

## Lecture 20

**Grover's search algorithm.** Grover's algorithm solves the problem of *unstructured search*. Suppose we have a function  $f: \{0, 1\}^n \rightarrow \{0, 1\}$ , and we are promised that either:

1.  $f(x) = 0$  for all  $x \in \{0, 1\}^n$ , or
2. there exists a unique  $x^* \in \{0, 1\}^n$  such that  $f(x^*) = 1$  and  $f(x) = 0$  for all  $x \neq x^*$ .

The goal is to determine which case holds, and if there is a marked element  $x^*$ , to find it. **Comment:** Think of  $f$  as evaluating a SAT formula.

Classically, you would need to query roughly  $2^n$  values in the worst case. Grover's algorithm can solve this problem using  $O(\sqrt{2^n})$  queries, giving a quadratic speedup.

**Fact 10.** In the classical query model, distinguishing these two cases requires  $\Omega(2^n)$  queries in the worst case, since you might need to check nearly all  $2^n$  possible inputs before finding  $x^*$  (or confirming no such  $x^*$  exists).

Recall from Lecture 13 that the quantum phase oracle for  $f: \{0, 1\}^n \rightarrow \{0, 1\}$  is given by

$$U_f |x\rangle |b\rangle = (-1)^{b \cdot f(x)} |x\rangle |b\rangle, \quad (112)$$

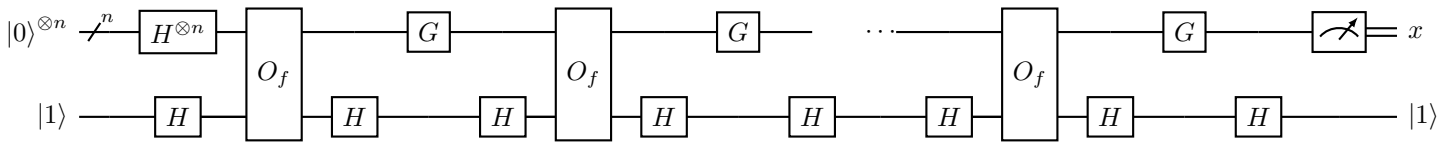
where  $x \in \{0, 1\}^n, b \in \{0, 1\}$ . The phase kickback trick shows that  $U_f = (\mathbb{1}_{2^n} \otimes H) O_f (\mathbb{1}_{2^n} \otimes H)$ .

**Proposition 9** (Grover's algorithm). *Given query access to  $f: \{0, 1\}^n \rightarrow \{0, 1\}$  where either  $f(x) = 0$  for all  $x$ , or there exists a unique  $x^* \in \{0, 1\}^n$  with  $f(x^*) = 1$ , Grover's algorithm can distinguish these cases and find  $x^*$  (if it exists) using*

$$O(\sqrt{2^n}) = O(2^{n/2}) \quad (113)$$

queries to  $f$ .

The quantum circuit for Grover's algorithm is:



where the pair  $(U_f, G)$  is repeated  $k \approx \frac{\pi}{4} \sqrt{2^n}$  times. The intermediate  $H$  gates on the ancilla could be deleted as they satisfy  $H^2 = \mathbb{1}$  – drawing them makes it clear where the  $U_f$ s come from.

*Proof.* Let  $N := 2^n$ . Let  $|\psi\rangle$  denote the  $N$ -dimensional quantum state

$$|\psi\rangle := H^{\otimes n} |0^n\rangle = \frac{1}{\sqrt{N}} \sum_{x \in \{0, 1\}^n} |x\rangle. \quad (114)$$

Let  $G \in \mathbb{C}^{N \times N}$  denote the *Grover diffusion operator*:

$$G := \mathbb{1}_N - 2|\psi\rangle\langle\psi|. \quad (115)$$

**Comment:** may discuss decomposing this operator into elementary quantum gates if there's time, else just Google or see, e.g., the first answer to this [StackExchange post](#).

Let  $V_f \in \mathbb{C}^{N \times N}$  denote the operation  $U_f$  implements on the first register when the ancilla qubit register is set to  $|1\rangle$ , i.e.,

$$V_f: |x\rangle \mapsto (-1)^{f(x)} |x\rangle \quad (116)$$

Then, the pair  $(U_f, G)$  forming each block implements the “Grover iteration” unitary  $GV_f$  on the first register. So suffices to analyze  $(GV_f)^k |\psi\rangle$ .

We analyze two cases:

1. **Case:**  $f(x) = 0$  for all  $x \in \{0, 1\}^n$ . In this case,  $V_f |x\rangle = |x\rangle$  for all  $x$ , so  $(GV_f)^k = G^k$ . Since  $G|\psi\rangle = -|\psi\rangle$ , we have  $G^k |\psi\rangle = (-1)^k |\psi\rangle$ . Therefore, measuring  $(GV_f)^k |\psi\rangle$  gives a uniformly random  $x \in \{0, 1\}^n$ , and we can verify that  $f(x) = 0$ , confirming this case.

2. **Case: there exists unique  $x^* \in \{0, 1\}^n$  with  $f(x^*) = 1$ .** Define the following quantum states:

$$|\psi_0\rangle := \frac{1}{\sqrt{N-1}} \sum_{x|f(x)=0} |x\rangle = \frac{1}{\sqrt{N-1}} \sum_{x \neq x^*} |x\rangle, \quad (117)$$

$$|\psi_1\rangle := |x^*\rangle. \quad (118)$$

These are normalized states:  $|\psi_0\rangle$  corresponds to the unmarked elements, and  $|\psi_1\rangle$  is the unique marked element. Note that  $\langle\psi_0|\psi_1\rangle = 0$  (orthogonal).

The initial state  $|\psi\rangle$  can be written as

$$|\psi\rangle = \sqrt{\frac{N-1}{N}} |\psi_0\rangle + \sqrt{\frac{1}{N}} |\psi_1\rangle = \cos(\theta) |\psi_0\rangle + \sin(\theta) |\psi_1\rangle, \quad (119)$$

where  $\theta := \arcsin(\sqrt{1/N}) \in (0, \pi/2]$ .

Analysis of Grover iteration  $GV_f$ :

We compute how  $GV_f$  acts on  $|\psi_0\rangle$  and  $|\psi_1\rangle$ :

$$\begin{aligned} GV_f |\psi_0\rangle &= G |\psi_0\rangle \quad (\text{since } V_f |\psi_0\rangle = |\psi_0\rangle) \\ &= |\psi_0\rangle - 2 |\psi\rangle \langle\psi|\psi_0\rangle \\ &= |\psi_0\rangle - 2 \cos(\theta) |\psi\rangle \\ &= |\psi_0\rangle - 2 \cos(\theta) (\cos(\theta) |\psi_0\rangle + \sin(\theta) |\psi_1\rangle) \\ &= (1 - 2 \cos^2(\theta)) |\psi_0\rangle - 2 \cos(\theta) \sin(\theta) |\psi_1\rangle \\ &= -\cos(2\theta) |\psi_0\rangle - \sin(2\theta) |\psi_1\rangle. \end{aligned}$$

Similarly,

$$\begin{aligned} GV_f |\psi_1\rangle &= -G |\psi_1\rangle \quad (\text{since } U_f |\psi_1\rangle = -|\psi_1\rangle) \\ &= -|\psi_1\rangle + 2 |\psi\rangle \langle\psi|\psi_1\rangle \\ &= -|\psi_1\rangle + 2 \sin(\theta) |\psi\rangle \\ &= 2 \sin(\theta) (\cos(\theta) |\psi_0\rangle + \sin(\theta) |\psi_1\rangle) - |\psi_1\rangle \\ &= 2 \sin(\theta) \cos(\theta) |\psi_0\rangle + (2 \sin^2(\theta) - 1) |\psi_1\rangle \\ &= \sin(2\theta) |\psi_0\rangle - \cos(2\theta) |\psi_1\rangle. \end{aligned}$$

Therefore,  $GV_f$  always maps the 2-dimensional subspace  $\text{span}(|\psi_0\rangle, |\psi_1\rangle)$  to itself. We can reduce the analysis to linear algebra in  $\mathbb{C}^2$  by working in the basis  $\{|\psi_0\rangle, |\psi_1\rangle\}$ .

In this basis,  $|\psi\rangle$  is represented as

$$\begin{pmatrix} \cos(\theta) \\ \sin(\theta) \end{pmatrix}, \quad (120)$$

and  $-GV_f$  is represented as the matrix

$$A := \begin{pmatrix} \cos(2\theta) & -\sin(2\theta) \\ \sin(2\theta) & \cos(2\theta) \end{pmatrix}. \quad (121)$$

This is the rotation matrix by angle  $2\theta$  anticlockwise!

**Comment:** Note that  $G = \mathbb{1}_N - 2|\psi\rangle\langle\psi|$  is a reflection about the hyperplane perpendicular to  $|\psi\rangle$ , while  $V_f = \mathbb{1}_N - 2|\psi_1\rangle\langle\psi_1|$  (check!) is a reflection about the hyperplane perpendicular to  $|\psi_1\rangle$ , so the above calculations also proves the mathematical fact that a product of two reflections is a rotation.