# Natural Language Models and Interfaces Part B, lecture 2

Ivan Titov

Institute for Logic, Language and Computation

# Today

- Parsing algorithms for CFGs
  - Recap, Chomsky Normal Form (CNF)
  - A dynamic programming algorithm for parsing (CKY)
  - Extension of CKY to support unary inner rules
- Parsing for PCFGs
  - Extension of CKY for parsing with PCFGs
- Parser evaluation (if we have time)

After this lecture you should be able to start working on the assignment step 2.1

## Parsing

Parsing is search through the space of all possible parses

**PCFG** model

e.g., we may want either any parse, all parses or the highest scoring parse (if PCFG):The probability by the

 $\underset{T \in G(x)}{\operatorname{arg\,max}} P(T)$ 

Set of all trees given by the grammar for the sentence *x* 

- Bottom-up:
  - One starts from words and attempt to construct the full tree
- Top-down
  - Start from the start symbol and attempt to expand to get the sentence

# CKY algorithm (aka CYK)

- Cocke-Kasami-Younger algorithm
  - Independently discovered in late 60s / early 70s
- ▶ An efficient bottom up parsing algorithm for (P)CFGs
  - can be used both for the recognition and parsing problems
- Very important in NLP (and beyond)

We will start with the non-probabilistic version

## Constraints on the grammar

The basic CKY algorithm supports only rules in the Chomsky Normal Form (CNF):

Unary preterminal rules (generation of words given PoS tags 
$$N \to telescope$$
,  $D \to the$ , ...)
$$C \to C_1 C_2$$
Binary inner rules (e.g.,  $S \to NP\ VP,\ NP \to D\ N$ )

- Any CFG can be converted to an equivalent CNF
  - Equivalent means that they define the same language

Makes linguists unhappy

- However (syntactic) trees will look differently
- It is possible to address it but defining such transformations that allows for easy reverse transformation

#### Transformation to CNF form

- What one need to do to convert to CNF form
  - lacktriangle Get rid of empty (aka epsilon) productions:  $C 
    ightarrow \epsilon$
  - Get rid of unary rules:  $C \rightarrow C_1$
  - N-ary rules:  $C \rightarrow C_1 \ C_2 \dots C_n \ (n > 2)$

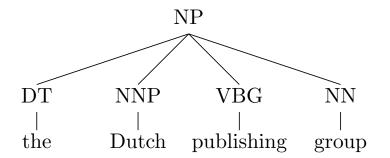
Crucial to process them, as required for efficient parsing

Generally not a problem as there are not empty production in the standard (postporcessed) treebanks

Not a problem, as our CKY algorithm will support unary rules

#### Transformation to CNF form: binarization

▶ Consider  $NP \rightarrow DT \ NNP \ VBG \ NN$ 



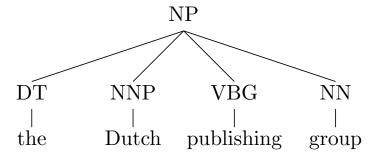
How do we get a set of binary rules which are equivalent?

$$NP \rightarrow DT X$$
 $X \rightarrow NNP Y$ 
 $Y \rightarrow VBG NN$ 

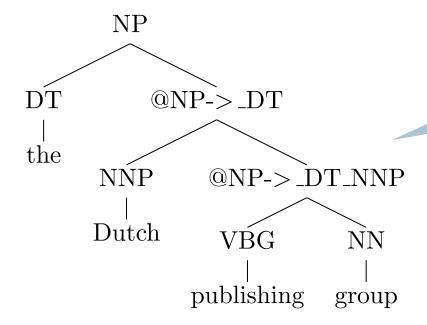
▶ A more systematic way to refer to new non-terminals

#### Transformation to CNF form: binarization

Instead of binarizing tules we can binarize trees on preprocessing:



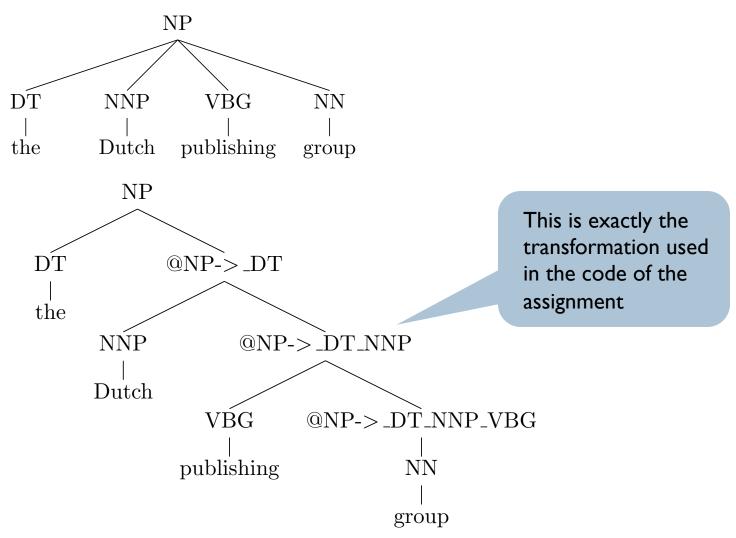
Also known as **lossless Markovization** in the context of PCFGs



Can be easily reversed on postprocessing

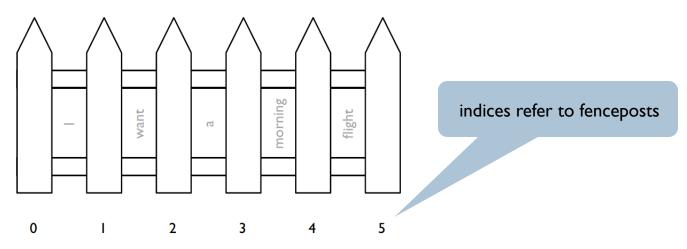
#### Transformation to CNF form: binarization

Instead of binarizing tules we can binarize trees on preprocessing:



[Some illustrations and slides in this lecture are from Marco Kuhlmann]

- We a given
  - a grammar  $G = (V, \Sigma, R, S)$
  - $m{v}$  a sequence of words  $m{w}=(w_1,w_2,\ldots,w_n)$
- lacktriangle Our goal is to produce a parse tree for w
- lacktriangle We need an easy way to refer to substrings of w



start symbol

span (i, j) refers to words between fenceposts i and j

## Key problems

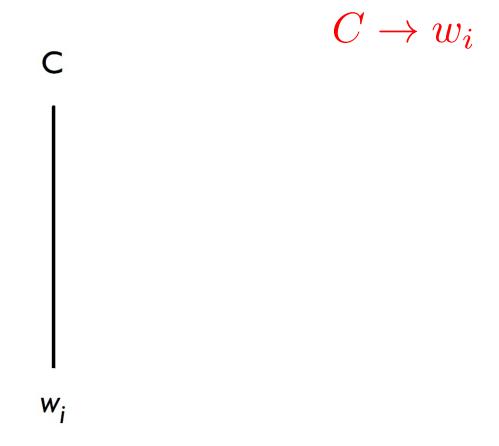
- Recognition problem: does the sentence belong to the language defined by CFG?
  - Is there a derivation which yields the sentence?
- Parsing problem: what is a derivation (tree) corresponding the sentence?
  - Probabilistic parsing: what is the most probable tree for the sentence?

# Parsing one word

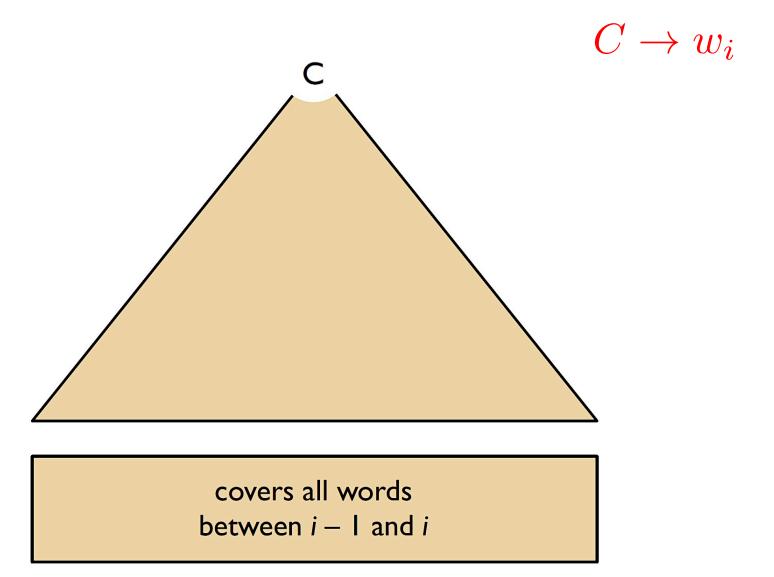
$$C \to w_i$$

 $w_i$ 

# Parsing one word

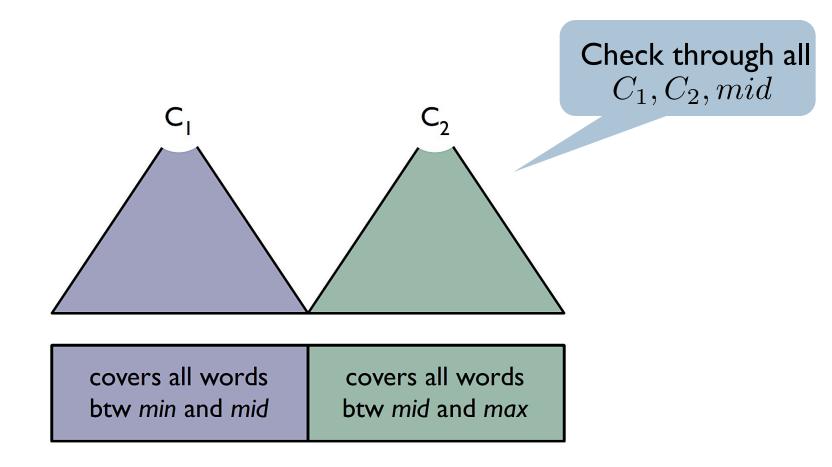


## Parsing one word



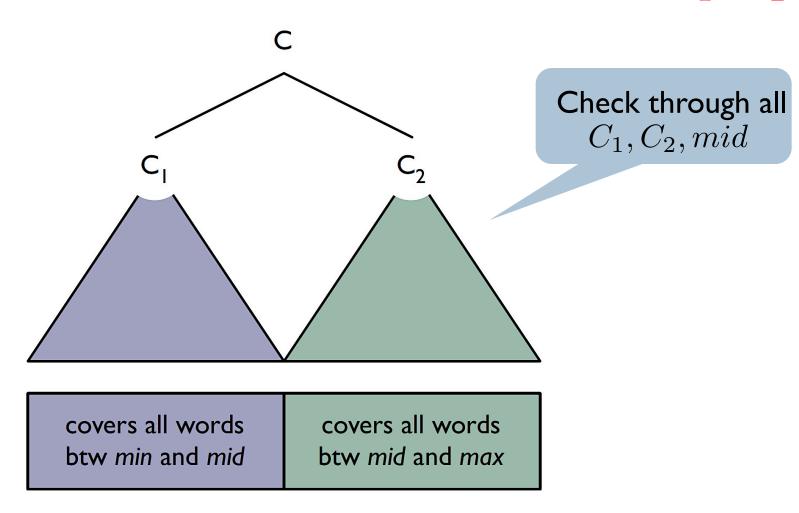
## Parsing longer spans

$$C \rightarrow C_1 \ C_2$$

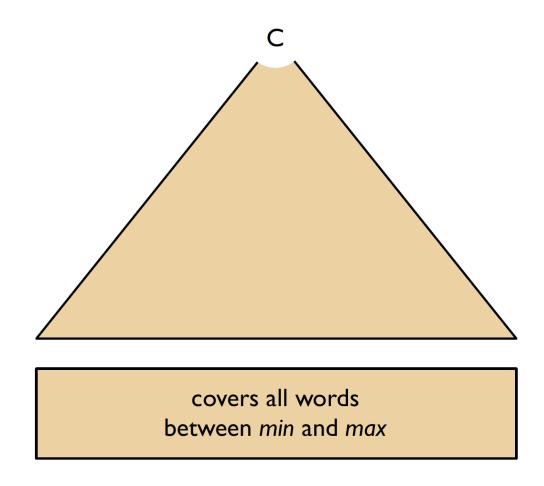


## Parsing longer spans





# Parsing longer spans



## Signatures

- Applications of rules is independent of inner structure of a parse tree
- We only need to know the corresponding span and the root label of the tree
  - lacksquare Its signature [min, max, C]

Also known as an edge

#### CKY idea

- Compute for every span a set of admissible labels (may be empty for some spans)
  - Start from small trees (single words) and proceed to larger ones
- ▶ When done, check if S is among admissible labels for the whole sentence, if yes the sentence belong to the language
  - That is if a tree with signature [0, n, S] exists
- Unary rules?

$$S \to NP VP$$

$$VP \rightarrow M \ V$$
  $VP \rightarrow V$ 

$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 
$$N \rightarrow lead$$
 
$$N \rightarrow poison$$

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

max = 1 max = 2 max = 3

S? min = 0 min = 1

 $S \to NP VP$ 

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$

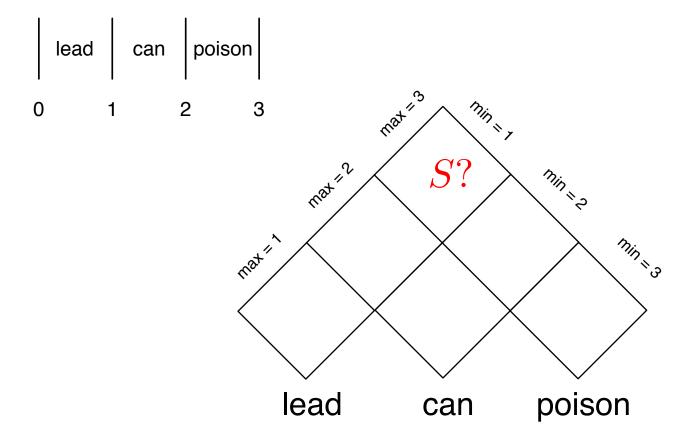
$$NP \to N NP$$

 $N \rightarrow can$   $N \rightarrow lead$   $N \rightarrow poison$ 

 $M \to can$   $M \to must$ 

Chart (aka parsing triangle)

$$V \to poison$$
  $V \to lead$ 



$$VP \to M V$$
 $VP \to V$ 

 $S \to NP \ VP$ 

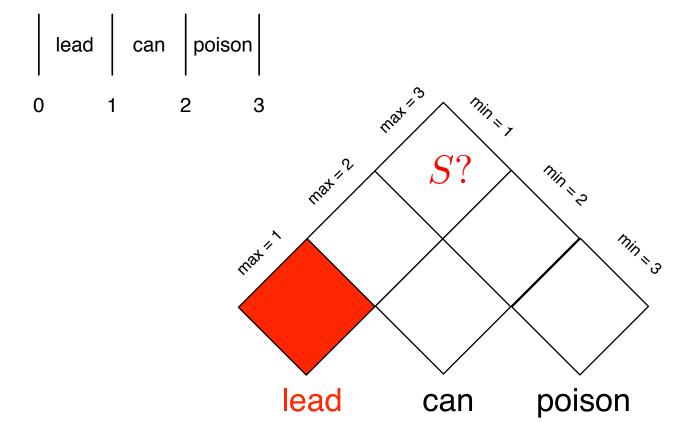
$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 $N \rightarrow lead$ 
 $N \rightarrow poison$ 

$$M \to can$$
  $M \to must$ 

$$V \to poison$$
 
$$V \to lead$$



$$VP \to M V$$
 $VP \to V$ 

 $S \to NP \ VP$ 

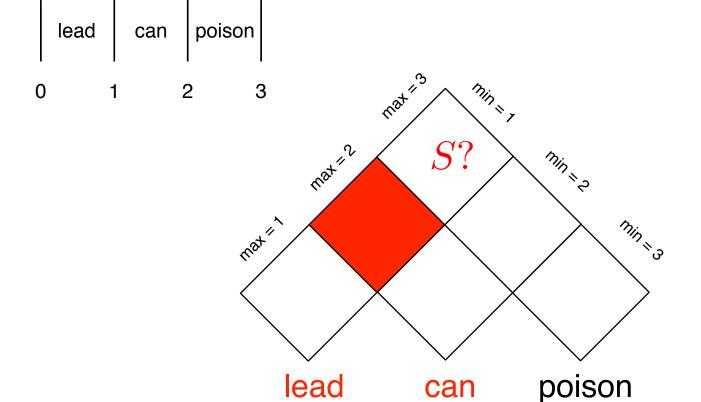
$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 $N \rightarrow lead$ 
 $N \rightarrow poison$ 

$$M \to can$$
  $M \to must$ 

$$V \to poison$$
 
$$V \to lead$$



$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

 $S \to NP \ VP$ 

$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 $N \rightarrow lead$ 
 $N \rightarrow poison$ 

$$M \to can$$
$$M \to must$$

$$V \to poison$$
 
$$V \to lead$$

max = 1 max = 2 max = 3

 $\min = 0$   $\begin{bmatrix} 1 & & 4 & & 6 \\ & S? & & \\ & & 2 & & 5 \\ \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\$ 

$$S \rightarrow NP VP$$

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$

$$NP \to N NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

max = 1 max = 2 max = 3

min = 0 ? 2 ? min = 1 3 ?

$$VP \rightarrow M \ V$$
  
 $VP \rightarrow V$ 

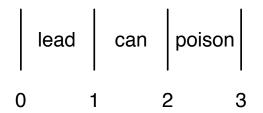
$$NP \to N$$

$$NP \to N NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison$$
 
$$V \rightarrow lead$$



max = 1 max = 2 max = 3 min = 0 min = 1 min = 2 min = 2 max = 3 min = 3

$$S \rightarrow NP VP$$

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

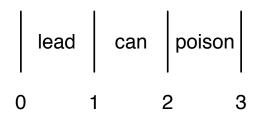
$$NP \to N$$

$$NP \to N NP$$

 $N \rightarrow can$   $N \rightarrow lead$   $N \rightarrow poison$ 

 $M \to can$   $M \to must$ 

 $V \rightarrow poison$   $V \rightarrow lead$ 



max = 1 max = 2 max = 3

$$S \to NP \ VP$$

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

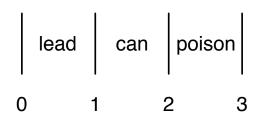
$$NP \to N$$

$$NP \to N NP$$

 $N \rightarrow can$   $N \rightarrow lead$   $N \rightarrow poison$ 

 $M \to can$   $M \to must$ 

 $V \rightarrow poison \\ V \rightarrow lead$ 

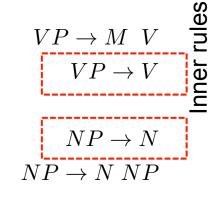


min = 0	$egin{bmatrix} {\sf 1} & N, V & \ NP, VP & \ \end{pmatrix}$		
min = 1		$ \begin{array}{c} 2 \\ N, M \\ NP \end{array} $	
min = 2			NP, VP

max = 2

max = 1

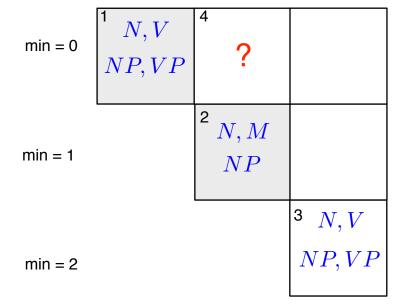
max = 3



 $N \rightarrow can$   $N \rightarrow lead$   $N \rightarrow poison$ 

 $M \to can$   $M \to must$ 

$$V \rightarrow poison \\ V \rightarrow lead$$



$$S \rightarrow NP VP$$

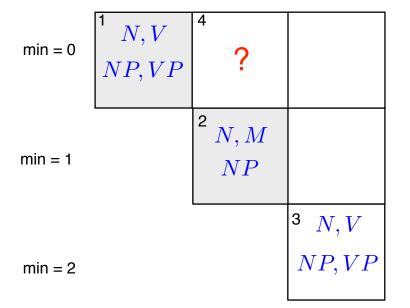
$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$
$$NP \to N NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$\begin{aligned} M \to can \\ M \to must \end{aligned}$$

$$V \to poison$$
 
$$V \to lead$$



$$S \to NP VP$$

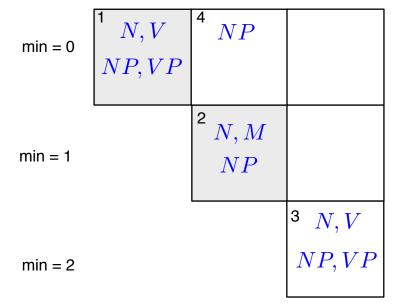
$$VP \to M V$$
 $VP \to V$ 

$$NP \to N$$
$$NP \to N \ NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$



$$S \to NP VP$$

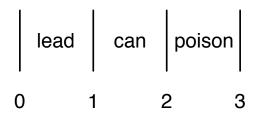
$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

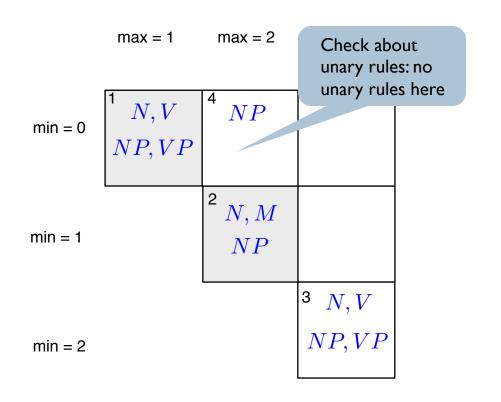
$$NP \to N$$
$$NP \to N \ NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$





$$S \to NP \ VP$$

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

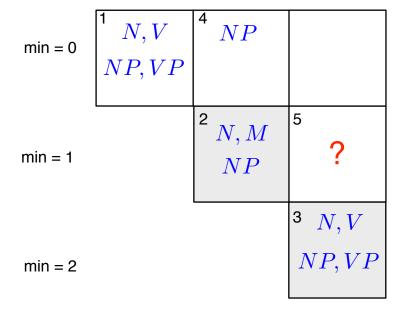
$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison$$
  $V \rightarrow lead$ 



$$S \rightarrow NP VP$$

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$

$$NP \to N NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison$$
  $V \rightarrow lead$ 

max = 1 max = 2 max = 3

$$\min = 0 \quad \begin{bmatrix} 1 & N, V & 4 & NP \\ NP, VP & & & \\ \end{bmatrix}$$

$$\min = 1 \quad \begin{bmatrix} 2 & N, M & 5S, VP, \\ NP & & NP \\ \end{bmatrix}$$

$$\min = 2 \quad \begin{bmatrix} 3 & N, V \\ NP, VP \\ \end{bmatrix}$$

 $S \rightarrow NP \ VP$ 

$$\begin{array}{c} VP \to M \ V \\ VP \to V \end{array}$$

$$NP \to N$$

$$NP \to N \ NP$$

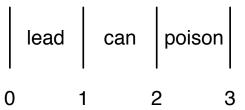
$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

 $VP \to M \ V$ 

#### CKY in action



 $VP \to V$  $NP \rightarrow N$  $NP \rightarrow N NP$ 

$$max = 1$$
  $max = 2$   $max = 3$ 

min = 0	$\begin{bmatrix} 1 & N, V \\ NP, VP \end{bmatrix}$	<sup>4</sup> NP		
min = 1		$egin{array}{c} 2 \\ N, M \\ NP \end{array}$	${5 \atop NP}$	
min = 2			NP, VP	

Check about unary rules: no unary rules here

 $N \to poison$  $M \to can$  $M \to must$ 

 $N \to can$ 

 $N \rightarrow lead$ 

$$V \rightarrow poison$$
  $V \rightarrow lead$ 

max = 1 max = 2 max = 3

 $\min = 0 \quad \begin{bmatrix} 1 & N, V \\ NP, VP \end{bmatrix} \quad \begin{bmatrix} 4 & NP \\ NP, VP \end{bmatrix} \quad \begin{bmatrix} 6 \\ ? \\ NM \\ NP \end{bmatrix} \quad \begin{bmatrix} 5 \\ S, VP, \\ NP \end{bmatrix}$   $\min = 1 \quad \begin{bmatrix} 3 & N, V \\ NP, VP \end{bmatrix}$   $\min = 2 \quad \begin{bmatrix} NP, VP \\ NP \end{bmatrix}$ 

 $S \rightarrow NP VP$ 

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

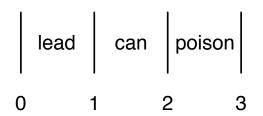
$$NP \to N$$

$$NP \to N NP$$

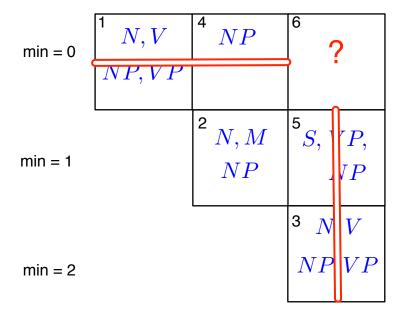
$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$



max = 1 max = 2 max = 3



 $S \to NP VP$ 

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$

$$NP \to N NP$$

$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \to poison$$
 
$$V \to lead$$

max = 2max = 1max = 3

mid = 1

min = 0	$ \begin{array}{c} 1 & N, V \\ NP, VP \end{array} $	$^{4}$ $NP$	$^{6}S$ , $NP$
min = 1		$ \begin{array}{c} 2 \\ N, M \\ NP \end{array} $	5S, VP, NP
			N, V
min = 2			NP, VP

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

 $S \to NP \ VP$ 

$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 
$$N \rightarrow lead$$
 
$$N \rightarrow poison$$

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

max = 1 max = 2 max = 3

mid = 2

min = 0	$\begin{bmatrix} 1 & N, V \\ NP, VP \end{bmatrix}$	$^{4}$ $NP$	$egin{array}{c} {f 6}_{m S},NP \ S(?!) \end{array}$
min = 1		$ \begin{array}{c c} 2 & N, M \\ NP & \end{array} $	5S, VP, NP
			$\left  egin{array}{ccc} {\sf 3} & N,V & \\ NP,VP & \end{array}  ight $
min = 2			

 $S \to NP \ VP$ 

$$VP \rightarrow M V$$
 $VP \rightarrow V$ 

$$NP \to N$$

$$NP \to N \ NP$$

$$N \rightarrow can$$
 
$$N \rightarrow lead$$
 
$$N \rightarrow poison$$

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

max = 1 max = 2 max = 3

mid = 2

$$\min = 0 \quad \begin{bmatrix} 1 & N, V & 4 & NP & 6 \\ NP, VP & & S(?!) & \\ \end{bmatrix}$$

$$\min = 1 \quad \begin{bmatrix} 2 & N, M & 5 \\ NP & & NP \\ \end{bmatrix}$$

$$3 \quad N \quad V$$

Apparently the sentence is ambiguous for the grammar: (as the grammar overgenerates)

$$VP \rightarrow M V$$
  
 $VP \rightarrow V$ 

 $S \to NP \ VP$ 

$$NP \to N$$

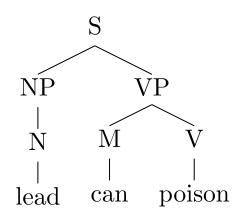
$$NP \to N NP$$

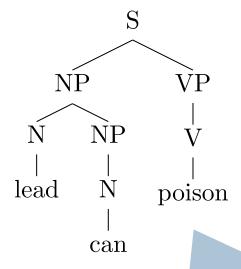
$$N \rightarrow can$$
  $N \rightarrow lead$   $N \rightarrow poison$ 

$$M \to can$$
 
$$M \to must$$

$$V \rightarrow poison \\ V \rightarrow lead$$

# **Ambiguity**





No subject-verb agreement, and poison used as an intransitive verb

Apparently the sentence is ambiguous for the grammar: (as the grammar overgenerates)

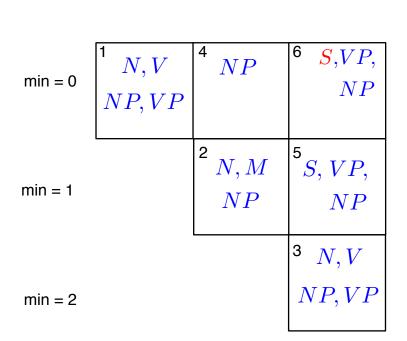
# **CKY** more formally

Here we assume that labels (C) are integer indices

▶ Chart can be represented by a Boolean array chart [min] [max] [C]

max = 3

- ightharpoonup Relevant entries have 0 < min < max  $\leq$  n
- chart[min][max][C] = true if the signature (min, max, C) is already added to the chart; false otherwise.



max = 1 max = 2

In the assignment code we use a class Chart but its access methods are similar

# Implementation: preterminal rules

```
for each wi from left to right
  for each preterminal rule C -> wi
    chart[i - 1][i][C] = true
```

## Implementation: binary rules

```
for each max from 2 to n

for each min from max - 2 down to 0

for each syntactic category C

for each binary rule C -> C1 C2

for each mid from min + 1 to max - 1

if chart[min][mid][C1] and chart[mid][max][C2] then

chart[min][max][C] = true
```

## Implementation: unary rules

```
for each max from 1 to n
                                             new bounds!
  for each min from max - 1 down to 0
   // First, try all binary rules as before.
    // Then, try all unary rules.
    for each syntactic category C
      for each unary rule C -> C1
        if chart[min][max][C1] then
          chart[min][max][C] = true
```

But we forgot something!

# Unary closure

What if the grammar contained 2 rules:

$$A \to B$$

$$B \to C$$

▶ But C can be derived from A by a chain of rules:

$$A \to B \to C$$

 One could support chains in the algorithm but it is easier to extend the grammar, to get the reflexive transitive closure

$$\begin{array}{ccc} A \to B & A \to A \\ B \to C & B \to B \\ A \to C & C \to C \end{array}$$

Convenient for programming reasons in the PCFG case

### Implementation: skeleton

```
// int n = number of words in the sequence

// int m = number of syntactic categories in the grammar

// int s = the (number of the) grammar's start symbol

boolean[][][] chart = new boolean[n + 1][n + 1][m]

// Recognize all parse trees built with with preterminal rules.

// Recognize all parse trees built with inner rules.

return chart[0][n][s]
```

## Algorithm analysis

Time complexity?

```
for each max from 2 to n

for each min from max - 2 down to 0

for each syntactic category C

for each binary rule C -> C1 C2

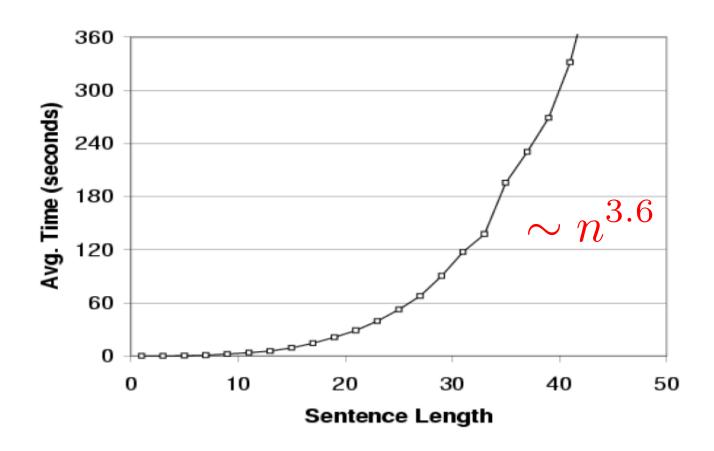
for each mid from min + 1 to max - 1
```

A few seconds for sentences under < 20 words for a non-optimized parser

- $heta(n^3|R|)$  , where |R| is the number of rules in the grammar
- There exist algorithms with better asymptotical time complexity but the `constant' makes them slower in practice (in general)

## Practical time complexity

Time complexity? (for the PCFG version)



[Plot by Dan Klein]

# Today

- Parsing algorithms for CFGs
  - Recap, Chomsky Normal Form (CNF)
  - A dynamic programming algorithm for parsing (CKY)
  - Extension of CKY to support unary inner rules
- Parsing for PCFGs
  - Extension of CKY for parsing with PCFGs
- Parser evaluation (if we have time)

## Probabilistic parsing

- We discussed the recognition problem:
  - check if a sentence is parsable with a CFG
- Now we consider parsing with PCFGs
  - Recognition with PCFGs: what is the probability of the most probable parse tree?
  - Parsing with PCFGs: What is the most probable parse tree?

#### Distribution over trees

- Let us denote by G(x) the set of derivations for the sentence x
- The probability distribution defines the scoring  $\,P(T)$  over the trees  $\,T\in G(x)\,$
- Finding the best parse for the sentence according to PCFG:

$$\underset{T \in G(x)}{\operatorname{arg\,max}} P(T)$$

### **CKY** with PCFGs

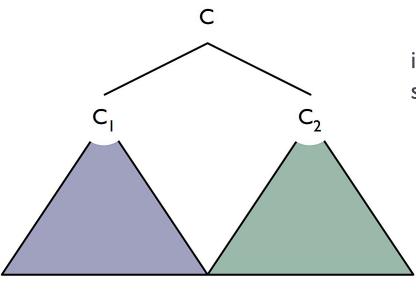
- ▶ Chart is represented by a double array chart [min] [max] [C]
  - It stores probabilities for the most probable subtree with a given signature
- chart [0] [n] [S] will store the probability of the most probable full parse tree

#### Intuition

covers all words

btw min and mid

$$C \rightarrow C_1 \ C_2$$



covers all words

btw mid and max

For every C choose  $C_1$ ,  $C_2$  and mid such that

$$P(T_1) \times P(T_2) \times P(C \to C_1 C_2)$$

is maximal, where  $T_1$  and  $T_2$  are left and right subtrees.

## Implementation: preterminal rules

```
for each wi from left to right
  for each preterminal rule C -> wi
    chart[i - 1][i][C] = p(C -> wi)
```

## Implementation: binary rules

```
for each max from 2 to n
  for each min from max - 2 down to 0
    for each syntactic category C
      double best = undefined
      for each binary rule C -> C1 C2
         for each mid from min + 1 to max - 1
           double t_1 = chart[min][mid][C_1]
           double t<sub>2</sub> = chart[mid][max][C<sub>2</sub>]
           double candidate = t_1 * t_2 * p(C \rightarrow C_1 C_2)
           if candidate > best then
             best = candidate
      chart[min][max][C] = best
```

# Unary (reflexive transitive) clo

The fact that the rule is composite needs to be stored to recover the true tree

$$A \rightarrow B \quad 0.1$$

$$B \rightarrow C \quad 0.2$$

$$A \rightarrow B \quad 0 \quad A \rightarrow A \quad 1$$

$$B \rightarrow C \quad 0.2 \quad B \rightarrow B \quad 1$$

$$A \rightarrow C \quad 0.2 \times 0.1 \quad C \rightarrow C \quad 1$$

Note that this is not a PCFG anymore as the rules do not sum to I for each parent

$$A \rightarrow B$$
 0.1  $A \rightarrow B$  0.1  $A \rightarrow A$  1   
  $B \rightarrow C$  0.2  $\Rightarrow$   $B \rightarrow C$  0.1  $B \rightarrow B$  1   
  $A \rightarrow C$  1. $e - 5$   $A \rightarrow C$  0.02  $C \rightarrow C$  1

What about loops, like:  $A \rightarrow B \rightarrow A \rightarrow C$ ?

## Recovery of the tree

- For each signature we store backpointers to the elements from which it was built (e.g., rule and, for binary rules, midpoint)
  - start recovering from [0, n, S]

- Be careful with unary rules
  - Basically you can assume that you always used an unary rule from the closure (but it could be the trivial one  $\,C \to C\,$  )

# Speeding up the algorithm (approximate search)

#### Basic pruning (roughly):

- For every span (i,j) store only labels which have the probability at most N times smaller than the probability of the most probable label for this span
- Check not all rules but only rules yielding subtree labels having non-zero probability

#### Coarse-to-fine pruning

Parse with a smaller (simpler) grammar, and precompute (posterior) probabilities for each spans, and use only the ones with non-negligible probability from the previous grammar

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#### Parser evaluation

Though has many drawbacks it is easier and allows to track state-of-the-art across many years

- Intrinsic evaluation:
  - Automatic: evaluate against annotation provided by human experts (gold standard) according to some predefined measure
  - Manual: ... according to human judgment
- Extrinsic evaluation: score syntactic representation by comparing how well a system using this representation performs on some task
  - E.g., use syntactic representation as input for a semantic analyzer and compare results of the analyzer using syntax predicted by different parsers.

## Standard evaluation setting in parsing

- Automatic intrinsic evaluation is used: parsers are evaluated against gold standard by provided by linguists
- ▶ There is a standard split into the parts:
  - training set: used for estimation of model parameters
  - development set: used for tuning the model (initial experiments)
  - test set: final experiments to compare against previous work

### Automatic evaluation of constituent parsers

The most standard measure; we will focus on it

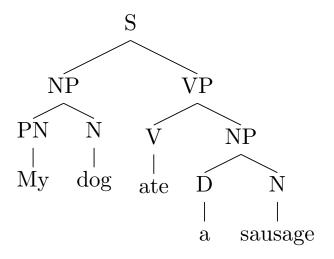
- Exact match: percentage of trees predicted correctly
- Bracket score: scores how well individual phrases (and their boundaries) are identified
- Crossing brackets: percentage of phrases boundaries crossing
- Dependency metrics: scores dependency structure corresponding to the constituent tree (percentage of correctly identified heads)

#### Brackets scores

Subtree signatures for CKY

- The most standard score is bracket score
- lacktriangle It regards a tree as a collection of brackets: [min, max, C]
- The set of brackets predicted by a parser is compared against the set of brackets in the tree annotated by a linguist
- Precision, recall and FI are used as scores

## Bracketing notation



▶ The same tree as a bracketed sequence

```
(S

(NP (PN My) (N Dog))

(VP (V ate)

(NP (D a ) (N sausage))

)
```

#### Brackets scores

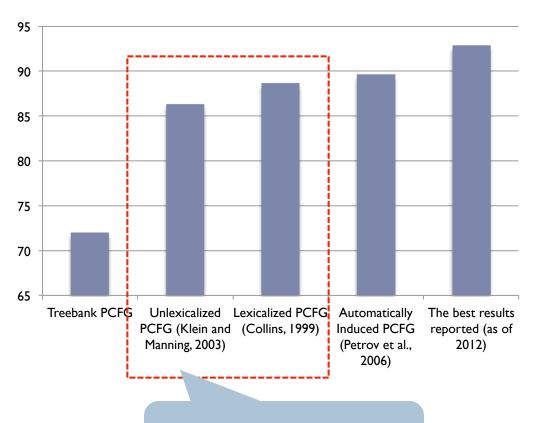
$$Pr = \frac{\text{number of brackets the parser and annotation agree on}}{\text{number of brackets predicted by the parser}}$$

$$Re = \frac{\text{number of brackets the parser and annotation agree on}}{\text{number of brackets in annotation}}$$

$$F1 = \frac{2 \times Pr \times Re}{Pr + Re}$$

Harmonic mean of precision and recall

### Preview: FI bracket score



We will introduce these models next time