

## LINKED LISTS (Background & Intro)

### Problem with Arrays

① How to implement Round-Robin Scheduling?

$P_0$	$P_1$	$P_2$	$P_3$	$P_4$	$P_5$
10	5	3	15	10	8

Time = 5

5	3	15	10	8	5
$P_1$	$P_2$	$P_3$	$P_4$	$P_5$	$P_0$

3	15	10	8	5
$P_2$	$P_3$	$P_4$	$P_5$	$P_0$

$P_1$  completes

15	10	8	5
$P_3$	$P_4$	$P_5$	$P_0$

$P_2$  completes

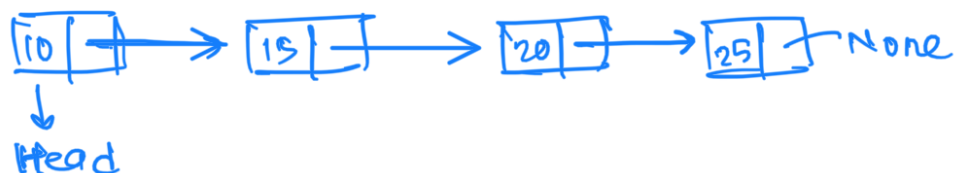
This becomes difficult with arrays, as you have to pick out item & insert in the end.

② Given a sequence of items, whenever we see an item  $x$  in the sequence we need to replace it with  $B$  instance of 'y'

I/P  $\rightarrow$  d e a x q r x p y

O/P  $\rightarrow$  d e a y y y y y q r y y y y y p y

### Linked Lists



Linked list structure

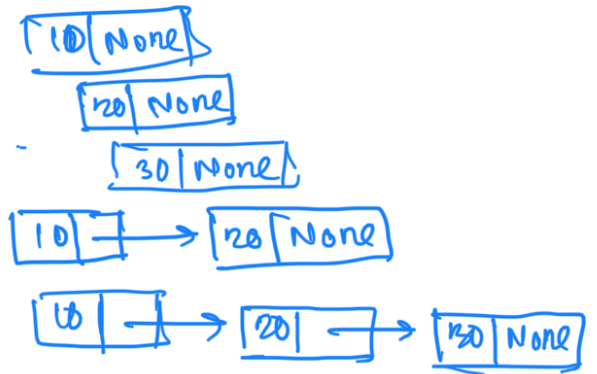
⇒ simple implementation



class Node :

```
def __init__(self, k):  
    self.key = k  
    self.next = None
```

```
temp1 = Node(10)  
temp2 = Node(20)  
temp3 = Node(30)  
temp1.next = temp2  
temp2.next = temp3  
head = temp1
```



shorter implementation

```
head = Node(10)  
head.next = Node(20)  
head.next.next = Node(30)
```

## Applications of LL

- 1) Insertion at end & begin are  $O(1)$ .
- 2) Insertions & deletions in middle are  $O(N)$ . If we have reference to previous node.
- 3) Worst case deletion from beginning is  $O(1)$ .
- 4) Round Robin implementation.
- 5) Merging two sorted linked list is faster than arrays.
- 6) Implementation of simple memory manager where we need to link free blocks.
- 7) Easier implementation of Queue & Deque data structures.

Traversal



```

head
def printList(head):
    curr = head
    while curr != None:
        print(curr.key, end = " ")
        curr = curr.next

```

### Insertion at the beginning

class Node:

```

def __init__(self, key):
    self.key = key
    self.next = None

def insertBegin(head, key):
    temp = Node(key)
    temp.next = head
    return temp

```

### Insertion at the end

```

def insertEnd(head, key):
    if head == None:
        return Node(key)
    curr = head
    while curr.next != None:
        curr = curr.next
    curr.next = Node(key)
    return head

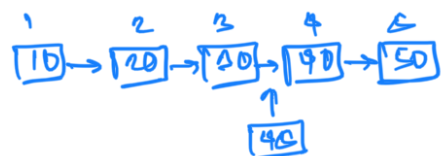
```

### Insertion at given position

```

def insertPos(head, data, pos):
    temp = Node(data)
    if pos == 1:
        temp.next = head
        return temp
    curr = head
    for i in range(pos-2):
        curr = curr.next
        if curr == None:

```



$\ominus \lfloor \min(pos, n) \rfloor$

```

return head
temp.next = curr.next
curr.next = temp
return head

```

### Delete first node

```

def deleteFirst(head):
    if head == None:
        return None
    else:
        return head.next

```

### Delete last node

```

def deleteLastNode(head):
    if head == None:
        return None
    if head.next == None:
        return None
    curr = head
    while curr.next.next != None:
        curr = curr.next
    curr.next = None
    return head

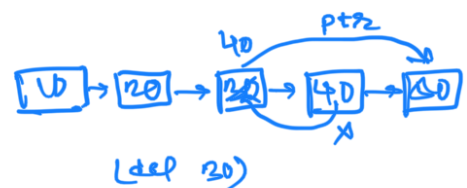
```

### Delete Node using pointer

```

def deleteNode(ptr):
    temp = ptr.next
    ptr.data = temp.data
    ptr.next = temp.next

```

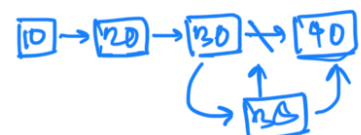


### Insertion in sorted list

```

def sortedInsert(head, val):
    temp = Node(val)
    if head == None:
        return temp
    elif val < head.data:
        temp.next = head

```



```

return temp
else:
    curr = head
    while curr.next != None & curr.next.data < val:
        curr = curr.next
    temp.next = curr.next
    curr.next = temp
    return head

```

## Middle of linked list

Brute  
Force

- Iterate through and count nodes.
- 2<sup>nd</sup> pass, iterate till count//2 & print

Optim  
Approach

- Two pointers, slow & fast

```
def printMiddle(head):
```

```
    if head == None:
```

```
        return
```

```
    slow = head
```

```
    fast = head
```

```
    while fast != None and fast.next != None:
```

```
        slow = slow.next
```

```
        fast = fast.next.next
```

```
    print(slow.data)
```