

MTH 9821 Numerical Methods for Finance I

Lecture 2

September 12, 2013

1 Cholesky Decomposition

Definition:

Given a symmetric positive definite nonsingular matrix A , there is a unique upper triangular matrix U with positive entries on the main diagonal such that

$$A = U^t U$$

Remark:

- ★ In Cholesky decomposition, we try to solve $\frac{n(n+1)}{2}$ unknowns from $\frac{n(n+1)}{2}$ equations (upper triangular part).
- ★ In LU decomposition, we try to solve n^2 unknowns from n^2 equations.

Lemma: A nonsingular symmetric matrix A has Cholesky decomposition, then A is spd.

Proof:

It's enough to show that $x^t A x > 0, \forall x \neq 0$.

Let U be a nonsingular upper triangular matrix such that $A = U^t U$.

$$x^t A x = x^t U^t U x = (Ux)^t Ux = \|Ux\|^2 \geq 0 \quad \forall x \text{ with } Ux \neq 0$$

Note: $Ux = 0 \Leftrightarrow x = 0$ since U is nonsingular.

Lemma: If A is spd, then there exists an upper triangular nonsingular matrix U such that $A = U^t U$ (and $U(i, i) > 0, \forall i = 1 : n$).

We can proof by solving $A = U^t U$ for a 4×4 example.

$$\underbrace{\begin{pmatrix} A(1,1) & A(1,2) & A(1,3) & A(1,4) \\ A(2,1) & A(2,2) & A(2,3) & A(2,4) \\ A(3,1) & A(3,2) & A(3,3) & A(3,4) \\ A(4,1) & A(4,2) & A(4,3) & A(4,4) \end{pmatrix}}_A = \underbrace{\begin{pmatrix} U(1,1) & 0 & 0 & 0 \\ U(1,2) & U(2,2) & 0 & 0 \\ U(1,3) & U(2,3) & U(3,3) & 0 \\ U(1,4) & U(2,4) & U(3,4) & U(4,4) \end{pmatrix}}_{U^t} \underbrace{\begin{pmatrix} U(1,1) & U(1,2) & U(1,3) & U(1,4) \\ 0 & U(2,2) & U(2,3) & U(2,4) \\ 0 & 0 & U(3,3) & U(3,4) \\ 0 & 0 & U(4,3) & U(4,4) \end{pmatrix}}_U$$

$$U(1,1)^2 = A(1,1) \implies U(1,1) = \sqrt{A(1,1)}$$

(Since A is spd, $A(1,1) = e_1^t A e_1 > 0$)

$$U(1,k) = \frac{A(1,k)}{U(1,1)} \quad \forall k = 1 : 4 \implies U(1,2:4) = \frac{A(1,2:4)}{U(1,1)} = \frac{A(1,2:4)}{\sqrt{A(1,1)}}$$

Denote $U(2:4, 2:4) = U_1$, we have

$$\begin{aligned}
U_1^t U_1 &= A(2:4, 2:4) - (U(1, 2:4))^t U(1, 2:4) \\
&= A(2:4, 2:4) - \frac{A(1,2:4)^t}{\sqrt{A(1,1)}} \frac{A(1,2:4)}{\sqrt{A(1,1)}} = \frac{A(1,2:4)^t A(1,2:4)}{A(1,1)} \\
\left(U_1^t U_1 = A(2:n, 2:n) - \frac{A(1,2:n)^t A(1,2:n)}{A(1,1)} \right)
\end{aligned}$$

In order to show that the matrix U such that $U^t U = A$ exists, it's equal to show that the matrix $U_1^t U_1$ is spd, given A is spd.

Proof is shown in Class Handout Lemma 11.4.

Pseudocode: *Cholesky Decomposition*

Input:

A = nonsingular symmetric positive definite matrix of size n

Output:

U = upper triangular matrix

such that $U^t U = A$

for $i = 1 : (n - 1)$

$$U(i, i) = \sqrt{A(i, i)};$$

for $k = i + 1 : n$

$$U(i, k) = \frac{A(i, k)}{U(i, i)};$$

end

for $j = (i + 1) : n$

for $k = j : n$

$$A(j, k) = A(j, k) - U(i, j)U(i, k);$$

end

end

end

$$U(n, n) = \sqrt{A(n, n)}$$

★ Operation Count of Cholesky Decomposition: $\frac{n^3}{3} + O(n^2)$
 ★ Operation saving comes from $k = j : n$, only need to update the upper triangular part.

$$\begin{aligned}
\text{Operation Count} &= \sum_{i=1}^{n-1} \left((n-i) + \sum_{j=i+1}^n \sum_{k=j}^n 2 \right) \\
&= \sum_{i=1}^{n-1} \left((n-i) + \sum_{j=i+1}^n 2(n-j+1) \right) \\
(\text{Let } p = n-j+1) &= \sum_{i=1}^{n-1} \left((n-i) + \sum_{p=1}^{n-i} p \right) \\
&= \sum_{i=1}^{n-1} [(n-1) + (n-i)^2 + (n-i)] \\
(\text{Let } q = n-i) &= \sum_{q=1}^{n-1} q^2 + 2q \\
&= \frac{(n-1)n(2n-1)}{6} + 2 \cdot \frac{n(n-1)}{2} \\
&= \frac{2n^3}{6} + O(n^2) = \frac{n^3}{3} + O(n^2)
\end{aligned}$$

Uniqueness: Cholesky decomposition of a spd matrix is unique.
Proof is similar to the uniqueness proof of LU decomposition.

2 Use of Cholesky Decomposition

Given a matrix A ,

$$A \begin{cases} \text{symmetric} & \longrightarrow \begin{cases} \text{Success} \\ \text{Fail} \longrightarrow \text{LU with row pivoting} \end{cases} \\ \text{asymmetric} & \longrightarrow \text{LU with row pivoting} \end{cases}$$

2.1 Using Cholesky Decomposition for Solving Linear System

Solve $\mathbf{Ax} = \mathbf{b}$

\mathbf{A} : $n \times n$ spd matrix given;

\mathbf{b} : $n \times 1$ vector given;

Find \mathbf{x} : $n \times 1$ vector

Thus,

$$\begin{aligned} Ax &= b \\ \iff U^t Ux &= b \\ \iff U^t y &= b, \quad Ux = y \end{aligned}$$

$\mathbf{x} = \text{linear_solve_LU_cholesky}(\mathbf{A}, \mathbf{b})$

$$\begin{aligned} U &= \text{Cholesky}(A) \\ y &= \text{forward_subst}(U^t, b) \\ x &= \text{backward_subst}(U, y) \\ (x &= \text{backward_subst}(U, \text{forward_subst}(U^t, b))) \end{aligned}$$

2.2 Solve p Linear Systems

$$Ax_i = b_i, \quad \forall i = 1 : p$$

$$\begin{aligned} U &= \text{Cholesky}(A); \\ \text{for } i &= 1 : p \\ y &= \text{forward_subst}(U^t, b_i) \\ x_i &= \text{backward_subst}(U, y) \\ \text{end} \end{aligned}$$

Remark:

- ★ Implement Cholesky before the for loop
- ★ Operation count: $2pn^2 + \frac{n^3}{3} + O(n^2)$

2.3 Least Square

A is an $m \times n$ matrix with $m > n$, it has linearly independent columns.

Least square solution:

$$\begin{aligned} \text{Solve } \underbrace{A^t A}_{m \times m \text{ spd matrix}} x &= A^t b \\ \min \|b - Ax\|, \quad x &\in \mathbb{R}^n \\ \mathbf{x} &= \text{linear_solve_cholesky}(\mathbf{A}^t \mathbf{A}, \mathbf{A}^t \mathbf{y}) \end{aligned}$$

2.4 Find Normal Random Sample with Correlating

- Find normal random variables X_1, X_2, \dots, X_n with spd covariance matrix A .

$$\begin{pmatrix} X_1 \\ \vdots \\ X_n \end{pmatrix} = U^t \begin{pmatrix} Z_1 \\ \vdots \\ Z_n \end{pmatrix}$$

– Linear Transformation Property: $Y = MX \Rightarrow \Sigma_Y = MM^t \Sigma_X$.

– n=2 case:

$$\begin{pmatrix} X_1 \\ X_2 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ \rho & \sqrt{1-\rho^2} \end{pmatrix} \begin{pmatrix} Z_1 \\ Z_2 \end{pmatrix} \Rightarrow \Sigma_X = \begin{pmatrix} 1 & \rho \\ \rho & 1 \end{pmatrix} = \Omega_X$$

- Given samples Z_1, \dots, Z_{2N} of Z , the pairs $(Z_{2i+1}, \rho Z_{2i+1} + \sqrt{1-\rho^2} Z_{2i+2})$, $i = 0 : (N-1)$ are independent and with correlation.

– Remark: This can be used for Monte Carlo simulation. Brownian motion cannot be used here because random variables generated have no correlation.

3 Cholesky Decomposition for Banded Matrices

A is a banded matrix of band m iff $A(j, k) = 0 \forall |j - k| > m$.

Lemma: A is spd of band $m \Rightarrow$ Cholesky factor U of A is also banded with band m . e.g., 8×8 matrix with band 3

$$\begin{pmatrix} x & x & x & x & 0 & 0 & 0 & 0 \\ x & x & x & x & x & 0 & 0 & 0 \\ x & x & x & x & x & x & 0 & 0 \\ x & x & x & x & x & x & x & 0 \\ 0 & x & x & x & x & x & x & x \\ 0 & 0 & x & x & x & x & x & x \\ 0 & 0 & 0 & x & x & x & x & x \\ 0 & 0 & 0 & 0 & x & x & x & x \end{pmatrix}$$

Remark:

- ★ In the above 8×8 matrix, only shaded entries need to be updated in this step.
- ★ m^2 entries need to be updated

Pseudocode: *Cholesky Decomposition for banded matrices*

Input:

A = nonsingular symmetric positive definite banded matrix of banded m and size n

Output:

U = upper triangular matrix

such that $U^t U = A$

```
for  $i = 1 : (n - 1)$ 
     $U(i, i) = \sqrt{A(i, i)}$ ;
    for  $k = i + 1 : \min(i + m, n)$ 
         $U(i, k) = \frac{A(i, k)}{U(i, i)}$ ;
    end
    for  $j = (i + 1) : \min(i + m, n)$ 
        for  $k = j : \min(i + m, n)$ 
             $A(j, k) = A(j, k) - U(i, j)U(i, k)$ ;
        end
    end
end
 $U(n, n) = \sqrt{A(n, n)}$ 
```

[Operation Count: nm^2]

4 Efficient Decomposition of a Tridiagonal Matrix

$$A = \begin{pmatrix} 2 & -1 & & & \\ -1 & \ddots & \ddots & & \\ & \ddots & \ddots & \ddots & \\ & & & -1 & 2 \end{pmatrix}$$

$$\begin{pmatrix} x & x & 0 & 0 & 0 & 0 \\ x & \boxed{x} & x & 0 & 0 & 0 \\ 0 & x & x & x & 0 & 0 \\ 0 & 0 & x & x & x & 0 \\ 0 & 0 & 0 & 0 & x & x \end{pmatrix}$$

Tridiagonal: only one entry needed to be updated, i.e., no difference for Cholesky and LU. But in LU, we have an advantage: 1's on main diagonal of L .

Method 1: `x = linear_solve_cholesky_tridiagonal_spd(A, b)`

```
U = Cholesky(A)
y = forward_subst(Ut, b)
x = backward_subst(U, y)
(x = backward_subst(U, forward_subst(Ut, b)))
```

Pseudocode: *Cholesky Decomposition for Solving spd Tridiagonal Linear System*

(Cholesky Decomposition)

for $i = 1 : (n - 1)$

$$U(i, i) = \sqrt{A(i, i)};$$

$$U(i, i + 1) = \frac{A(i, i + 1)}{U(i, i)};$$

$$A(i + 1, i + 1) = A(i + 1, i + 1) - U(i, i + 1)^2;$$

end

$$U(n, n) = \sqrt{A(n, n)}$$

(Forward Substitution)

$$y(1) = \frac{b(1)}{U(1, 1)};$$

for $i = 2 : n$

$$y(i) = \frac{b(i) - U(i, i - 1)y(i - 1)}{U(i, i)}$$

end

(Backward Substitution)

$$x(n) = \frac{y(n)}{U(n, n)};$$

for $i = (n - 1) : 1$

$$x(i) = \frac{y(i) - U(i, i + 1)x(i + 1)}{U(i, i)}$$

end

[Operation Count: $4n + 3n + 3n = 10n + O(1)$]

Method 2: $\mathbf{x} = \text{linear_solve_LU_no_pivoting}(\mathbf{A}, \mathbf{b})$

```
[L,U] = lu(A)
y = forward_subst(L, b)
x = backward_subst(U, y)
(x = backward_subst(U, forward_subst(L, b)))
```

Remark:: We can do LU without pivoting for spd tridiagonal matrices, since all eigenvalue is greater than 0, then all leading principal minor is greater than 0.

Pseudocode: *LU Decomposition for Solving spd Tridiagonal Linear System*

(LU Decomposition no Pivoting)

```
for i = 1 : (n - 1)
    U(i, i) = A(i, i), U(i, i + 1) = A(i, i + 1);
    L(i + 1, i) =  $\frac{A(i+1, i)}{U(i, i)}$ , L(i, i) = 1;
    A(i + 1, i + 1) = A(i + 1, i + 1) - L(i + 1, i)U(i, i + 1);
end
U(n, n) = A(n, n), L(n, n) = 1;
```

(Forward Substitution)

```
y(1) =  $\frac{b(1)}{L(1,1)}$  = b(1);
for i = 2 : n
    y(i) =  $\frac{b(i) - L(i, i-1)y(i-1)}{L(i, i)}$  = b(i) - L(i, i - 1)y(i - 1)
end
```

(Backward Substitution)

```
x(n) =  $\frac{y(n)}{U(n, n)}$ ;
for i = (n - 1) : 1
    x(i) =  $\frac{y(i) - U(i, i+1)x(i+1)}{U(i, i)}$ 
end
```

[Operation Count: $3n + 2n + 3n = 8n + O(1)$]

Method 3: Direct Calculation of Cholesky Factor

For

$$B_N = \begin{pmatrix} 2 & -1 & & & \\ -1 & \ddots & \ddots & & \\ & \ddots & \ddots & \ddots & \\ & & & -1 & 2 \end{pmatrix}$$

Cholesky factor

$$U_N(i, i) = \sqrt{\frac{i+1}{i}};$$

$$U_N(i, i+1) = -\sqrt{\frac{i}{i+1}} = -\frac{1}{U_N(i, i)};$$

★ Operation Count for Direct Calculating Cholesky Factor $3n + n = 4n$

★ Operation Count for Method 3 $4n + 3n + 3n = 10n + O(1)$, i.e., no saving than implementing Cholesky decomposition