

Longest Increasing Subsequence

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1 Bottom-up Dynamic Programming solution

Let's define $L(i)$ as the length of the longest strictly increasing subsequence ending at index i . The recurrence formula for the longest strictly increasing subsequence is given by:

$$L(i) = 1 + \max_{\substack{j < i \\ \text{arr}[j] < \text{arr}[i]}} L(j)$$

This equation states that the length of the longest increasing subsequence ending at index i is 1 plus the maximum length obtained by considering all indices j less than i , where the corresponding element $\text{arr}[j]$ is less than $\text{arr}[i]$.

Complexity :

$$T(n) = \mathcal{O}(n^2)$$

$$M(n) = \mathcal{O}(n)$$

```
1 class Solution {
2     private int max(int[] L) {
3         int maxLength = Integer.MIN_VALUE;
4         for (final int length : L) {
5             maxLength = Math.max(maxLength, length);
6         }
7         return maxLength;
8     }
9
10    public int lengthOfLIS(int[] nums) {
```

```

11  int n = nums.length;
12  int[] L = new int[n];
13  // Initialize the array with minimum length 1 for each index
14  Arrays.fill(L, 1);
15
16  // Iterate to fill in the values of L(i) using the recurrence relation
17  for (int i = 1; i < n; i++) {
18      for (int j = 0; j < i; j++) {
19          if (nums[i] > nums[j]) {
20              L[i] = Math.max(L[i], L[j] + 1);
21          }
22      }
23  }
24  // Find the maximum value in the array L
25  return max(L);
26 }
27 }

```

Listing 1: $\mathcal{O}(n^2)$ DP solution

2 Backtracking Algorithm

```

1  public class Solution {
2      private List<List<Integer>> generateSubsequences(int[] arr) {
3          List<List<Integer>> allSubsequences = new ArrayList<>();
4          generateSubsequencesHelper(arr, 0, new ArrayList<>(), allSubsequences);
5          return allSubsequences;
6      }
7
8      private void generateSubsequencesHelper(int[] arr, int index, List<Integer> current,
9          List<List<Integer>> allSubsequences) {
10         if (index == arr.length) {
11             // Base case: add the current subsequence to the result
12             allSubsequences.add(new ArrayList<>(current));
13             return;
14         }
15         // Exclude the current element
16         generateSubsequencesHelper(arr, index + 1, current, allSubsequences);
17         // Include the current element
18         current.add(arr[index]);
19         generateSubsequencesHelper(arr, index + 1, current, allSubsequences);
20         // Backtrack to exclude the current element
21         current.removeLast();
22     }
23
24     private boolean isStrictlyIncreasing(List<Integer> list) {
25         for (int i = 1; i < list.size(); i++) {
26             if (list.get(i) <= list.get(i - 1)) {
27                 return false;
28             }
29         }
30         return true; // Strictly increasing
31     }
32
33     public int lengthOfLIS(int[] nums) {
34         List<List<Integer>> allSubsequences = generateSubsequences(nums);
35         int max = 1;

```

```

35     for (List<Integer> subsequence : allSubsequences) {
36         if (isStrictlyIncreasing(subsequence)) {
37             max = Math.max(max, subsequence.size());
38         }
39     }
40     return max;
41 }
42 }

```

Listing 2: $\mathcal{O}(2^n)$ Backtracking Algorithm

3 DP with Binary Search

```

1  import java.util.Arrays;
2
3  public class Solution {
4      public int lengthOfLIS(int[] nums) {
5          if (nums == null || nums.length == 0) {
6              return 0;
7          }
8
9          int[] dp = new int[nums.length];
10         int len = 0;
11
12         for (int num : nums) {
13             int index = Arrays.binarySearch(dp, 0, len, num);
14             if (index < 0) {
15                 index = -(index + 1);
16             }
17             dp[index] = num;
18             if (index == len) {
19                 len++;
20             }
21         }
22
23         return len;
24     }
25 }

```

Listing 3: $\mathcal{O}(n \log n)$ DP with Binary Search

4 Recursive Algorithm

```

1  public class Solution {
2      private int max;
3      public int lengthOfLISHelper(int[] arr, int n) {
4          // Base case: if there is only one element, the LIS length is 1
5          if (n == 1) {
6              return 1;
7          }
8          int currResult;
9          int maxEnding = 1;
10
11         for (int i = n - 1; i > 0; i--) {
12             // Recursively calculate LIS for previous elements

```

```

13     currResult = lengthOfLISSHelper(arr, i);
14
15     // Check if the current element can be included in the increasing subsequence
16     if ((arr[i - 1] < arr[n - 1]) && (currResult + 1 > maxEnding)) {
17         maxEnding = currResult + 1;
18     }
19 }
20 max = Math.max(maxEnding, max);
21 return maxEnding;
22 }
23
24 public int lengthOfLISS(int[] nums) {
25     lengthOfLISSHelper(nums, nums.length);
26     return max;
27 }
28 }

```

Listing 4: $\mathcal{O}(2^n)$ Recursive Algorithm