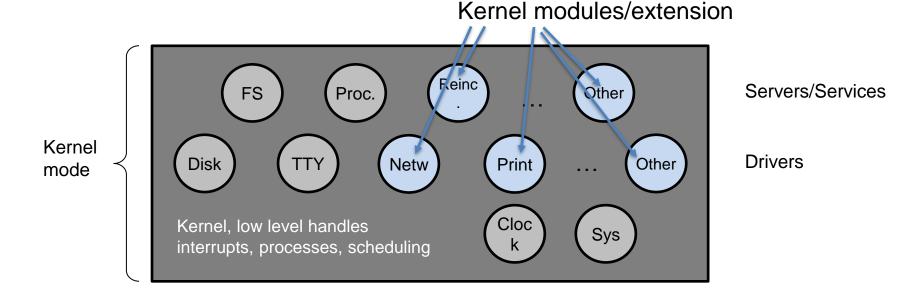


OS Kernel Functionality Extensibility

- Compile Time
 - Select what components to compile in
- Execution Time
 - Loadable Kernel Modules/Extensions
 - Microkernel: load the code in user space (any)
 - Monolithic: load the code in kernel space (Solaris, Linux, etc.)
 - Program Injection (event-based)
 - o eBPF/BPF
 - WebAssembly



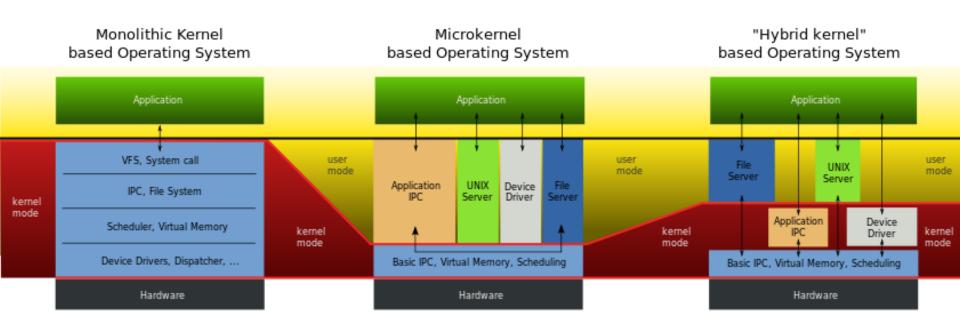
Loadable Kernel Modules (Monolithic)



- Core services in the kernel, others dynamically loaded
- Advantages for monolithic OSes
 - Convenient: no need for rebooting for newly added modules
 - Efficient: no need for message passing unlike microkernel
 - Flexible: any module can call any other module unlike layered model

Hybrid OS Design

- Many different approaches
 - Key idea: exploit the benefits of monolithic and microkernel designs
 - Windows, Xnu/Darwin, DragonFly BSD, ...
- Extensibility via kernel modules



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Summary

- Fundamental distinction between user and privileged modes supported by most hardware
- OS design has been an evolutionary process of trial and error
- Successful OS designs have run the spectrum from monolithic, to layered, to micro kernels
- The role and design of an OS are still evolving
- There is no "ideal" OS structure