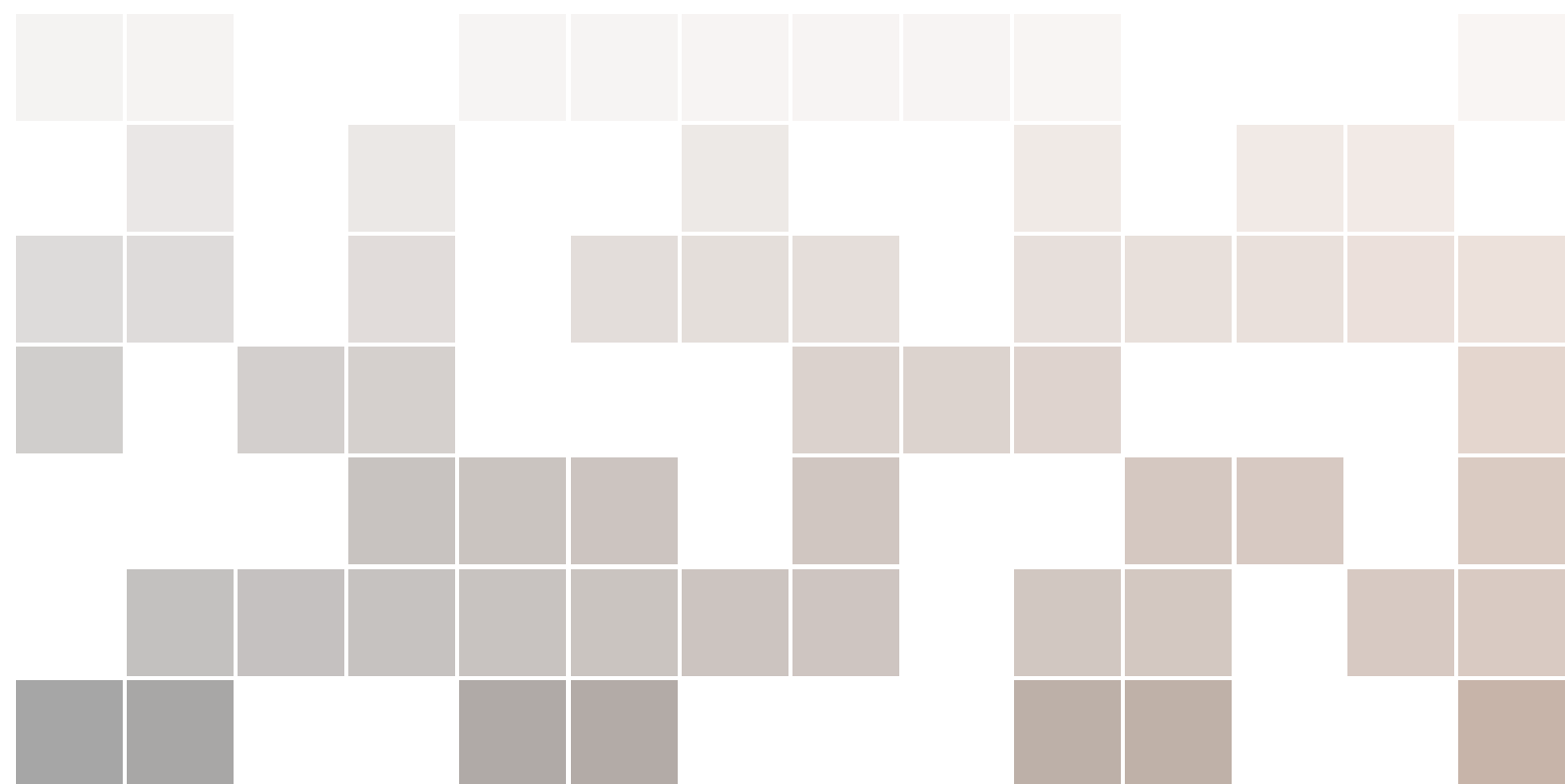


Representation and Reasoning for Intelligent Systems

Stefan Klaus



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First printing, March 2013

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Constraint Satisfaction Problem

Basic terms about CSP:

Arc Consistency Algorithm(AC-3)

Path consistency

Backtracking

Improving backtracking efficiency

Citation

Lists

Numbered List

Bullet Points

Descriptions and Definitions

1. Text Chapter

1.1 Constraint Satisfaction Problem

Constraint satisfaction problems(CSP) are mathematical problems defined as a set of objects whose state must satisfy a number of constraints or limitations. One example of a CSP is the game “Sudoku”

1.2 Basic terms about CSP:

- an assignment is to assign values to some or all variables
- an assignment that does not violate any constraints is called a consistent assignment. With values from domain D_i
- A complete assignment is one in which each variable is assigned
- a partial assignment is one that assigns value to some of the variables

Constraint graph

Binary CSP: each constraint relates at most two variables, e.g: WASA constraint graph: nodes are variables(e.g. region WA), arcs show constraints(e.g. WASA)

Varieties of Variables

Discrete variables

- Finite Domains
- boolean CSPs, include: boolean satisfiability(NP - complete)
- Sudoku
- Infinite Domains(Integers, Strings, etc.)
- job scheduling, variables are start/end days for each job, need a constraint language, e.g.
 $StartJob_1 + 5 \leq StarJob_3$

Continuous variables

- start /end times for Hubble Telescope observations
- Unary constraints involve a single variable
 - $SA \neq Green$

- binary constraints involve pairs of variables
 - $SA \neq WA$
- Higher-order constraints involve 3 or more variables
 - cryptarithmic column constraints
- Preferences(soft) constraints
 - red is better than green
 - CSPs with preference are often with optimization search algorithms constraints optimization problems

Node Consistency

If a node is node-consistent if all the value's domain satisfy the variable's unary constraints.

Example:

$D = \text{Red, Green, Blue}$

Variable X

X dislikes green, then X starts with $D = \text{Red, Green, Blue}$, and becomes node consistent after eliminating Green. X is node consistent with the reduced domain, $D = \text{Red, Blue}$.

Arc Consistency

X_i is arc-consistent with respect to another variables X_j if for every value in X_i 's current domain D_i there is one value in X_j 's domain D_j that satisfies the binary constraint on the arc (X_i, X_j)

Example:

Given two variables X_i, X_j with values in $0, 1, 2, \dots, 9$ and constraints $(0,0), (1,1), (2,4), (3,9)$.

To make X_i arc-consistent with respect to X_j , we reduce X_i 's domain to $0, 1, 2, 3$.

To make X_j arc-consistent with respect to X_i , we reduce X_j 's domain to $0, 1, 4, 9$.

1.3 Arc Consistency Algorithm(AC-3)

function AC-3(*csp*) **returns** false if an inconsistency is found and true otherwise

inputs: *csp*, a binary CSP with components (X, D, C)

local variables: *queue*, a queue of arcs, initially all the arcs in *csp*

while *queue* is not empty **do**

$(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\text{queue})$

if REVISE(*csp*, X_i, X_j) **then**

if size of $D_i = 0$ **then return** false

for each X_k **in** $X_i.\text{NEIGHBORS} - \{X_j\}$ **do**

 add (X_k, X_i) to *queue*

return true

function REVISE(*csp*, X_i, X_j) **returns** true iff we revise the domain of X_i

revised \leftarrow false

for each x **in** D_i **do**

if no value y in D_j allows (x, y) to satisfy the constraint between X_i and X_j **then**

 delete x from D_i

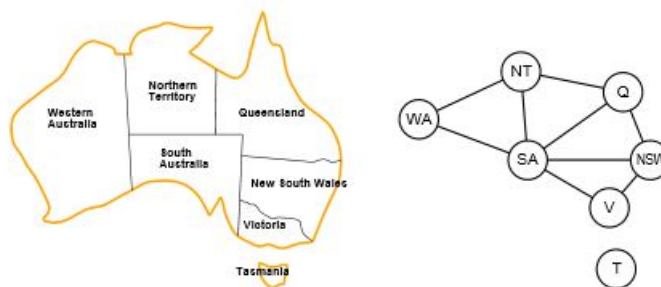
revised \leftarrow true

return *revised*

1. initially let a queue contain all arcs

2. remove an arc (X_i, X_j) from the queue and make the variable X_i arc-consistent to X_j
 - (a) IF X_i domain D_i is unchanged, then check the next arc in the queue
 - (b) IF X_i 's domain D_i is revised(smaller), then add all arcs (X_k, X_i) in the queue
 - (c) If X_i 's domain D_i is empty, then CSP no solution
3. Keep checking all arcs in the queue until the queue is empty

1.4 Path consistency



A two variable set X_i, X_j is path consistency with respect to a third variable X_m if for every assignment $X_i = a, X_j = b$ consistent with the constraints on X_i, X_j , where is an assignment to X_m , that satisfies the constraints on X_i, X_m and X_m, X_j

Example:

Can we color the Australia map with two colors?

Make the set WA, SA path consistent with respect to NT?

Assignments: WA = blue, SA = red or WA =blue,SA=red

But no assignment exists for NT

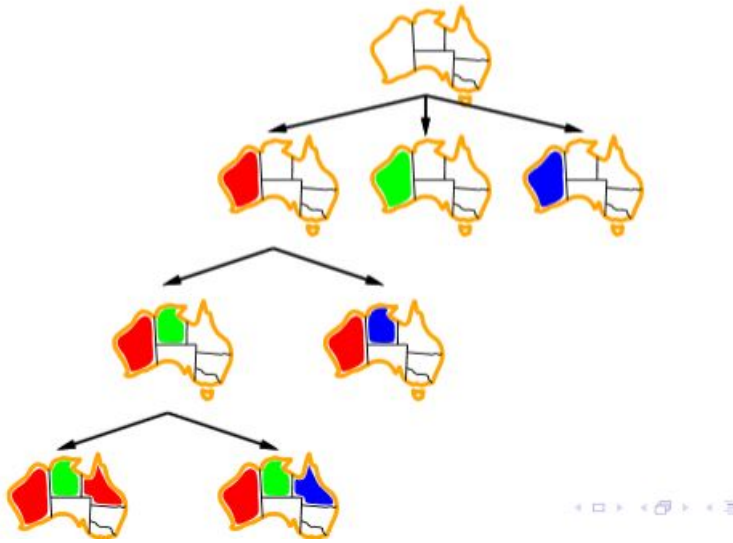
Constraint propagation

- constraint propagation is a specific type of inference
 - use the constraints to reduce the number of legal values for a variable, which in turn can reduce the legal values for another variable, and so on
- node consistency
- arc consistency
- path consistency
- k-consistency: k variables involved
- Global consistency

Limits

- Indeed AC-3 works for the easiest Sudoku puzzles
- slightly harder ones can be solved by PC-2, but at a greater computational cost: there are 255,960 different path constraints to consider in a Sudoku puzzle
- To solve the hardest puzzles and to make efficient progress, we will have to be more clever

1.5 Backtracking



1. Select an unassigned value
2. assign values
3. Depth-first search
4. If an inconsistency is detected, then BACKTRACK returns failure, causing the previous call to try another value

Backtracking search is a depth-first search for CSPs

- choose values for one variable at a time
- backtrack when a variables has no legal left to assign

Backtracking search is the basic uninformed algorithm for CSP. Can solve n-queens for $n \approx 25$.

```

function BACKTRACKING-SEARCH(csp) returns a solution, or failure
    return BACKTRACK({ }, csp)

function BACKTRACK(assignment, csp) returns a solution, or failure
    if assignment is complete then return assignment
    var  $\leftarrow$  SELECT-UNASSIGNED-VARIABLE(csp)
    for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
        if value is consistent with assignment then
            add { var = value } to assignment
            inferences  $\leftarrow$  INFERENCE(csp, var, value)
            if inferences  $\neq$  failure then
                add inferences to assignment
                result  $\leftarrow$  BACKTRACK(assignment, csp)
                if result  $\neq$  failure then
                    return result
            remove { var = value } and inferences from assignment
    return failure

```

1.5.1 Improving backtracking efficiency

General-purpose methods can give huge gains in speed:

1. which variable should be assigned next?
2. in what order should values be tried?
3. can we detect inevitable failure early?

4. Can we take advantage of problem structure?

Minimum remaining values

Minimum remaining values (MRV): choose the variable with the fewest legal values.

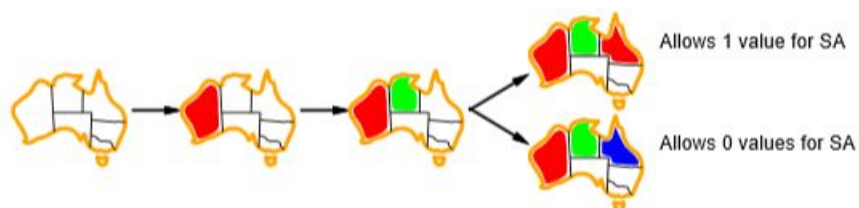
Degree heuristic

- Tie-breaker among MRV variables
- Degree heuristic: choose the variable with the most constraints on remaining variables



Least constraining value

Given a variable, choose the least constraining value: the one that rules out the fewest values in the remaining variables.



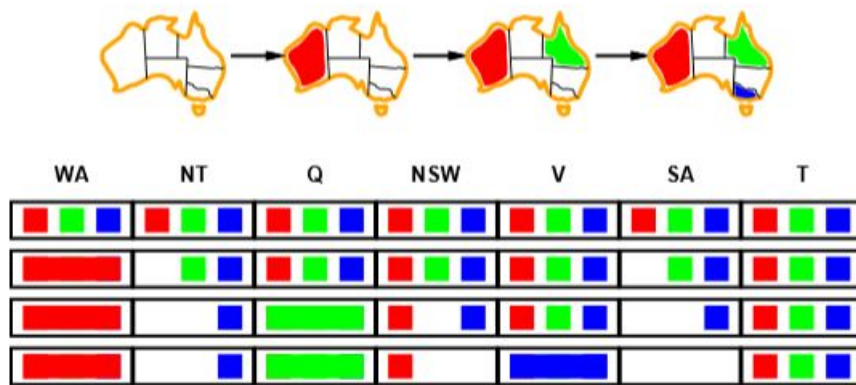
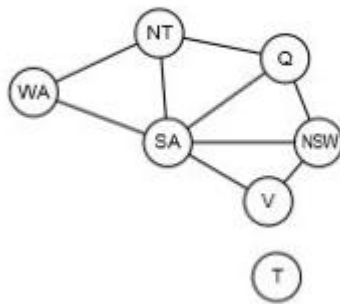
Combining

these heuristics makes 1000 queens feasible.

Forward checking

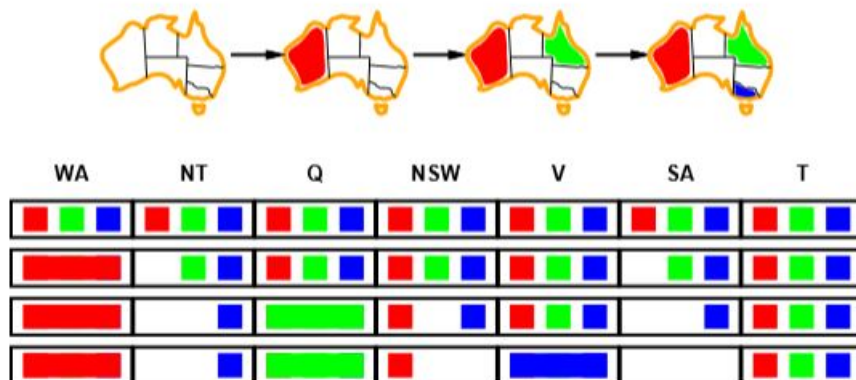
Idea: Keep track of remaining legal values for unassigned variables.

Terminates search when any variable has no legal values.



Constraint propagation

Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:



NT and SA cannot both be blue!

Constraint propagation repeatedly enforces constraints locally.

Arc consistency

Simplest form of propagation makes each arc consistent.

$X \rightarrow Y$ is consistent if for every value x of X there is some allowed y

1.6 Citation

This statement requires citation [Smi12]; this one is more specific [Smi13, page 122].

1.7 Lists

Lists are useful to present information in a concise and/or ordered way¹.

1.7.1 Numbered List

1. The first item
2. The second item
3. The third item

1.7.2 Bullet Points

- The first item
- The second item
- The third item

1.7.3 Descriptions and Definitions

Name Description

Word Definition

Comment Elaboration

¹Footnote example...

Theorems

Several equations
Single Line

Definitions

Notations

Remarks

Corollaries

Propositions

Several equations
Single Line

Examples

Equation and Text
Paragraph of Text

Exercises

Problems

Vocabulary

2. In-text Elements

2.1 Theorems

This is an example of theorems.

2.1.1 Several equations

This is a theorem consisting of several equations.

Theorem 2.1.1 — Name of the theorem. In $E = \mathbb{R}^n$ all norms are equivalent. It has the properties:

$$||\mathbf{x}|| - ||\mathbf{y}|| \leq ||\mathbf{x} - \mathbf{y}|| \quad (2.1)$$

$$||\sum_{i=1}^n \mathbf{x}_i|| \leq \sum_{i=1}^n ||\mathbf{x}_i|| \quad \text{where } n \text{ is a finite integer} \quad (2.2)$$

2.1.2 Single Line

This is a theorem consisting of just one line.

Theorem 2.1.2 A set $\mathcal{D}(G)$ is dense in $L^2(G)$, $|\cdot|_0$.

2.2 Definitions

This is an example of a definition. A definition could be mathematical or it could define a concept.

Definition 2.2.1 — Definition name. Given a vector space E , a norm on E is an application, denoted $||\cdot||$, E in $\mathbb{R}^+ = [0, +\infty[$ such that:

$$||\mathbf{x}|| = 0 \Rightarrow \mathbf{x} = \mathbf{0} \quad (2.3)$$

$$||\lambda \mathbf{x}|| = |\lambda| \cdot ||\mathbf{x}|| \quad (2.4)$$

$$||\mathbf{x} + \mathbf{y}|| \leq ||\mathbf{x}|| + ||\mathbf{y}|| \quad (2.5)$$

2.3 Notations

Notation 2.1. Given an open subset G of \mathbb{R}^n , the set of functions φ are:

1. Bounded support G ;
2. Infinitely differentiable;

a vector space is denoted by $\mathcal{D}(G)$.

2.4 Remarks

This is an example of a remark.

R The concepts presented here are now in conventional employment in mathematics. Vector spaces are taken over the field $\mathbb{K} = \mathbb{R}$, however, established properties are easily extended to $\mathbb{K} = \mathbb{C}$.

2.5 Corollaries

This is an example of a corollary.

Corollary 2.5.1 — Corollary name. The concepts presented here are now in conventional employment in mathematics. Vector spaces are taken over the field $\mathbb{K} = \mathbb{R}$, however, established properties are easily extended to $\mathbb{K} = \mathbb{C}$.

2.6 Propositions

This is an example of propositions.

2.6.1 Several equations

Proposition 2.6.1 — Proposition name. It has the properties:

$$||\mathbf{x}| - |\mathbf{y}|| \leq |\mathbf{x} - \mathbf{y}| \quad (2.6)$$

$$||\sum_{i=1}^n \mathbf{x}_i|| \leq \sum_{i=1}^n ||\mathbf{x}_i|| \quad \text{where } n \text{ is a finite integer} \quad (2.7)$$

2.6.2 Single Line

Proposition 2.6.2 Let $f, g \in L^2(G)$; if $\forall \varphi \in \mathcal{D}(G)$, $(f, \varphi)_0 = (g, \varphi)_0$ then $f = g$.

2.7 Examples

This is an example of examples.

2.7.1 Equation and Text

■ **Example 2.1** Let $G = \{x \in \mathbb{R}^2 : |x| < 3\}$ and denoted by: $x^0 = (1, 1)$; consider the function:

$$f(x) = \begin{cases} e^{|x|} & \text{si } |x - x^0| \leq 1/2 \\ 0 & \text{si } |x - x^0| > 1/2 \end{cases} \quad (2.8)$$

The function f has bounded support, we can take $A = \{x \in \mathbb{R}^2 : |x - x^0| \leq 1/2 + \varepsilon\}$ for all $\varepsilon \in]0; 5/2 - \sqrt{2}[$. ■

2.7.2 Paragraph of Text

■ **Example 2.2 — Example name.** Nam dui ligula, fringilla a, euismod sodales, sollicitudin vel, wisi. Morbi auctor lorem non justo. Nam lacus libero, pretium at, lobortis vitae, ultricies et, tellus. Donec aliquet, tortor sed accumsan bibendum, erat ligula aliquet magna, vitae ornare odio metus a mi. Morbi ac orci et nisl hendrerit mollis. Suspendisse ut massa. Cras nec ante. Pellentesque a nulla. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Aliquam tincidunt urna. Nulla ullamcorper vestibulum turpis. Pellentesque cursus luctus mauris.

■

2.8 Exercises

This is an example of an exercise.

Exercise 2.1 This is a good place to ask a question to test learning progress or further cement ideas into students' minds.

■

2.9 Problems

Problem 2.1 What is the average airspeed velocity of an unladen swallow?

2.10 Vocabulary

Define a word to improve a students' vocabulary.

Vocabulary 2.1 — Word. Definition of word.

3. Presenting Information

3.1 Table

Treatments	Response 1	Response 2
Treatment 1	0.0003262	0.562
Treatment 2	0.0015681	0.910
Treatment 3	0.0009271	0.296

Table 3.1: Table caption

3.2 Figure

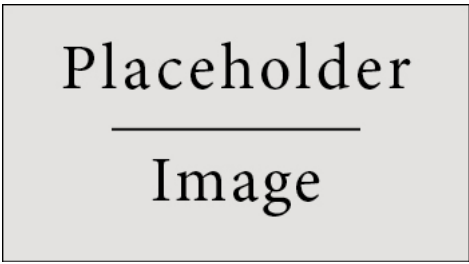


Figure 3.1: Figure caption



Bibliography

Books

[Smi12] John Smith. *Book title*. 1st edition. Volume 3. 2. City: Publisher, Jan. 2012, pages 123–200 (cited on page 10).

Articles

[Smi13] James Smith. “Article title”. In: 14.6 (Mar. 2013), pages 1–8 (cited on page 10).

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