Random Number Generation

CMS 380 Simulation and Stochastic Modeling

Pseudorandom Number Generators

High-quality random numbers are the building blocks of simulation programs. Every modern programming environment includes a method for generating, at minimum, a uniformly distributed random value in the range (0,1).

First, consider the problem of generating a "true" random number on a computer. If the "random" value is being generated by a program, it can't *really* be truly random, because there must be some underlying deterministic code calculating the value. Therefore, programs that calculate streams of random numbers are better termed *pseudorandom number generators* (PRNGs). A good PRNG is one that can produce long streams of values that have the same statistical properties as true random values.

It is possible to generate "true" random values if you sample from a random physical process. Various services have been created that do things like sample from the atmospheric background radiation, which is supposed to be effectively random. It is also possible to generate random values by sampling from human user input.

The Linear Congruential Generator

A linear congruential pseudorandom number generator has the form

$$X_{n+1} = aX_n + c \mod m$$

It generates a stream of values starting from an initial starting value X_0 , which is called the *seed*.

Here's an example of the sequence produced when $a=5,\,c=1,\,m=16,$ and $X_0=0.$

```
X_0
                                seed
                                               0
                                               1
X_1
              (5 \cdot 0 + 1) \mod 16
X_2
              (5 \cdot 1 + 1) \mod 16
                                               6
X_3
              (5 \cdot 5 + 1) \mod 16
                                              15
X_4
            (5 \cdot 15 + 1) \mod 16
                                              12
X_5
             (5 \cdot 12 + 1) \mod 16
                                              13
                                               2
X_6
             (5 \cdot 13 + 1) \mod 16
X_7
              (5 \cdot 2 + 1) \mod 16
                                              11
X_8
            (5 \cdot 11 + 1) \mod 16
                                               8
X_9
              (5 \cdot 8 + 1) \mod 16
                                               9
X_{10}
              (5 \cdot 9 + 1) \mod 16
                                              14
X_{11}
           (5 \cdot 14 + 1) \mod 16
                                               7
X_{12}
              (5 \cdot 7 + 1) \mod 16
                                               4
X_{13}
              (5 \cdot 4 + 1) \mod 16
                                               5
X_{14}
              (5 \cdot 5 + 1) \mod 16
                                              10
       = (5 \cdot 10 + 1) \mod 16
                                               3
X_{15}
X_{16}
              (5 \cdot 3 + 1) \mod 16
                                               0
```

 $X_16 = X_0$, so the sequence would repeat from that point forward.

Period. An LCG can only produce values in the range 0 to m-1, because that's the output range of the mod operation. The generator above has *full period*: it cycles through all 16 possible values for m=16 before repeating. Having a long period is a good property for a PRNG, because a simulation to generate long sequences of random values without worrying about repetition. Having a large m is necessary, but the relationship between a, c, and m is also important. You can do a little more reading and find out about some of the theorems governing their relationship if you're interested.

Parameters. Choosing the values of a, c, and m is a challenging problem. It requires careful testing and validation to find a stream that produces both a long period and a stream of random numbers with good randomness properties. For example, Java's Random class uses $m=2^{48}$, a=25214903917, and c=11.

Because parameterizing a PRNG is hard, **you should never write your own to use in production code**, unless there's some reason you have no other choice. It's too easy to come up with an implementation that seems superficially correct, but ends up being biased or inaccurate in some way. There was, in fact, a famous bug in a C function called RANDU, which had the formula

$$X_{n+1} = 65539 \cdot X_n \mod 2^{31}$$

This turns out to produce random number streams that are highly correlated, meaning that there is a dependence between successive values, which makes them definitely not random enough for valid simulation results. In fact, if the initial seed is odd (which it was in implementations), the formula would only generate odd numbers! When the RANDU bug was finally discovered, it required re-validation of some simulation results that could no longer be trusted due to the low quality of its random number streams.

Uniform Random Variables. One last point: the LCG generates integers in [0, m-1], but we often want floating point values in (0,1). If the generator has a good stream and m is large, then:

$$U_n = \frac{X_n}{m}$$

is a good approximation of a uniformly distributed random variable in (0,1).

The Inverse CDF Method

Suppose that you have implemented a good quality PRNG and can use it to generate those uniformly distributed random values in (0,1). What do you do when you want, say, an exponentially distributed random variable in a simulation?

The answer: you need to *transform* a uniformly distributed random number (which you can get from your PRNG) into an exponentially distributed random number. This is the problem of *random variate generation*.

Recall that the CDF of the exponential distribution is

$$F_X(x) = P(X \le x) = 1 - e^{-\lambda x}$$

 $F_X(x)$ is a probability, so it must be between 0 and 1.

Here's the strategy:

- Let u be uniform random value from the PRNG.
- Set $F_X(x) = u$ and solve for x.
- The resulting formula will give you a way to generate random values drawn from the distribution with F_X as its CDF.

Applying this strategy to the exponential CDF yields:

$$1 - e^{-\lambda x} = u$$
$$e^{-\lambda x} = 1 - u$$

Take the natural logarithm of both sides to undo the exponentiation:

$$-\lambda x = \ln(1 - u)$$

The final formula is

$$x = -\frac{\ln(1-u)}{\lambda}$$

This can be simplified a little bit: if u is uniformly distributed in (0,1), then so is 1-u, so it's sufficient to use

$$x = -\frac{\ln u}{\lambda}$$

Variates calculated by this method will be exponentially distributed with parameter λ . Here's some Python code.

from math import random, log

def randExp(mu):

"""Generate an exponential random variate with parameter mu.

Input: mu the parameter of the distribution

Output: a random variate x ~ exp(mu)

Tip: use mu as the parameter because lambda is a Python keyword (it's used to create anonymous functions)

return -log(random()) / mu

That's it. Python's log method calculates the natural logarithm if you use it with only one input argument (it can take the base of the log as an optional second argument).

The Rayleigh distribution is a continuous distribution that has applications to physics and aerodynamics. It has one parameter, σ , which is called the scale. It's CDF is

$$F_X(x) = 1 - e^{\frac{-x^2}{2\sigma^2}}$$

Use the Inverse CDF method to derive a function for generating Rayleigh distributed random variates.

Set $F_X(x) = u$ and solve for x:

$$1 - e^{\frac{-x^2}{2\sigma^2}} = u$$

$$e^{\frac{-x^2}{2\sigma^2}} = 1 - u$$

$$\frac{-x^2}{2\sigma^2} = \ln(1 - u)$$

$$x^2 = -2\sigma^2 \ln(1 - u)$$

$$x = \sqrt{-2\sigma^2 \ln(1 - u)}$$

The last equation can be simplified a little bit by again replacing 1-u with an equivalent u and bringing σ outside the square root:

$$x = \sigma \sqrt{-2 \ln u}$$

Note that it isn't a problem to have -2 under the square root. In fact, because u<1, $\ln u$ must be negative, so multiplying by -2 makes the entire term under the radical positive.