## **TECHIE DELIGHT </>**

Q

PRACTICE FAANG Interview Prep

Data Structures & Algorithms ➤

## Find the shortest path in a maze

Given a <u>maze</u> in the form of a binary rectangular matrix, find the shortest path's length in the maze from a given source to a given destination. The path can only be constructed out of cells having value 1, and at any moment, we can only move one step in one of the four directions.

The valid moves are:

**Go Top:** 
$$(x, y) \longrightarrow (x - 1, y)$$

**Go Left:** 
$$(x, y) \longrightarrow (x, y - 1)$$

**Go Down:** 
$$(x, y) \longrightarrow (x + 1, y)$$

**Go Right:** 
$$(x, y) \longrightarrow (x, y + 1)$$

For example, consider the following binary matrix. If source = (0, 0) and destination = (7, 0)

5), the shortest path from source to destination has length 12.

$$[\ 1\ 1\ 1\ 1\ 1\ 0\ 0\ 1\ 1\ 1\ ]$$

$$[\;0\;\;0\;\;1\;\;0\;\;1\;\;1\;\;1\;\;0\;\;0\;\;1\;]$$

$$[\ 1\ 0\ 1\ 1\ 1\ 0\ 1\ 1\ 0\ 1]$$

$$[\;0\;\;0\;\;0\;\;1\;\;0\;\;0\;\;0\;\;1\;\;0\;\;1\;]$$

This website uses cookies. By using this site you agree to the use of cookies, our policies, copyright terms and other conditions. Read our Privacy Policy.

```
[0 1 1 1 1 1 1 1 0 0]
[1 1 1 1 1 0 0 1 1 1]
[0 0 1 0 0 1 1 0 0 1]
```

## Practice this problem

To find the maze's shortest path, search for all possible paths in the maze from the starting position to the goal position until all possibilities are exhausted. We can easily achieve this with the help of backtracking. The idea is to start from the given source cell in the matrix and explore all four paths possible and recursively check if they will lead to the destination or not. Then update the minimum path length whenever the destination cell is reached. If a path doesn't reach the destination or explored all possible routes from the current cell, backtrack. To make sure that the path is simple and doesn't contain any cycles, keep track of cells involved in the current path in a matrix, and before exploring any cell, ignore the cell if it is already covered in the current path.

Following is the C++, Java, and Python implementation of the idea:

```
C++
Java
Python
1
     import sys
2
3
4
     # Check if it is possible to go to (x, y) from the current position.
5
     # function returns false if the cell is invalid, has value 0 or alrea
6
     def isSafe(mat, visited, x, y):
7
         return 0 \le x \le len(mat) and 0 \le y \le len(mat[0]) and \setminus
8
                 not (mat[x][y] == 0 \text{ or } visited[x][y])
9
10
11
     # Find the shortest possible route in a matrix `mat` from source cell
     # to destination cell `dest`.
12
13
14
     # `min dist` stores the length of the longest path from source to a d
15
     # found so far, and `dist` maintains the length of the path from a so
```

This website uses cookies. By using this site you agree to the use of cookies, our policies, copyright terms and other conditions. Read our Privacy Policy.

```
19
20
         # if the destination is found, update `min dist`
21
         if (i, j) == dest:
22
             return min(dist, min dist)
23
24
         # set (i, j) cell as visited
25
         visited[i][j] = 1
26
27
         # go to the bottom cell
28
         if isSafe (mat, visited, i + 1, j):
29
             min dist = findShortestPath(mat, visited, i + 1, j, dest, min
30
31
         # go to the right cell
         if isSafe(mat, visited, i, j + 1):
32
33
             min dist = findShortestPath(mat, visited, i, j + 1, dest, min
34
35
         # go to the top cell
36
         if isSafe(mat, visited, i - 1, j):
37
             min dist = findShortestPath(mat, visited, i - 1, j, dest, min
38
39
         # go to the left cell
         if isSafe(mat, visited, i, j - 1):
40
41
             min dist = findShortestPath(mat, visited, i, j - 1, dest, min
42
43
         # backtrack: remove (i, j) from the visited matrix
44
         visited[i][j] = 0
45
46
         return min dist
47
48
49
     # Wrapper over findShortestPath() function
50
     def findShortestPathLength(mat, src, dest):
51
52
         # get source cell (i, j)
53
         i, j = src
54
         \# get destination cell (x, y)
55
56
         x, y = dest
57
58
         # base case
59
         if not mat or len(mat) == 0 or mat[i][j] == 0 or mat[x][y] == 0:
60
             return -1
61
         # `M × N` matrix
62
63
         (M, N) = (len(mat), len(mat[0]))
64
65
         \# construct an `M 	imes N` matrix to keep track of visited cells
         visited = [[False for in range(N)] for in range(M)]
66
67
68
         min dist = findShortestPath(mat, visited, i, j, dest)
69
70
         if min dist != sys.maxsize:
71
             return min dist
72
         else:
72
              raturn = 1
```

This website uses cookies. By using this site you agree to the use of cookies, our policies, copyright terms and other conditions. Read our Privacy Policy.

```
76
     if name == ' main ':
77
78
         mat = [
79
              [1, 1, 1, 1, 1, 0, 0, 1, 1, 1],
80
              [0, 1, 1, 1, 1, 1, 0, 1, 0, 1],
81
              [0, 0, 1, 0, 1, 1, 1, 0, 0, 1],
82
              [1, 0, 1, 1, 1, 0, 1, 1, 0, 1],
83
              [0, 0, 0, 1, 0, 0, 0, 1, 0, 1],
84
              [1, 0, 1, 1, 1, 0, 0, 1, 1, 0],
85
              [0, 0, 0, 0, 1, 0, 0, 1, 0, 1],
86
              [0, 1, 1, 1, 1, 1, 1, 1, 0, 0],
87
              [1, 1, 1, 1, 1, 0, 0, 1, 1, 1],
88
              [0, 0, 1, 0, 0, 1, 1, 0, 0, 1]
89
         ]
90
91
         src = (0, 0)
92
         dest = (7, 5)
93
94
         min dist = findShortestPathLength(mat, src, dest)
95
96
         if min dist !=-1:
97
             print("The shortest path from source to destination has lengt
98
         else:
99
             print("Destination cannot be reached from source")
100
```

Download Run Code

## Output:

The shortest path from source to destination has length 12

The time complexity of the above backtracking solution will be higher since all paths need to be traveled. However, since it is the shortest path problem, Breadth-first search (BFS) would be an ideal choice. If BFS is used to solve this problem, we travel level by level. So the destination node's first occurrence gives us the result, and we can stop our search there. The BFS approach is discussed here.

- Backtracking, Matrix
- Maze, Medium, Recursive

This website uses cookies. By using this site you agree to the use of cookies, our policies, copyright terms and other conditions. Read our Privacy Policy.

Techie Delight © 2021 All Rights Reserved. Privacy Policy Contact us Useful Tools

Online Compiler/IDE

This website uses cookies. By using this site you agree to the use of cookies, our policies, copyright terms and other conditions. Read our Privacy Policy.