

*. Computer Networks

Unit -5 [The Link Layer and Local Area Network]

*. Link Layer Services:

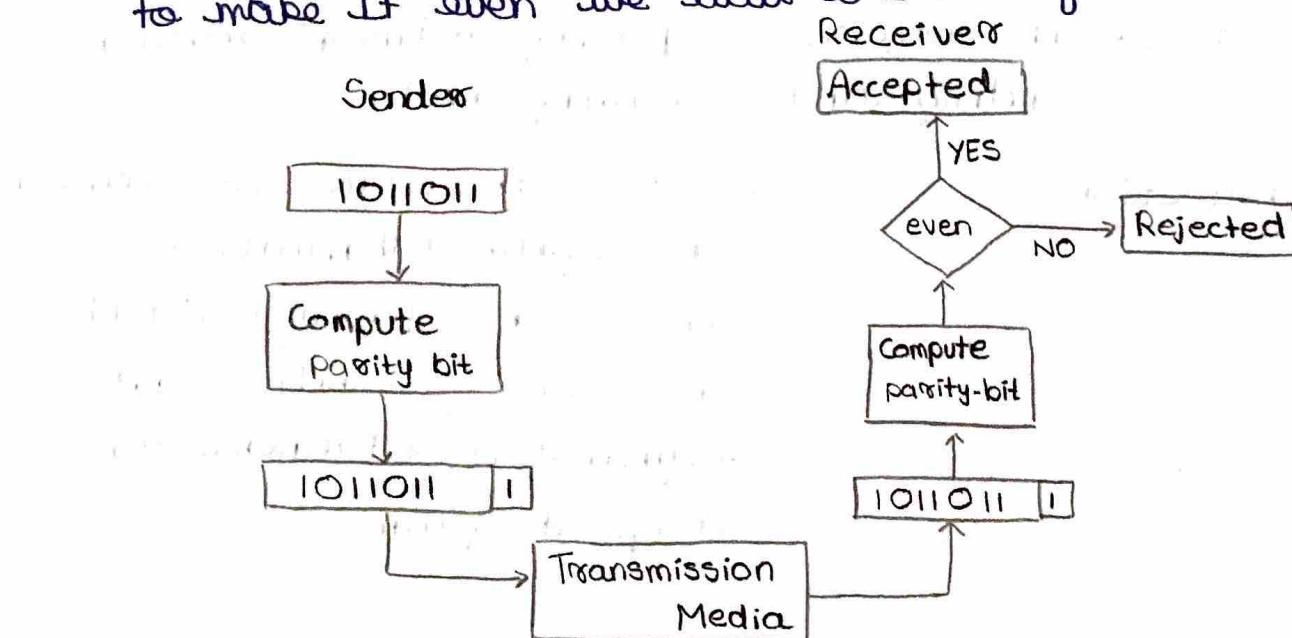
-) Framing - it encapsulates the datagrams into frames by adding header and trailer information.
-) Link Access - a medium access control (MAC) addresses is used in frame headers to identify source and destination. It is different from IP address.
-) Reliable Delivery - this layer protocol provides reliable delivery service i.e. it guarantees to move each network-layer datagram across the link without error. It can be achieved with acknowledgements and retransmissions.
-) Flow Control - it requires a pacing between adjacent sending and receiving nodes.
-) Error Detection & Correction - here the errors are caused by signal attenuation or noise. The receiver detects the errors & sender send signal for retransmission or drops the frame.

*. Error ~~Detection~~ Detection & correction technique:
- there are three techniques for error detection,

-) Parity Check.
-) Checksum Method.
-) Cyclic Redundancy Checks [CRC].

• Parity Check:

- In this, a redundant bit called parity bit is added to every data so that the no. of 1's in the data becomes even.
- For example, if even parity is used and no. of 1's is even then one bit of value 0 is added. This way the no. of 1's remains even.
- For example Similarly, if the no. of 1's is odd, to make it even we add one bit of value 1.



- the receiver counts the no. of 1's in a frame. If the count is even, the frame is considered to be not-corrupted and is accepted and if the count is odd, then the frame is not accepted.

•) Checksum:

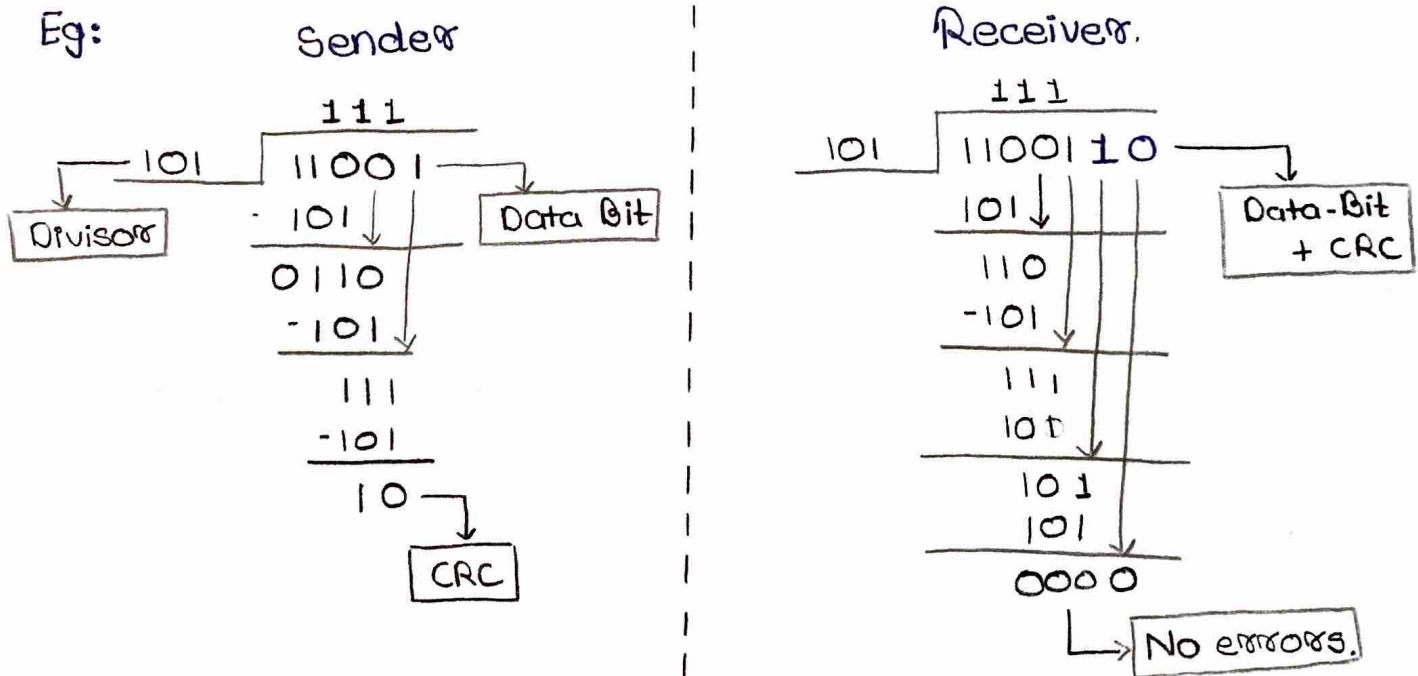
- Data is divided into k segments each of m bits.
- At sender's side, the segments are added using 1's complement to get the sum. This sum is complemented to get the checksum.
- The checksum segment is sent along with the data segments.
- At receiver's side, all received segments are added using 1's complement to get complemented sum.
- If the result is zero, the received data is accepted and if the result is non zero, then the data is corrupted.

$$\begin{array}{r} 10110011 \\ 10101011 \\ \hline 1.01011110 \\ \quad \quad \quad \searrow 1 \\ \hline 01011111 \\ 01011010 \\ \hline 10111001 \\ \quad \quad \quad \searrow 1 \\ \hline 11010101 \\ \hline 101101110 \\ \quad \quad \quad \searrow 1 \\ \hline 01101111 = \text{Sum} \\ 10010000 = \end{array}$$

* Cyclic Redundancy Checks [CRC]:

- o) CRC is the most powerful and easy to implement technique.
- o) CRC is based on binary division
- o) In CRC, a sequence of redundant bits are added at the end of the data so that the resulting data unit becomes exactly divisible by a second, predetermined binary no.
- o) At the destination, the incoming data is divided by the same no.
- o) If there is no remainder [i.e $r=0$], the data is assumed to be correct and is accepted.
- o) A remainder indicates that the data has been damaged in transit & is rejected.

Eg:



* Multiple Access Links and Protocols:

There are two types of links,

- A point to point link consists of a single sender at one end of the link and a single receiver at the other end of the link.
- A broadcast link, can have multiple sending and receiving nodes all connected to the same, single shared channel.

• Multiple Access Problem:

- In broadcast link all nodes are capable of transmitting frames, more than two nodes can transmit frames at the same time.
- When this happens, all of the nodes receive multiple frames at the same time; i.e., the transmitted frames collide at all receivers.
- Thus, all the frames involved in the collision are lost.
- Therefore, if many nodes want to transmit frames frequently, then many transmissions will result in collisions & the broadcast channel is wasted.
- In order to ensure that the broadcast channel performs well, it is necessary to coordinate the transmissions.
- And this coordination job is done by Multiple Access Protocol.

* Multiple Access Protocol:

There are 3 categories of multiple access protocol;

- Channel Partitioning Protocol:

- divide channel into smaller parts
- allocate parts to node for exclusive use
- Eg:- TDMA - Time division multiple access
FDMA - Frequency division multiple access
CDMA - Code division multiple access

- Random Access Protocol:

- channel not divided and allow collisions.
- "recovers" from collisions.
- Eg:- MAC - Medium access control
 - ALOHA
 - slotted ALOHA
 - CSMA, CSMA/CD, CSMA/CA.

- Taking-Turns Protocols:

- nodes take turns but nodes with more to send can take longer turns.
- Eg:- Polling
Token passing.

*. TDMA : Time Division, Multiple Access

- Suppose the channel supports N nodes and that the transmission rate of the channel is R bps.
- TDMA divides time into time frames and further divides each time frame into N time slots.
- Each time slot is then assigned to one of the N nodes.
- Whenever a node has a packet to send, it transmits the packet's bits during its assigned time slot in the TDMA frame.
- TDMA is attractive as it eliminates collisions and is perfectly fair.
- But there are two major drawbacks:
 1. A node is limited to an avg. rate of R/N bps. even when it is the only node with packets to send.
 2. A node must always wait for its turn, even when it is the only node with a frame to send.

*. FDMA: Frequency Division Multiple Access

- FDM divides the channels into different frequencies and assigns each frequency to one of the nodes.
- FDM thus creates N smaller channels out of the single.
- FDM shares both the advantages & disadvantages of TDM
- It avoids collisions & divides the bandwidth fairly among the nodes.

*. CDMA : Code Division Multiple Access.

- CDMA assigns a different code to each node. while TDM and FDM assign time and frequencies respectively
- Each node then uses its unique code to encode the data bits it sends
- If the codes are chosen carefully, CDMA networks have the property that different nodes can transmit simultaneously.
- Their respective receivers correctly receive a sender's encoded data bits in spite of interfering transmissions by other nodes.

*. Pure ALOHA Protocol:

- It allows user to transmit whenever they have data to be sent
- Senders wait to see if a collision occurred (after whole message has been sent)
- If collision occurs, each station involved waits a random ~~size~~ amount of time then tries again.
- Systems in which multiple users share a common channel in a way that can lead to conflicts are known as contention systems.
- Whenever two frames try to occupy the channel at the same time, there will be a collision and both will be grabbed.
- If the first bit of a new frame overlaps with with the last bit of the frame, which almost finished, both frames will be totally destroyed and both will have to be retransmitted later.
- The throughput of the Pure ALOHA is max. when the frames are of uniform length.
- The formula to calculate the throughput of the Pure ALOHA is:

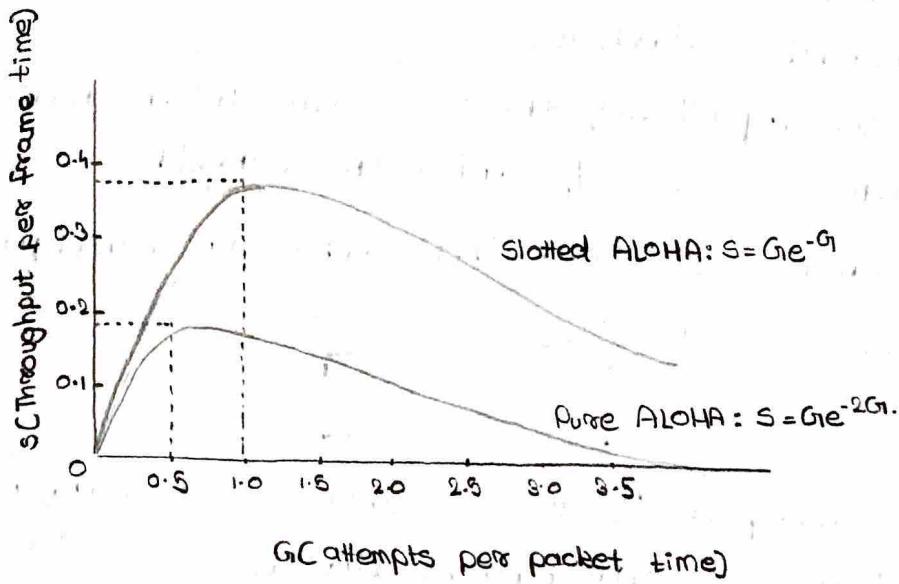
$$S = G_1 e^{-2G_1}$$

- The throughput is max when $G_1 = 1/2$ which is 18% of the total transmitted data.

* Slotted ALOHA:

- It was invented to improve the efficiency of Pure ALOHA as chances of collision were high in Pure ALOHA.
- The time of the shared channel is divided into discrete intervals called slots.
- The stations can send frequency a frame only at the beginning of the slot and only one frame is sent in each slot.
- If any station is not able to place the frame onto the channel at the beginning, then the station has to wait until the beginning of next time slot.
- The formula to calculate the throughput of the Slotted ALOHA is:

$$S = G e^{-G}$$



•) Advantages of Slotted ALOHA:

- A single active node can continuously transmit at full rate of channel.
- Slotted ALOHA is also highly decentralized.
- It is also an simple protocol.

•) Disadvantages of Slotted ALOHA:

- when there are multiple active nodes, a fraction of the slot will have collision
- idle slots
- clock synchronization.

*. ARP [Address Resolution Protocol]:

-) ARP is a protocol for mapping an IP address to a physical machine address [MAC - Media access control] that is recognized in the local network.
-) For example, in IP v4 - an address is 32-bit long. whereas the addresses for attached devices are 48-bit long.
-) A table called ARP cache, is used to maintain a correlation between each MAC address & its corresponding IP address.
-) ARP provides the protocol rules for making this correlation and providing address conversion in both directions.