

Danny Tan
Dt1462
HW#9

1)

a)

```
insert ( x, Node) {  
  if t is null  
    set Node to a new Node with element c  
  if x is less than the element of Node  
    do recursion of left subtree of Node  
  else if t is greater than the element of Node  
    do recursion of right subtree of Node  
}
```

b)

```
rangedPrint (low , high, Node) {  
  if Node is null  
    return
```

```
  if the element of Node is greater than low  
    do recursion of left subtree of Node
```

```
  if element is in between low and high  
    print out its element
```

```
  if the element of Node is less than high  
    do recursion of right subtree of Node
```

```
}
```

c)

```
stringy () {  
  set num to number of Nodes ( the size of tree)  
  do an inorder traversal on the root and store all the Nodes into a queueOfNodes  
  loop through the queue using num  
  insert the node in the front of queue onto root  
  pop that Node  
}
```

d)

```
averageNodeDepth() {  
  return depthSum / numOfNodes  
}
```

```

depthSum(Node, depth) {
  if Node is null
  return 0
endif
return depth + depthSum for left subtree+ depthSum for right subtree
}

```

```

numOfNodes (Node) {
  if Node is null
  return 0
endif
return 1 + numOfNodes for left subtree + numOfNodes for right subtree
}

```

2) Average case and worst case is $O(n)$.

3)

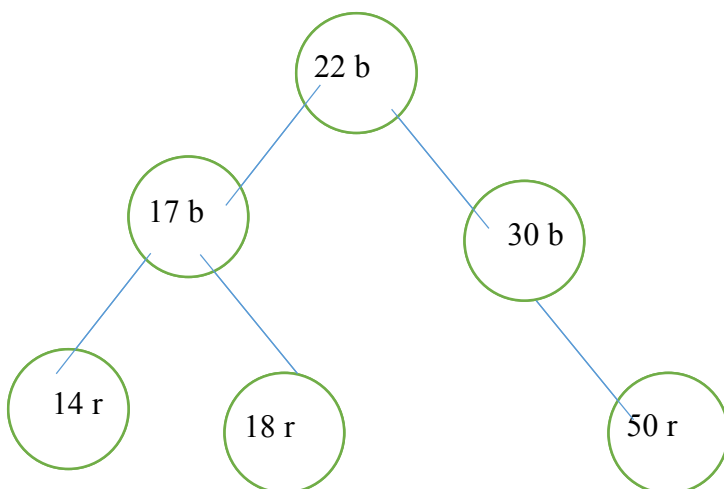
There is nothing wrong with the code. However the $t \neq \text{nullptr}$ is not necessary.

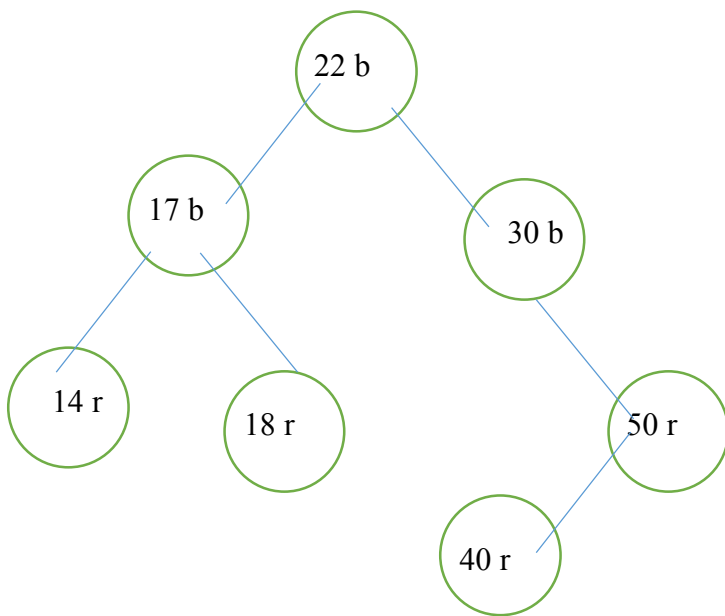
4)

`new Node{ x, nullptr, nullptr, 1 }` part will make the program not work correctly because at the end of the program, the size is increment again. So the total size of the inserted node, a leaf node will be 2 but that is not correct because its size is suppose to be 1

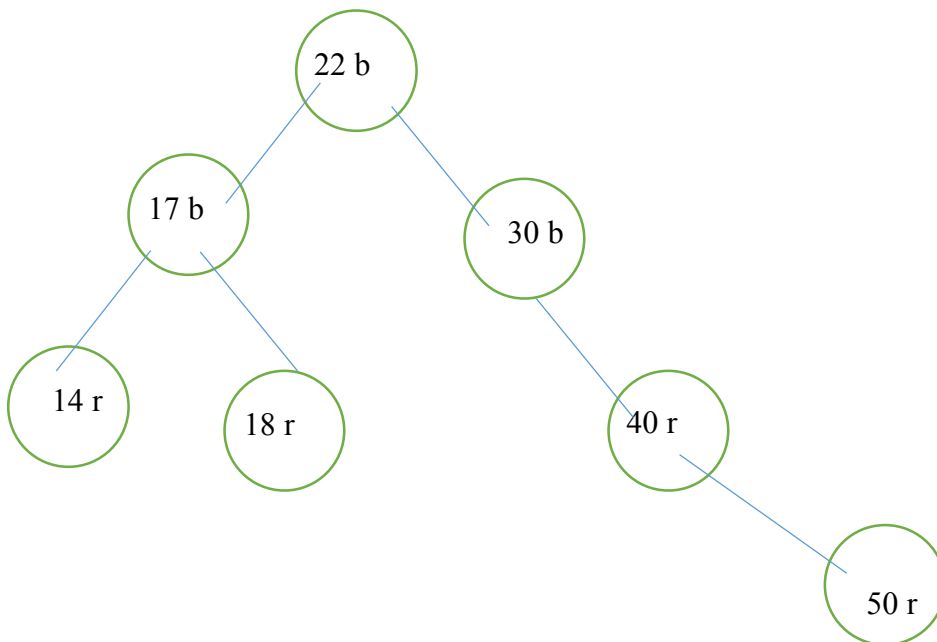
The correct code will be `new Node{ x, nullptr, nullptr, 0 }`

5)

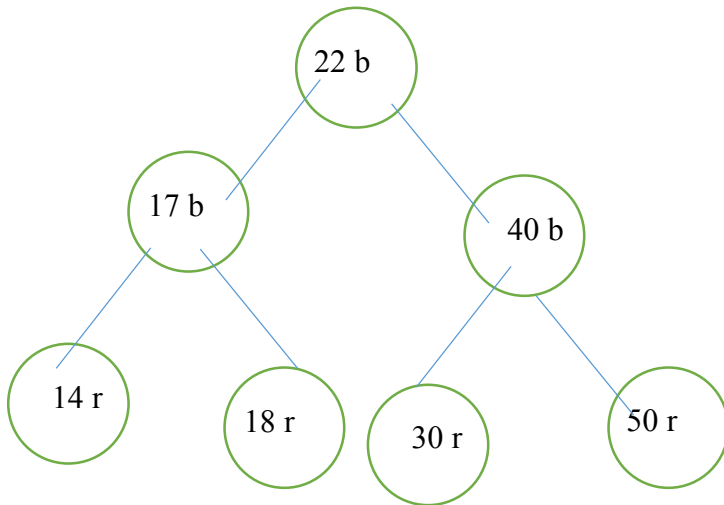




Double red problem case 3. Right Rotate



Double red problem case 2. Left rotate



7)

```
template <class Comparable>
void RedBlackTree::rightRotateRecolor( Node * & k2 ) {
    Node *k1 = k2->left;
    K2->left = k1->right;
    K1->left = k2;
    K2 = k1;
```

8) In the worst case scenario, double red problem case 1 2 and 3 is used. Case 1 does not change the pointer. Case 2 and case 3 both changes 3 pointer. So the maximum pointer change is 6.