Write a simple C code for 2 dimensional array with dynamic memory allocation.

```
int i, j, **pi;
pi = (int **) malloc(sizeof(int*)*10);
if (pi == NULL) exit(1);

for (i = 0; i < 10; i++)
  pi[i] = (int *)malloc(sizeof(int)*20);
  // *(pi+i)</pre>
```

 Write a C program that prints out the integer values of x, y, z in ascending order.

```
if (x < y) {
        if (y < z) printf ("%d %d %d\n", x, y, z);
        else if (x < z) printf ("%d %d %d\n", x, z, y);
        else printf ("%d %d %d\n", z, x, y);
} else {
        if (x < z) printf ("%d %d %d\n", y, x, z);
        else if (y < z) printf ("%d %d %d\n", y, z, x);
        else printf ("%d %d %d\n", z, y, x);
}</pre>
```

Show that the following statements are correct.

$$-5n^2 - 6n = O(n^2) \rightarrow n \ge 0, 5n^2 - 6n \le 5n^2$$

$$- \sum_{i=0}^{n} i^2 = O(n^3) \to n \ge 0, \frac{n(n+1)(2n+1)}{6} \le n^3$$

$$-n^3 + 10^6 n^2 = O(n^3) \rightarrow n \ge 10^6$$
, $n^3 + 10^6 n^2 \le 2n^3$

$$- \frac{6n^3}{\log n + 1} = O(n^3) \to n \ge 1, \frac{6n^3}{\log n + 1} \le 6n^3$$

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Let length[i] be the desired length of row i of a two dimensional array. Write a function similar to make2dArray() to create a two dimensional array such that row i has length[i] elements.

```
int i, j, **pi, noRow;
noRow = 5;
pi = (int **) malloc(sizeof(int*)*noRow);
if (pi == NULL) exit(1);

for (i = 0; i < noRow; i++)
  pi[i] = (int *)malloc(sizeof(int)*length[i]);</pre>
```

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Develop a structure to represent each of the following geometric objects: rectangle, triangle, and circle.

```
typedef struct {
  int x, y;
} coordinate;

typedef struct {
  int no_point;
  coordinate p[4];
} polygon;
```

Name and student ID Implement Remove() and Attach() in slide 21 using polynomial representation in slide 22. #define MAX DEGREE 101 typedef struct { int degree; float coef[MAX DEGREE]; } polynomial; $A(x) = 3x^6 + 2x^5 + 41$ 3 6 0 0 0 0 41

degree

coef[]

```
polynomial Attach(polynomial *p, float c, int e) {
   int diff;
   if (e <= p->degree) p->coef[p->degree+1-e] += c;
   else {
        diff = e - p->degree;
        for (i = e + 1; i >= diff + 1; i--)
           p->coef[i] = p->coef[i-diff];
        p->degree = e;
        p->coef[1] = c;
Polynomial Remove (polynomial *p, int e) {
   int i = 1:
   p->coef[p->degree+1-e] = 0;
   if (e == p->degree) {
        while (p->coef[i] == 0) i++;
        for (j = 1; j \le p-)degree + 1 - i; j++)
           p->coef[j] = p->coef[j+i];
        p->degree -= i;
```

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We have a 500 x 500 matrix. Find the smallest number of nonzero elements such that the memory space of 2-D array representation becomes less than that of the sparse matrix representation.

<Solution>

2D array size = 500 x 500 x 4

Sparse matrix size = 3 x 4 x n

where n is the number of nonzero elements.

Therefore, the sparse matrix size becomes larger When $n \ge 83,334$.

```
Name and student ID
Rewrite the following codes without rowTerms[] to save memory.
Hint: Reuse startingPos[].
for (i=1; i <= numTerms; i++) {rowTerms[a[i].col]++;}</pre>
startingPos[0] = 1;
for(i = 1; i < numCols; i++) {</pre>
        startingPos[i]=startingPos[i-1]+rowTerms[i-1];}
<Solution>
for(i=1; i<= numTerms; i++) {startingPos[a[i].col]++;}</pre>
temp0 = startingPos[0];
startingPos[0] = 1;
for (i = 1; i < numCols; i++) {
       temp1 = startingPos[i];
        startingPos[i]=startingPos[i-1]+temp0;
        temp0 = temp1;
```

```
Name and student ID
Write the string remove function to remove j characters beginning
   from i in string s.
<Solution>
void stringremove (char *s, int i, int j) {
   len = strlen(s);
   if (len < i+j) {
        fprintf (stderr, "position is out of bounds\n");
        exit(1);
   for (x = i; x < len - j; x++)
        s[x] = s[x+j];
   s[x] = (0);
```

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Write a function strnchar() that takes a string and a character as input, and returns the number of occurrences of the character in the string.

```
int strnchar(char *s, char p) {
    int i, n, p_count;
    n = strlen(s);
    p_count = 0;
    for (i = 0; i < n; i++) {
        if (s[i] == p) p_count ++;
    }
    return p_count;
}</pre>
```

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Compute the failure function for each of the following patterns.

(a) aaaab: -1 0 1 2 -1

(b) ababaa: -1 -1 0 1 2 0

(c) abaabaab: -1 -1 0 0 1 2 3 1

```
Name and student ID
Rewrite pmatch() in slide 12 using for loop instead of while loop.
int pmatch (char* string, char* pat)
{
  int i=0; int j=0;
  int lens=strlen(string); int lenp=strlen(pat);
  for (i=0, j=0; i < lens && j < lenp;) {
    if (string[i] == pat[j]) { i++; j++; }
    else if (j == 0) i++;
    else j = failure[j-1]+1;
  return (j == lenp ? i-lenp : -1);
```

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```
Suppose a slightly modified maze problem such that we can move to only
  four directions (North, East, South, and West). Then, how do we need
  to modify path() ?
       while (dir < 4 && !found) {
           /* move in direction dir*/
           nextRow = row + move[dir].vert;
           nextCol = col + move[dir].horiz;
           if (nextRow==EXIT ROW && nextCol==EXIT COL)
              found = TRUE;
           else if ( !maze[nextRow][nextCol]
                         && !mark[nextRow][nextCol]) {
              mark[nextRow][nextCol]) = 1;
              position.row = row; position.col = col;
              position.dir = ++dir;
              push(position);
              row = nextRow; col = nextCol; dir = 0;
           else ++dir;
        } /* while (dir < 4 & !found)</pre>
```

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Evaluate the following postfix expressions. (assume that all the operands are single digit)

(a)
$$7 \ 3 \ 2 + * \ 3 - = 32$$

(b)
$$3 \ 4 \ 5 \ * \ 4 \ / \ + = 8$$

$$(c) 9 3 * 4 5 * + 3 - = 44$$

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Write the postfix form of the following expressions:

(a)
$$a * b * c \Rightarrow a b * c *$$

$$(b) -a + b - c + d \Rightarrow -a b + c - d +$$

$$(c) a * -b + c \Rightarrow a -b * c +$$

$$(d)(a + b) * d + e / (f + a * d) + c => a b + d * e f a d * + / c +$$

Name and student ID Modify delete() function in the previous slide so that it takes only two arguments, first and x without trail. void delete(listPointer *first, listPointer x) { listPointer trail; if (*first == x) *first = (*first)->link; // deleting the first node else { trail = first; for (;trail->link != x; trail = trail->link); trail->link = x->link; free(x);

Name and student ID In slide 17, modify addq() with the assumption that we have only front not rear. void addq (int i, element item) { /* insert an item at the rear of queue i */ queuePointer temp, rear; MALLOC(temp, sizeof(*temp)); temp->data = item; temp->link = NULL; if (front[i]) { for (rear = front[i]; rear->link != NULL; rear = rear->link); rear->link = temp; else front[i] = temp;

```
Name and student ID

In the previous slide, modify erase() using for loop instead of while().

void erase(polyPointer* ptr) {
   polyPointer temp;
   for ( ; *ptr != NULL; ) {
       temp = *ptr;
       *ptr = (*ptr)->link;
       free(temp);
   }
}
```

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In the previous slide, explain how to modify cpadd() if there is no special node indicating the first node. You do not need to submit the actual code.

>> If there is no special node (with exponent '-1') in a circular list, we have to check if we reach the first node whenever we move to the next node.

```
Name and student ID

Implement a function, invertedCopyList(), for a given list, to copy the list, invert it, and return.

listPointer invertedCopyList(listPointer ptr){
}

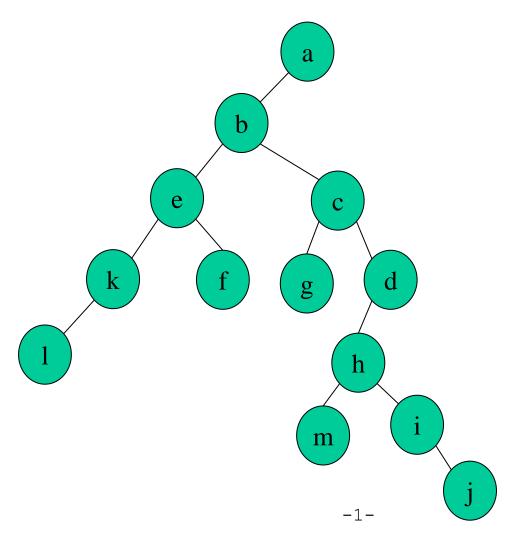
>> You can easily implement by slightly modifying invert().
```

```
Name and student ID
   count = 0;
   for (i = 0; i < n; i++)
       if (out[i]) {
          count ++;
          printf("\nNew class: %5d", i);
          out[i] = FALSE; x = seq[i]; top = NULL;
          for (;;) {
              while (x) {
                j = x->data;
                 if (out[j]) {
                  printf(%5d", j); out[j] = FALSE;
                  y=x->link; x->link=top; top=x; x=y;
                } else x = x-\lambda;
              if (!top) break;
              x = seq[top->data];
              top = top->link;
          } /* for (;;) */
   } /* for (i) */
   printf("%d\n", count);
 } /* main() */
```

>> Since out[i] is initialized as TRUE for 0<= i <n, count simply becomes 'n'.

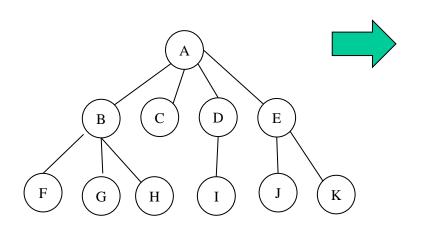
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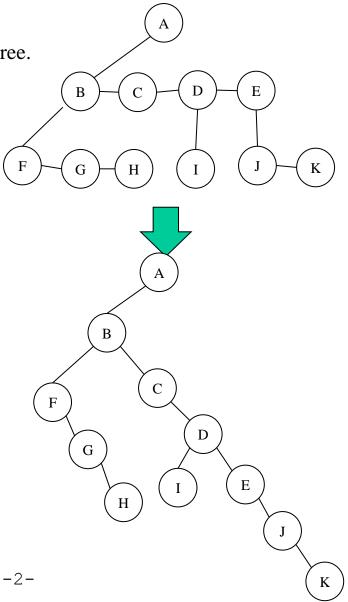
Draw the tree given as the list (a(b(e(k(l),f),c(g,d(h(m,i(j))))))).



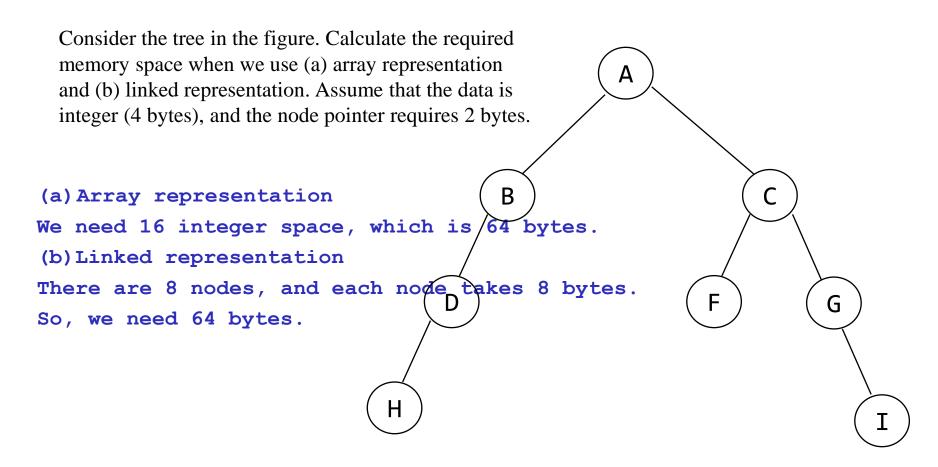
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Convert the tree in the figure into a binary tree.





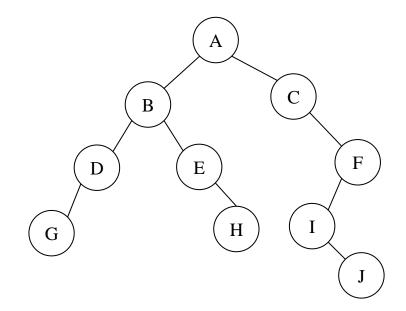
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Write out the inorder, preorder, postorder, and level-order traversals for the given binary tree.

Inorder: G D B E H A C I J F
Preorder: A B D G E H C F I J
Postorder: G D H E B J I F C A
Level-order: A B C D E F G H I J



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Write a function treeAdd() which returns the sum of data in a tree.

```
int treeAdd(treePointer tree) {
  int sum = 0;

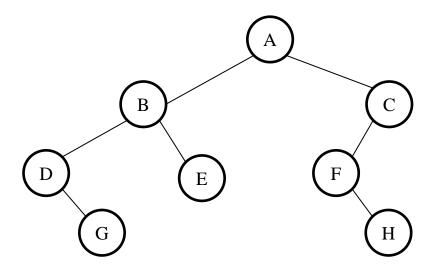
if (ptr) {
  sum = treeAdd(ptr->leftChild);
  sum += ptr->data;
  sum += treeAdd(ptr->rightChild);
  }
  return sum;
}
```

```
typedef struct node* treePointer;
typedef struct node {
   int data;
   treePointer leftChild, rightChild;
};
```

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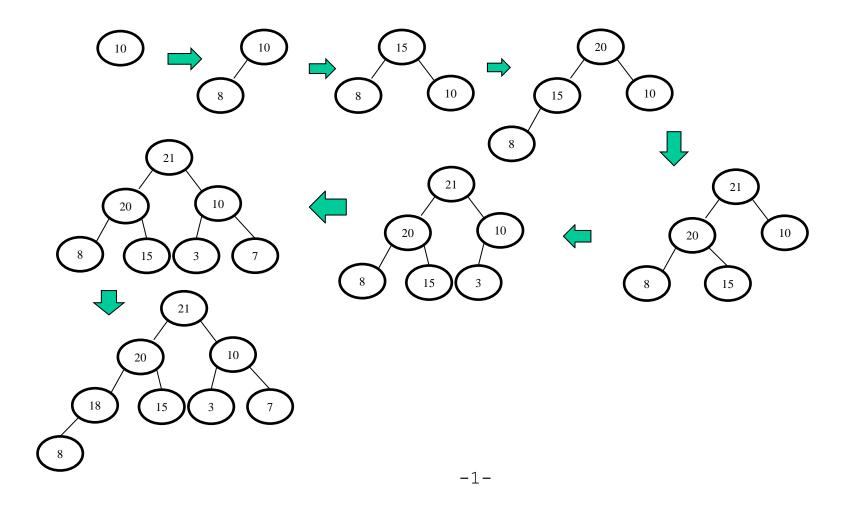
Rebuild a binary tree from the following traversal results:

Inorder traversal: DGBEAFHC Postorder traversal: GDEBHFCA



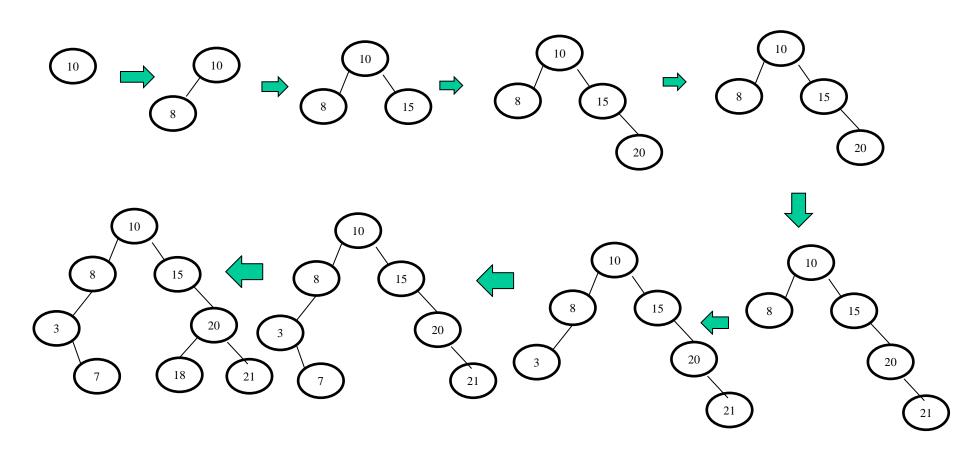
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Draw the max heap after inserting 10, 8, 15, 20, 21, 3, 7, 18 from an empty heap.

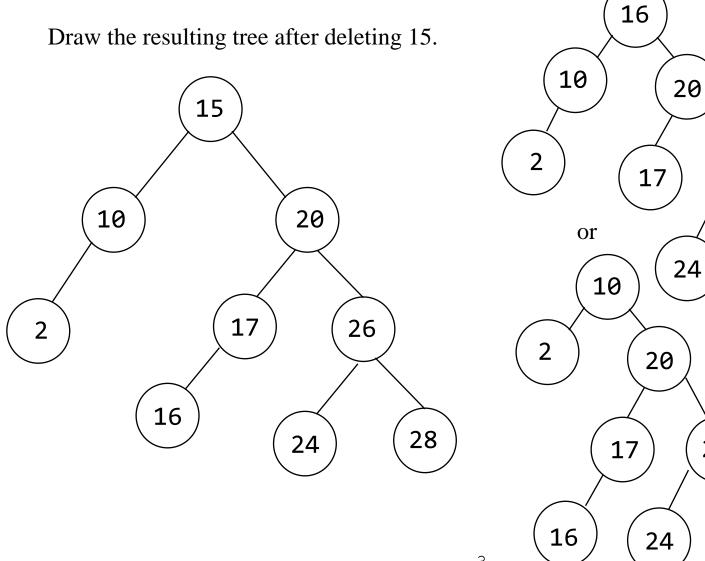


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Draw the BST after inserting 10, 8, 15, 20, 21, 3, 7, 18 from an empty BST.



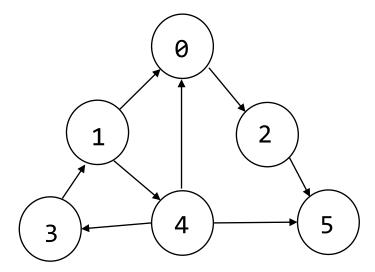
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Name and student ID

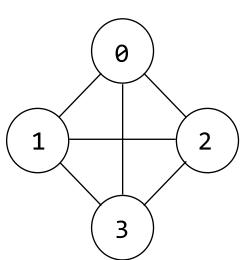
Draw the strongly connected components of the following graph.

Since there is a path between all the pairs in the graph, the graph is strongly connected.



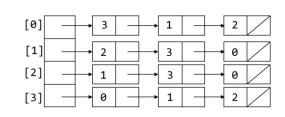
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Calculate the required memory to represent the following graph using adjacency matrix, adjacency list, and sequential representation, respectively.



Adjacency matrix: $4 \times 4 \times 4$ byte = 64 bytes

- Adjacency list: 8 + 72 = 80 bytes
- 4 Head pointers (4 x 2 bytes)
- 12 nodes (12 \times (4 + 2) bytes)



Sequential representation: $(n + 2xe + 1) \times 4$ byte

- n = 4, e = 6
- $(4 + 2x6 + 1) \times 4$ byte = 68 bytes

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Show the adjacency multilist for the following graph (slide 32).

