CMU 17-651 Models of Software Systems Fall 2016

Homework #3: Proof

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INSTRUCTIONS: before completing your homework, make sure you have considered the following:

- You are allowed to use only primitive inference rules.
- Remember to include vertical lines to represent the scope any assumptions.
- Remember that all assumptions in your proofs must be discharged.
- Double-check that your line references are correct when applying inference rules.

I. Proof: Propositional Logic

In this section *use only the primitive inference rules* of propositional calculus.

1. Provide derivations for each of the following, using natural deduction:

a. $(p \wedge q) \wedge r \vdash p \wedge (q \wedge r)$

b. $p \wedge q, p \Rightarrow s, q \Rightarrow t \vdash s \wedge t$

c. $q \Rightarrow \neg p, p \wedge q \vdash r$

1.	$q \Rightarrow \neg p$	premise	
2.	$p \wedge q$	premise	
3.	p	$\wedge - elim, 2$	
4.	$\neg p$	$modusponens, 1,$	
5,	$\neg p$	$\Rightarrow - elim$	
6.	$\neg r$	assumption	l
7.	p	Copy from 3	l
8.	$\neg p$	Copy from 4	l
4.	r	$\neg -$ *elim*, 6-8	

d. $p \wedge q \vdash p \Rightarrow q$

1.	$p \wedge q$	premise
2.	p	$\wedge - elim$,1
3.	q	$\wedge - elim$,1
4.	$p \Rightarrow q$	$\Rightarrow - intro$,2,3

e. $\neg\neg q \vdash q \vee r$

1.	$\neg\neg q$	premise
2.	$\neg q$	$\neg -$ *elim*, 1
3.	q	$\neg -$ *elim*, 2
4.	$q \vee r$	$\vee - intro, 3$

f. $p \Rightarrow (q \wedge r), \neg p \Rightarrow r, p \vee \neg p \vdash r$

II. Proof: Predicate Logic

In this section *use only the primitive inference rules* of predicate calculus.

2. Show using natural deduction:

   a. $\forall x : T \bullet P(x) \wedge Q(x) \dashv\vdash (\forall x : T \bullet P(x)) \wedge (\forall y : T \bullet Q(y))$

   b. $\exists x : T \bullet P(x) \vee Q(x) \dashv\vdash (\exists x : T \bullet P(x)) \vee (\exists x : T \bullet Q(x))$

   (NOTE: $p \dashv\vdash q$ is a shorthand for "$p \vdash q$ and $q \vdash p$." That is, for $p \dashv\vdash q$ you need to show two separate derivations: one for $p \vdash q$ and another for $q \vdash p$.)