

Answer Set Programming for Scheduling at Bosch

T. Eiter¹, T. Geibinger¹, N. Musliu¹, J. Oetsch¹, P. Skocovsky^{1,2}, D. Stepanova²

¹Vienna University of Technology (TU Wien)

²Bosch center for AI, CR/PJ-AI-R26

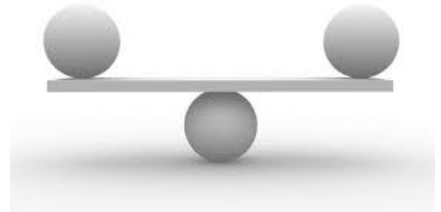
Pain Point 1: Solution Robustness

■ Independence

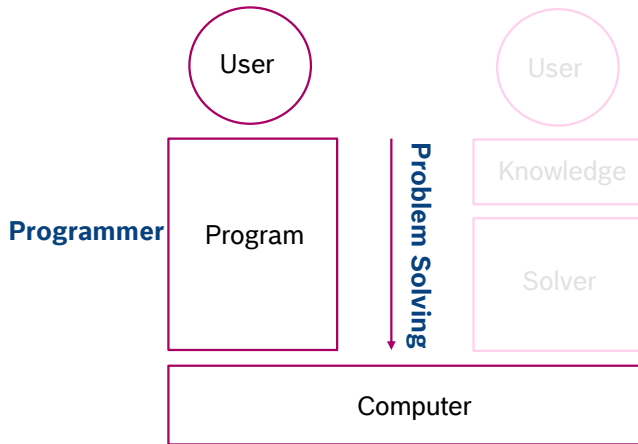
- Depending on commercial software is problematic
 - Additional costs
 - Issues if the company is sold, disappears, changes contracting conditions...

■ Reliability

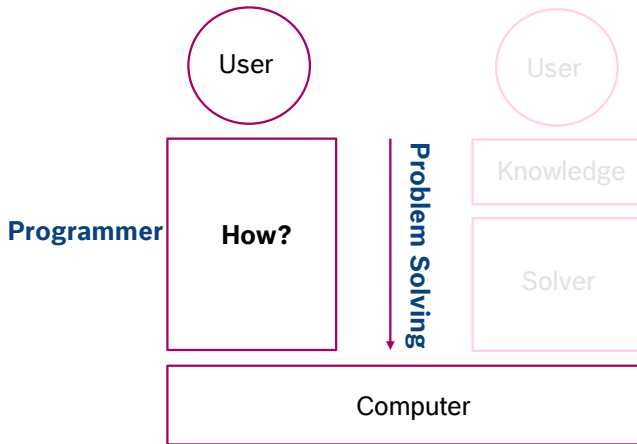
- Software should be
 - Efficient
 - Maintainable
 - Ideally, optimality guarantees



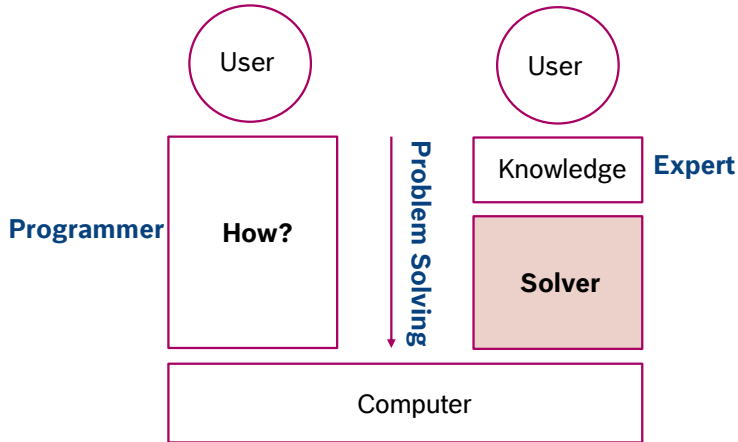
Traditional VS Knowledge-driven Problem Solving



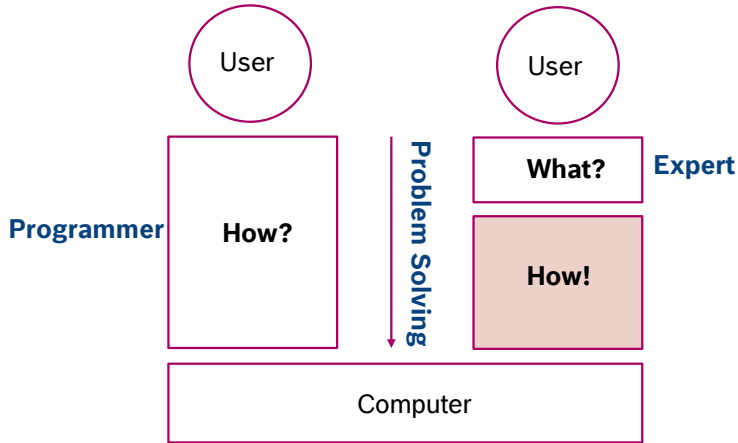
Traditional VS Knowledge-driven Problem Solving



Traditional VS Knowledge-driven Problem Solving



Traditional VS Knowledge-driven Problem Solving



Knowledge-driven Problem Solving

- + Transparency
- + Flexibility
- + Maintainability
- + Reliability

- + Generality
- + Efficiency
- + Optimality
- + Availability

Knowledge

Expert

Solver

Knowledge-driven Problem Solving

- + Transparency
- + Flexibility
- + Maintainability
- + Reliability

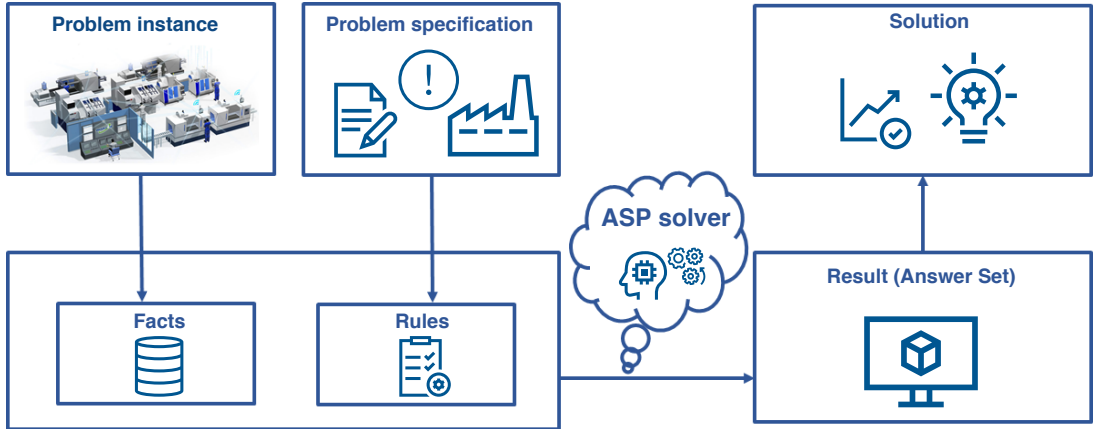
- + Generality
- + Efficiency
- + Optimality
- + Availability

Knowledge

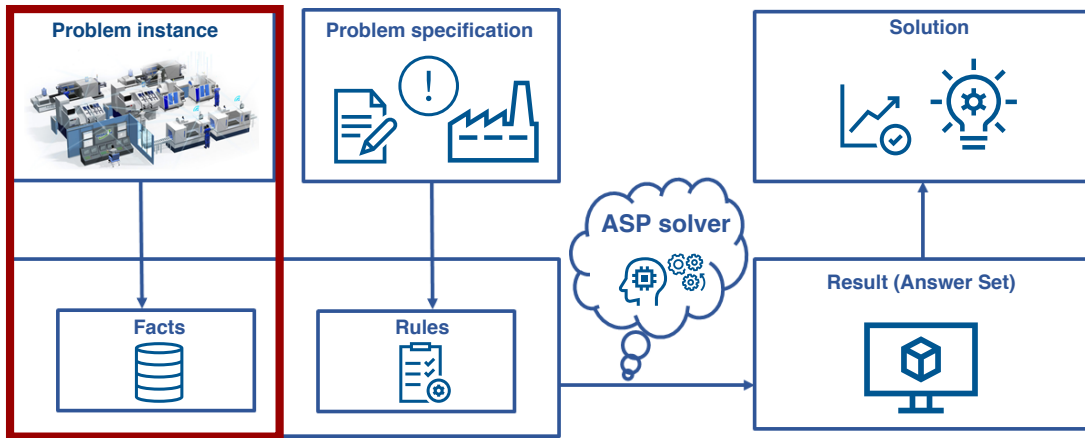
Expert

Solver

Answer Set Programming



Facts



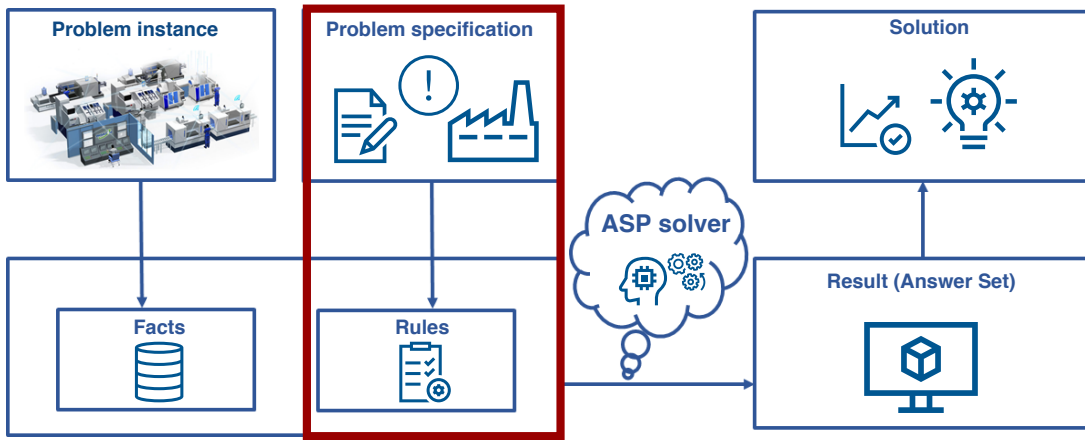
Facts

Example

Machine **m1** is capable of processing the job **j6** with impact of type “**r**”, whose duration is **50** minutes, due time is in **180** minutes and whose weight on machine **m1** is **67**.

- `machine(m1)`
- `job(j6)`
- `impact(j6, r)`
- `...`
- `weight(j6, m1, 67)`

Rules



Rules

Example

Job must be assigned to a single machine that is capable of performing the job.

```
■ 1 { assigned(J,M) : machine(M), capable(M,J) } 1 :- job(J).
```

Rules

Example

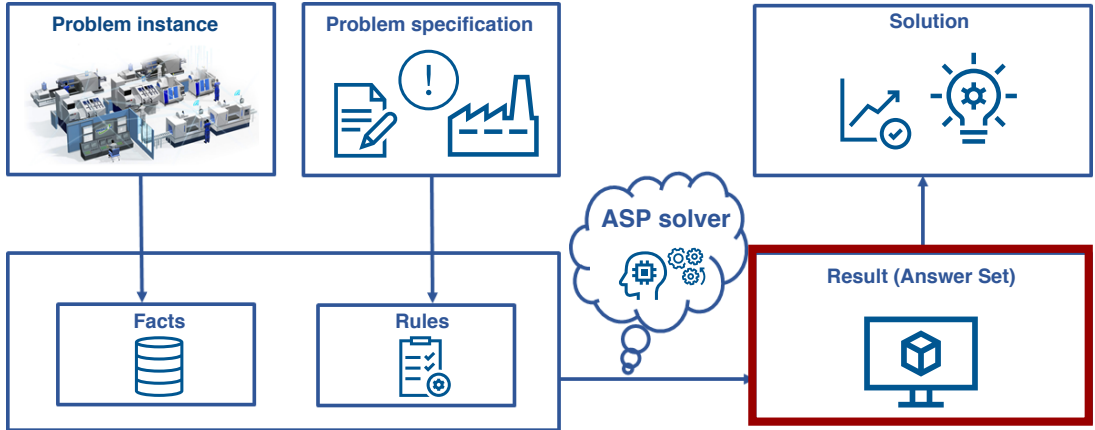
Job must be assigned to a single machine that is capable of performing the job.

- `1 { assigned(J,M) : machine(M), capable(M,J) } 1 :- job(J).`

A job has to start after its release time.

- `&diff{ 0 - start(J) } <= -R :- job(J), assigned(J,M), release(J,M,R).`

ASP Output (Answer Set)



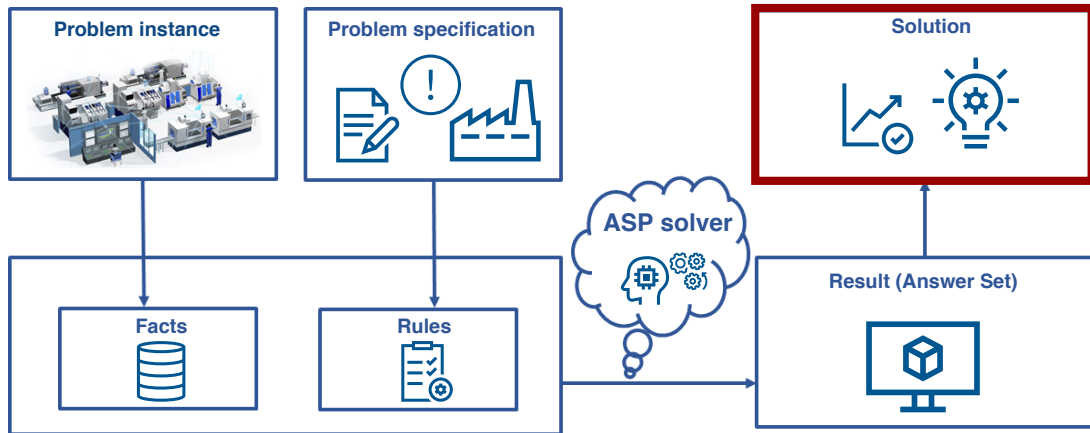
ASP Output (Answer Set)

Example

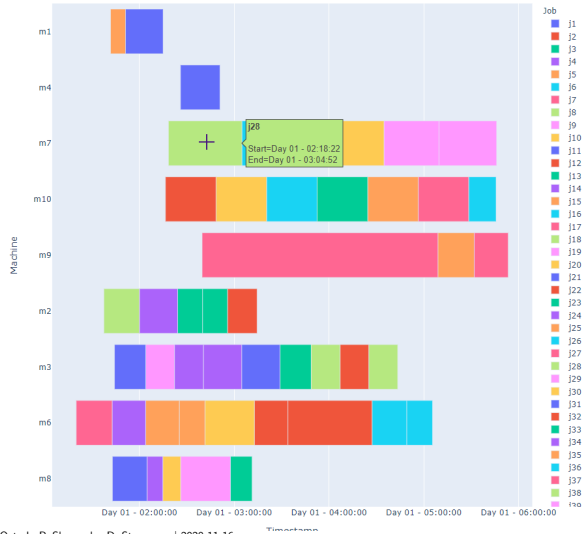
```
$ clingo-dl -heuristic=domain -minimize-variable=makespan asp_encoding.lp
Optimization: 115835
:
Optimization: 93707
Answer: 11
assigned(j1, m6), ..., assigned(j98,m5), assigned(j99,m14)

OPTIMUM FOUND
Models : 11
Time : 6.68s (Solving: 3.03s 1st Model: 3.33s Unsat: 0.02s)
CPU Time : 5.718s
```

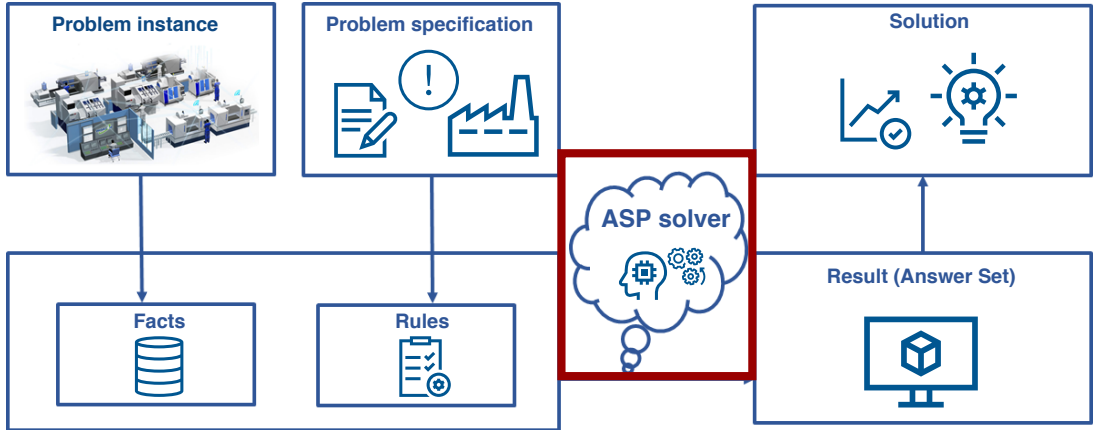

Solution



Solution

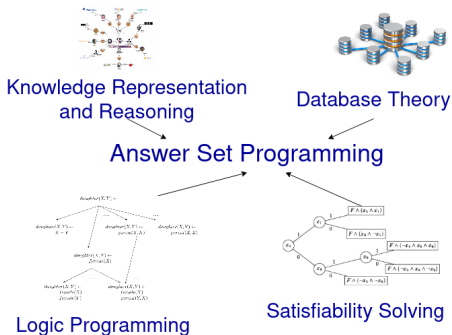


ASP Solvers



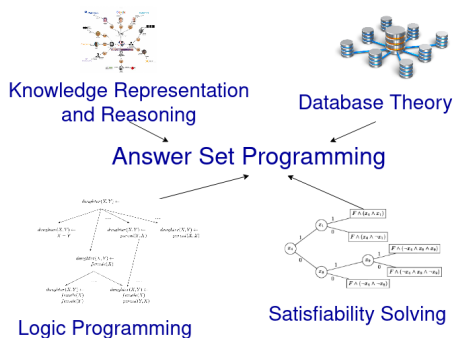
ASP: Research, Solvers and Applications

Academic research



ASP: Research, Solvers and Applications

Academic research



ASP solvers

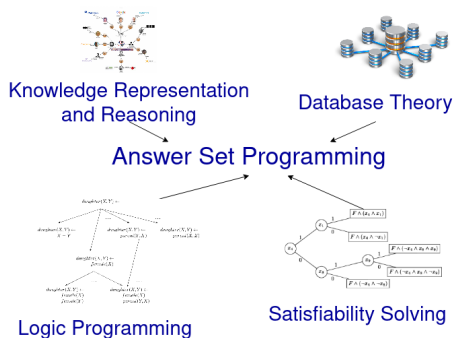
Clingo

- Open source (MIT license)
- Well-supported



ASP: Research, Solvers and Applications

Academic research



ASP solvers

Clingo

- Open source (MIT license)
- Well-supported



Industrial applications

SIEMENS
Ingenuity for life



Railroad interlocking configurators



Routing driverless transport vehicles in car assembly

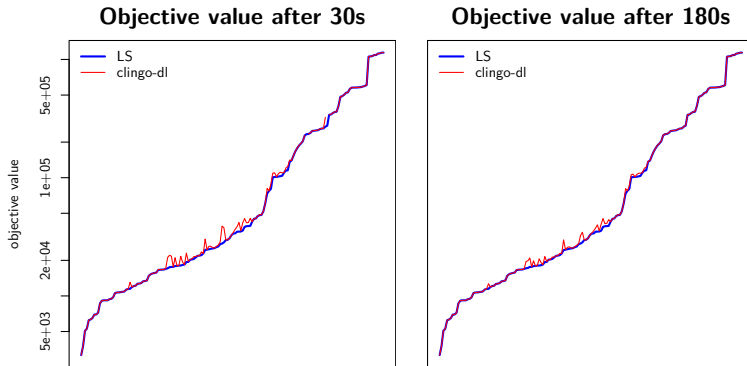


SBB CFF FFS



Train scheduling for Swiss railroad

Comparison to Local Solver on Implantation Rtp Problem



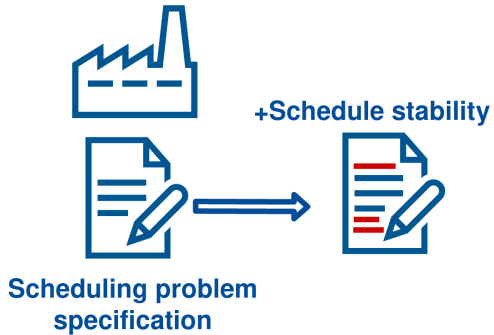
Objective values of local solver and clingo-dl achieved on 162 real snapshots

Pain Point 2: Solution Adaptability

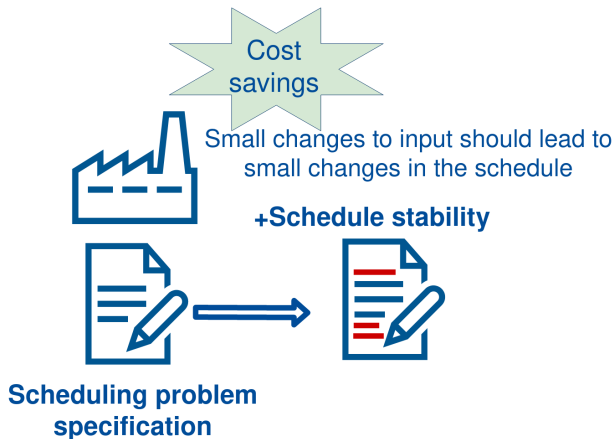


**Scheduling problem
specification**

Pain Point 2: Solution Adaptability

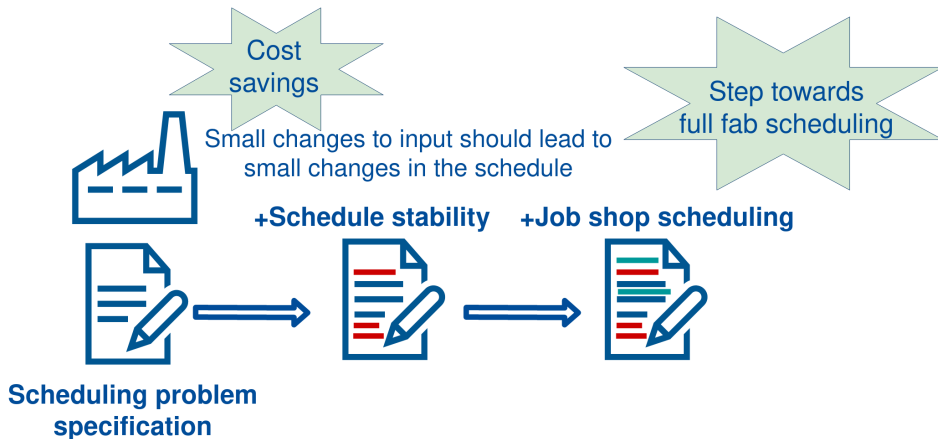


Pain Point 2: Solution Adaptability

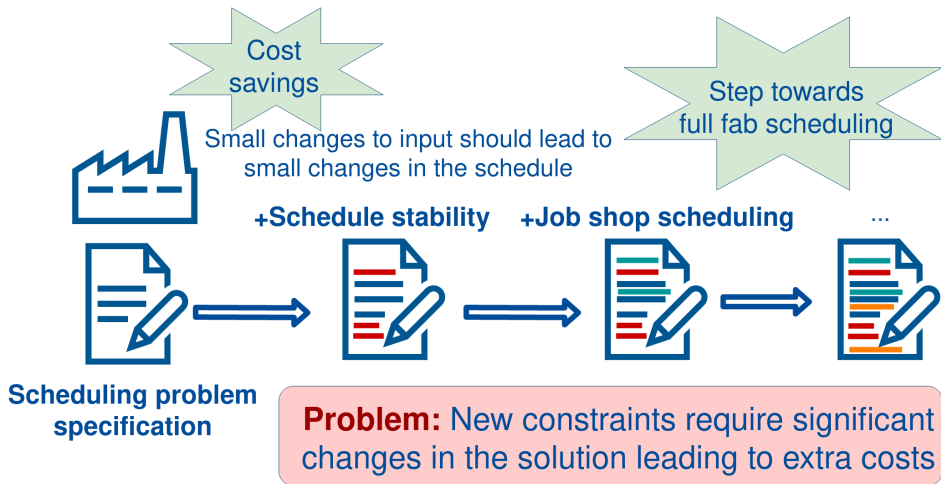


...

Pain Point 2: Solution Adaptability

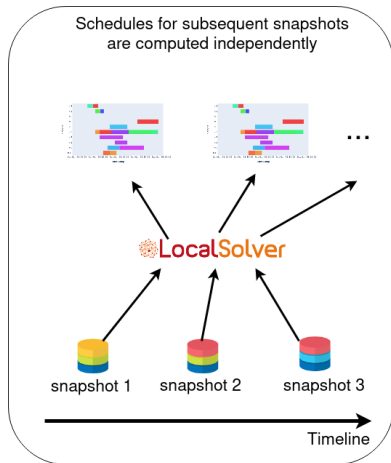


Pain Point 2: Solution Adaptability

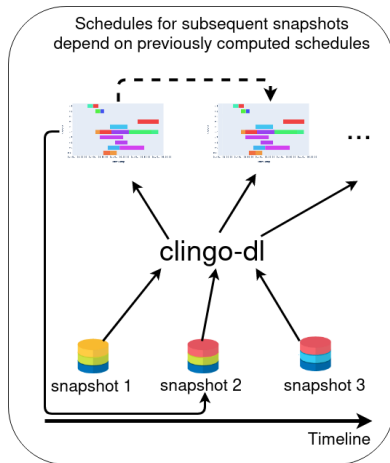


Schedule Stability

Current solution

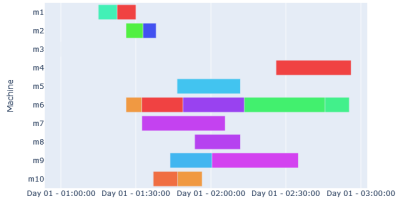


ASP-based solution

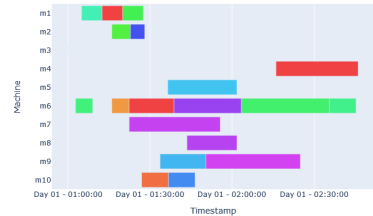
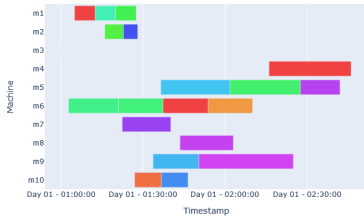
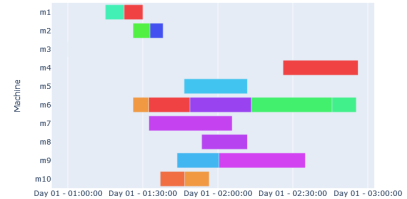


Schedule Stability

Current solution (unstable)



ASP-based solution (stable)



Job Shop Scheduling

Example

ASP scheduling encoding

```
before(J1,J2,M) | before(J2,J1,M) :- assigned(J1,M), assigned(J2,M), J1 < J2.
:- before(J1,J2,M), before(J2,J3,M), not before(J1,J3,M).

first(J1,M) :- assigned(J1,M), #count { J2 : before(J2,J1,M) } = 0.
1 { next(J1,J2,M) : before(J1,J2,M) } 1 :- assigned(J2,M), not first(J2,M).
:- next(J1,J2,M), before(J1,J3,M), before(J3,J2,M).
1 { assigned(J,M) : machine(M), capable(M,J) } 1 :- job(J).

earliest(J,M,T) :- job(J), machine(M),
    T = #max { 0 ; R : release(J,M,R) ;
    F : lastJobFinished(M,F) ;
    A : availableFrom(M,A) }.

&diff{ 0 - compl(J) } <= -(T+D) :- assigned(J,M),
    duration(J,M,D),
    earliest(J,M,T).

&diff{ 0 - compl(J1) } <= -(T+D+S) :- first(J1,M), duration(J1,M,D),
    histJob(M,J2), earliest(J1,M,T),
    setupTime(M,J2,J1,S).

&diff{ compl(J2) - compl(J1) } <= -(D+S) :- before(J2,J1,M),
    next(J3,J1,M),
    setupTime(M,J3,J1,S),
    duration(J1,M,D).

&diff{ 0 - compl(J1) } <= -(T+D+S) :- next(J2,J1,M),
    setupTime(M,J2,J1,S),
    duration(J1,M,D),
    earliest(J1,M,T).

&diff{ compl(J) - 0 } <= T :- assigned(J,M), availableTo(M,T).
:- next(J2,J1,M), forbidden(J2,J1).
:- first(J1,M), histJob(M,J2), forbidden(J2,J1).

&diff{ compl(J) - makespan } <= 0 :- job(J).
&diff{ 0 - makespan } <= -F :- machine(M), lastJobFinished(M,F).
```

ASP job shop scheduling encoding

```
operation(J,1..N) :- job(J,N).
before(J1,J2,M) | before(J2,J1,M) :- assigned(J1,M), assigned(J2,M), J1 < J2.
:- before(J1,J2,M), before(J2,J3,M), not before(J1,J3,M).

first(J1,M) :- assigned(J1,M), #count { J2 : before(J2,J1,M) } = 0.
1 { next(J1,J2,M) : before(J1,J2,M) } 1 :- assigned(J2,M), not first(J2,M).
:- next(J1,J2,M), before(J1,J3,M), before(J3,J2,M).
1 { assigned(+(J,0),M) : machine(M), capable(M,J) } 1 :- job(J)operation(J,0).

earliest(J,M,T) :- job(J)job(J,_), machine(M),
    T = #max { 0 ; R : release(J,M,R) ;
    F : lastJobFinished(M,F) ;
    A : availableFrom(M,A) }.

&diff{ 0 - compl(J,0) } <= -(T+D) :- assigned(+(J,0),M),
    duration(+(J,0),M,D),
    earliest(J,M,T).

&diff{ 0 - compl(J1,01) } <= -(T+D+S) :- first(J1,M), duration(+(J1,01),M,D),
    histJob(M,J2), earliest(J1,M,T),
    setupTime(M,J2,J1,S).

&diff{ compl(J2,02) - compl(J1,01) } <= -(D+S) :- before(+(J2,02),+(J1,01),M),
    next(+(J3,03),+(J1,01),M),
    setupTime(M,J3,J1,S),
    duration(+(J1,01),M,D).

&diff{ 0 - compl(J1,01) } <= -(T+D+S) :- next(+(J2,02),+(J1,01),M),
    setupTime(M,J2,J1,S),
    duration(+(J1,01),M,D),
    earliest(J1,M,T).

&diff{ compl(J,0) - 0 } <= T :- assigned(+(J,0),M), availableTo(M,T).
:- next(+(J2,_),+(J1,_),M), forbidden(J2,J1).
:- first(+(J1,_),M), histJob(M,J2), forbidden(J2,J1).

&diff{ compl(J,0) - makespan } <= 0 :- job(J)operation(J,0).
&diff{ 0 - makespan } <= -F :- machine(M), lastJobFinished(M,F).

&diff{ (J,01) - (J,02) } <= -D :- assigned((J,01),M), duration((J,01),M,D),
    operation(J,02), 02 = 01+1.
```

Heuristics

Example

Scenario:

- Weighted shortest processing time (wpts) heuristic:
schedule job j before k if $d_j/w_j < d_k/w_k$, where
 - d_j is duration of j
 - w_j is weight of j

ASP encoding of wpts heuristic to guide search:

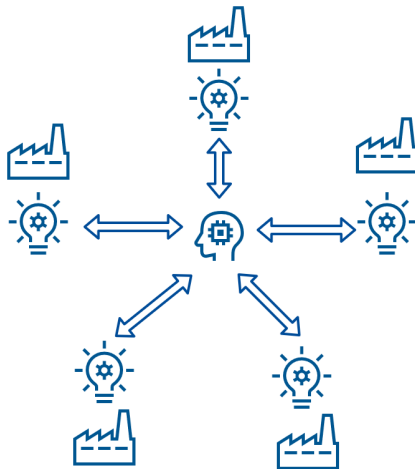
- `wspt(J,R) :- duration(J,D), weight(J,W), R = D*1000 / W.`
`#heuristic next(J,K) : wspt(J,RJ), wspt(K,RK), RJ > RK. [-1,sign]`

Summary

- Pain point: solution robustness
 - Clingo-dl: years of research, open source
 - ASP encoding for implantation work center
 - Performance on par with local solver
- Pain point: solution adaptability
 - Computation of stable schedules
 - Extension to job shop

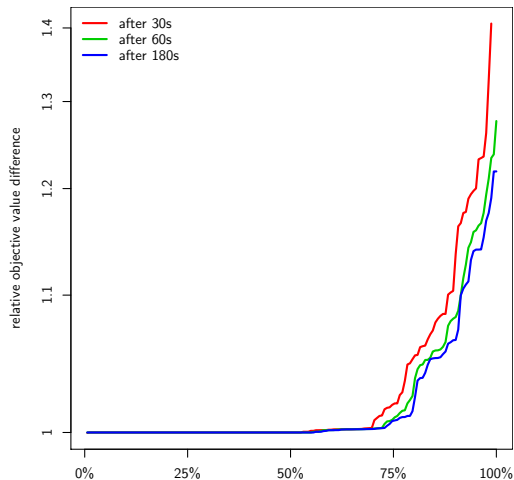
Summary

- Pain point: solution robustness
 - Clingo-dl: years of research, open source
 - ASP encoding for implantation work center
 - Performance on par with local solver
- Pain point: solution adaptability
 - Computation of stable schedules
 - Extension to job shop
- Further steps:
 - Realization of various optimization functions
 - Full fab scheduling
 - Learn heuristics using ML (e.g., re-opt, R34), encode as ASP rules



Distribution of Relative Objective Value Difference

- Open source ASP solver clingo-dl performs on par with Local Solver



Job Shop Scheduling

Example (ASP scheduling job shop encoding)

```
% Every job comprises of N operations: 1..N  
operation(J,1..N) :- job(J,N).
```

Job Shop Scheduling

Example (ASP scheduling job shop encoding)

```
% Every job comprises of N operations: 1..N
```

```
operation(J,1..N) :- job(J,N).
```

```
% Every job must be assigned to exactly one machine capable of performing  
the job
```

```
1{assigned(J,M):machine(M), capable(M,J)}1 :- job(J).
```

Job Shop Scheduling

Example (ASP scheduling job shop encoding)

```
% Every job comprises of N operations: 1..N
```

```
operation(J,1..N) :- job(J,N).
```

```
% Every operation of every job must be assigned to exactly one machine  
capable of performing the job
```

```
1{assigned(J(J,0),M):machine(M), capable(M,J)}1 :- job(J)operation(J,0).
```

Job Shop Scheduling

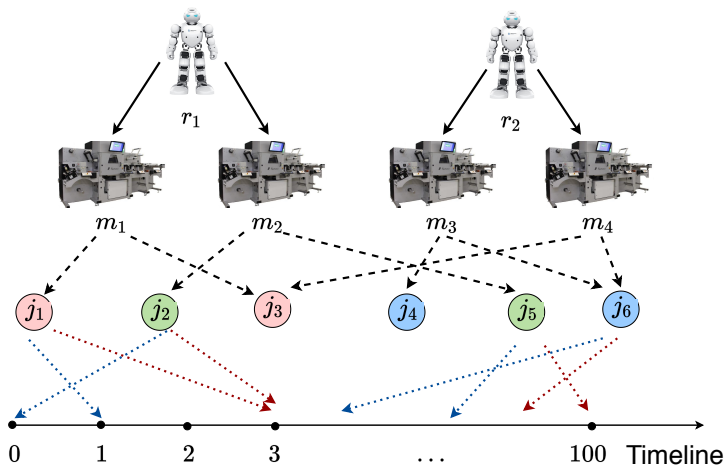
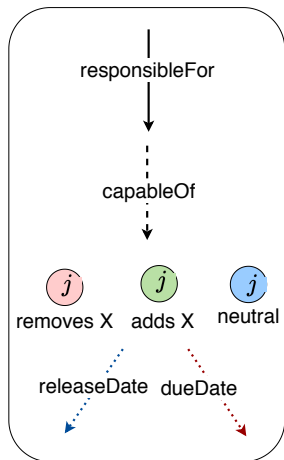
Example (ASP scheduling job shop encoding)

```
% Every job comprises of N operations: 1..N
operation(J,1..N) :- job(J,N).

% Every operation of every job must be assigned to exactly one machine
capable of performing the job
1{assigned(J(J,0),M):machine(M), capable(M,J)}1 :- job(J)operation(J,0).

% Operations in a job must not overlap
&diff{ (J,01) - (J,02) } <= -D :- assigned((J,01),M), duration((J,01),M,D),
                                operation(J,02), 02 = 01+1.
```

Semiconductor Production at Reutlingen



Scheduling Constraints

- Each job has:
 - machine-specific duration
 - machine-specific weight
- Each machine has:
 - availability time period
 - limit on available time capacity per job type
 - setup times: time needed to switch from job j_i to j_k
 - forbidden sequences of jobs
- Each robot has
 - limit on the number of jobs that can be assigned to machines served by the robot
- Objectives:
 - minimize the maximum of the completion time of the last job performed on a machine
 - minimize total weighted completion time

Comparison to CPLEX

- ASP is strictly more expressive than CPLEX
- ASP is more transparent, convenient to use, programs are easy to extend/modify
- Some implementations of ASP solvers even use CPLEX at the backend
- The implementation that we are working on is more efficient¹.

¹Tomi Janhunen, Roland Kaminski, Max Ostrowski, Sebastian Schellhorn, Philipp Wanko, Torsten Schaub: Clingo goes linear constraints over reals and integers. Theory and Practice of Logic Programming, 17(5-6): 872-888 (2017)