

**Problem 1**

(1.C.1) Consider the choice structure  $(\mathcal{B}, C(\cdot))$  with  $\mathcal{B} = (\{x, y\}, \{x, y, z\})$  and  $C(\{x, y\}) = \{x\}$ . Show that if  $(\mathcal{B}, C(\cdot))$  satisfies the weak axiom, then we must have  $C(\{x, y, z\}) = \{x\}, = \{z\}$ , or  $= \{x, z\}$ .

**Solution.** By the definition of WARP: if for any  $B$ , such that  $x, y \in B$  and  $B \subset \mathcal{B}$  it is true that  $x \in C(B)$ , then for all other  $B'$ , such that  $x, y \in B'$  the following should be satisfied:

$$x \in B'$$

Let's suppose, that  $y \in C(\{x, y, z\})$ , then:

$$\begin{cases} y \in \{x, y, z\} & \text{by condition} \\ y \in \{x, y\} & \text{by condition} \\ y \in C(\{x, y, z\}) & \text{by assumption} \end{cases} \xrightarrow{\text{WARP}} y \in C(\{x, y\}) \quad (1)$$

But conditions state, that  $y \notin C(\{x, y\}) = \{x\}$ , thus we got contradiction.

**Problem 2**

(1.D.3) Let  $X = \{x, y, z\}$ , and consider the choice structure  $(\mathcal{B}, C(\cdot))$  with

$$\mathcal{B} = \{\{x, y\}, \{y, z\}, \{x, z\}, \{x, y, z\}\}$$

and  $C(\{x, y\}) = \{x\}$ ,  $C(\{y, z\}) = \{y\}$  and  $C(\{x, z\}) = \{z\}$ . Show that  $(\mathcal{B}, C(\cdot))$  must violate the weak axiom.

**Solution.**

If the WARP is not violated:

1. From the  $C(\{x, y\}) = \{x\}$  it should follow, that  $y \notin C(\{x, y, z\})$
2. From the  $C(\{y, z\}) = \{y\}$  it should follow, that  $z \notin C(\{x, y, z\})$
3. From the  $C(\{x, z\}) = \{z\}$  it should follow, that  $x \notin C(\{x, y, z\})$

As we can see from the points above,  $C(\{x, y, z\}) = \emptyset$  which is contradictory to the definition of choice function, thus WARP is violated.

**Problem 3**

Let  $X = \mathbb{R}$ , and define a pair of binary relations,  $R_1$  and  $R_2$ , on this set as follows: For any  $x, y \in X$ :

$$x R_1 y \iff x \geq y - 1$$

$$x R_2 y \iff x \geq y + 1$$

For  $i = 1, 2$ , determine if  $R_i$  is complete, transitive or asymmetric. Find a simple expression for the symmetric part  $R_i^\sim$ , which is defined as  $x R_i^\sim y \iff [x R_i y \text{ and } y R_i x]$ . Explain your answer.

**Solution.****Completeness.**

The binary relation  $R_i$  is complete if the following holds:

$$\forall x, y \in X, \text{ we have that either } (x R_i y) \text{ or } (y R_i x)$$

In other words, if  $X \setminus A = \emptyset$ , where  $A$  – set of points, such that  $(x R_i y)$  or  $(y R_i x)$ , then relation is complete.

1.  $i = 1$

$A = \{y \geq x - 1\} \cup \{y \leq x + 1\}$ .  $A$  is obviously a union of two rays, with non-empty intersection at  $y \in [x - 1, x + 1]$  thus it covers the whole  $\mathbb{R} \implies R_1$  is complete.

2.  $i = 2$

$A = \{y \leq x - 1\} \cup \{y \geq x + 1\}$ .  $A$  doesn't include points  $x, y \in X$ , such that  $y \in (x - 1, x + 1) \Rightarrow R_2$  is not complete.

### Assymetry.

Lets start with  $i = 1$ . If asymmetry holds, then the following should be true:

$$\begin{aligned} x R_1 y &\Rightarrow \neg(y R_1 x) \\ x \geq y - 1 &\Rightarrow \neg(y \geq x - 1) \\ y \leq x + 1 &\Rightarrow \neg(y \geq x - 1) \\ y \leq x + 1 &\not\Rightarrow y < x - 1 \end{aligned}$$

We see that  $y \leq x + 1$  doesn't imply  $y < x - 1$  (for example points  $x, y \in X$  such that  $x = y$ ), therefore  $R_1$  is not  
By the analogy for  $i = 2$ :

$$\begin{aligned} x R_2 y &\Rightarrow \neg(y R_2 x) \\ x \geq y + 1 &\Rightarrow \neg(y \geq x + 1) \\ y \leq x - 1 &\Rightarrow \neg(y \geq x + 1) \\ y \leq x - 1 &\Rightarrow \neg(y \geq x + 1) \\ y \leq x - 1 &\Rightarrow y < x + 1 \end{aligned}$$

We see that  $y \leq x - 1$  implies  $y < x + 1$ , therefore  $R_2$  is assymetric.

### Transitivity.

Lets take  $x, y, z \in X$ . From transitivity the following should apply:

1.  $i = 1$

$$\begin{aligned} (x R_1 y) \wedge (y R_1 z) &\Rightarrow x R_1 z \\ (x R_1 y) \wedge (y R_1 z) &\Leftrightarrow \begin{cases} y \leq x + 1 \\ z \leq y + 1 \end{cases} \Rightarrow z \leq x + 2 \not\Rightarrow x \geq z - 1 \Leftrightarrow x R_1 z \end{aligned}$$

Thus relation is not transitive.

2.  $i = 2$

$$\begin{aligned} (x R_2 y) \wedge (y R_2 z) &\Rightarrow x R_2 z \\ (x R_2 y) \wedge (y R_2 z) &\Leftrightarrow \begin{cases} y \leq x - 1 \\ y \geq z + 1 \end{cases} \Rightarrow z \leq x - 2 \Rightarrow x \geq z + 1 \Leftrightarrow x R_1 z \end{aligned}$$

Thus relation is transitive.

### Simmetric part.

So we define  $R_i^{\sim}$  as  $[(x R_i y) \wedge (y R_i x)]$

1.  $i = 1$

$$\begin{aligned} R_1^{\sim} &\Leftrightarrow [(x R_1 y) \wedge (y R_1 x)] \\ &\Leftrightarrow \begin{cases} y \leq x + 1 \\ x \leq y + 1 \end{cases} \Rightarrow 0 \leq 2 \quad \forall x, y \in X \\ &\Leftrightarrow x, y \in X \end{aligned}$$

2.  $i = 2$

$$\begin{aligned} R_2 &\iff [(x R_2 y) \wedge (y R_2 x)] \\ &\iff \begin{cases} x \geq y + 1 \\ y \geq x + 1 \end{cases} \implies 0 \geq 2 \quad \forall x, y \in X \\ &\iff x, y \notin X \end{aligned}$$

#### Problem 4

Consider a binary relation  $P$  on a set  $X$  that is asymmetric and negatively transitive. Define a further binary relation  $\succsim_p$  on  $X$  as follows:

$$x \succsim_p y \iff \neg(y P x)$$

Show that  $\succsim_p$  is a complete and transitive binary relation.

#### Solution.

*Completeness.*

$\forall x, y \in X$ , we have that either  $(x P y)$  or  $(y P x)$

Lets notice that  $x \succsim_p y \iff \neg(y P x)$  and  $y \succsim_p x \iff \neg(x P y)$ , applying asymmetry we get:

$$x \succsim_p y \iff \neg(y P x) \tag{2}$$

$$x \succsim_p y \iff x P y \tag{3}$$

So from point (3) we see, that  $x \succsim_p y [x P y]$  is actually the opposite for  $x \prec_p y [\neg(x P y)]$  and therefore relations are complete.

*Transitivity.*

$\forall x, y, z \in X$ , we have that if  $(x \succsim_p y)$  and  $(y \succsim_p z)$  then  $(x \succsim_p z)$

$$\begin{cases} \neg(y P x) \\ \neg(z P y) \end{cases} \xrightarrow{\text{neg. trans}} \neg(z P x) = \neg(x P z) \quad \forall x, y, z \in X$$

*Strict part.*

$x \succ_p y \xrightarrow{\text{definition}} \forall x, y \in X$ , we have that  $(x \succsim_p y)$  and not  $(y \succsim_p x)$

$$\begin{aligned} &\neg(y P x) \wedge x P y \\ &x P y \wedge x P y \iff x P y \end{aligned}$$

Therefore  $x \succ_p y \xrightarrow{\text{definition}} P$ .

#### Problem 5

Consider a choice structure  $(\mathcal{B}, C)$  on a set  $X$ . Suppose that any subset of  $X$  that has two elements belongs to  $\mathcal{B}$  (in addition to other sets, possibly). Let  $\succsim^*$  denote the revealed preference relation derived from  $C$ , as we denoted in class. Show that if  $C$  satisfies WARP, then for any  $x, y \in X$ , we have  $x \succsim^* y$  iff  $x \in C\{x, y\}$

**Solution.** Firstly, lets show, that there is rational revealed preference relation  $\succsim^*$ . The preference is rational once it satisfies transitivity and completeness:

1. *Completeness.* Since  $\mathcal{B}$  includes all pairs of elements in  $X$ , then  $\forall x, y \in X$  either  $C(\{x, y\})$  is either  $\{x\}$  or  $\{y\}$  or  $\{x, y\}$ . Thus, all elements from  $X$  can be sorted by preference.
2. *Transitivity.* Let  $x \succsim^* y$  and  $y \succsim^* z$  and lets consider  $B = \{x, y, z\}$ . If  $\succsim^*$  is rational, then  $x \in C(B)$ , because by the transitivity  $x \succsim^* z$ . Because of the completeness,  $C(B) \neq \emptyset$ . Lets suppose  $C(\cdot)$  is such that  $y \in C(B)$ , then, by the WARP and  $x \succsim^* y$ , the  $x$  must be in  $C(B)$ . Suppose instead  $z \in C(B)$ . Then, by the WARP and  $\succsim^*$  we conclude that  $y \in C(B)$ . From the previous logic, we conclude that  $x$  is also in  $C(B)$ .

Now, let's consider rationalized choice structure  $C^*(B, \succsim)$  as

$$C^*(B, \succsim) = \{x \in B : x \succsim y, \forall y \in B\} \text{ for all } B \text{ in } \mathcal{B}$$

and let's prove, that  $C(B) = C^*(B, \succsim)$  for all  $B$  in  $\mathcal{B}$ . First, let  $x \in C(B)$ , then  $x \succsim^* y, \forall y \in B$ . Therefore  $x \in C^*(B, \succsim^*)$ , which means, that  $C(B) \subset C^*(B, \succsim^*)$ . From the other hand, if  $x \in C^*(B, \succsim^*)$ , then  $x \succsim^* y$  for all  $y \in B$ . So, for each  $y \in B$  there must be some set  $B_y \in \mathcal{B}$ , which satisfies  $x, y \in B_y$  and  $x \in C(B_y)$ . Due to WARP,  $x \in C(B)$ , therefore  $C^*(B, \succsim) \subset C(B)$ . From this point it can be concluded, that  $C(B) = C^*(B, \succsim^*)$ .

And as  $X$  consists of all pairs of elementary objects, thus there is unique rationalizing preference relation  $\succsim^*$ , therefore we can get from choice structure and WARP to rational preference relation, or in other words

$$x \succsim^* y \iff x \in C\{x, y\}$$