

Homework #3: Proof

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INSTRUCTIONS: before completing your homework, make sure you have considered the following:

- You are allowed to use only primitive inference rules.
- Remember to include vertical lines to represent the scope any assumptions.
- Remember that all assumptions in your proofs must be discharged.
- Double-check that your line references are correct when applying inference rules.

I. Proof: Propositional Logic

In this section *use only the primitive inference rules* of propositional calculus.

1. Provide derivations for each of the following, using natural deduction:

- a. $p \wedge q, p \Rightarrow s, q \Rightarrow t \vdash s \wedge t$
1. $p \wedge q$ premise
 2. $p \Rightarrow s$ premise
 3. $q \Rightarrow t$ premise
 4. p \wedge -elim, 1
 5. s modus ponens, 2,4
 6. q \wedge -elim, 1
 7. t modus ponens, 3,6
 8. $s \wedge t$ \wedge -intro, 5,7

- b. $q \Rightarrow \neg p, p \wedge q \vdash r$
1. $q \Rightarrow \neg p$ premise
 2. $p \wedge q$ premise
 3. q \wedge -elim, 2
 4. $\neg p$ modus ponens, 1
 5. p \wedge -elim, 2
 6. $\neg r$ assumption 1
 7. $\neg p$ Copy from 4 1
 8. p Copy from 5 1
 4. r \neg -elim, 6-8

- c. $p \wedge q \vdash p \Rightarrow q$
1. $p \wedge q$ assumption 1
 2. p \wedge -elim, 1 1
 3. q \wedge -elim, 1 1
 4. $p \Rightarrow q$ \Rightarrow -intro, 1-3, 2, 3



- d. $\neg\neg q \vdash q \vee r$
- | | | |
|-----------------|-----------------------|---|
| 1. $\neg\neg q$ | assumption | 1 |
| 2. $\neg q$ | $\neg - elim, 1$ | 1 |
| 3. q | $\neg - elim, 2$ | 1 |
| 4. $q \vee r$ | $\vee - intro, 1 - 3$ | |



- e. $p \Rightarrow (q \wedge r), \neg p \Rightarrow r, p \vee \neg p \vdash r$
- | | | | |
|---------------------------------|----------------------|---|---|
| 1. $\neg p \Rightarrow r$ | premise | | |
| 2. $p \vee \neg p$ | assumption | 1 | |
| 3. p | $\vee - elim, 2$ | 1 | |
| 4. $p \Rightarrow (q \wedge r)$ | assumption | 1 | 1 |
| 5. $(q \wedge r)$ | modus ponens, 4, 3 | 1 | 1 |
| 6. r | $\wedge - elim, 2-5$ | | |



II. Proof: Predicate Logic

In this section *use only the primitive inference rules* of predicate calculus.

2. Show using natural deduction:

- a. $\forall x : T \bullet P(x) \wedge Q(x) \dashv\vdash (\forall x : T \bullet P(x)) \wedge (\forall y : T \bullet Q(y))$

$$\forall x : T \bullet P(x) \wedge Q(x) \vdash (\forall x : T \bullet P(x)) \wedge (\forall y : T \bullet Q(y))$$

- | | | |
|--|---------------------------|-------|
| 1. $\forall x : T \bullet P(x) \wedge Q(x)$ | assumption | 1 |
| 2. $x \in T$ | assumption | 1 1 |
| 3. $P(x) \wedge Q(x)$ | $\forall - elim, 1$ | 1 1 |
| 4. $P(x)$ | $\wedge - elim, 3$ | 1 1 |
| 5. $Q(x)$ | $\wedge - elim, 3$ | 1 1 |
| 6. $y \in T$ | assumption | 1 1 1 |
| 7. $Q(y)$ | transforming x to y, 5, 6 | 1 1 1 |
| 8. $\forall y : T \bullet Q(y)$ | $\forall - intro 6, 7$ | 1 1 |
| 9. $\forall x : T \bullet P(x)$ | $\forall - intro 2, 4$ | 1 1 |
| 10. $\forall x : T \bullet P(x) \wedge (\forall y : T \bullet Q(y))$ | $\wedge - intro 8, 9$ | |



$$(\forall x : T \bullet P(x)) \wedge (\forall y : T \bullet Q(y)) \vdash \forall x : T \bullet P(x) \wedge Q(x)$$

- | | | |
|---|--------------------------|-------|
| 1. $(\forall x : T \bullet P(x)) \wedge (\forall y : T \bullet Q(y))$ | assumption | 1 |
| 2. $\forall x : T \bullet P(x)$ | | 1 |
| 3. $\forall y : T \bullet Q(y)$ | | 1 |
| 4. $a \in T$ | assumption | 1 1 |
| 5. $b \in T$ | assumption | 1 1 1 |
| 6. $P(a)$ | $\forall - elim, 2, 4$ | 1 1 1 |
| 7. $Q(b)$ | $\forall - elim, 3, 5$ | 1 1 1 |
| 8. $P(a) \wedge Q(b)$ | $\wedge - intro, 6, 7$ | 1 1 1 |
| 9. $\forall x : T \bullet P(x) \wedge Q(x)$ | $\forall - intro, 4 - 8$ | |

- b. $\exists x : T \bullet P(x) \vee Q(x) \dashv\vdash (\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x))$

$$\exists x : T \bullet P(x) \vee Q(x) \vdash (\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x))$$

1. $\exists x : T \bullet P(x) \vee Q(x)$	assumption	1
2. $a \in T \wedge P(a) \vee Q(a)$	assumption	1 1
3. $P(a) \vee Q(a)$	$\wedge -elim, 2$	1 1
4. $P(a)$	$\vee -elim, 3$	1 1
5. $Q(a)$	$\vee -elim, 3$	1 1
6. $a \in T$	$\wedge -elim, 2$	1 1
6. $\exists x : T \bullet P(x)$	$\exists -intro, 6, 4$	1 1
7. $\exists x : T \bullet Q(x)$	$\exists -intro, 6, 5$	1 1
8. $(\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x))$	$\vee -intro, 6, 7$	



$(\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x)) \vdash \exists x : T \bullet P(x) \vee Q(x)$		
1. $(\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x))$	assumption	1
2. $\exists x : T \bullet P(x)$	$\vee -elim, 1$	1
3. $\exists x : T \bullet Q(x)$	$\vee -elim, 1$	1
4. $a \in T \wedge P(a)$	assumption	1 1
5. $P(a)$	$\wedge -elim, 4$	1 1
6. $b \in T \wedge Q(b)$	assumption	1 1 1
7. $Q(b)$	$\wedge -elim, 6$	1 1 1
8. $b \in T$	$\wedge -elim, 6$	1 1 1
9. $P(a) \vee Q(b)$	$\vee -intro, 2 - 4$	1 1 1
10. $\exists x : T \bullet P(x) \vee Q(x)$	$\exists -intro, 8, 9$	1 1
11. $\exists x : T \bullet P(x) \vee Q(x)$	$\exists -elim, 2, 4, 10$	
12. $\exists x : T \bullet P(x) \vee Q(x)$	$\exists -elim, 3, 6 - 10$	

(NOTE: $p \dashv\vdash q$ is a shorthand for “ $p \vdash q$ and $q \vdash p$.” That is, for $p \dashv\vdash q$ you need to show two separate derivations: one for $p \vdash q$ and another for $q \vdash p$.)