

rate = R bps

distance = m meters

propagation delay = m/s

packet size = L bits

a. $d_{prop} = \frac{m}{s}$

b. $d_{trans} = \frac{L}{R}$

c. end-to-end = $d_{prop} + d_{trans} = \frac{m}{s} + \frac{L}{R}$

d. $t = 0$

@ $t = d_{trans}$

At d_{trans} , the last bit should just be leaving Host A.

e. $d_{prop} < d_{trans}$

$t = d_{trans}$

The first bit should already be at host B.

f. $s = 2.5 \times 10^4$ m/s

$L = 100$ bits

$R = 28$ Kbps

$d_{prop} = d_{trans}$

$\frac{d}{2.5 \times 10^4 \text{ m/s}} = \frac{100 \text{ bits}}{28000 \text{ b/s}} = 0.00357 \text{ s}$

$d = 89.286 \text{ m}$

2. # RTT

index.html w/ 5 jpeg

a. Non-persistent connections require 2 RTT per object.

In this case, non-persistent connections require 12 RTT

b. With parallel connections, the index.html file requires 2 RTT, 2 for 4 jpeg and 2 for final jpeg. Total # of RTT = 6

c. Persistent connections require 1 RTT per object plus 1 for startup

In this case, the connection require 7 RTT

d. 1 setup + 1 index.html + 1 jpegs = 3 RTT's

3.3 There will be 1000 server-side sockets and 100 port #. There are 1000 sockets for each TCP connection and 1 additional server-side port # ready for a new connection.

4. when $d_{\text{queue}} = 0$

a. $130 \mu = \frac{m}{2 \times 10^8} + \frac{8(1024)}{10^9} \Rightarrow m = 24.4 \text{ km}$

b. $d_{\text{end to end}} = 142 \mu$

$d_{\text{end to end}} - (d_{\text{prop}} + d_{\text{trans}}) = d_{\text{queue}} = 142 \mu - 130 \mu = 12 \mu$

time without queuing delay to transmit: $\frac{8(1024)}{10^9} = 8.19 \mu$

$\Rightarrow \frac{12}{8.19} = 1.47$ packets are queued on average.

$142 = \frac{1000 \cdot 1024}{1.66/s}$