

Functors Applicatives

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1 Functors

Functors increase the level of generality in Haskell. Think of them as functions over a range of parameterised types such as lists, trees.

Functors have both an **infix notation** where $(\<\$>) = \text{fmap}$, e.g. $\text{fmap } (>1) [1,2,3]$ and a **prefix notation** where $\text{fmap } (g.h) = \text{fmap } g . \text{fmap } h$, e.g. $(>1) \<\$> (\text{Just } 1)$

Functor Laws

1. $\text{fmap id} = \text{id}$
2. $\text{fmap } (g.h) = \text{fmap } g . \text{fmap } h$

Application of Composition $(f.g) x, f.g x, f g x, f g x$

Structural Induction Suppose S is some **recursively defined structure** (e.g. [a] or Tree a) that has **substructure** (e.g. sublist, subtree) and there is **partial ordering**, e.g. length or number of nodes.

Structural induction implies if..

1. P is true for **all minimum structures**, and (Base Case)
2. $P(x')$ is true for x' any **immediate substructure** of x (Induction)

2 Applicatives

Functors can map a function over each element in a structure. Suppose we wish to generalise the idea to allow functions with **any number of arguments** to be mapped.

Applicative Laws

1. $\text{pure id } \<*> x = x$ *mapping identity has no effect*
2. $\text{pure } (g x) = \text{pure } g \<*> \text{pure } x$ *pure distribution*
3. $x \<*> \text{pure } y = \text{pure } (-> g y) \<*> x$ *effectful function disregards evaluation order*
4. $x \<*> (y \<*> z) = (\text{pure } (.) \<*> x \<*> y) \<*> z$ *associativity*

Corollary $g \<\$> x = \text{pure } g \<*> x$

Effectful Programming Applicative functors abstract the idea of applying **pure functions** to **effectful arguments**, where the effects depend on the underlying functor. Effects may be: possibility of failure, having many ways to succeed or I/O actions.

Applicative Style

Note that: $\text{fmap } g x = \text{pure } g \<*> x$, so this means

$\text{pure } g \<*> x1 \<*> x2 \<*> \dots \<*> xn$

is same as

$g \<\$> x1 \<*> x2 \<*> \dots \<*> xn$