

Homework #1

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Question 1

- a) The auction is a Bayesian game where each player $i \in \{1, \dots, N\}$ observes type $v_i \sim F(x)$ and chooses a continuous action $b_i \geq 0$ to maximize their expected payoff, $\mathbb{E}[u_i]$, where:

$$u_i(v_i, b_i, b_{-i}) = \begin{cases} v_i - b_i, & b_i > b_j \forall j \neq i \\ \frac{1}{k}(v_i - b_i), & b_i \in A = \max\{b_1, \dots, b_N\}, \text{ where } |A| = k \\ 0 & b_i < \max\{b_1, \dots, b_N\} \end{cases}$$

On any continuous distribution, the probability that any two draws are exactly equal is zero. Therefore, the knife's-edge case is excluded from the remainder of the analysis for brevity.

- b) Player i 's expected payoff is:

$$\mathbb{E}[u_i] = (v_i - b_i) \left[\prod_{j \neq i} \Pr(b_i > b_j) \right] + 0 \left[1 - \prod_{j \neq i} \Pr(b_i > b_j) \right] = (v_i - b_i) \left[\prod_{j \neq i} \Pr(b_i > b_j) \right]$$

Assume that all players bet according to a linear function of their valuation: $b_j = d + cv_j$. Then, player i 's expected payoff becomes:

$$\begin{aligned} \mathbb{E}[u_i] &= (v_i - b_i) \left[\prod_{j \neq i} \Pr(b_i > d + cv_j) \right] = (v_i - b_i) \left[\prod_{j \neq i} \Pr\left(v_j < \frac{b_i - d}{c}\right) \right] \\ &= (v_i - b_i) \left[\prod_{j \neq i} \left(\frac{b_i - d}{c}\right)^a \right] = (v_i - b_i) \left(\frac{b_i - d}{c}\right)^{Na} \end{aligned}$$

Player i 's optimal bet, then, can be derived from the first-order condition of their expected payoff:

$$\frac{v_i Na}{c} \left(\frac{b_i - d}{c}\right)^{Na-1} - \left(\frac{b_i - d}{c}\right)^{Na} - \frac{b_i Na}{c} \left(\frac{b_i - d}{c}\right)^{Na-1} = 0$$

$$Nav_i - b_i + d - b_i Na = 0$$

$$b_i = \left(\frac{Na}{1 + Na} \right) v_i + \frac{d}{Na + 1}$$

Assuming every player uses the same formula, then, we can solve:

$$c = \frac{Na}{Na + 1} d = \frac{d}{Na + 1} \Rightarrow d = 0$$

Thus, the Bayesian Nash equilibrium is for all players to bet:

$$b = \left(\frac{Na}{Na + 1} \right) v$$

- c)
- d) The limit of b as $a \rightarrow \infty$ is v , so the payoff of the winner converges to zero. Therefore, bidding becomes more competitive as a increases.
- e)

Question 2

- a)
- b)
- c)
- d)

Question 3

- a) The auction is a Bayesian game where each player $i \in \{1, 2, 3\}$ observes type $v_i \sim U[0, 1]$ and chooses a continuous action $b_i \geq 0$ to maximize their expected payoff, $\mathbb{E}[u_i]$, where, for $i \neq j \neq k$:

$$u_i(v_i, b_i, b_j, b_k) = \begin{cases} v_i - b_i, & b_i > b_j \geq b_k \\ \frac{1}{2}(v_i - b_i), & b_i = b_j > b_k \\ \frac{1}{3}(v_i - b_i), & b_i = b_j = b_k \\ 0 & b_i < b_j \leq b_k \end{cases}$$

On any continuous distribution, the probability that any two draws are exactly equal is zero. Therefore, the knife's-edge cases are excluded from the reminder of the analysis for brevity.

- b)

- c) To determine expected revenue, we must first determine the expected value of the third-highest valuation. Begin by solving the pdf of the third-highest price, f^3 , using the order statistic pdf formula:

$$f^3(x) = \frac{n!(n-3)!}{2!} x^2(1-x)^{n-3} = \left(\frac{n(n-1)(n-2)}{2} \right) x^2(1-x)^{n-3}$$

Then, the expected revenue is the expected value of the third-lowest bid:

$$\begin{aligned} R_s &= \frac{n(n-1)(n-2)}{2} \int_0^1 v^2(1-v)^{n-3} dv = \frac{n(n-1)(n-2)}{2} \int_1^0 (1-2y+y^2)y^{n-3} dy \\ &= \frac{n(n-1)(n-2)}{2} \int_1^0 y^{n-3} - 2y^{n-2} + y^{n-1} dy \end{aligned}$$

Thus, the expected revenue of the seller is:

$$\mathbb{E} \left[\frac{n-1}{n-2} v \right] = \frac{n-1}{(n-2)(n+1)}$$

- d)