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# Spatial organization of genes along mammalian chromosomes

EPFL Bachelor Project – Computer Science – Spring 2016

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Lab : Computational Systems Biology Lab (Felix Naef)

Why we are doing this research

# Goals of the research

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01

**One:**

Genes coding for functional products are placed along each chromosome, the function of its specific positioning **being yet to be understood**

02

**Two:**

Are genes placed in **clusters** along the chromosomes ?

03

**Three:**

If they are, it would suggest a **functional role of the specific positioning of genes along a chromosome**

The idea

## Partitioning algorithm

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- How to identify the clusters and partition the genes efficiently ?
- From an  $O(2^{n-1})$  problem to an  $O(N^2)$  using dynamic programming

Simple example

## Partitioning algorithm

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- A length  $n$  of a metal rod. Example for  $n = 4$ . So  $2^3 = 8$  possible solutions
- A table of prices  $p_i$  for rods of lengths  $i = 1, \dots, n$



Simple example

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price $p_i$	1	5	8	9	10	17	17	20	24	30

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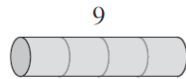
- **Objective** : Decide how to cut the rod into pieces and maximize / minimize the price.

Simple example

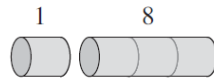
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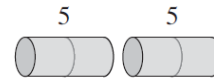
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(a)



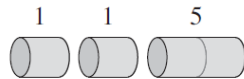
(b)



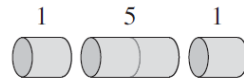
(c)



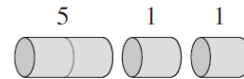
(d)



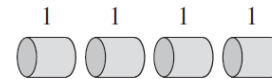
(e)



(f)



(g)



(h)

Simple example

## Partitioning algorithm

---

- If an optimal solution cuts the rod into  $k$  pieces, for some  $1 \leq k \leq n$ , then an optimal decomposition
- $n = i_1 + i_2 + \dots + i_k$



Simple example

## Partitioning algorithm

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- $n = i_1 + i_2 + \dots + i_k$
- of the rod into pieces of lengths  $i_1, i_2, \dots, i_k$  provides maximum revenue
- $r_n = p_{i_1} + p_{i_2} + \dots + p_{i_k} p_{i_1}$

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CUT-ROD( $p, n$ )

```
1  if  $n == 0$ 
2      return 0
3   $q = -\infty$ 
4  for  $i = 1$  to  $n$ 
5       $q = \max(q, p[i] + \text{CUT-ROD}(p, n - i))$ 
6  return  $q$ 
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Simple example

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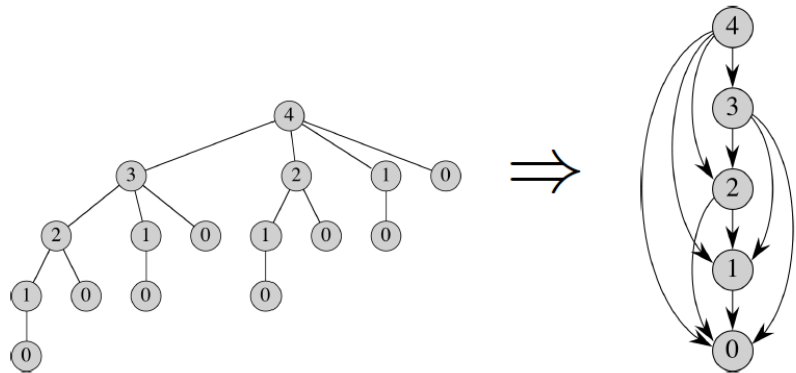
BOTTOM-UP-CUT-ROD( $p, n$ )

```
1  let  $r[0..n]$  be a new array
2   $r[0] = 0$ 
3  for  $j = 1$  to  $n$ 
4       $q = -\infty$ 
5      for  $i = 1$  to  $j$ 
6           $q = \max(q, p[i] + r[j - i])$ 
7       $r[j] = q$ 
8  return  $r[n]$ 
```

Simple example

## Partitioning algorithm

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Simple example

# Partitioning algorithm

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- Algorithm only returns the optimal revenue, but we want the whole solution path.

Simple example

## Partitioning algorithm

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- **Approach** : Each cell of the stored table corresponds to a decision: the location of the leftmost cut. Store the decision corresponding to every cell in a separate table

Simple example

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- **Approach** : Each cell of the stored table corresponds to a decision: the location of the leftmost cut. Store the decision corresponding to every cell in a separate table

EXTENDED-BOTTOM-UP-CUT-ROD( $p, n$ )

```
1  let  $r[0..n]$  and  $s[0..n]$  be new arrays
2   $r[0] = 0$ 
3  for  $j = 1$  to  $n$ 
4       $q = -\infty$ 
5      for  $i = 1$  to  $j$ 
6          if  $q < p[i] + r[j - i]$ 
7               $q = p[i] + r[j - i]$ 
8               $s[j] = i$ 
9       $r[j] = q$ 
10 return  $r$  and  $s$ 
```

Simple example

## Partitioning algorithm

- Algorithm only returns the optimal revenue, but we want the whole solution path.
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**PRINT-CUT-ROD-SOLUTION**( $p, n$ )

```
1  ( $r, s$ ) = EXTENDED-BOTTOM-UP-CUT-ROD( $p, n$ )
2  while  $n > 0$ 
3      print  $s[n]$ 
4       $n = n - s[n]$ 
```

$i$	0	1	2	3	4	5	6	7	8	9	10
$r[i]$	0	1	5	8	10	13	17	18	22	25	30
$s[i]$	0	1	2	3	2	2	6	1	2	3	10



Real partitioning

## Partitioning algorithm

---

We are actually doing quite the same with our RNA-Seq data

- Where **n** is the number of genes of a chromosome we want to partition
- Genes are ordered in order to follow a cluster logic
- We try to minimize the score instead of maximizing it

Real partitioning

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What is the equivalent to the **prices** of the rod cutting algorithm ?

## Model selection

---

- The «price» is computed for each block of genes by :
- **$\min\{\text{score of model 1}, \text{score of model 2}\}$**
- Where the score is the residue (square of the errors) + a penalty.
- Model 1 is the flat model and model 2 is a circadian (rhythmic model) :
- Residue of a block :  $\sum_{i=1}^n (M_i - \hat{Y}_{ji})^2$ , for  $j = 1, 2$  (for model 1 and 2), M is the matrix of RNA-Seq data and  $\hat{Y}_j$  is the predictor matrix

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- Flat model :  $\hat{Y} = \textit{mean}(M_i)$
- Circadian model :  $\hat{Y} = \alpha \sin(2\pi f) + \beta \cos(2\pi f)$

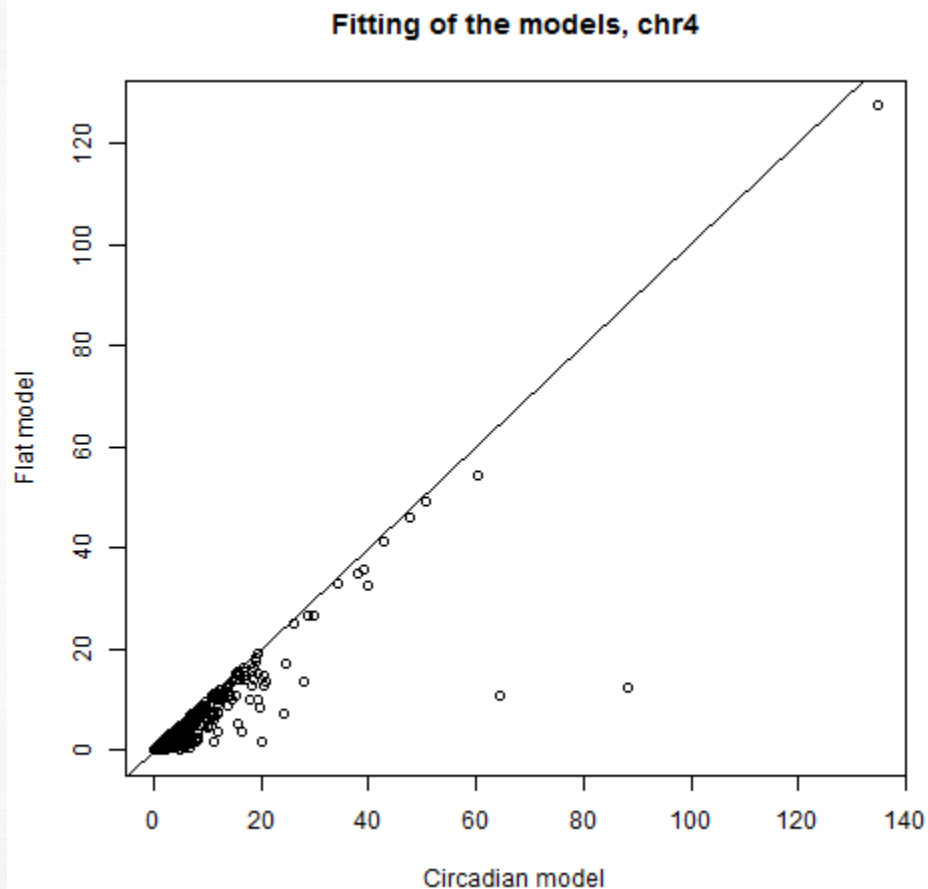


Real partitioning

## Model selection

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Circadian model always better or equal than the flat model !



Real partitioning

## Model selection

---

- Have to introduce a penalty : Score = residue + penalty
- $\text{Penalty} = \sigma^2 * (k + 1) * \log(N)$
- Where  $m$  is the number of parameters,  $N$  the total number of data points and  $\sigma$  a customizable value.

## Model selection

---

- Have to introduce a penalty :  $\text{Score} = \text{residue} + \text{penalty}$
- $\text{Penalty} = \sigma^2 * (k + 1) * \log(N)$
- Where  $\mathbf{m}$  is the number of parameters,  $\mathbf{N}$  the total number of data points and  $\sigma$  a customizable value.
- For example, for chromosome 7 containing 1139 chromosomes and 24 time points, the penalty for model 2 would be:
- $\text{Penalty} = \sigma^2 * ((1139 + 2) + 1) * \log(1139 * 24)$

Real partitioning

## Model selection

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- Instead of using an abstract sigma value, we compute a value of sigma given **the percentage of expected circadian blocks**
- We use the inverse cumulative function to find it :

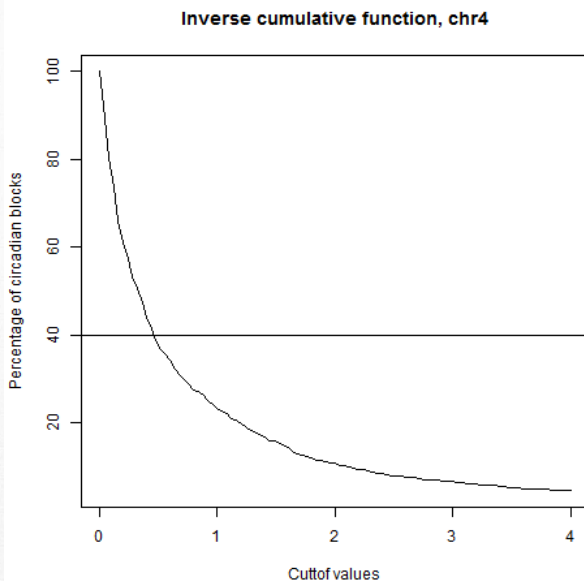


Real partitioning

## Model selection

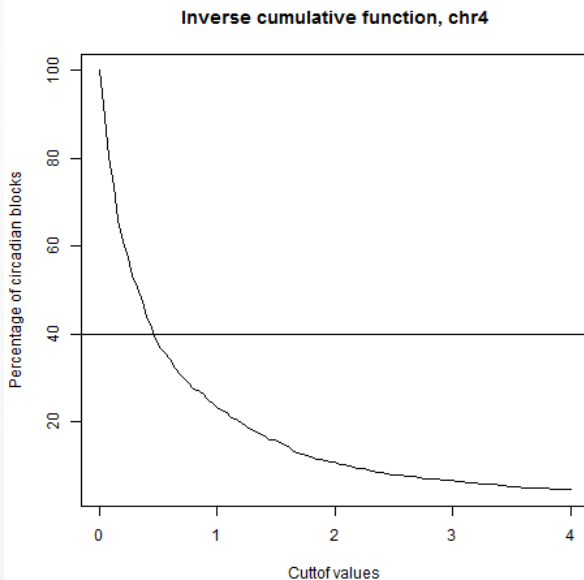
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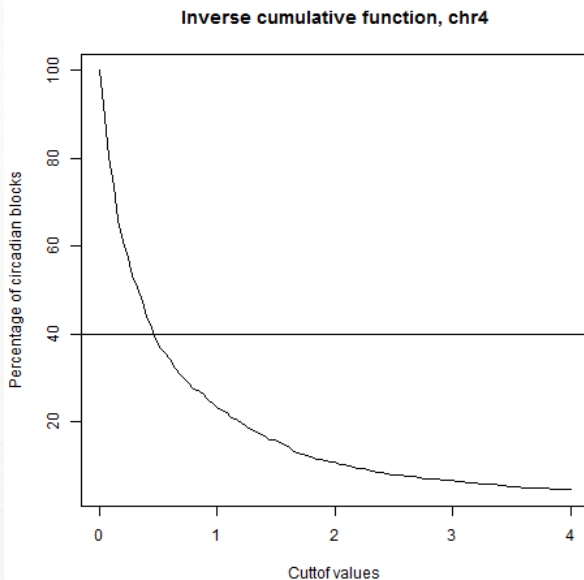
- Instead of using an abstract sigma value, we compute a value of sigma given **the percentage of expected circadian blocks**
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- **Cutoff** =  $\text{quantile}(\text{res.1} - \text{res.2}, 1 - x)$ ,
- Where **res.1** and **res.2** are the residues of the flat model and circadian model and **x** is the expected percentage of circadian blocks



Real partitioning

## Model selection

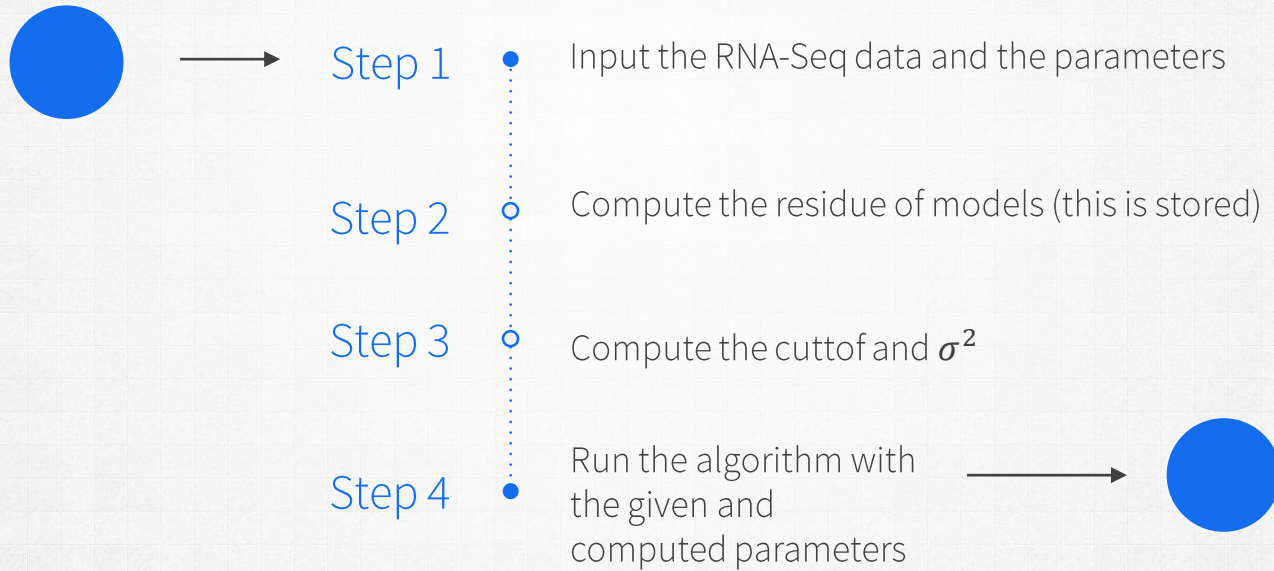
- Instead of using an abstract sigma value, we compute a value of sigma given the percentage of expected circadian blocks
- We use the inverse cumulative function to find it :
- Cutoff 40% = 0.4621029
- $\sigma^2 = \frac{cutoff}{2 * (\log(N))} = 0.02289724$



Complete algorithm

# Pipeline of computation

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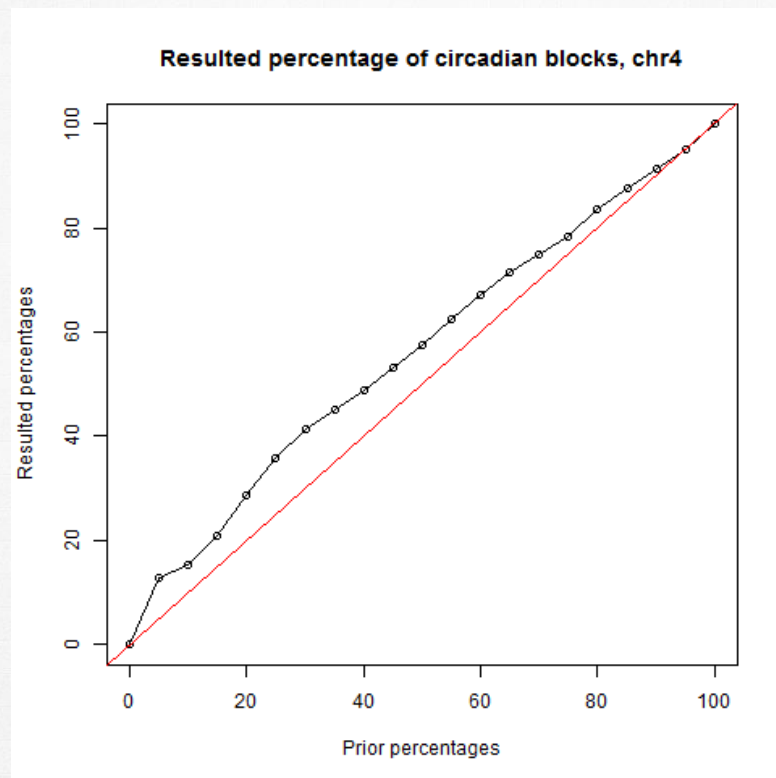




Results

# Resulted percentages of circadian blocks

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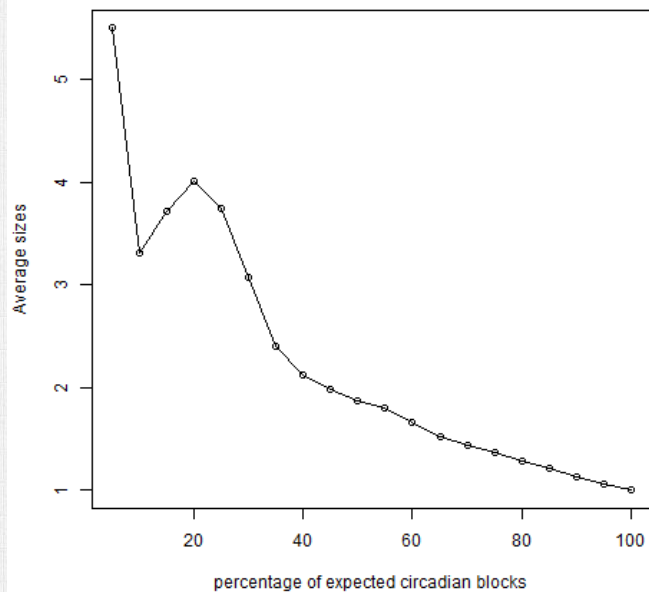


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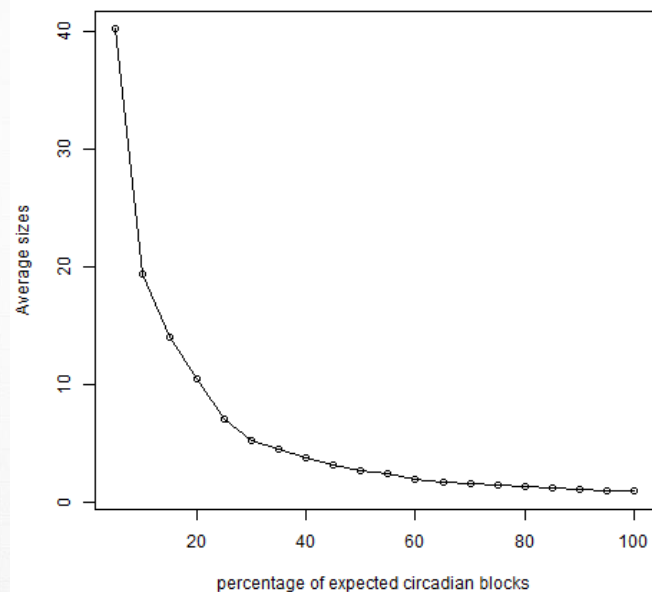
# Average size of partitions

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Average size of circadian partitions chr4

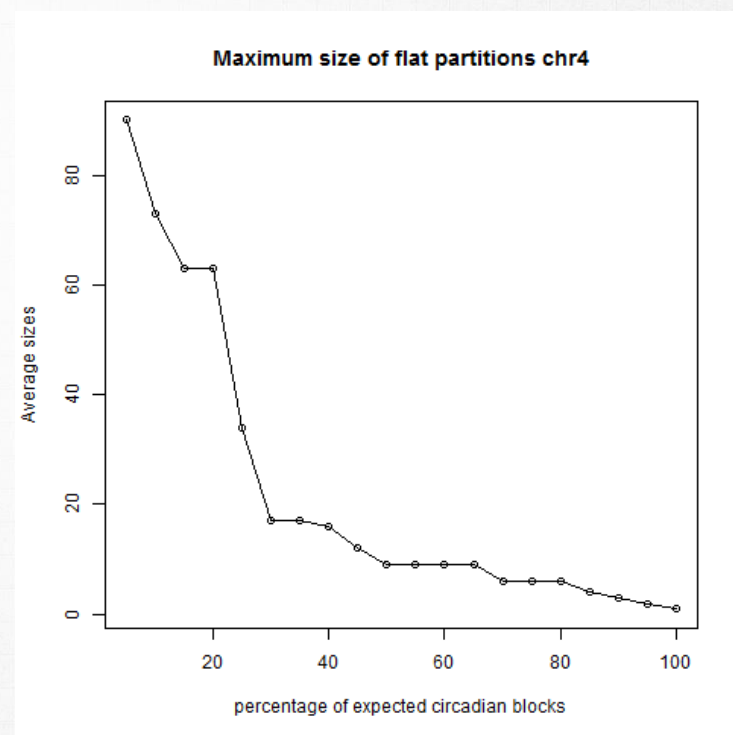
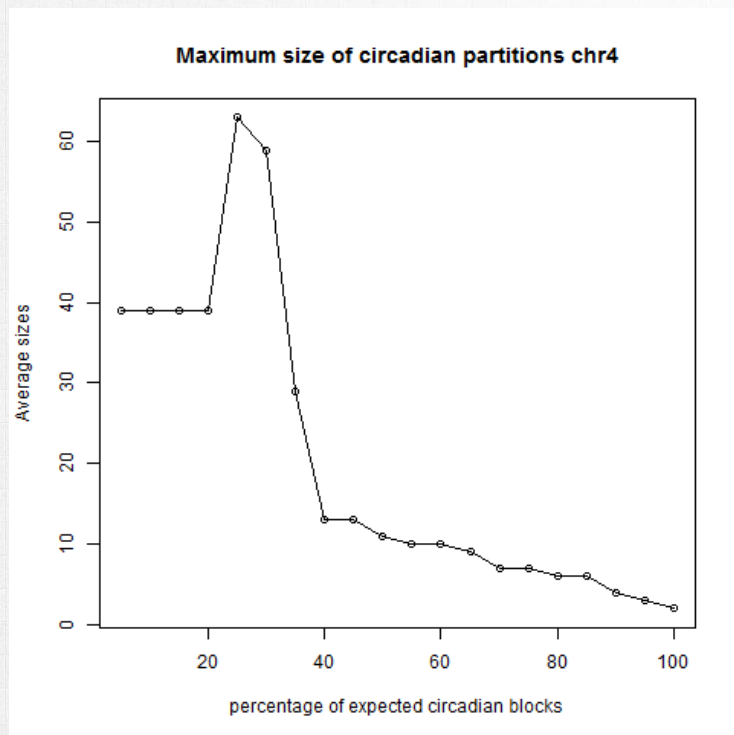


Average size of flat partitions chr4



Results

# Maximum sizes of partitions



Results

## Do clusters of genes actually exist ?

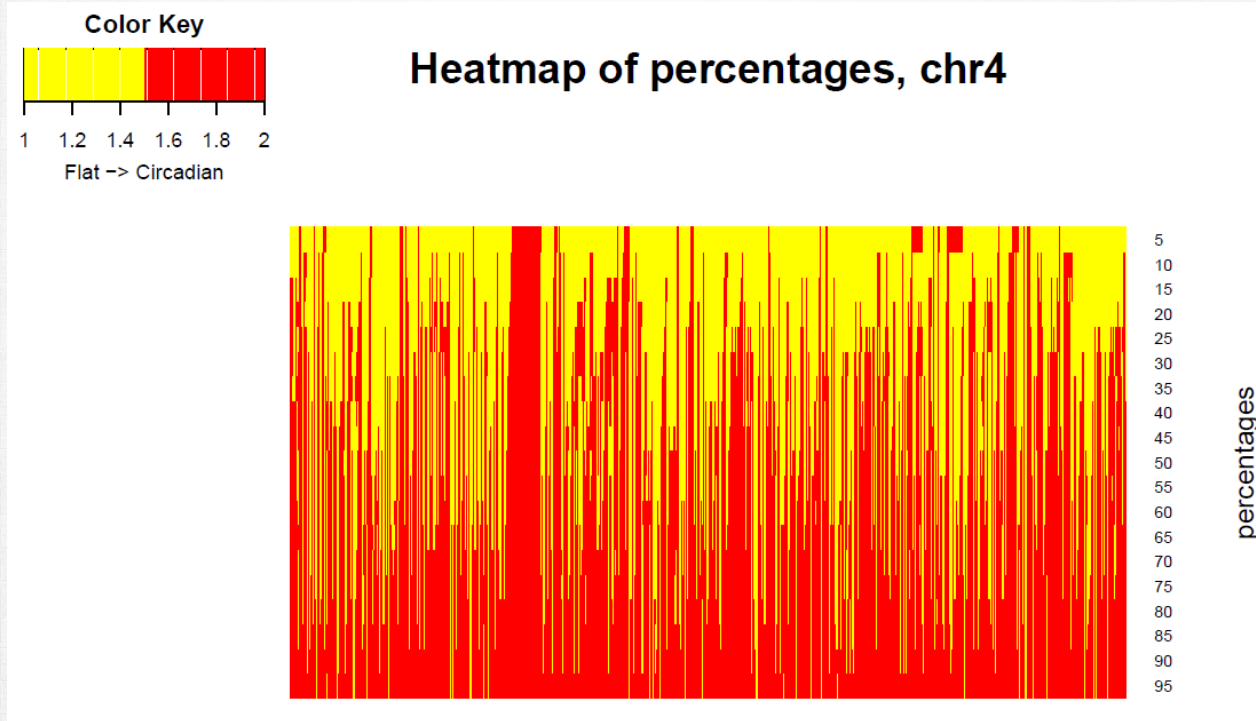
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- We ran the algorithm for chromosomes with randomly permuted genes
- 50 times each
- Comparison of sizes of circadian partitions



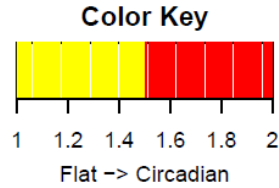
Results

## Heatmap percentages 5 to 95

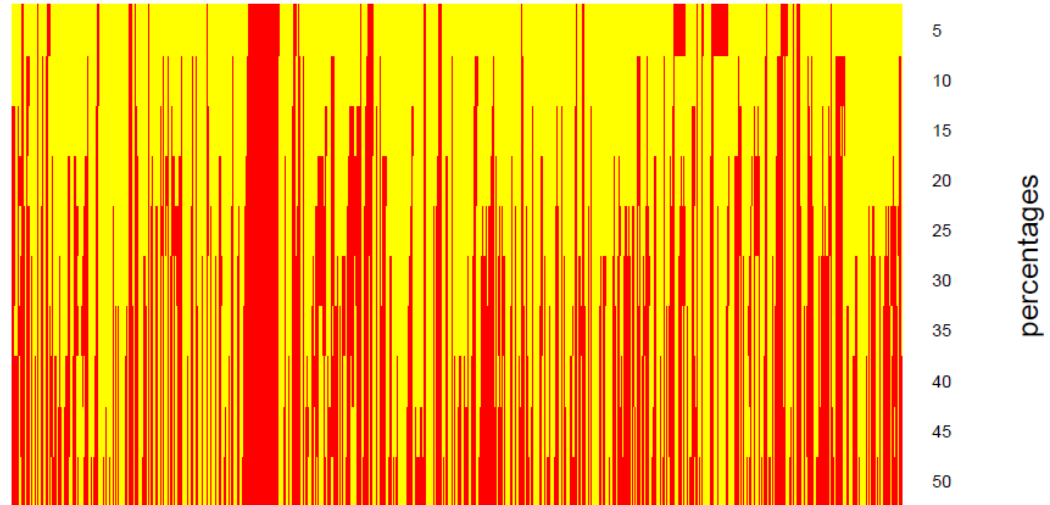


Results

## Heatmap percentages 5 to 50

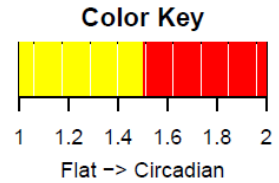


**Heatmap of percentages, chr4**

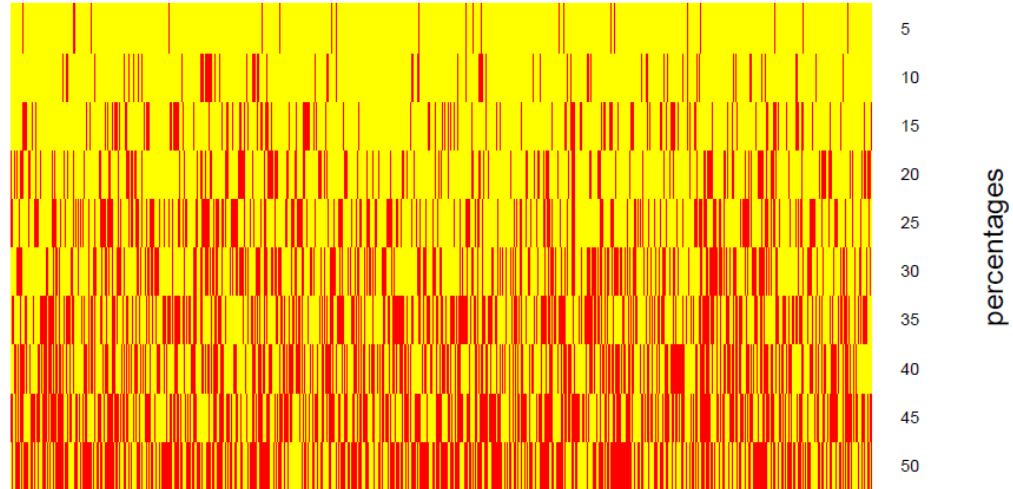


Results

# Heatmap randomized chromosome

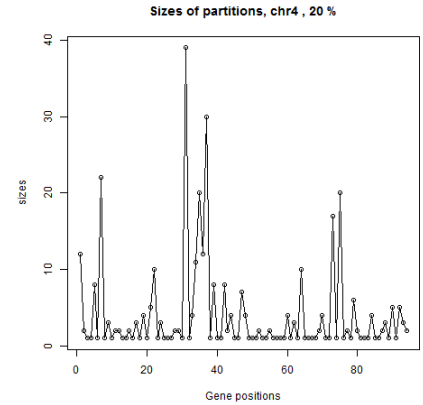
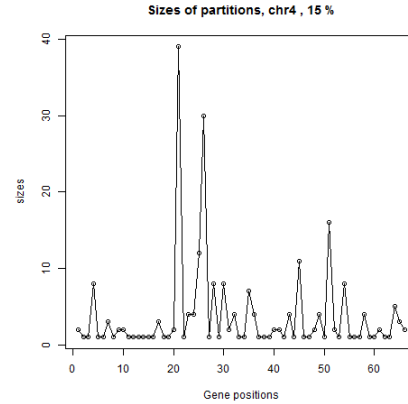
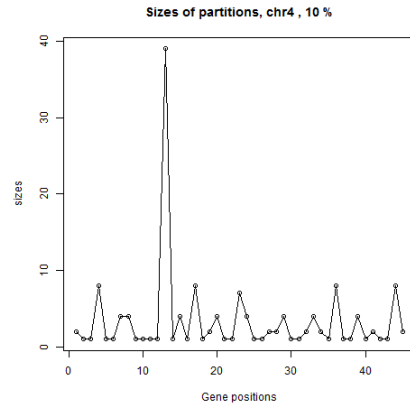
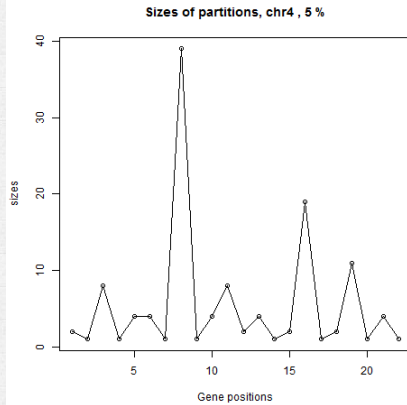


**Heatmap of percentages, random, chr4**



Results

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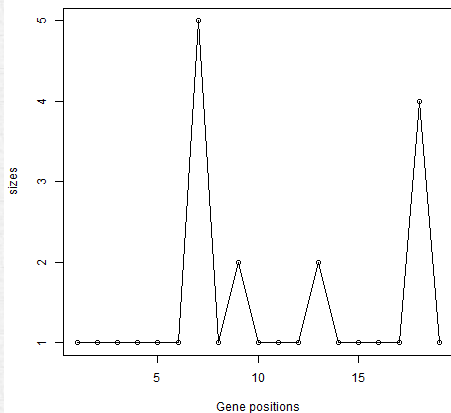




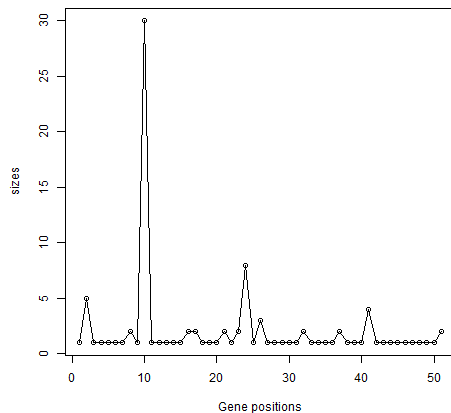
Results

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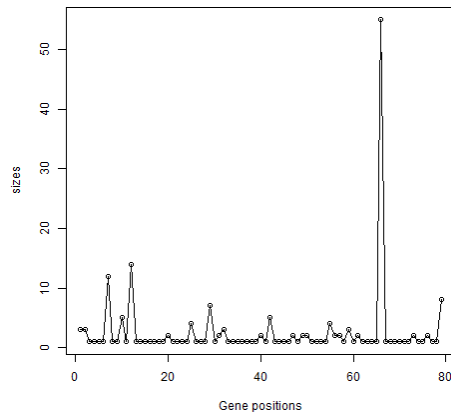
Sizes of randomized partitions, chr4 , 5 %



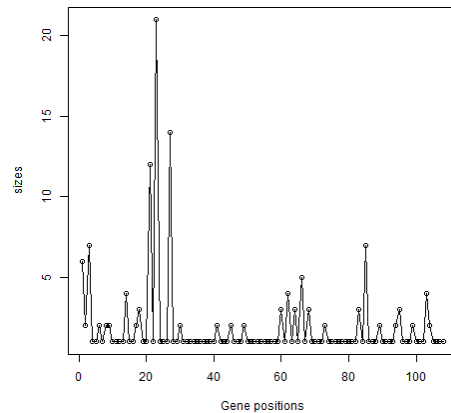
Sizes of randomized partitions, chr4 , 10 %



Sizes of randomized partitions, chr4 , 15 %



Sizes of randomized partitions, chr4 , 20 %



A final review

# CONCLUSIONS

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- Seems like the positioning is not random (heatmap)
- But the partitioning seems doesn't seem optimal → Penalty is the problem ?
- Seems Genes are actually placed in clusters along the chromosome

A final review

# FUTURE RESEARCH

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- Find a better way to penalize models
- Increase number of models
- Use different data with to do the partitioning (we can prove that genes are positioned in clusters using different data)
- Use the randomized data to prove that clusters do exist

THANK YOU