Name: TeamGVP

College : Gayatri vidya parishad.

Contestants:

PVSM Praveen Kumar

S Bhargava Ram

S Raghuveer Sharma

Index :		• Linear Equation Solver [15]
Query Algo's :		Mobius Function [15]
Segment Tree Lazy [Min/Max/Sum in]		 Chinese Remainder Theorem [16]
Range	[3]	 Binomial Coefficients DP [16]
BIT with Binary Search	[3]	• Longest Common Substring [16]
BIT with Range Updates	[4]	0
Fenwick Tree	[4]	Geometry:
 Union Find 	[5]	Point ,Operator Overloading Func [16]ClockWise Cclockwise Rotations [16]
 Longest Increasing Sequence 	[5]	
Segment Tree without Lazy	[5]	Project Point on Line assingPoints [16]Project Point on LineSegment [16]
Matrix Exponentation	[6]	 Project Point on LineSegment [16] Distance of point to LineSegment [16]
 Square Root Decomposition 	[6]	• Parallel Lines Checkup [17]
		• Collinear Lines [17]
Graphs:		• LineSegment Intersection [17]
 Dijkstra 	[7]	• Points on SameSide of a Line [17]
 Floyd Warshall 	[8]	 Compute Center of Circle with 3 Points
 Euclidian Path 	[8]	 Determine Point Inside a Polygon [17]
 Single Source Shortest Path 	[8]	• Determine Point on Polygon [17]
 Fast Dijkstra 	[9]	• Circle Circle Intersection [17]
 BFS,DFS,Biparatite Checking 	[9]	• Convex Hull of N Points [18]
 Kruskal's, Prim's MST 	[10]	Utilities:
 Flood Fill 	[11]	• STL C++ [19,20]
 Connected Components 	[11]	• Bit Manipulations [21]
• LCA	[11]	• String Tokenizer in C++ [21]
 LCA Using RMQ 	[12]	• I/O Manip in C++ [21]
		• Java Big Integer [21]
Mathematics:		• Python Fast IO Methods [21]
 Modular Mul,Exp 	[12]	• Python Inbuilt Functions [21]
 Fermet Prime Test 	[13]	Missellaneaus
 Extended Euclid 	[13]	Miscellaneous:
 Miller Rabbin Test 	[13]	• KMP String Matching [22]
 Euler Mod Inverse 	[13]	Longest Common Prefix [22]Rectangle DP Sum [23]
 Pollard Rho Factorization 	[13]	
 Highest Prime Factor 	[13]	 Stable Marriage Problem [23] Longest Alternating Sequence [24]
 Prime Factorization 	[14]	• Longest Increasing Subsequence [24]
 Slow Prime Tests 	[14]	• Z-Algorithm [24]
 Sieve, Segmented Sieve 	[14]	• Top Down DP [24]
 ModInverse DP 	[15]	• Manacher's Algorithm [25]
• Euler Totient	[15]	5 [-]

Segment Tree Lazy: (Range Max/Min)

```
const int N=5;
#define MAX (1+(1<<6)) // Why? :D
#define inf 0x7fffffff
int arr[N+25];
int tree[MAX];
int lazy[MAX];
void build_tree(int node, int a, int b) {
   if(a > b) return;
   if(a == b) { tree[node] = arr[a]; return; }
   build_tree(node^*2, a, (a+b)/2);
   build_tree(node*2+1, 1+(a+b)/2, b
   tree[node] = max(tree[node*2],
   tree[node*2+1]);
}
void update_tree(int node, int a, int b, int i, int j, int
value) {
   if(lazy[node] != 0) {
                tree[node] += lazy[node];
                if(a != b) {
         lazy[node*2] += lazy[node];
   lazy[node*2+1] += lazy[node];
      }
      lazy[node] = 0;
   if(a > b \mid\mid a > j \mid\mid b < i) return;
   if(a \ge i \&\& b \le j) {
      tree[node] += value;
      if(a != b) {
         lazy[node*2] += value;
         lazy[node*2+1] += value;
      }
      return;
   update_tree(node*2, a, (a+b)/2, i, j, value);
   update_tree(1+node*2, 1+(a+b)/2, b, i, j, value);
   tree[node] = max(tree[node*2],
   tree[node*2+1]);
}
```

```
int query_tree(int node, int a, int b, int i, int j) {
   if(a > b \mid\mid a > j \mid\mid b < i) return -inf; //+inf in case
Min
   if(lazy[node] != 0) {
      tree[node] += lazy[node];
      if(a != b) {
          lazy[node*2] += lazy[node];
   lazy[node*2+1] += lazy[node];
      lazy[node] = 0;
   if(a \ge i \&\& b \le j)return tree[node];
   int q1 = query_tree(node*2, a, (a+b)/2, i, j);
   int q2 = query_tree(1+node*2, 1+(a+b)/2, b, i, j)
   int res = max(q1, q2)
   return res;
}
//update_tree(1, 1, N, L, R, x);
//query_tree(1, 1, N, L, R);
Note: For sum in range query in lazy propagation.
    1. Change tree[node] += lazy[node] to
        tree[node] += (b-a+1)*lazy[node]
    2. Change max(q1,q2) to q1+q2
    3. Out of range : -INF,+INF to 0.
```

BIT With Binary Search:

}

// Binary indexed tree supporting binary search.

```
struct BIT {
  int n;
  vector<int> bit:
// BIT can be thought of as having entries f[1], ..., f[n]
//0 Initialized.
  BIT(int n):n(n), bit(n+1) {}
    // returns f[1] + ... + f[idx-1] precondition idx <= n+1
  int read(int idx) {
     idx--;
     int res = 0:
     while (idx > 0) { res += bit[idx];
        idx = idx \& -idx;
     } return res;
// \text{ returns f}[\text{idx1}] + ... + \text{f}[\text{idx2-1}] \text{ precondition idx1} <= \text{idx2} <=
  int read2(int idx1, int idx2) {
     return read(idx2) - read(idx1);
```

```
// adds val to f[idx]
// precondition 1 <= idx <= n (there is no element 0!)</pre>
  void update(int idx, int val) {
    while (idx \leq n) { bit[idx] += val;
       idx += idx \& -idx; }
  // returns smallest positive idx such that read(idx) >= target
  int lower_bound(int target) {
    if (target <= 0) return 1;
    int pwr = 1; while (2*pwr \le n) pwr*=2;
    int idx = 0; int tot = 0;
    for (; pwr; pwr >= 1) {
       if (idx+pwr > n) continue;
       if (tot + bit[idx+pwr] < target) {</pre>
         tot += bit[idx+=pwr]; } }
    return idx+2:
}
  // returns smallest positive idx such that read(idx) > target
  int upper_bound(int target) {
    if (target < 0) return 1;
    int pwr = 1; while (2*pwr \le n) pwr*=2;
    int idx = 0; int tot = 0;
    for (; pwr; pwr >>= 1) {
       if (idx+pwr > n) continue;
      if (tot + bit[idx+pwr] <= target) {</pre>
         tot += bit[idx+=pwr];
      }
    return idx+2;
  }
};
```

BIT with Range Updates:

```
// BIT with range updates, inspired by Petr Mitrichev

struct BIT {
   int n;
   vector<int> slope;
   vector<int> intercept;

// BIT can be thought of as having entries f[1], ..., f[n]

// which are 0-initialized
   BIT(int n): n(n), slope(n+1), intercept(n+1) {}

// returns f[1] + ... + f[idx-1] precondition idx <= n+1</pre>
```

```
int query(int idx) {
    int m = 0, b = 0;
    for (int i = idx-1; i > 0; i -= i\&-i) {
       m += slope[i];
       b += intercept[i];
    return m*idx + b; }
// adds amt to f[i] for i in [idx1, idx2)
// precondition 1 \le idx1 \le idx2 \le n+1 (you can't //
update element 0)
void update(int idx1, int idx2, int amt) {
    for (int i = idx1; i \le n; i += i\&-i) {
       slope[i] += amt;
       intercept[i] -= idx1*amt;
    for (int i = idx2; i \le n; i += i\&-i) {
       slope[i] -= amt;
       intercept[i] += idx2*amt;
    }
};
```

```
struct FenwickTree (Normal BIT)
{
    typedef ll T;
    vector<T> v;
    FenwickTree(int n): v(n, 0) {}
    void add(int i, T x) {
        for(; i < (int)v.size(); i |= i+1) v[i] += x;
    }
    T sum(int i) { //[0, i)
        T r = 0;
        for(-- i; i >= 0; i = (i & (i+1)) - 1) r += v[i];
        return r;
    }
    T sum(int left, int right) { //[left, right)
        return sum(right) - sum(left);
    }
};
```

BIT Example:

};

int tree [(1 << LOGSZ)+1];

```
int N = (1 << LOGSZ);
        // add v to value at x
void set(int x, int v) {
 while(x \le N) { tree[x] += y; x += (x \& -x); }
}
        // get cumulative sum up to and including x
int get(int x) {
 int res = 0;
 while(x) { res += tree[x]; x -= (x \& -x); }
 return res;
}
        // get largest value with cumulative sum
less than or equal to x;
        // for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
 while(mask && idx < N) {
  int t = idx + mask;
  if(x \ge tree[t]) 
   idx = t;
   x = tree[t];
  mask >>= 1;
 }
 return idx;
struct UnionFind
{
   vector<int> data;
   void init(int n) { data.assign(n+1, -1); }
   bool unionSet(int x, int y) \{x = root(x); y = root(x)\}
root(y);
   if(x != y) \{ if(data[y] < data[x]) swap(x, y);
        data[x] += data[y]; data[y] = x; }
      return x != y;
   bool findSet(int x, int y) { return root(x) ==
root(y); }
   int root(int x) { return data[x] < 0? x: data[x]
        = root(data[x]); }
   int size(int x) { return -data[root(x)]; }
```

```
vector<int> longestIncreaseSequence(const
vector<int>& a) (NlogN) Binary Search + Greedy
   const int n = a.size();
   vector<int> A(n, INF);
   vector<int> id(n);
   for(int i = 0; i < n; ++i) {
      id[i] = lower_bound(all(A), a[i]) - A.begin();
      A[id[i]] = a[i];
   int m = *max_element(id.begin(), id.end());
   vector<int> b(m+1);
   for(int i = n-1; i >= 0; --i)
      if(id[i] == m) b[m--] = a[i];
   return b;
}
struct SegTree (Without Lazy.)
   vector<int> Tree:
       void init(int Arr[],int N) {
       Tree.assign(4*N,0); Build(Arr,1,N,1); }
   void Build(int Arr[],int L,int R,int POS)
    if(L==R) { Tree[POS]=Arr[L]; return; }
      int Mid=(L+R)/2;
               Build(Arr,L,Mid,2*POS);
               Build(Arr,Mid+1,R,2*POS+1);
    Tree[POS]=max(Tree[POS*2],Tree[POS*2+1]);
   }
   int Query(int L,int R,int QL,int QR,int POS)
     if(QL<=L && QR>=R) return Tree[POS];
     if (QL>R || QR<L) return -1;
     int Mid=(L+R)/2; return max(
               Query(L, Mid, QL,QR, 2*POS),
               Query(Mid+1,R,QL,QR,
               2*POS+1));
   }
```

```
void Update(int L,int R,int QL,int QR,int POS,int
VALUE)
{
         if(L>R || QL>R || QR<L) { return; }
         if(L=R) { if(QL==L && QR==R)
             Tree[POS]=VALUE; return; }
         int Mid=(L+R)/2;
             Update(L,Mid,QL,QR,2*POS,VALUE);
             Update(Mid+1,R,QL,QR,2*POS+1,VALUE);
             Tree[POS]=max(Tree[POS*2],Tree[POS*2+1]);
        }
};
Seg.Update(1,N,L,R,1,VALUE));
Seg.Query(1,N,L,R,1);</pre>
```

Matrix Exponentation:

```
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
const int N=4:
ll MOD=1e9+7;
void Copy(ll A[N][N],ll B[N][N]) {
  for(int i=0;i< N;i++)
    for(int j=0;j<N;j++)
      A[i][j]=B[i][j];
void Mul(ll A[N][N],ll B[N][N]) {
     ll C[N][N];
          for(int i=0;i< N;i++){
            for(int j=0;j<N;j++){
              C[i][i]=0;
              for(int k=0;k<N;k++) {
C[i][j] = (C[i][j] + (ll)A[i][k]*(ll)B[k][j])%MOD;
              } } }
     Copy(A,C);
}
void Power(ll A[N][N],int Exp) {
ll Ans[4][4]=\{\{1,0,0,0\},\{0,1,0,0\},\{0,0,1,0\},\{0,0,0,1\}\};
  while(Exp){
         if(Exp&1) Mul(Ans,A);
         Mul(A,A);
         Exp=Exp>>1;
      }
      Copy(A,Ans);
}
```

```
ll Solve(int K)
  if(K==0)
              return 0;
  if(K==1)
              return 1;
  if(K==2)
              return 2:
  ll A[4][4] = \{\{1,1,1,0\},\{1,0,0,0\},\{0,0,1,1\},\{0,0,0,1\}\};
  Power(A,K-1);
 return
((A[0][0]*1+A[0][1]*0+A[0][2]*1+A[0][3]*1)
%MOD);
}
The above code is for f(n) = f(n-1)+f(n-2) + n-1
                     | 1 1 1 0 | | F(N-1) |
|F(n)|
| F(n-1) |
                     | 1 0 0 0 | | F(N-2) |
                     | 0 0 1 1 | | N-1 |
| N |
  1 |
                     | 0 0 0 1 | | 1
```

Square Root Decomposition:

```
#include <bits/stdc++.h>
using namespace std;
   // Variable to represent block size. This is made
global
   // so compare() of sort can use it.
int block:
   // Structure to represent a query range
struct Query{
   int L, R;
}:
   // Function used to sort all queries so that all queries
   // of same block are arranged together and within a
block.
   // queries are sorted in increasing order of R values.
bool compare(Query x, Query y){
   // Different blocks, sort by block.
  if (x.L/block != y.L/block)
    return x.L/block < y.L/block;
   // Same block, sort by R value
  return x.R < y.R;
// Prints sum of all query ranges. m is number of queries
// n is size of array a \square.
void queryResults(int a[], int n, Query q[], int m){
  // Find block size
  block = (int)sqrt(n);
  // Sort all queries so that queries of same blocks
```

```
// are arranged together.
  sort(q, q + m, compare);
  // Initialize current L, current R and current sum
  int currL = 0, currR = 0;
  int currSum = 0;
  // Traverse through all queries
for (int i=0; i < m; i++){
    // L and R values of current range
    int L = q[i].L, R = q[i].R;
    // Remove extra elements of previous range. For
    // example if previous range is [0, 3] and current
    // range is [2, 5], then a[0] and a[1] are subtracted
    while (currL < L){
      currSum -= a[currL];
      currL++;
    }
    // Add Elements of current Range
    while (currL > L){
      currSum += a[currL-1];
      currL--;
    }
    while (currR <= R){
      currSum += a[currR];
      currR++;
    // Remove elements of previous range. For example
    // when previous range is [0, 10] and current range
    // is [3, 8], then a[9] and a[10] are subtracted
    while (currR > R+1){
      currSum -= a[currR-1];
      currR--;
// Print sum of current range
cout << "Sum of [" << L << ", " << R << "] is " << currSum
<< endl:
} //end for loop
}//End queryResults
// Driver program
int main()
  int a[] = \{1, 1, 2, 1, 3, 4, 5, 2, 8\};
  int n = sizeof(a)/sizeof(a[0]);
  Query q[] = \{\{0, 4\}, \{1, 3\}, \{2, 4\}\};
  int m = sizeof(q)/sizeof(q[0]);
  queryResults(a, n, q, m);
  return 0;
}
```

Graph Algorithms:

Dijkstra:

```
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
// This function runs Dijkstra's algorithm for single source
// shortest paths. No negative cycles allowed!
// Running time: O(|V|^2)
// INPUT: start, w[i][j] = cost of edge from i to j
// OUTPUT: dist[i] = min weight path from start to i
       prev[i] = previous node on the best path from the
//
            start node
//
void Dijkstra (const VVT &w, VT &dist, VI &prev, int
start){
 int n = w.size();
 VI found (n);
 prev = VI(n, -1);
 dist = VT(n, 1000000000);
 dist[start] = 0;
while (start != -1){
  found[start] = true;
  int best = -1:
  for (int k = 0; k < n; k++) if (!found[k]){
   if (dist[k] > dist[start] + w[start][k]){
    dist[k] = dist[start] + w[start][k];
    prev[k] = start;
   if (best == -1 || dist[k] < dist[best]) best = k;
  start = best;
```

FlovdWarshall:

```
// This function runs the Floyd-Warshall algorithm for all-
// shortest paths. Also handles negative edge weights.
Returns true
// if a negative weight cycle is found.
// Running time: O(|V|^3)
// INPUT: w[i][j] = weight of edge from i to j
// OUTPUT: w[i][j] = shortest path from i to j
// prev[i][j] = node before j on the best path starting at i
bool FloydWarshall (VVT &w, VVI &prev){
 int n = w.size(); prev = VVI (n, VI(n, -1));
 for (int k = 0; k < n; k++){ for (int i = 0; i < n;
i++){
   for (int j = 0; j < n; j++){ if (w[i][j] > w[i][k] +
w[k][j]){
      w[i][j] = w[i][k] + w[k][j];
      prev[i][j] = k; } } }
 // check for negative weight cycles
 for(int i=0;i< n;i++)
  if (w[i][i] < 0) return false;
 return true;
}
```

Euclidean Path:

```
struct Edge;
typedef list<Edge>::iterator iter;

struct Edge{
  int next_vertex;
  iter reverse_edge;
Edge(int next_vertex) :next_vertex(next_vertex) {}
};
const int max_vertices = ??;
int num_vertices;
list<Edge> adj[max_vertices]; // adjacency list
vector<int> path;
void find_path(int v){
    while(adj[v].size() > 0) {
    int vn = adj[v].front().next_vertex;
    adj[vn].erase(adj[v].front().reverse_edge);
    adj[v].pop_front();
```

```
find_path(vn);
}
path.push_back(v);
}
void add_edge(int a, int b){
  adj[a].push_front(Edge(b));
  iter ita = adj[a].begin();
  adj[b].push_front(Edge(a));
  iter itb = adj[b].begin();
  ita->reverse_edge = itb;
  itb->reverse_edge = ita;
}
```

Single Source Shortest Path:

```
struct Node{
 int node, cost;
 // i.e Node o is higher up in priority in the queue
 bool operator<(Node const &o) const {</pre>
  return cost > o.cost; }
};
vector<Node> AdjList[MAXN];
int dist[MAXN];
void ShortestPath(int src){
   Node start, curr, next;
   start.node = src:
   start.cost = 0;
   memset(dist, -1, sizeof(dist));
         // Initialize distance
   priority_queue<Node> q;
         // Initialize queue
   q.push(start);
   while(!q.empty()){
      curr = q.top();
      q.pop();
      if(dist[curr.node] == -1){
          dist[curr.node] = curr.cost;
          for(int i = 0; i < AdjList[curr.node].size();</pre>
i++){
             if(dist[AdjList[curr.node][i].node] != -
1) continue;
         next.node = AdjList[curr.node][i].node;
         next.cost = curr.cost +
AdjList[curr.node][i].cost;
          q.push(next);
```

```
}
}
scanf("%d %d %d",&x,&y,&c);
AdjList[x].push_back({y,c});
AdjList[y].push_back({x,c});
ShortestPath(src);
```

Fast Dijkstra:

```
typedef pair<int, int> PII;
vector< vector<PII> > edges(N);
edges[i].push_back(make_pair(dist, vertex));
        // note order of arguments here
  // use priority queue in which top element has the
"smallest" priority
priority_queue<PII, vector<PII>, greater<PII> > Q;
vector<int> dist(N, INF), dad(N, -1);
Q.push(make_pair(0, s));
dist[s] = 0;
 while (!Q.empty()) {
       PII p = Q.top();
       Q.pop();
       int here = p.second;
       if (dist[here] != p.first) continue;
   //vector<PII>::iterator it
   for (auto it=edges[here].begin();
                it!=edges[here].end(); it++) {
    if (dist[here] + it->first < dist[it->second]) {
          dist[it->second] = dist[here] + it->first;
          dad[it->second] = here;
         Q.push(make_pair(dist[it->ss], it->ss));
    }
         //ss = second
}
```

BFS:

```
vector<int> Graph[V];
int Color[V]; // {0,-1,1} = {White,Grey,Black} =
{Unvisited,Discovered,Explored}
int Dist[V],Pre[V];
//Set Dist[],Pre[] = INF , Color[],Dtime[],Ftime[] = 0
```

```
void BFS(int s){
    queue <int> q;    q.push(s);
    Color[s] = -1;    Dist[s]=Pre[s]=0;
    while(!q.empty()){        int u = q.front(); q.pop();
        for(int i=0;i<Graph[u].size();i++){
            int v = Graph[u][i];        if(Color[v]==0){
                Color[v] = -1;        Dist[v] = Dist[u]+1;
               Pre[v] = u; q.push(v);        }
        } Color[u] = 1;
    }
    return;
}</pre>
```

DFS:

```
void DFS(); void DFS_VISIT(int u);
list<int> TopoSort;
void DFS(int N){
   for(int u=1;u\leq N;u++){
    if(Color[u]==0){
      DFS_VISIT(u);
      }
}}
void DFS_VISIT(int u){
  Time++;
  DTime[u] = Time; Color[u] = -1;
  for(int i=0;i<Graph[u].size();i++){</pre>
    int v = Graph[u][i]; if(Color[v]==0){
       Pre[v]=u; DFS_VISIT(v); }
  Color[u]=1; Time++;
  FTime[u]=Time; TopoSort.push_front(u);
```

Biparatite Graph

```
vector<int> Graph[VMax];
int Color[VMax];
int Visited[VMax];
```

```
bool isBipartite(int s)
   Color[s] = 1;
   queue <int> q;
   q.push(s); Visited[s]=1;
   while (!q.empty()){
          int u = q.front();
          Visited[u]=1;
          q.pop();
          for (int i=0; i<Graph[u].size();i++) {
             int v = Graph[u][i];
                if (Color[v] == -1){
                Color[v] = 1 - Color[u];
                q.push(v);
                Visited[v]=1;
                else if (Color[v] == Color[u])
                    return false:
         }
   return true;
```

Kruskal:

```
//Kruskal
// inside int main()
vector< pair<int, ii> > EdgeList; // (weight, two
vertices) of the
edge
 for (int i = 0; i < E; i++) {
  scanf("%d %d %d", &u, &v, &w); // read the
triple: (u, v, w)
  EdgeList.push_back(make_pair(w, ii(u, v))); }
  sort(EdgeList.begin(), EdgeList.end()); // sort by
edge weight
O(ElogE)
  // note: pair object has built-in comparison function
int mst_cost = 0;
UnionFind UF(V);
                     // all V are disjoint sets initially
 for (int i = 0; i < E; i++) { // for each edge, O(E)
  pair<int, ii> front = EdgeList[i];
  if (!UF.isSameSet(front.second.first,
front.second.second)){
    mst_cost += front.first;
                                  // add the weight of
e to MST
   UF.unionSet(front.second.first,
front.second.second);
```

```
}

// note: the runtime cost of UFDS is very light
// note: the number of disjoint sets must eventually be 1
for a valid MST

printf("MST cost = %d (Kruskal's)\n", mst_cost);
```

PRIMS(ADDING VERTICES-SPANNING TREE)--(O(V+E)logV):

```
#include <iostream>
#include <vector>
#include <queue>
#include <functional>
#include <utility>
using namespace std;
const int MAX = 1e4 + 5;
typedef pair<long long, int> PII;
bool marked[MAX];
vector <PII> adj[MAX];
long long prim(int x){
priority_queue<PII, vector<PII>, greater<PII> > Q;
  long long minimumCost = 0;
  PII p;
  Q.push(make_pair(0, x));
  while(!Q.empty()) {
    // Select the edge with minimum weight
    p = Q.top();
    Q.pop();
    x = p.second;
    // Checking for cycle
    if(marked[x] == true)
      continue;
    minimumCost += p.first;
    marked[x] = true;
    for(int i = 0;i < adj[x].size();++i)
     y = adj[x][i].second;
      if(marked[y] == false)
        Q.push(adj[x][i]);
```

```
}
  return minimumCost;
int main(){
  int nodes, edges, x, y;
  long long weight, minimumCost;
  cin >> nodes >> edges;
  for(int i = 0; i < edges; ++i){
    cin >> x >> y >> weight;
    adj[x].push_back(make_pair(weight, y));
    adj[y].push_back(make_pair(weight, x));
  }
  // Selecting 1 as the starting node
  minimumCost = prim(1);
  cout << minimumCost << endl;</pre>
  return 0;
}
```

Flood fill:

```
int dr[] = \{1,1,0,-1,-1,-1,0,1\};
// trick to explore an implicit 2D grid
int dc[] = \{0,1,1,1,0,-1,-1,-1\};
// S,SE,E,NE,N,NW,W,SW neighbors
int floodfill(int r, int c, char c1, char c2) {
// returns the size of CC
  if (r < 0 || r >= R || c < 0 || c >= C) return 0;
// outside grid
  if (grid[r][c] != c1) return 0;
// does not have color c1
  int ans = 1;
// adds 1 to ans because vertex (r, c) has c1 as its color
  grid[r][c] = c2;
// now recolors vertex (r, c) to c2 to avoid .cycling!
  for (int d = 0; d < 8; d++)
   ans += floodfill(r + dr[d], c + dc[d], c1, c2);
  return ans; // the code is neat due to dr[] and dc[]
}
```

Connected components

```
numCC = 0;
dfs_num.assign(V, UNVISITED);
// sets all vertices' state to UNVISITED
```

```
for (int i = 0; i < V; i++)
// for each vertex i in [0..V-1]
if (dfs_num[i] == UNVISITED)
// if vertex i is not visited yet
    printf("CC %d:", ++numCC), dfs(i), printf("\n");</pre>
```

LCA (Least Common Ancestor)

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
// children[i] contains the children of node i
// A[i][j] is the 2^j-th ancestor of node i, or -1 if that
ancest or does not exist
int A[max_nodes][log_max_nodes+1];
// L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int L[max_nodes];
int lb(unsigned int n){
   if(n==0)
      return -1;
   int p = 0:
   if (n >= 1 << 16) \{ n >>= 16; p += 16; \}
   if (n >= 1 << 8) \{ n >= 8; p += 8; \}
   if (n >= 1 << 4) \{ n >= 4; p += 4; \}
   if (n >= 1 << 2) \{ n >>= 2; p += 2; \}
   if (n >= 1 << 1) \{ p += 1; \}
   return p;
void DFS(int i, int l){
   L[i] = l;
   for(int j = 0; j < children[i].size(); j++)
      DFS(children[i][j], l+1);
}
int LCA(int p, int q){
   // ensure node p is at least as deep as node q
   if(L[p] < L[q])
      swap(p, q);
   // "binary search" for the ancestor of node p
situated on the same level as q
   for(int i = log_num_nodes; i >= 0; i--)
      if(L[p] - (1 << i) >= L[q])
```

```
p = A[p][i];
      if(p == q)
         return p;
   // "binary search" for the LCA
   for(int i = log_num_nodes; i >= 0; i--)
      if(A[p][i] != -1 && A[p][i] != A[q][i])  {
         p = A[p][i];
         q = A[q][i];
      }
   return A[p][0];
}
int main(int argc,char* argv[])
{
// read num_nodes, the total number of nodes
log_num_nodes=lb(num_nodes);
for(int i = 0; i < num\_nodes; i++)
{
   int p;
   // read p, the parent of node i or -1 if node i is
the root
   A[i][0] = p;
   if(p!=-1)
      children[p].push_back(i);
   else
      root = i;
}
// precompute A using dynamic programming
for(int j = 1; j \le \log_n num_nodes; j++)
   for(int i = 0; i < num_nodes; i++)
      if(A[i][j-1]!=-1)
         A[i][j] = A[A[i][j-1]][j-1];
      else
         A[i][j] = -1;
   // precompute L
   DFS(root, 0);
   return 0;
}
```

LCA using RMO

```
//LCA(U,V) = MIN(E[H[U]....H[V])
int L[2*MAX_N], E[2*MAX_N], H[MAX_N], idx;
void dfs(int cur, int depth) {
```

```
H[cur] = idx;
E[idx] = cur;
L[idx++] = depth;
for (int i = 0; i < children[cur].size(); i++) {
    dfs(children[cur][i], depth+1);
    E[idx] = cur; // backtrack to current node
    L[idx++] = depth;
}

yoid buildRMQ(){
    idx= 0;
    memset(H,-1, sizeof H);
    dfs(0, 0);
    // we assume that the root is at index 0
}</pre>
```

<u>Arithematic:</u>

Modular Multiplication:

```
ull MulMod(ull A,ull B,ull Mod) (user if doubt.)
{
  ull x=0,y=A%Mod;
  while(B>0)
  {
    if(B&1)
        x=(x+y)%Mod;
    y=(y*2)%Mod;
    B=B/2;
  }
  return x%Mod;
}
```

Modular Exponent:

Fermet Prime Test:

```
bool FermatPrimeTest(ll N,ll Iterations)
{
   if(N==1)
      return 0;
   srand(time(NULL));

for(ll i=0;i<Iterations;i++)
   {
      ll a=rand()%(N-1)+1;
      if(PowerMod(a,N-1,N)!=1)
       return 0;
   }
   return 1;
}</pre>
```

Extended Euclid:

```
int d, x, y;
void extendedEuclid(int A, int B) {
  if(B == 0) {
     d = A;
     x = 1;     y = 0;
  }
  else {
     extendedEuclid(B, A%B);
     int temp = x;
     x = y;
     y = temp - (A/B)*y;
  }
}
extendedEuclid(16, 10);
cout << "The GCD of 16 and 10 is " << d << endl;
cout << "Coefficients x and y are "<< x << "and ";
cout << y << endl;</pre>
```

Mod Multiplicative Inverse: (IF M is composite)

```
int d,x,y;
int modInverse(int A, int M) {
    extendedEuclid(A,M);
    return (x%M+M)%M; //x may be negative
}
```

Miller Rabbin Test:

```
bool MillerRabinTest(ll N,ll Iterations)
{
    if(N<2) return 0; if(N==2) return 1;
    if(N%2==0) return 0;
    ll M = N-1;
    while(M%2==0) M=M/2;
    for(int i=0;i<Iterations;i++){
        ll a=rand()%(N-1)+1, K=M;
        ll Mod = ModExp(a,K,N);
        while(K!=N-1 && Mod!=1 && Mod!=N-1){
            Mod = MulMod(Mod,Mod,N);
            K=K*2;
        }
        if(Mod!=N-1 && K%2==0) return 0;
    }
    return 1;
}</pre>
```

Pollard Rho Factorization:

```
Il pollardRho(ll n) {
    if(n<=1) return -1;
    if(n%2==0)
    return 2;
    srand (time(NULL));
    ll x, y, g=1, a;
    x = rand() % n + 1;
    y = x;
    a = rand() % n + 1;
    while(g==1) {
        x = ((x*x) + a)%n;
        y = ((y*y) + a)%n;
        y = ((y*y) + a)%n;
        g = gcd(abs(x - y), n);
    }
    return g;}</pre>
```

Highest Prime Factor:

```
bool Prime[MAX+25];
int HPF[MAX+25]; //Highest Prime Factor
void sieve(){
  int i,j;
  memset(Prime,1,sizeof(Prime));
  Prime[0]=Prime[1]=0;
```

```
for(i=2;i<=sqrt(MAX);i+=2) Prime[i]=0;
    HPF[i] =2;
    for(i=3;i<=sqrt(MAX);i+=2) {
        if(Prime[i]){
            HPF[i]=i;
            K=i;
            for(j=i*i;j<=MAX;j+=i){
                Prime[j]=0;
            }
        }
    }
}</pre>
```

Prime Factorization:

```
//returns 60 = (2,2)(3,1)(5,1)
auto ExpCal(ll N)
{
    vector< pair<int,int> > Fac;
    ll NN = N,Pow,prime;
    while(NN>1){
        Pow=0;
        prime=HPF[NN];
        while(NN%prime==0){
            NN=NN/prime;
            Pow++;
        }
        Fac.push_back(make_pair(prime,Pow));
    }
    return Fac;
}
```

Small Number Prime Tests:

```
bool PrimeI(ll x){
    if(x==2 || x==3) return 1;
    if(x<=1 || x%2==0 || x%3==0) return 0;
    if(x%6!=1 && x%6!=5) return 0;
    ll i=5;
    while(i*i<=x){
        if(x%i==0 || x%(i+2)==0) return 0;
        i+=6;
    }
    return 1;
}</pre>
```

```
bool PrimeII(ll x){
    if(x<=1)return 0;
    if(x<=3)return 1;
    if(x%6==1||x%6==5){
        ll y=sqrt(x);
        for(ll i=2;i<=y;i++)
            if(x%i==0)
        return 0;
    return 1;
    }
    return 0;
}</pre>
```

```
#define MAX 100000005

#define setval(a,val) memset(a,val,sizeof(a))
bitset<MAX+25>Prime;

//bool Prime[MAX] gave RunTime error in
//HackerRank but accepted in Codechef and Spoj.
//Have to Check this in Practice Contest.
```

Sieve of Erathothesis:

```
void Sieve(){
  int i,j,POS; Prime[0]=Prime[1]=1;
  for(i=2;i<=180;i++)
    if(!Prime[i])
    for(j=2;(POS=j*i)<=32000;j++)
        Prime[POS]=1;
}</pre>
```

Segmented Sieve:

```
void SegmentSieve(int L,int R){
   int lim = sqrt(R);
   int i, j;
   for (i = 2; i <= lim; i++) {
      if (!Prime[i]) {
        for (j = L - L%i; j <= R; j += i)
            if (j>=L && !Prime[j] && j!=i)
            Prime[j] = 1;
      }
   }
}
If(R>200)
SegmentSieve(L,R);
```

Modular Inverse DP:

```
void DpModInv(int N){
  int i;
  ModInverse[1]=1;
  for(i=2;i<N;i++){
      ModInverse[i]=(-ll)(MOD/i)*
      (ll)ModInverse[MOD%i]) % MOD + MOD;
  }
}</pre>
```

Mobius Function Precompute:

Totient Function:

```
#include <stdlib.h>
// This code took less than 0.5s to calculate with
MAX = 10^7
#define MAX 10000000
int phi[MAX];
bool pr[MAX];
void totient(){
  for(int i = 0; i < MAX; i++){
    phi[i] = i;
    pr[i] = true;
}

for(int i = 2; i < MAX; i++)
  if(pr[i]){
    for(int j = i; j < MAX; j+=i){
        pr[j] = false;
}</pre>
```

```
phi[j] = phi[j] - (phi[j] / i);
}
pr[i] = true;
}
```

The sum of all values of Totient Function of all divisors of N is equal to N.

$$\varphi(p^k) = p^k - p^{k-1} = p^{k-1}(p-1)$$
 // p is prime.

```
\begin{split} \varphi(n) &= \varphi(p_1^{k_1}) \varphi(p_2^{k_2}) \cdots \varphi(p_r^{k_r}) \\ &= p_1^{k_1} \left( 1 - \frac{1}{p_1} \right) p_2^{k_2} \left( 1 - \frac{1}{p_2} \right) \cdots p_r^{k_r} \left( 1 - \frac{1}{p_r} \right) \\ &= p_1^{k_1} p_2^{k_2} \cdots p_r^{k_r} \left( 1 - \frac{1}{p_1} \right) \left( 1 - \frac{1}{p_2} \right) \cdots \left( 1 - \frac{1}{p_r} \right) \\ &= n \left( 1 - \frac{1}{p_1} \right) \left( 1 - \frac{1}{p_2} \right) \cdots \left( 1 - \frac{1}{p_r} \right). \end{split}
```

GCD(A,B) == 1 then $Totient(A) \times Totient(B) = Totient(A \cdot B)$.

Linear Equation Solver:

```
// finds all solutions to ax = b (mod n)
VI mod_linear_equation_solver(int a, int b, int n) {
    int x, y;
    VI solutions;
    int d = extended_euclid(a, n, x, y);
    if (!(b%d)) {
        x = mod (x*(b/d), n);
        for (int i = 0; i < d; i++)
            solutions.push_back(mod(x + i*(n/d), n));
    }
    return solutions;
}</pre>
```

Chinese Remainder Theorem 1:

```
// Chinese remainder theorem (special case): find z such
that
// z % x = a, z % y = b. Here, z is unique modulo M =
lcm(x,y).
// Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int
b) {
   int s, t;
   int d = extended_euclid(x, y, s, t);
   if (a%d!= b%d) return make_pair(0, -1);
   return make_pair(mod(s*b*x+t*a*y,x*y)/d,
   x*y/d);
}
```

Chinese Remaider Theorem 2:

```
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
    PII ret = make_pair(a[0], x[0]);
    for (int i = 1; i < x.size(); i++) {
        ret = chinese_remainder_theorem(ret.first, ret.second, x[i], a[i]);
        if (ret.second == -1) break;
    }
    return ret;
}</pre>
```

Longest Common Substring:

```
int LCSubStr(char *X, char *Y, int m, int n){
    // Create a table to store lengths of longest
    common suffixes of
    // substrings. Notethat LCSuff[i][j] contains
length of longest
    // common suffix of X[0..i-1] and Y[0..j-1]. The
first row and
    // first column entries have no logical meaning,
they are used only
    // for simplicity of program
    int LCSuff[m+1][n+1];
```

```
int result = 0; // To store length of the longest
common substring
    /* Following steps build LCSuff[m+1][n+1] in
bottom up fashion. */
    for (int i=0; i<=m; i++) {
        for (int j=0; j<=n; j++){
            if (i == 0 || j == 0)
                LCSuff[i][j] = 0;
        else if (X[i-1] == Y[j-1]){
               LCSuff[i][j] = LCSuff[i-1][j-1] + 1;
            result = max(result, LCSuff[i][j]);
        }
        else LCSuff[i][j] = 0;
    }
} return result;
}</pre>
```

```
int binomialCoeff(int n, int k)
{
  int C[n+1][k+1];  int i, j;
  for (i = 0; i <= n; i++){
    for (j = 0; j <= min(i, k); j++){
      if (j == 0 || j == i)
            C[i][j] = 1;
      else
            C[i][j] = C[i-1][j-1] + C[i-1][j];
    }
  }
  return C[n][k];</pre>
```

Geometry:

```
double INF = 1e100;
double EPS = 1e-12;

struct PT {
  double x, y;
  PT() {}
  PT(double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y) {}
  PT operator + (const PT &p) const { return
  PT(x+p.x, y+p.y); }
```

```
PT operator - (const PT &p) const { return PT(x-
p.x, y-p.y); }
PT operator * (double c) const { return PT(x*c,
y*c ); }
PT operator / (double c) const { return PT(x/c,
y/c ); }
};
double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
os << "(" << p.x << "," << p.y << ")";
}
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
return PT(p.x*cos(t)-p.y*sin(t),
p.x*sin(t)+p.y*cos(t));
}
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
}
// project point c onto line segment through a and
b
// if the projection doesn't lie on the segment,
returns closest vertex
PT ProjectPointSegment(PT a, PT b, PT c) {
        double r = dot(b-a,b-a);
        if (fabs(r) < EPS) return a;
        r = dot(c-a, b-a)/r;
        if (r < 0) return a;
        if (r > 1) return b;
        return a + (b-a)*r;
}
// compute distance from c to segment between a
and b
double DistancePointSegment(PT a, PT b, PT c) {
```

```
return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
}
// determine if lines from a to b and c to d are
parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
}
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
         && fabs(cross(a-b, a-c)) < EPS
         && fabs(cross(c-d, c-a)) < EPS;
}
// determine if line segment from a to b intersects
with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
      if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
         dist2(b, c) < EPS \mid\mid dist2(b, d) < EPS) return true;
      if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-a, d-b) > 0 \&\& dot
b, d-b) > 0
        return false;
      return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
false:
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
}
// determine if c and d are on same side of line
passing through a and b
bool OnSameSide(PT a, PT b, PT c, PT d) {
  return cross(c-a, c-b) * cross(d-a, d-b) > 0;
}
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  b=(a+b)/2;
  c=(a+c)/2;
  return ComputeLineIntersection(b,
b+RotateCW90(a-b), c, c+RotateCW90(a-c));
}
```

```
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0:
 for (int i = 0; i < p.size(); i++){
  int j = (i+1)\%p.size();
  if ((p[i].y \le q.y \& q.y \le p[j].y ||
   p[j].y \le q.y & q.y \le p[i].y) & 
   q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) /
(p[j].y - p[i].y)
   c = !c;
 }
 return c;
}
// determine if point is on the boundary of a
bool PointOnPolygon(const vector<PT> &p, PT q)
 for (int i = 0; i < p.size(); i++)
  if (dist2(ProjectPointSegment(p[i],
p[(i+1)\%p.size()], q), q) < EPS)
   return true:
  return false;
}
// compute intersection of circle centered at a with
radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b,
double r, double R) {
 vector<PT> ret;
 double d = sqrt(dist2(a, b));
 if (d > r + R || d + min(r, R) < max(r, R)) return ret;
 double x = (d*d-R*R+r*r)/(2*d);
 double y = sqrt(r*r-x*x);
 PT v = (b-a)/d;
 ret.push_back(a+v*x + RotateCCW90(v)*y);
 if (y > 0)
  ret.push_back(a+v*x - RotateCCW90(v)*y);
 return ret;
}
```

Convex Hull:

```
// A C++ program to find convex hull of a set of points.
// http://www.geeksforgeeks.org/orientation-3-
ordered-points/
// for explanation of orientation()
#include <iostream>
#include <stack>
#include <stdlib.h>
using namespace std;
struct Point{ int x, y; };
// A globle point needed for sorting points with
reference
// to the first point Used in compare function of
qsort()
Point p0;
// A utility function to find next to top in a stack
Point nextToTop(stack<Point> &S){
  Point p = S.top(); S.pop();
  Point res = S.top(); S.push(p);
  return res:
}
// A utility function to swap two points
int swap(Point &p1, Point &p2){
  Point temp = p1; p1 = p2;
  p2 = temp;
// A utility function to return square of distance
// between p1 and p2
int distSq(Point p1, Point p2){
  return (p1.x - p2.x)*(p1.x - p2.x) +
     (p1.y - p2.y)*(p1.y - p2.y);
// To find orientation of ordered triplet (p, q, r).
// The function returns following values
// 0 --> p, q and r are colinear
// 1 --> Clockwise
// 2 --> Counterclockwise
int orientation(Point p, Point q, Point r){
  int val = (q.y - p.y) * (r.x - q.x) -
        (q.x - p.x) * (r.y - q.y);
```

```
if (val == 0) return 0; // colinear
  return (val > 0)? 1: 2; // clock or counterclock wise
}
// A function used by library function qsort() to sort
an array of
// points with respect to the first point
int compare(const void *vp1, const void *vp2){
 Point *p1 = (Point *)vp1;
 Point p2 = (Point *)vp2;
 // Find orientation
 int o = orientation(p0, *p1, *p2);
 if (o == 0)
   return (distSq(p0, *p2) >= distSq(p0, *p1))? -1:1;
 return (o == 2)? -1: 1;
}
// Prints convex hull of a set of n points.
void convexHull(Point points[], int n){
 // Find the bottommost point
 int ymin = points[0].y, min = 0;
 for (int i = 1; i < n; i++){
  int y = points[i].y;
  // Pick the bottom-most or chose the left
  // most point in case of tie
  if ((y < ymin) \mid | (ymin == y \&\&
     points[i].x < points[min].x))
    ymin = points[i].y, min = i;
 // Place the bottom-most point at first position
 swap(points[0], points[min]);
 // Sort n-1 points with respect to the first point.
 // A point p1 comes before p2 in sorted ouput if p2
 // has larger polar angle (in counterclockwise
 // direction) than p1
 p0 = points[0];
 qsort(&points[1], n-1, sizeof(Point), compare);
 // If two or more points make same angle with p0,
 // Remove all but the one that is farthest from p0
 // Remember that, in above sorting, our criteria was
 // to keep the farthest point at the end when more
 // one points have same angle.
 int m = 1; // Initialize size of modified array
 for (int i=1; i<n; i++){
    // Keep removing i while angle of i and i+1 is same
```

```
// with respect to p0
    while (i < n-1 && orientation(p0, points[i],
                     points[i+1]) == 0
     i++;
    points[m] = points[i];
    m++; // Update size of modified array
 // If modified array of points has less than 3 points,
 // convex hull is not possible
 if (m < 3) return;
 // Create an empty stack and push first three points
 // to it.
 stack<Point>S;
 S.push(points[0]);
 S.push(points[1]);
 S.push(points[2]);
 // Process remaining n-3 points
 for (int i = 3; i < m; i++){
   // Keep removing top while the angle formed by
   // points next-to-top, top, and points[i] makes
   // a non-left turn
   while (orientation(nextToTop(S), S.top(), points[i])
!= 2)
     S.pop();
   S.push(points[i]);
 // Now stack has the output points, print contents
of stack
 while (!S.empty()){
    Point p = S.top();
    cout << "(" << p.x << ", " << p.y <<")" << endl;
    S.pop();
 }
}
// Driver program to test above functions
int main(){
  Point points[] = \{\{0, 3\}, \{1, 1\}, \{2, 2\}, \{4, 4\},
             \{0, 0\}, \{1, 2\}, \{3, 1\}, \{3, 3\}\};
  int n = sizeof(points)/sizeof(points[0]);
  convexHull(points, n);
  return 0;}
```

```
Utilities:
                                                                   5. s.substr(pos)
                                                                                      s.substr(pos,length)
                                                                   6. s.swap(s2)
                                                                                      7. s.find(string) s.find(string,pos)
queue < type > q;
1. q.push(element)
                      2. q.pop()
                                                                   8. s.replace(pos,len,string) 9. we can sort a string
3. q.front()
                      4. q.back()
                                                                   char array to string --> string str(ch_arr);
5. q.empty()(!)
                      6. q.size()
                                                                   string to char array --> str.c_str()
7. q.swap(q2)
                                                                   vector < type > vi;
stack <int> sk;
                                                                   1. vi.assign(val,no.of)
                                                                                               vi.assign(it first,it last)
                                                                                      vi.assign(arr, arr+N)
1. sk.empty() (!)
                      2. sk.pop()
                                                                   2. vi.clear()
3. sk.push(element) 4. sk.top()
                                                                                         4. vi.push_back()
                                                                   3. vi.empty()
                                                                                  (!)
5. sk.size()
                      6. sk.swap(sk2)
                                                                   5. vi.pop_back()
                                                                   6. vi.erase(it)
                                                                                         vi.erase(it first,it last)
priority_queue < type , vector < type > , cmp > pq;
                                                                   7. vi.front()
                                                                                      8. vi.back()
// cmp :
            //opposite
                                                                   9. vi.insert(it pos,val,no.of) vi.insert(it pos,it
struct cmp {
                                                                   first.itlast)
   bool operator()(const type &aa,const type &bb){
                                                                               vi.insert(it pos,arr, arr+N)
         return aa.ff < b.ff;
                                                                   10. operator " = " ( vector1 = vector2 )
      }
                                                                   11. operator " [] " (vi[i]) 12. vi.begin()
};
                                                                   13. vi.end() (!)
                                                                                     vector<typ>::iterator it
                                                                   14. vi.rbegin()
                                                                                      vector<typ>::reverse_iterator it
                                                                   15. vi.rend()
1. pq.push(element) 2. pq.pop()
                                                                                         16. vi.size()
                                                                   17. vi.swap(vii) 18. lower_bound(all(vi),element)
3. pq.top()
                      4. pq.empty() (!)
                                                                   19. upper_bound(all(vi),element) 20. vi.find()
5. pq.size()
                      6. pq.swap(pq2)
                                                                   21. set<int> uniq(vi.begin(),vi.end()) // uniq will
                                                                   contain unique elements
1. stoi(string)
                   //string to integer
2. to_string(integer) // integer to string
3. // normal :
   fill(a,0); fill(a,-1); fill(a,0x3f); //INF
                                                                   // comparision :
                                                                   bool cmp(type aa,type bb){
4. //bool:
   fill(a,0); fill(a,1);
                                                                      return aa.ff < bb.ff;
                                                                                               //direct
5.
                                                                  }
set_intersection(first.begin(),first.end(),second.begi
n(),s econd.end(),inserter(ans,ans.begin()));
                                                                   map < key, value > mp;
                                                                                               //different elements
                                                                   multimap < key, value > mmp;
6. // erase from back
                                                                                                     //same elements
minheap.erase(--minheap.end()); // c++ 4.3.2
                                                                   unordered_map < key, value > u_mp;
minheap.erase(std::prev(minheap.end())); // c++11
                                                                   unordered_multimap < key , value > um_mp;
7. setbase, setfill, setw, setprecision
                                                                   1. operator " = "
                                                                                         2. operator " [] "
8.numeric limits:
                                                                   3. mp.begin()
                                                                                         4. mp.end()
                                                                                                            // same to
   numeric_limits<int>::max()
                                                                   rbegin() and rend()
   numeric_limits<int>::min() ... etc..
                                                                   5. mp.insert(element)
                                                                                               6. mp.empty() (!)
                                                                   7. mp.size()
                                                                                      8. mp.erase(it)
                                                                      mp.erase(it first,it last)
string s;
1. operator " + " 2. operator " = "
                                                                   9. mp.clear()
                                                                                         10. mp.count(element)
3. operator " [] " 4. s.size()
                                                                   11. mp.find(element)
                                                                                               //!= mp.end()
```

// returns

```
iterators
13. mp.upper_bound(element) 14. mp.swap(mp2)
set < type, cmp > st;
multiset < type, cmp > mst;
unordered_set < type, cmp > u_st;
unordered_multiset < type, cmp > um_st;
1. st.begin()
                  2. st.end() (!)
3. st.size()
                  4. st.empty()
5. st.find() //!= st.end() 6. st.insert(type)
7. st.erase(it)
                  st.erase(it first,it last)
8. st.lower_bound(element)
                               O(logN)
9. st.upper_bound(element)
                               O(logN)
10 st.swap(st2)
                      11. st.clear()
12. st.count(element)
// cmp:
   struct cmp {
      bool operator()(const int &aa, const int &bb)
{
         return aa < bb;
      }
   };
//inbuit cmp functions.
greater<int>()
   ex: sort (arr, arr+5, std::greater<int>())
less<int>()
deque < type > dq;
1. operator " = "
2. operator " [] "
3. dq.clear()
4. dq.size()
5. dq.begin()
6. dq.end()
7. dq.erase(it)
                  dq.erase(it first,it last)
8. dq.push_back(element)
9. dq.pop_back()
10. dq.push_front(element)
11. dq.pop_front()
12. dq.empty()
Set a bit
Use the bitwise OR (|) operator.
```

12. mp.lower_bound(element)

```
number |= 1 << x; // Sets bit x
```

Clear a bit

```
Use the bitwise AND (&) and NOT (\sim) operators. number &= \sim(1 << x); // Clears bit x number &= !(1 << x); // Also works
```

Toggle a bit

```
Use the bitwise XOR (^) operator.
number ^= 1 << x; // Toggles bit x
```

Check a bit

```
Use the bitwise AND (&) operator.
bit = number & (1 << x); // Store value of bit x
```

```
#define isOn(S, j) (S & (1 << j))
#define setBit(S, j) (S |= (1 << j))
#define clearBit(S, j) (S &= \sim(1 << j))
#define toggleBit(S, j) (S ^= (1 << j))
#define lowBit(S) (S & (-S))
#define setAll(S, n) (S = (1 << n) - 1)
#define modulo(x, N) ((x) & (N - 1))
// returns x % N, where N is a power of 2

#define isPowerOfTwo(x) ((x & (x - 1)) == 0)
#define nearestPowerOfTwo(x)
((int)pow(2.0, (int)((log((double)x) / log(2.0)) + 0.5)))
```

String Tokenizer:

```
setfill - cout << setfill ('x') << setw (5); cout << 77 << endl;
----- Prints xxx77
setprecision - cout << setprecision (4) << f << endl;
----- Prints x.xxxx
```

• <u>lava</u>: <u>Big Integer</u>

```
import java.math.BigInteger
BigInteger sum = BigInteger.ZERO; //zero Constant
BigInteger two = new BigInteger("2");
sum = sum.add(V); //addition
BigInteger x = BigInteger(sc.next(), b);
// the second parameter is base.
```

BigInteger gcd_pq = p.gcd(q); //p and q are also //BigIntegers

Methods:

```
abs(x) clearBit(p) modInverse(m)
add(x) compareTo(x) modPower(e,m)
and(x) divide(x)nextProbablePrime(x,Random)
bitCount(x)gcd(x) probablePrime(x,Random r)
bitLength(x) mod(m)
```

Python:

Fast IO:

from sys import stdin, stdout

```
T = int(stdin.readline())
n,q = map(int,stdin.readline().split())
stdout.write(str(ans)+" ")
stdout.write("\n")
Int('string',base) oct(number) hex(number)
    bin(number) isinstance(data/variable,datatype)
```

Catalan numbers :

```
Cat(m) = ((2m \times (2m-1))/((m+1) \times m)) \times Cat(m-1).
```

```
Import itertools print (list((itertools.product("abc",repeat=3)))) ('a', 'a', 'a'), ('a', 'a', 'b'), ('a', 'a', 'c'), ('a', 'b', 'a').......]
```

```
print (list((itertools.permutations("abc",r=2))))
[('a', 'b'), ('a', 'c'), ('b', 'a'), ('b', 'c'), ('c', 'a'), ('c', 'b')]
```

```
Similarily,
combinations('ABCD', 2)
combinations with replacement('ABCD', 2)
```

Dynamic Programming:

Knuth-Morris-Prath algorithm

```
#define MAX_N 100010
char T[MAX_N], P[MAX_N];
// T = text, P = pattern
int b[MAX_N], n, m;
// b = back table, n = length of T, m = length of P
// call this before calling kmpSearch()
void kmpPreprocess() {
int i = 0, j = -1; b[0] = -1; // starting values
 while (i < m) \{ // \text{ pre-process the pattern string } P
 while (i \ge 0 \&\& P[i] != P[i]) i = b[i];
           // different, reset j using b
  i++; j++; // if same, advance both pointers
 b[i] = j; // observe i=8, 9, 10,11,12,13 with j = 0, 1, 2, 3, 4,
} // in the example of P = "SEVENTY SEVEN" above
void kmpSearch() {
     // this is similar as kmpPreprocess(), but on string T
     int i = 0, j = 0; // starting values
     while (i < n) \{ // \text{ search through string } T
     while (i \ge 0 \&\& T[i] != P[i]) i = b[i];
     // different, reset j using b
     i++; j++; // if same, advance both pointers
     if (j == m) {
```

printf("P is found at index %d in T\n", i - j); j = b[i]; // prepare j for the next possible match

Longest Common Prefix

}

```
void computeLCP() {
  int i, L;
  Phi[SA[0]] = -1; // default value
  for (i = 1; i < n; i++) // compute Phi in O(n)
     Phi[SA[i]] = SA[i-1];</pre>
```

// a match found when j == m

Rectangle Sum DP

```
void BuildSum(){
    for(int i = 1; i <= NAX; i++)
        for(int j = 1; j <= NAX; j++)
            Sum[i][j]+=Sum[i][j-1]+Arr[i][j];
    for(int j=1;j<=NAX;j++)
        for(int i=1;i<=NAX;i++)
            Sum[i][j]=Sum[i-1][j]+Sum[i][j];
    return;
}
Sum of Rectangle (x1,y1) and (x2,y2) = Sum[x2][y2] -
Sum[x1-1][y2] - Sum[x2][y1-1] + Sum[x1-1][y1-1].</pre>
```

Stable Marriage Problem:

```
// Gale-Shapley algorithm for the stable marriage problem.
// madj[i][j] is the jth highest ranked woman for man i.
// fpref[i][j] is the rank woman i assigns to man j.
// Returns a pair of vectors (mpart, fpart), where mpart[i]
gives the partner of man i, and fpart is analogous

pair<vector<int>, vector<int> > stable_marriage
(vector<vector<int> >& madj, vector<vector<int> >&
fpref) {
    int n = madj.size();
    vector<int> mpart(n, -1), fpart(n, -1);
    vector<int> midx(n);
    queue<int> mfree;
    for (int i = 0; i < n; i++) {
        mfree.push(i);
    }
}</pre>
```

```
while (!mfree.empty()) {
    int m = mfree.front(); mfree.pop();
    int f = madj[m][midx[m]++];
    if (fpart[f] == -1) {
        mpart[m] = f; fpart[f] = m;
    } else if (fpref[f][m] < fpref[f][fpart[f]]) {
        mpart[fpart[f]] = -1; mfree.push(fpart[f]);
        mpart[m] = f; fpart[f] = m;
    } else {
        mfree.push(m);
    }
}
return make_pair(mpart, fpart);
}</pre>
```

Longest Alternating Sequence:

```
#include<bits/stdc++.h>
using namespace::std;
typedef long long II;
II N;
II Arr[6000];
II lis[6000];
II ALTSEQ(II Arr[], II N){
  II i, j, Max = 0;
  for (i = 0; i < N; i++)
     lis[i] = 1;
  for (i = 1; i < N; i++)
     for (j = 0; j < i; j++)
       if (abs(Arr[i]) > abs(Arr[j]) && lis[i] < lis[j] + 1
              && 1LL*Arr[i]*Arr[j] < 0)
          lis[i] = lis[j] + 1;
  for (i = 0; i < N; i++)
     if (Max < lis[i])
       Max = lis[i];
  return Max;
Input: 1 2 -2 -3 5 -7 -8 10
Output: 5 (1-25-810)
```

Kadane Algorithm:

```
#include<iostream>
using namespace std;
int maxSubArraySum(int a[], int size){
   int max_so_far = a[0];
   int curr max = a[0];
   int curStartIndex = 0;
   int maxStartIndex;
   int maxEndIndex;
   for (int i = 1; i < size; i++){
      curr_max = max(a[i], curr_max+a[i]);
      if(curr_max > max_so_far){
          max_so_far = curr_max;
          maxStartIndex = curStartIndex;
          maxEndIndex = i;
   if(curr max < 0){
       curStartIndex = i+1;
   }
}
   cout << "Max Subarray starts at : " <<
   maxStartIndex << " and ends at :" << maxEndIndex
   << " and sum is : " << max_so_far << endl;
   return max so far;
}
```

Longest Increasing SubSequence (N^2) DP:

```
int increasingSubsequece(int n) {
  int maxLength = 1, bestEnd = 0;
  dp[0] = 1;
  prev[0] = -1;
  for (int i=1; i<n; ++i) {
    dp[i] = 1;
    prev[i] = -1;
    for (int j=i-1; j>=0; --j) {
       if (dp[j] + 1 > dp[i] && a[j] < a[i]) {
         dp[i] = dp[j] + 1;
         prev[i] = j;
       }
    }
    if (dp[i] > maxLength) {
       bestEnd = i;
       maxLength = dp[i];
```

```
}
}
// constructing the longest increasing subsequence
from prev[] array
/*
    vector <int> lis;
    for (int j=0; j<maxLength; ++j)
    {
        lis.push_back(a[bestEnd]);
        bestEnd = prev[bestEnd];
    }
    reverse(lis.begin(), lis.end());
    for (int i=0; i<lis.size(); ++i)
        printf("%d ", lis[i]);
    printf("\n");
    */
    return maxLength;
}</pre>
```

Z- Algorithm: O(N)

```
Pattern: "abc"
Text: "xabcabzabc"
vector<int> z_function(string s, int n) {
   vector<int> z(n);
   for (int i=1, l=0, r=0; i<n; ++i) {
       if (i <= r) {
           z[i] = min(r-i+1, z[i-l]);
       while (i+z[i] < n \&\& s[z[i]] == s[i+z[i]]) \{
           ++z[i];
       if (i+z[i]-1 > r) {
           l = i, r = i + z[i] - 1;
       }
   }
   return z;
z = z_function("abc$xabcabzabc",14);
z = [0\ 0\ 0\ (0)\ 0\ 3\ 0\ 0\ 2\ 0\ 0\ 3\ 0\ 0]
where len(pat) = z[i], it matches from that position.
```

Manacher's Algorithm [Longest Palindrome Substr]

#include <stdio.h>
#include <string.h>
char text[100];

```
void findLongestPalindromicString(){
      int N = strlen(text);
      if(N == 0) return;
       N = 2*N + 1; //Position count
      int L[N]; //LPS Length Array
       L[0] = 0; L[1] = 1;
      int C = 1; //centerPosition
      Int R = 2; //centerRightPosition
      int i = 0; //currentRightPosition
      int iMirror; //currentLeftPosition
      int expand = -1, diff = -1;
      int maxLPSLength = 0;
      int MaxLPSCenterPosition = 0;
      int start = -1,end = -1;
  //Uncomment it to print LPS Length array
  //printf("%d %d ", L[0], L[1]);
          for (i = 2; i < N; i++) {
    //get currentLeftPosition iMirror for
currentRightPosition i
    iMirror = 2*C-i;
//Reset expand - means no expansion required
    expand = 0; diff = R - i;
//If currentRightPosition i is within
//centerRightPosition R
    if(diff > 0) {
       if(L[iMirror] < diff) // Case 1
         L[i] = L[iMirror];
       else if(L[iMirror] == diff && i == N-1)
         L[i] = L[iMirror];
       else if(L[iMirror] == diff && i < N-1)
       {
            L[i] = L[iMirror];
            expand = 1; // expansion required
       else if(L[iMirror] > diff)
         L[i] = diff;
         expand = 1; // expansion required
       }
    }
    else{
       L[i] = 0;
       expand = 1; // expansion required
```

```
//Attempt to expand palindrome centered at
currentRightPosition i
//Here for odd positions, we compare characters and
//if match then increment LPS Length by ONE
//If even position, we just increment LPS by ONE
without
//any character comparison
    if (expand == 1){
       while (((i + L[i]) < N \&\& (i - L[i]) > 0) \&\&
         (((i + L[i] + 1) \% 2 == 0) | |
         (\text{text}[(i + L[i] + 1)/2] == \text{text}[(i-L[i]-1)/2])))
         L[i]++;
       }
    }
    if(L[i] > maxLPSLength){// Track maxLPSLength
       maxLPSLength = L[i];
       maxLPSCenterPosition = i;
    }
// If palindrome centered at currentRightPosition i
// expand beyond centerRightPosition R,adjust
// centerPosition C based on expanded palindrome.
    if (i + L[i] > R){
       C = i;
       R = i + L[i];
    //Uncomment it to print LPS Length array
    //printf("%d ", L[i]);
  }
  //printf("\n");
  start = (maxLPSCenterPosition - maxLPSLength)/2;
  end = start + maxLPSLength - 1;
  //printf("start: %d end: %d\n", start, end);
  printf("LPS of string is %s : ", text);
  for(i=start; i<=end; i++)</pre>
    printf("%c", text[i]);
  printf("\n");
```