

# Context-Free Grammar

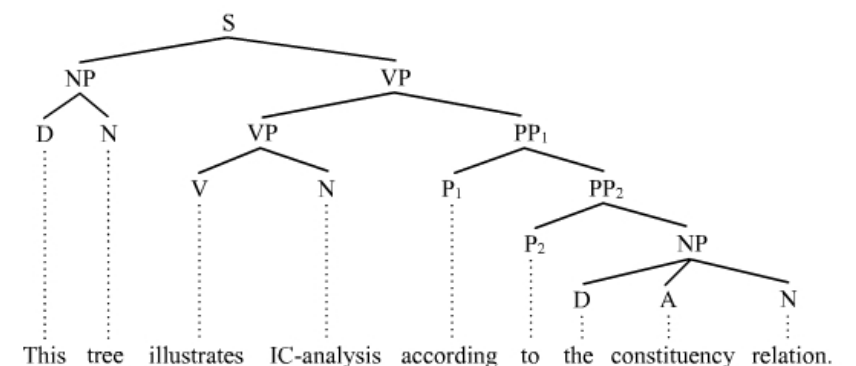
COMP90042

Natural Language Processing

Lecture 14



THE UNIVERSITY OF  
MELBOURNE



# Recap

- Center embedding
  - ▶ The cat loves Mozart
  - ▶ The cat **the dog chased** loves Mozart
  - ▶ The cat **the dog the rat bit chased** loves Mozart
  - ▶ The cat **the dog the rat the elephant admired bit chased** loves Mozart
- Cannot be captured by regular expressions ( $S^nV^n$ )
- **Context-free grammar!**

# Basics of Context-Free Grammars

- **Symbols**

- ▶ **Terminal**: word such as *book*
- ▶ **Non-terminal**: syntactic label such as NP or VP
- ▶ Convention to use upper and lower-case to distinguish, or else “quotes” for terminals

- **Productions (rules)**

$$W \rightarrow X Y Z$$

- ▶ Exactly one non-terminal on left-hand side (LHS)
- ▶ An ordered list of symbols on right-hand side (RHS)
  - can be **Terminals** or **Non-terminals**

- **Start symbol: S**

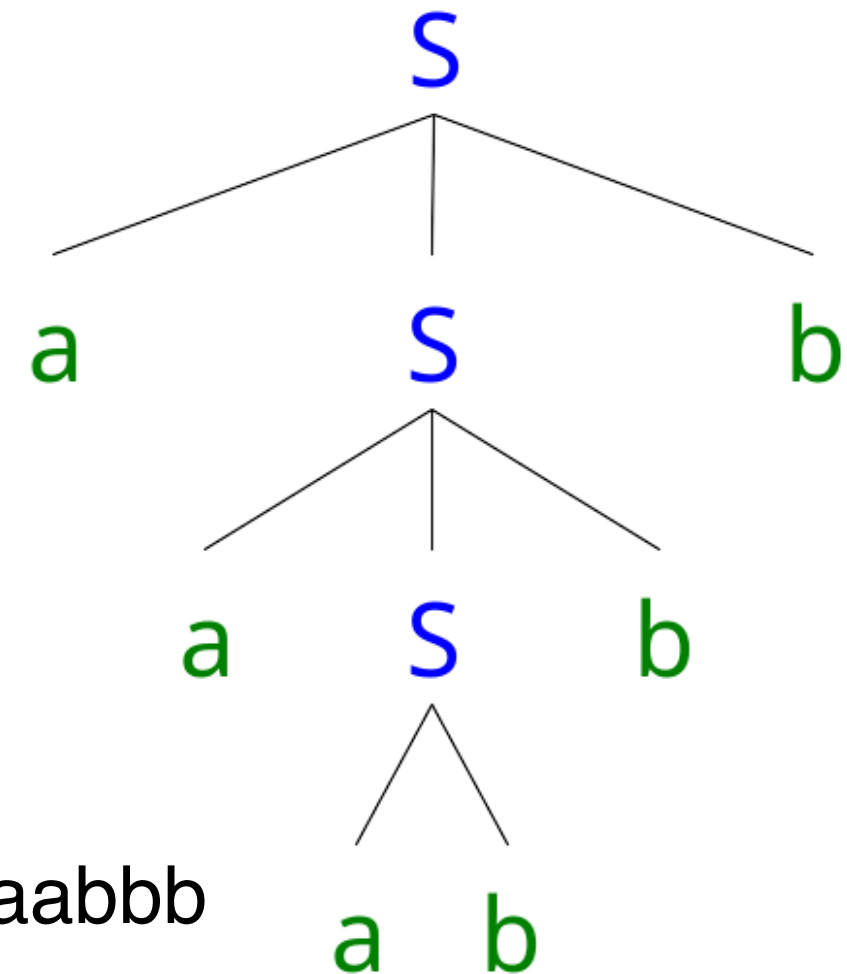
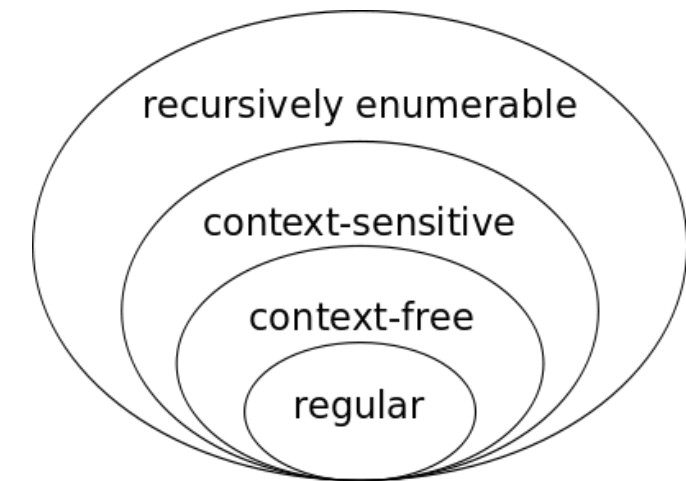
# Why “Context Free”

$$W \rightarrow X Y Z$$

- Production rule depends only on the LHS (and not on ancestors, neighbours)
  - ▶ Analogous to Markov chain
  - ▶ Behaviour at each step depends only on current state

# Context-Free vs. Regular

- Context-free languages more general than Regular languages
  - ▶ Allows recursive nesting
- CFG for  $a^n b^n$ :
  - ▶  $S \rightarrow a S b$
  - ▶  $S \rightarrow a b$



Parse tree for  $aaabbbb$

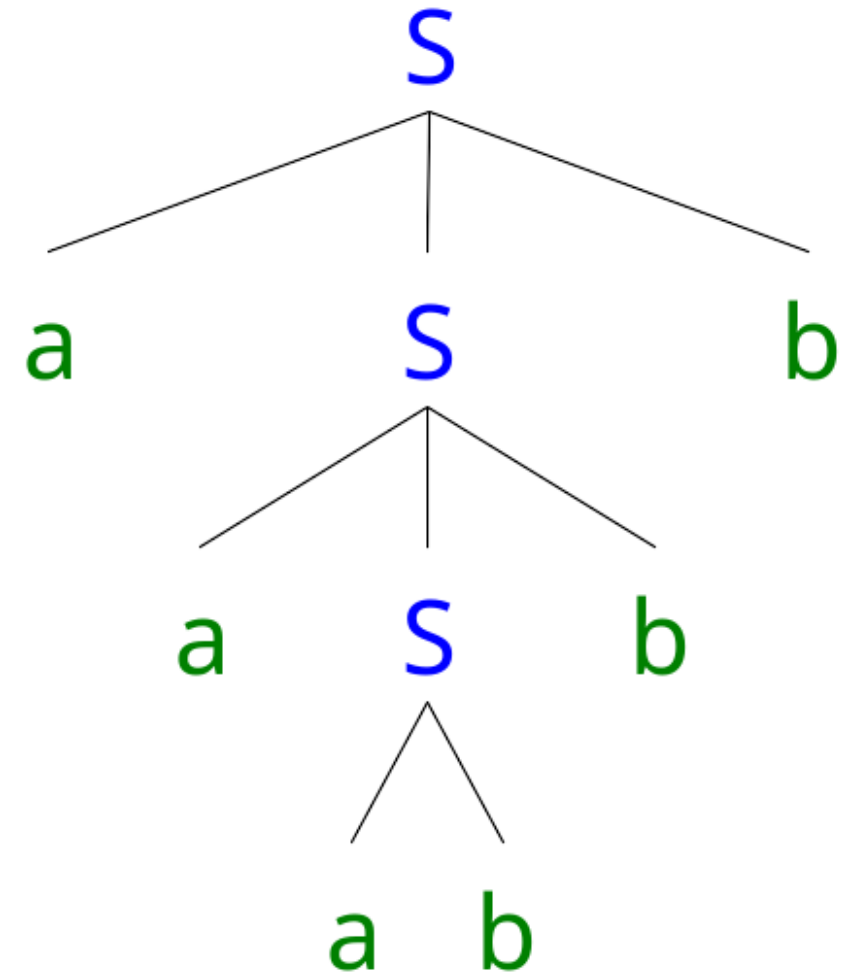
# Parsing

- Given a string and production rules
- Produce a valid parse tree

►  $S \rightarrow a S b$

►  $S \rightarrow a b$

aaabbbb



# What This Means?

- If English can be represented with CFG:
  - ▶ we can build a “parser” to automatically judge whether a sentence is grammatical!
- But is natural language context-free?
- Not quite: cross-serial dependencies ( $a^m b^n c^m d^n$ )

Swiss-German:

...de Karl d'Maria em Peter de Hans laat hilfe larne schwume

English:

...Charles lets Mary help Peter to teach John to Swim

# But...

- Context-free representations strike a good balance:
  - ▶ CFG cover most syntactic patterns
  - ▶ CFG parsing is computationally efficient
- We use CFG to describe a core fragment of English syntax



# Syntactic Constituents

- Sentences are broken into **constituents**
  - ▶ word sequence that function as a **coherent unit** for linguistic analysis
- Constituents have certain key properties:
  - ▶ movement
  - ▶ substitution
  - ▶ coordination

# Movement

- Constituents can be moved around sentences
  - ▶ Abigail gave [her brother] [a fish]
  - ▶ Abigail gave [a fish] to [her brother]
- Contrast: [gave her], [brother a]

# Substitution

- Constituents can be substituted by other phrases of the same type
  - ▶ Max thanked [his older sister]
  - ▶ Max thanked [her]
- Contrast: [Max thanked], [thanked his]

# Coordination

- Constituents can be conjoined with coordinators like *and* and *or*
  - ▶ [Abigail] and [her young brother] brought a fish
  - ▶ Abigail [bought a fish] and [gave it to Max]
  - ▶ Abigail [bought] and [greedily ate] a fish
- Contrast: [brother brought], [bought a]

# Constituents and Phrases

- Once we identify constituents, we use **phrases** to describe them
- Phrases are determined by their **head word**:
  - ▶ noun phrase: her younger **brother**
  - ▶ verb phrase: greedily **ate** it
- We can use CFG to formalise these intuitions

# A Simple CFG for English

Terminal symbols: *rat, the, ate, cheese*

Non-terminal symbols: S, NP, VP, DT, VBD, NN

Productions:

$S \rightarrow NP VP$

$NP \rightarrow DT NN$

$VP \rightarrow VBD NP$

$DT \rightarrow the$

$NN \rightarrow rat$

$NN \rightarrow cheese$

$VBD \rightarrow ate$

# Generating Sentences with CFGs

Always start with S (the sentence/start symbol)

**S**

Apply a rule with S on LHS ( $S \rightarrow NP VP$ ), i.e substitute RHS

**NP VP**

Apply a rule with NP on LHS ( $NP \rightarrow DT NN$ )

**DT NN VP**

Apply rule with DT on LHS ( $DT \rightarrow the$ )

***the* NN VP**

Apply rule with NN on LHS ( $NN \rightarrow rat$ )

***the rat* VP**

In each step we  
rewrite the left-most  
non-terminal

# Generating Sentences with CFGs

Apply rule with VP on LHS ( $VP \rightarrow VBD\ NP$ )

***the rat*** VBD NP

Apply rule with VBD on LHS ( $VBD \rightarrow ate$ )

***the rat ate*** NP

Apply rule with NP on LHS ( $NP \rightarrow DT\ NN$ )

***the rat ate*** DT NN

Apply rule with DT on LHS ( $DT \rightarrow the$ )

***the rat ate the*** NN

Apply rule with NN on LHS ( $NN \rightarrow cheese$ )

***the rat ate the cheese***

No non-terminals  
left, we're done!

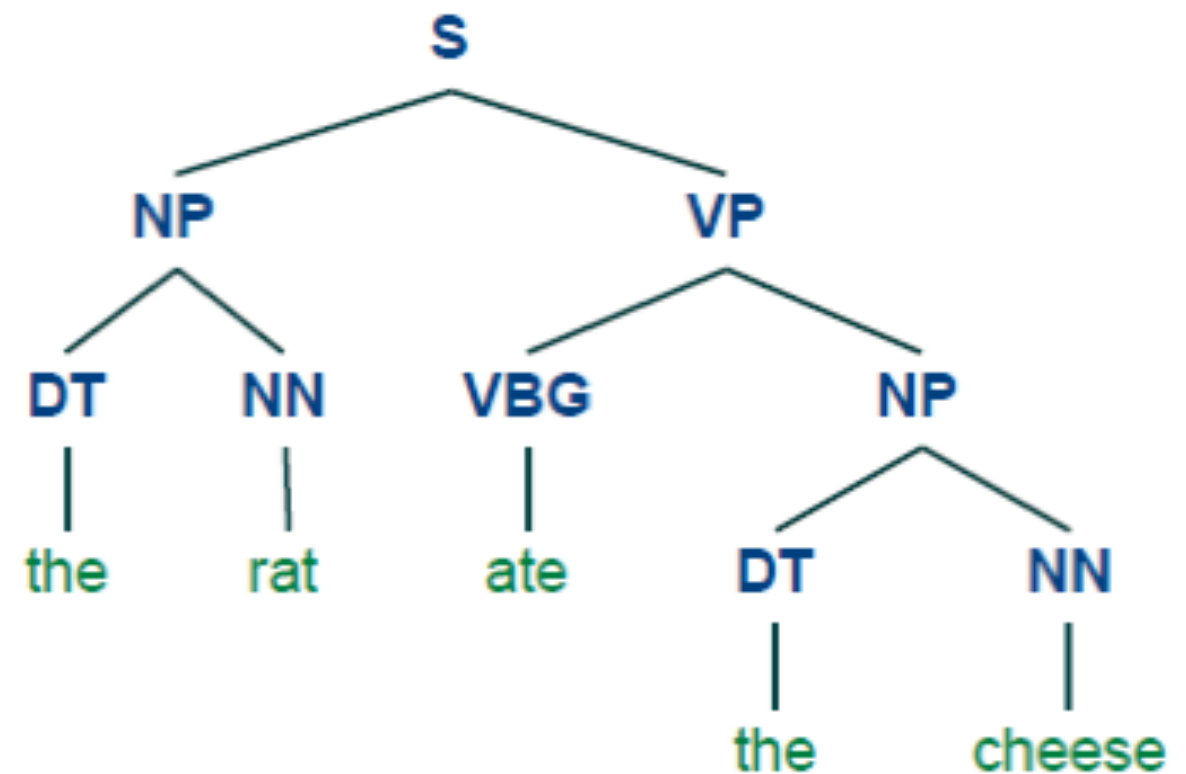




# CFG Trees

- Generation corresponds to a syntactic tree
- Non-terminals are internal nodes
- Terminals are leaves

(S (NP (DT the)  
         (NN rat))  
  (VP (VBG ate)  
      (NP (DT the)  
          (NN cheese)))))



- Parsing is the **reverse** process

# A CFG for Arithmetic Expressions

$$S \rightarrow S \text{ OP } S \mid \text{NUM}$$

$$\text{OP} \rightarrow + \mid - \mid \times \mid \div$$

$$\text{NUM} \rightarrow \text{NUM DIGIT} \mid \text{DIGIT}$$

$$\text{DIGIT} \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$$

- $S$  = starting symbol
- ' $\text{OP}$ ' = operator OR
- Recursive, NUM and  $S$  can produce themselves

# Parsing

- Is '4' a valid string?

$$S \rightarrow S \text{ OP } S \mid \text{NUM}$$

$$\text{OP} \rightarrow + \mid - \mid \times \mid \div$$

$$\text{NUM} \rightarrow \text{NUM DIGIT} \mid \text{DIGIT}$$

$$\text{DIGIT} \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$$

S  
|  
Num  
|  
Digit  
|  
4

# Parsing

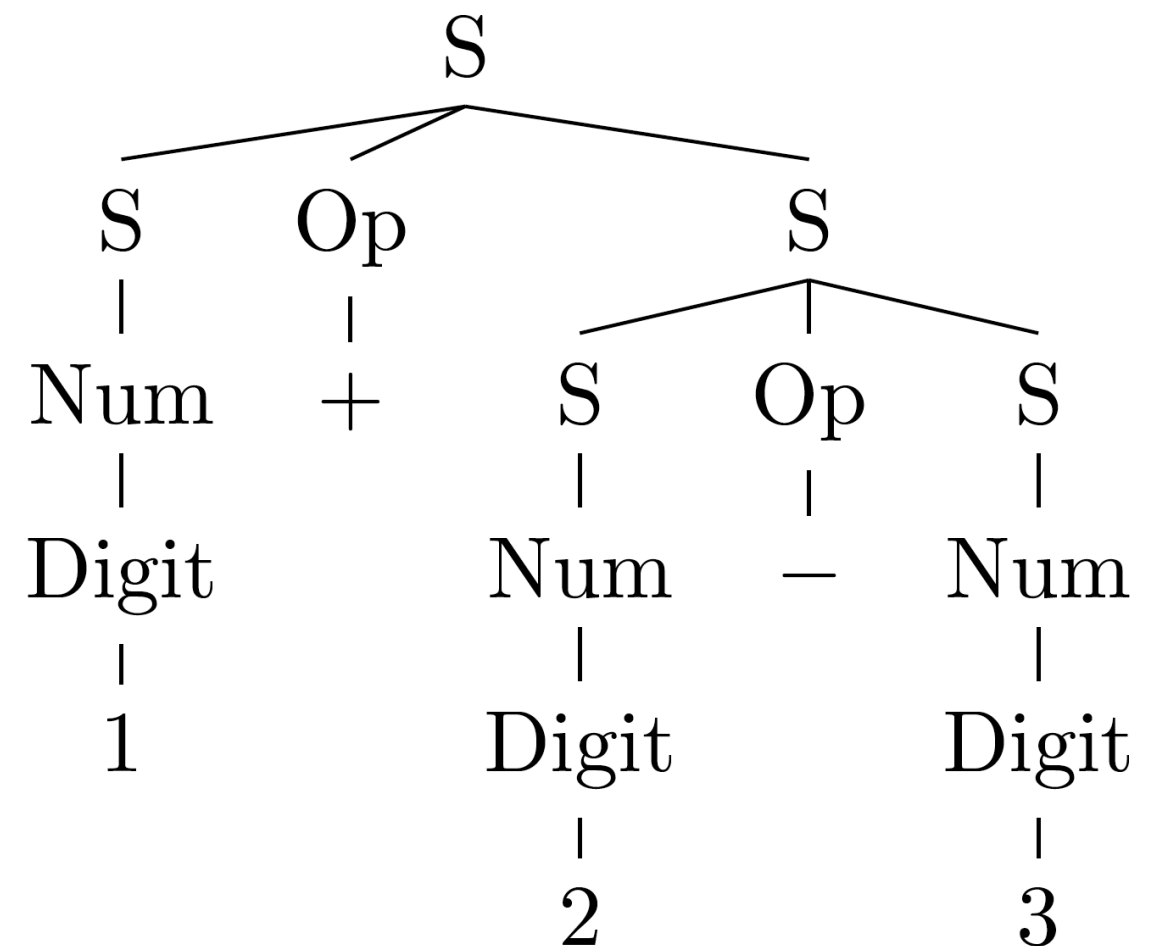
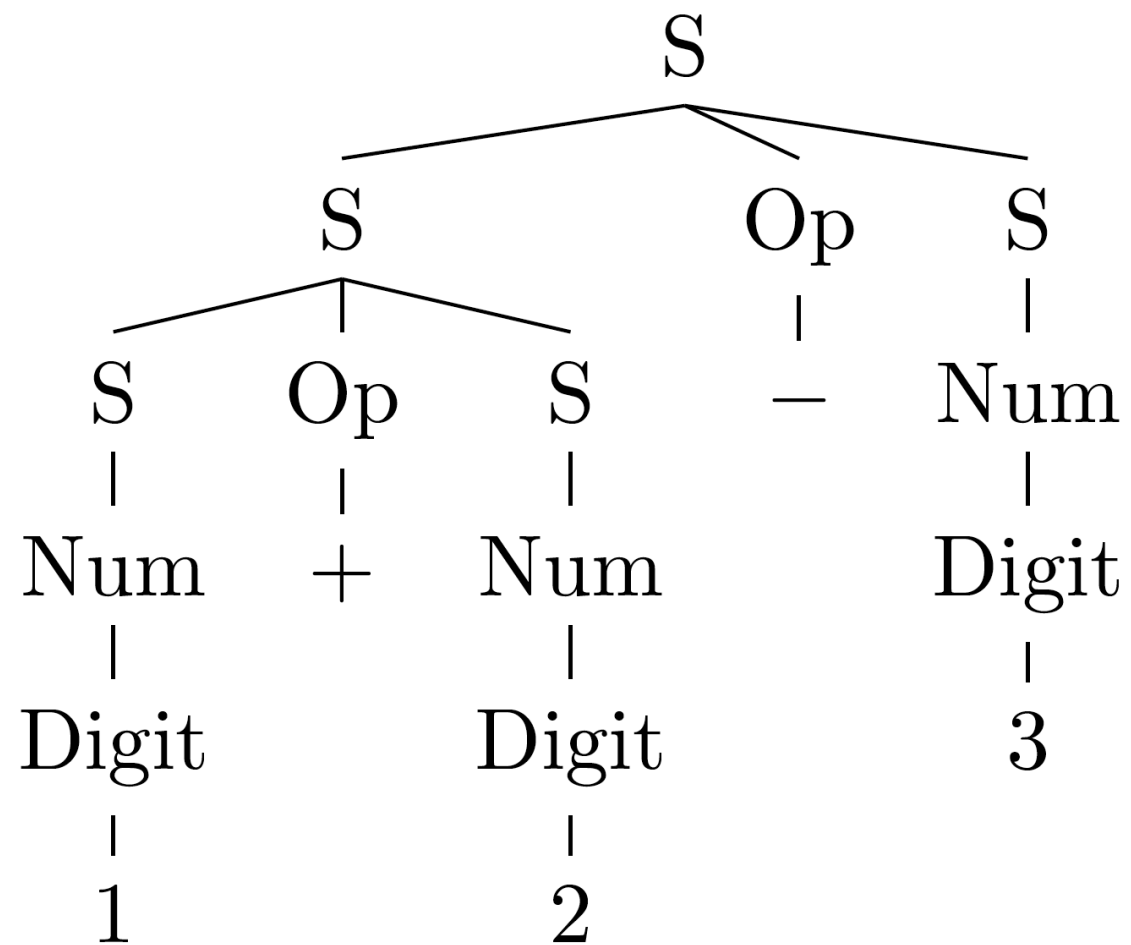
$$S \rightarrow S \text{ OP } S \mid \text{NUM}$$

$$\text{OP} \rightarrow + \mid - \mid \times \mid \div$$

$$\text{NUM} \rightarrow \text{NUM DIGIT} \mid \text{DIGIT}$$

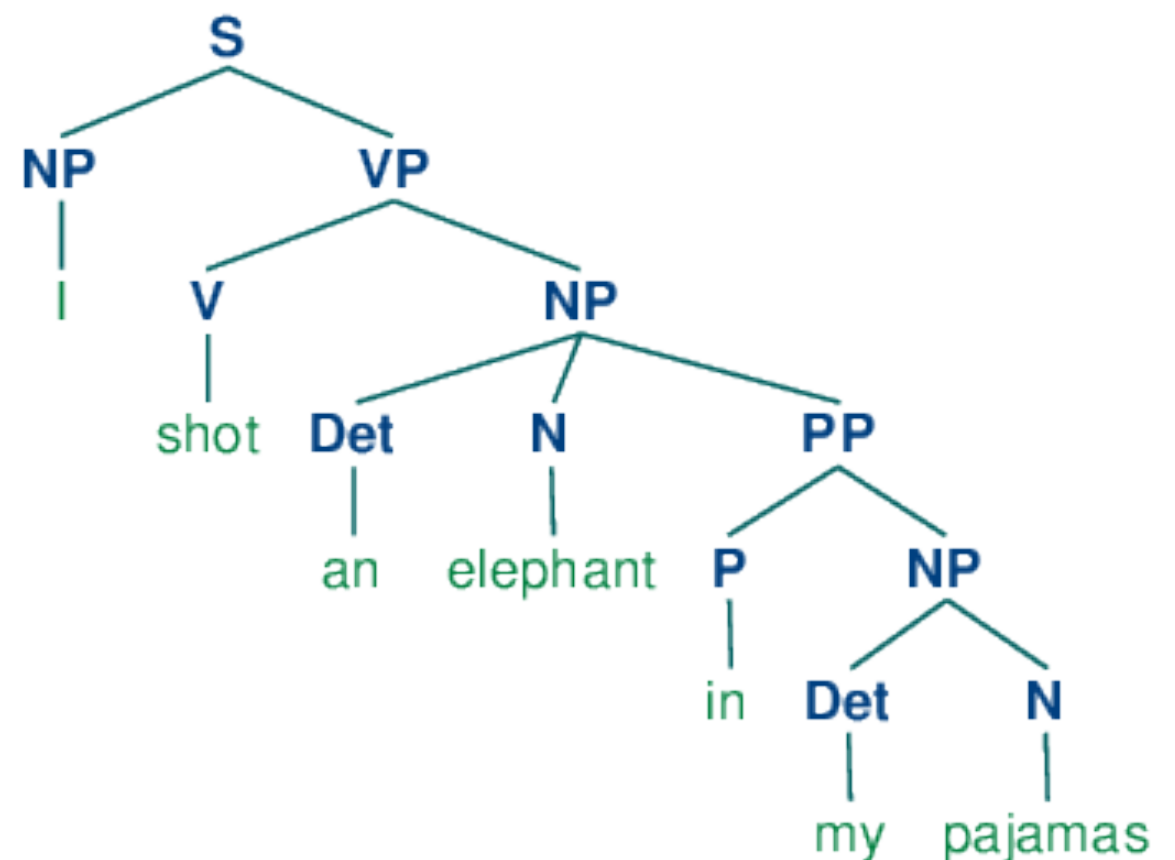
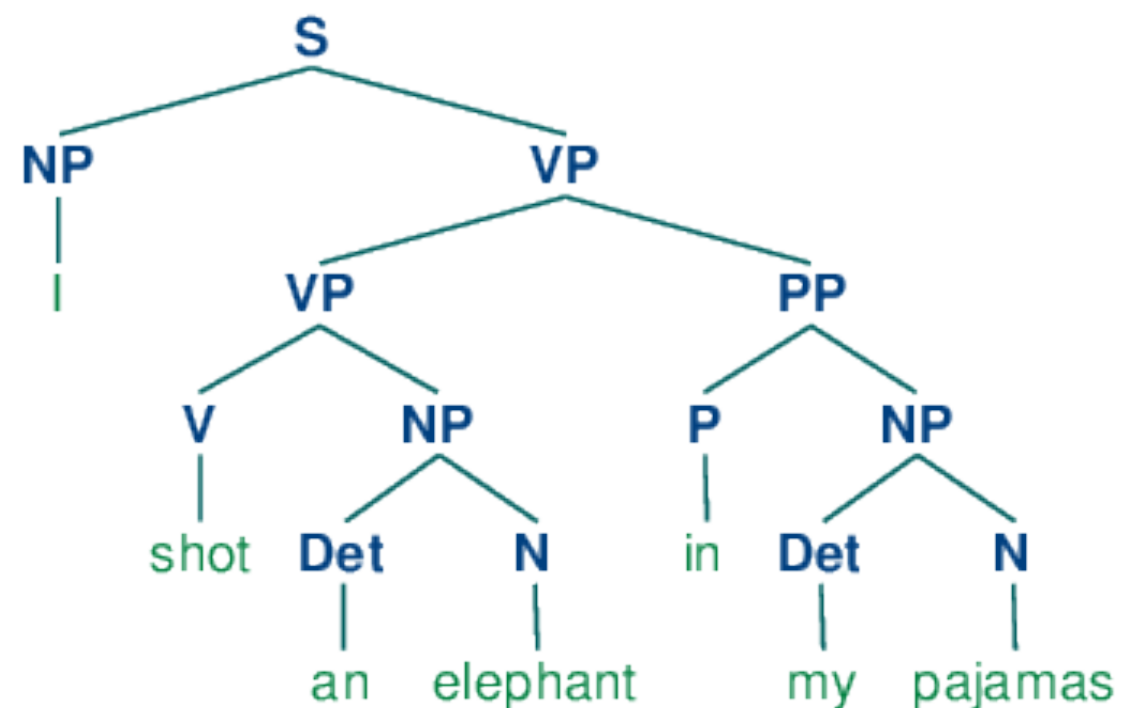
$$\text{DIGIT} \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$$

- Is '1+2-3' a valid string?



# Parse Ambiguity

- Often more than one tree can describe a string
- “*While hunting in Africa, I shot an elephant in my pajamas. How he got into my pajamas, I don't know.*”  
*Animal Crackers* (1930)



# Parsing CFG

# CYK Algorithm

- Bottom-up approach to parsing in CFG
- Tests whether a string is valid given a CFG, without enumerating all possible parses
- Core idea: form small constituents first, and merge them into larger constituents
- Requirement: CFGs must be in **Chomsky Normal Forms**

# Convert to Chomsky Normal Form

- Change grammar so all rules of form:
  - ▶  $A \rightarrow B C$
  - ▶  $A \rightarrow a$
- Convert rules of form  $A \rightarrow B c$  into:
  - ▶  $A \rightarrow B X$
  - ▶  $X \rightarrow c$



# Convert to Chomsky Normal Form

- Convert rules  $A \rightarrow B C D$  into:
  - ▶  $A \rightarrow B Y$
  - ▶  $Y \rightarrow C D$
  - ▶ E.g.,  $VP \rightarrow VP NP NP$   
for ditransitive cases, “*sold [her] [the book]*”
- $X, Y$  are new symbols we have introduced

# CNF (cont)

- CNF disallows unary rules,  $A \rightarrow B$ .
- Imagine  $NP \rightarrow S$ ; and  $S \rightarrow NP \dots$  leads to infinitely many trees with same yield.
- Replace RHS non-terminal with its productions
  - E.g convert  $A \rightarrow B$ ,  $B \rightarrow 1$ ,  $B \rightarrow 2$  into:
  - $A \rightarrow 1$ ,  $A \rightarrow 2$

# The CYK Parsing Algorithm

- Convert grammar to Chomsky Normal Form (CNF)
- Fill in a parse table
- Use table to derive parse
- S in top right corner of table = success!
- Convert result back to original grammar

we	eat	sushi	with	chopsticks
[0,1]	[0,2]	[0,3]	[0,4]	[0,5]
	[1,2]	[1,3]	[1,4]	[1,5]
		[2,3]	[2,4]	[2,5]
			[3,4]	[3,5]
				[4,5]

$S \rightarrow NP VP$

$NP \rightarrow NP PP$

$PP \rightarrow IN NP$

$VP \rightarrow V NP$

$VP \rightarrow VP PP$

$NP \rightarrow we$

$NP \rightarrow sushi$

$NP \rightarrow chopsticks$

$IN \rightarrow with$

$V \rightarrow eat$

we	eat	sushi	with	chopsticks
NP [0,1]	[0,2]	[0,3]	[0,4]	[0,5]
	V [1,2]	[1,3]	[1,4]	[1,5]
		NP [2,3]	[2,4]	[2,5]
			IN [3,4]	[3,5]
				NP [4,5]

S → NP VP

NP → NP PP

PP → IN NP

VP → V NP

VP → VP PP

NP → we

NP → sushi

NP → chopsticks

IN → with

V → eat

we	eat	sushi	with	chopsticks
NP [0,1]	∅ [0,2]	[0,3]	[0,4]	[0,5]
	V [1,2]	[1,3]	[1,4]	[1,5]
		NP [2,3]	[2,4]	[2,5]
			IN [3,4]	[3,5]
				NP [4,5]

S → NP VP

NP → NP PP

PP → IN NP

VP → V NP

VP → VP PP

NP → we

NP → sushi

NP → chopsticks

IN → with

V → eat

we	eat	sushi	with	chopsticks
NP [0,1]	∅ [0,2]	[0,3]	[0,4]	[0,5]
	V [1,2]	VP [1,3]	[1,4]	[1,5]
		NP [2,3]	[2,4]	[2,5]
			IN [3,4]	[3,5]
				NP [4,5]

S → NP VP

NP → NP PP

PP → IN NP

VP → V NP

VP → VP PP

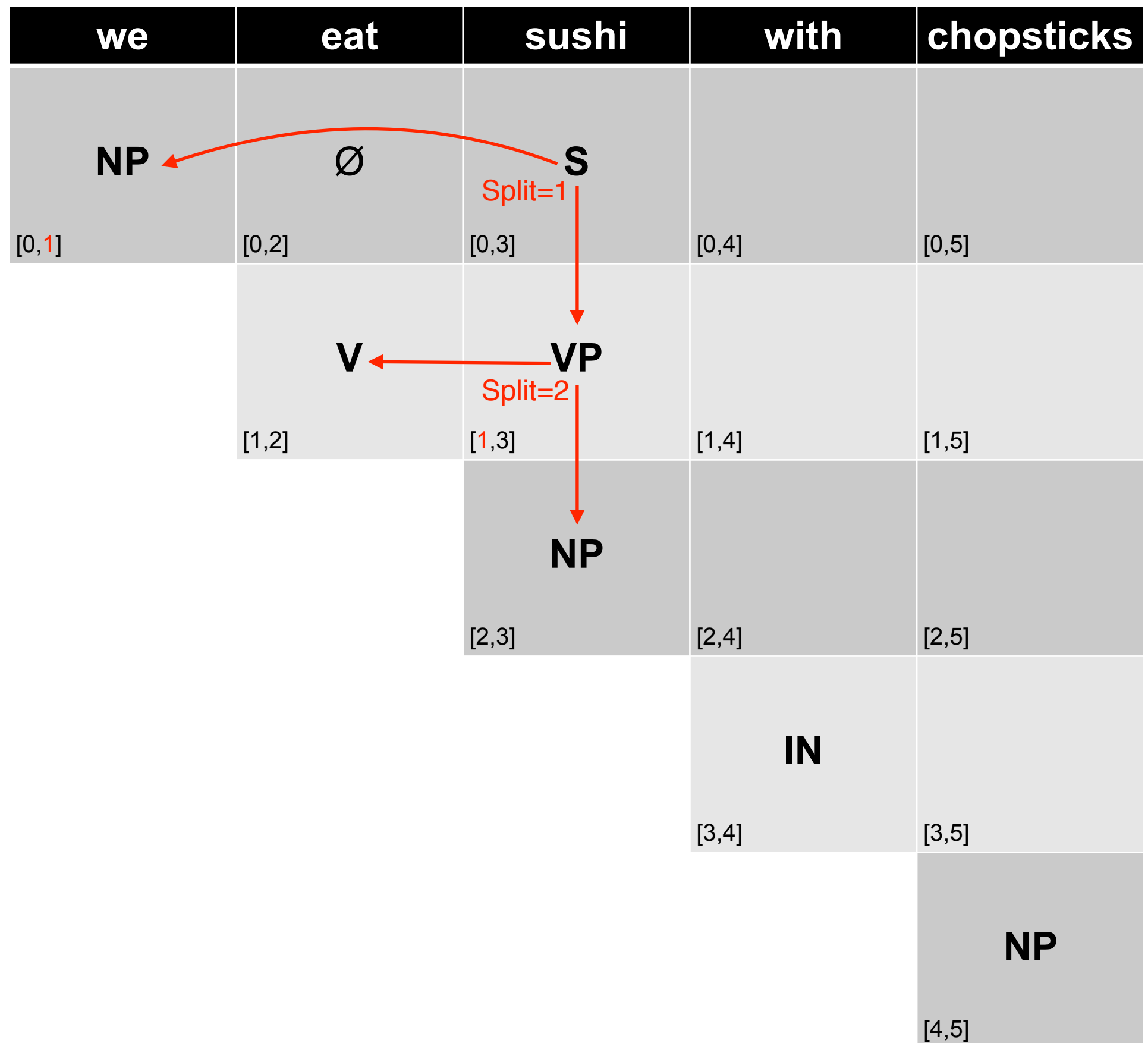
NP → we

NP → sushi

NP → chopsticks

IN → with

V → eat



**S → NP VP**

NP → NP PP

PP → IN NP

VP → V NP

VP → VP PP

NP → we

NP → sushi

NP → chopsticks

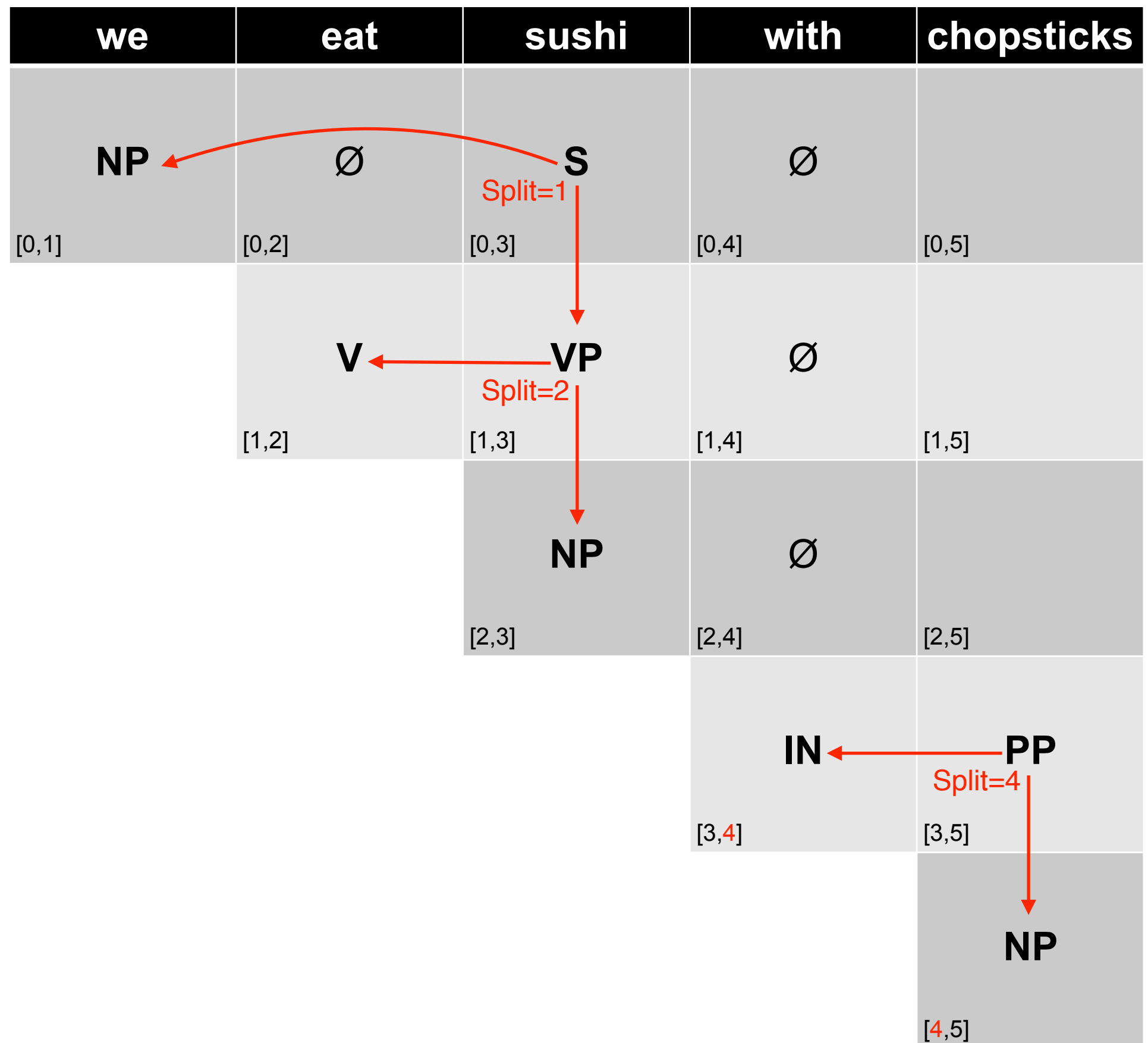
IN → with

V → eat

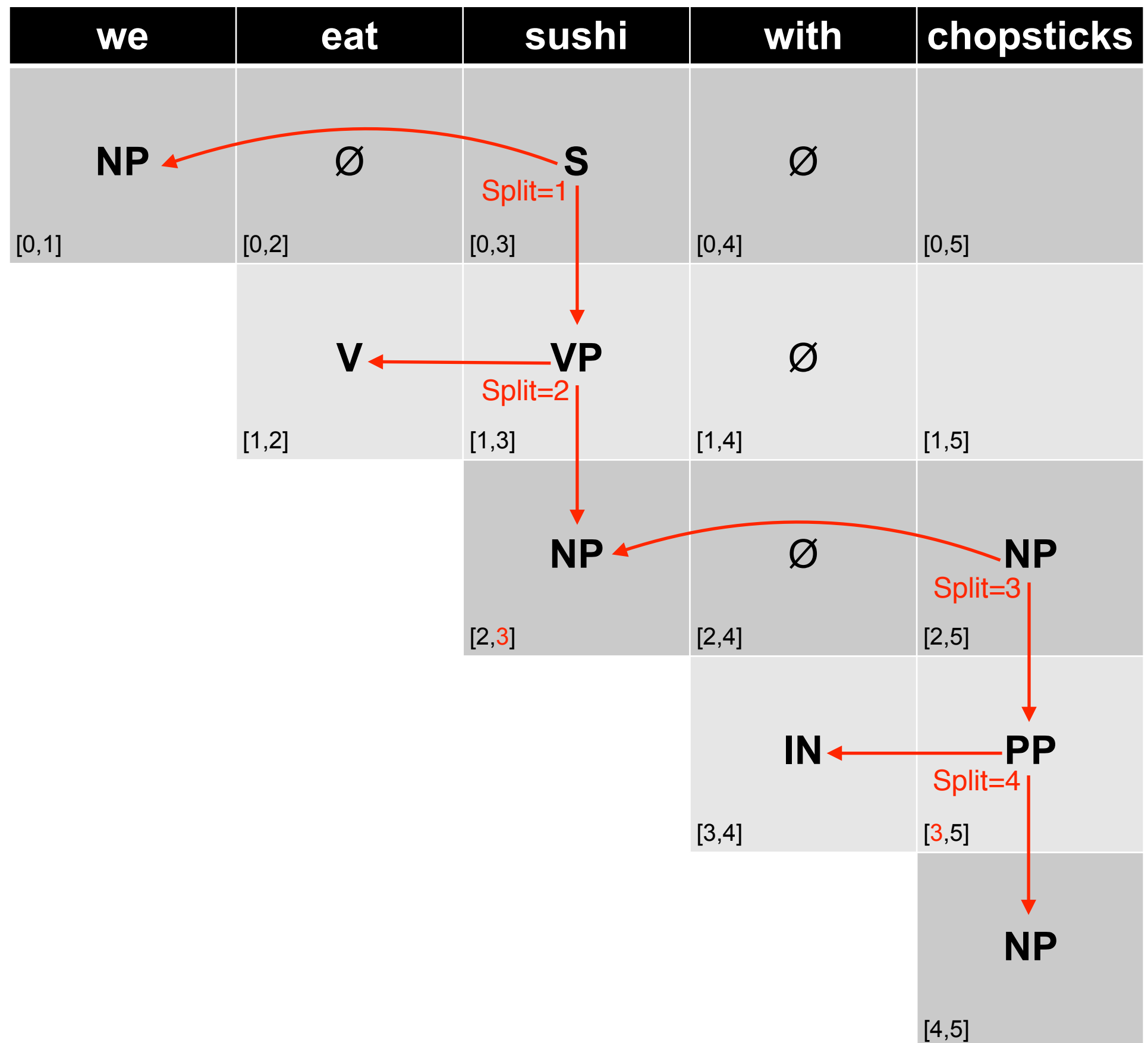


we	eat	sushi	with	chopsticks
NP [0,1]	∅ [0,2]	S [0,3]	∅ [0,4]	∅ [0,5]
	V [1,2]	VP [1,3]	∅ [1,4]	∅ [1,5]
		NP [2,3]	∅ [2,4]	∅ [2,5]
			IN [3,4]	∅ [3,5]
				NP [4,5]

$S \rightarrow NP VP$   
 $NP \rightarrow NP PP$   
 $PP \rightarrow IN NP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow VP PP$   
 $NP \rightarrow we$   
 $NP \rightarrow sushi$   
 $NP \rightarrow chopsticks$   
 $IN \rightarrow with$   
 $V \rightarrow eat$



$S \rightarrow NP VP$   
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$S \rightarrow NP VP$

$NP \rightarrow NP PP$

$PP \rightarrow IN NP$

$VP \rightarrow V NP$

$VP \rightarrow VP PP$

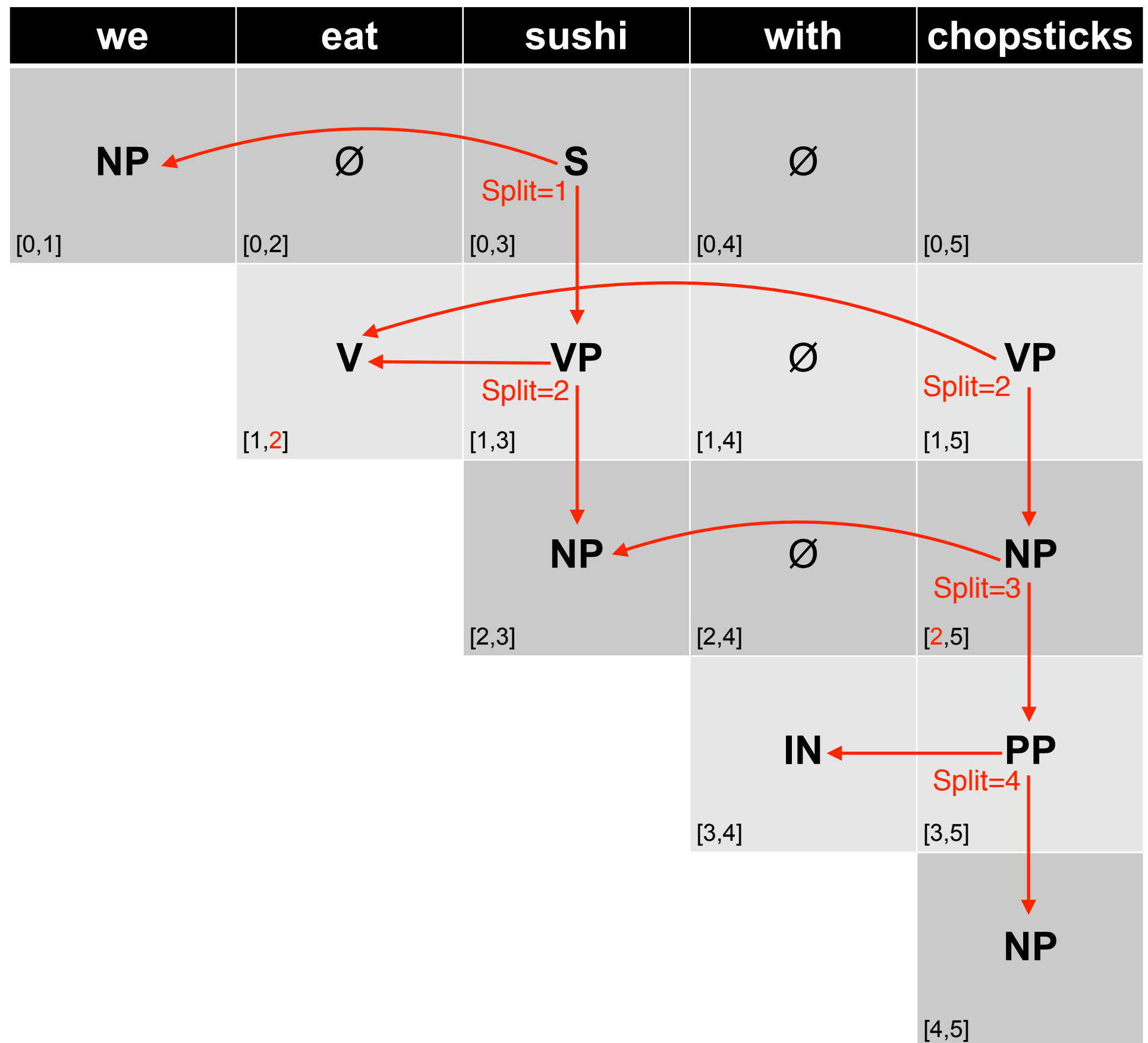
$NP \rightarrow we$

$NP \rightarrow sushi$

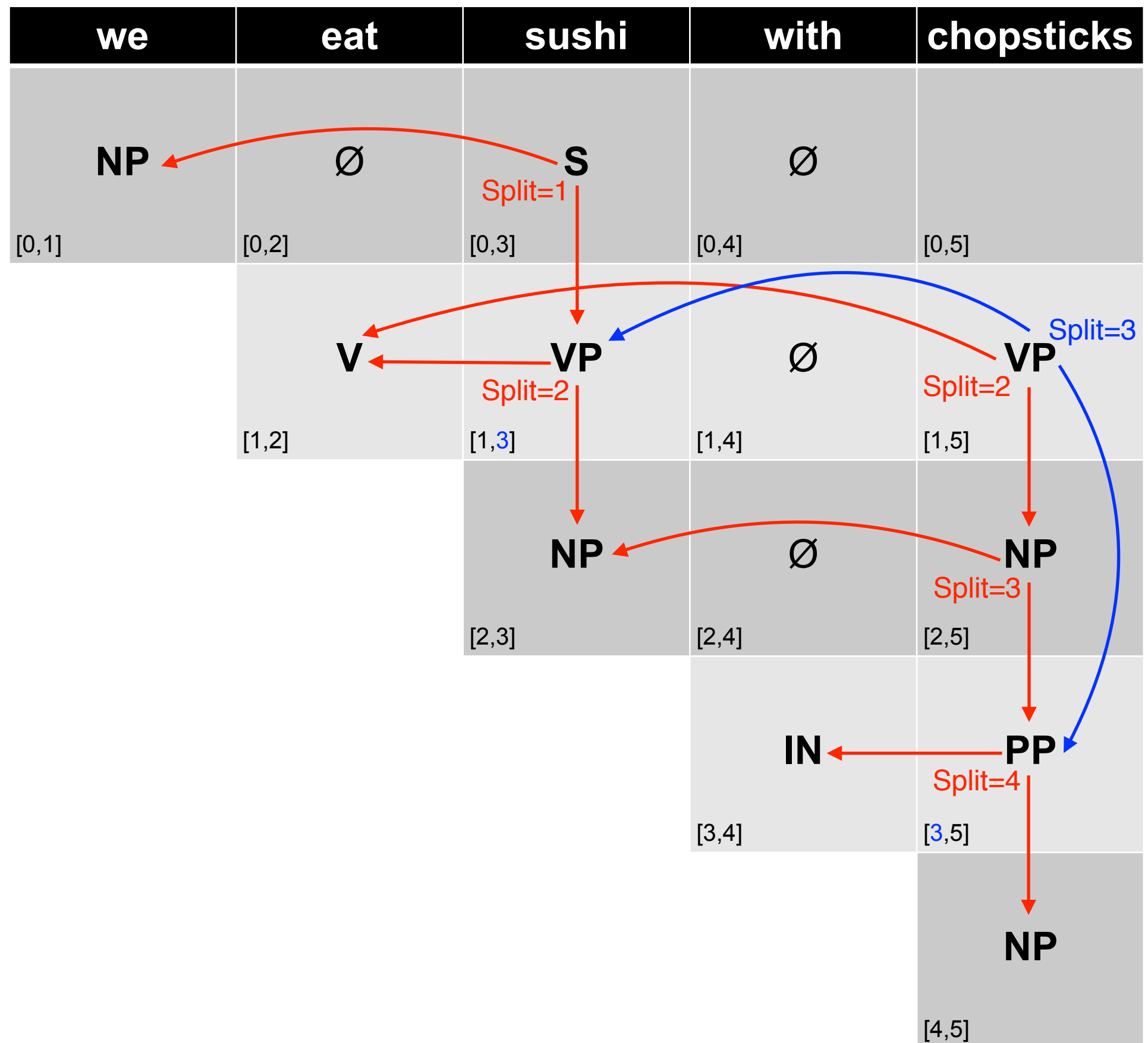
$NP \rightarrow chopsticks$

$IN \rightarrow with$

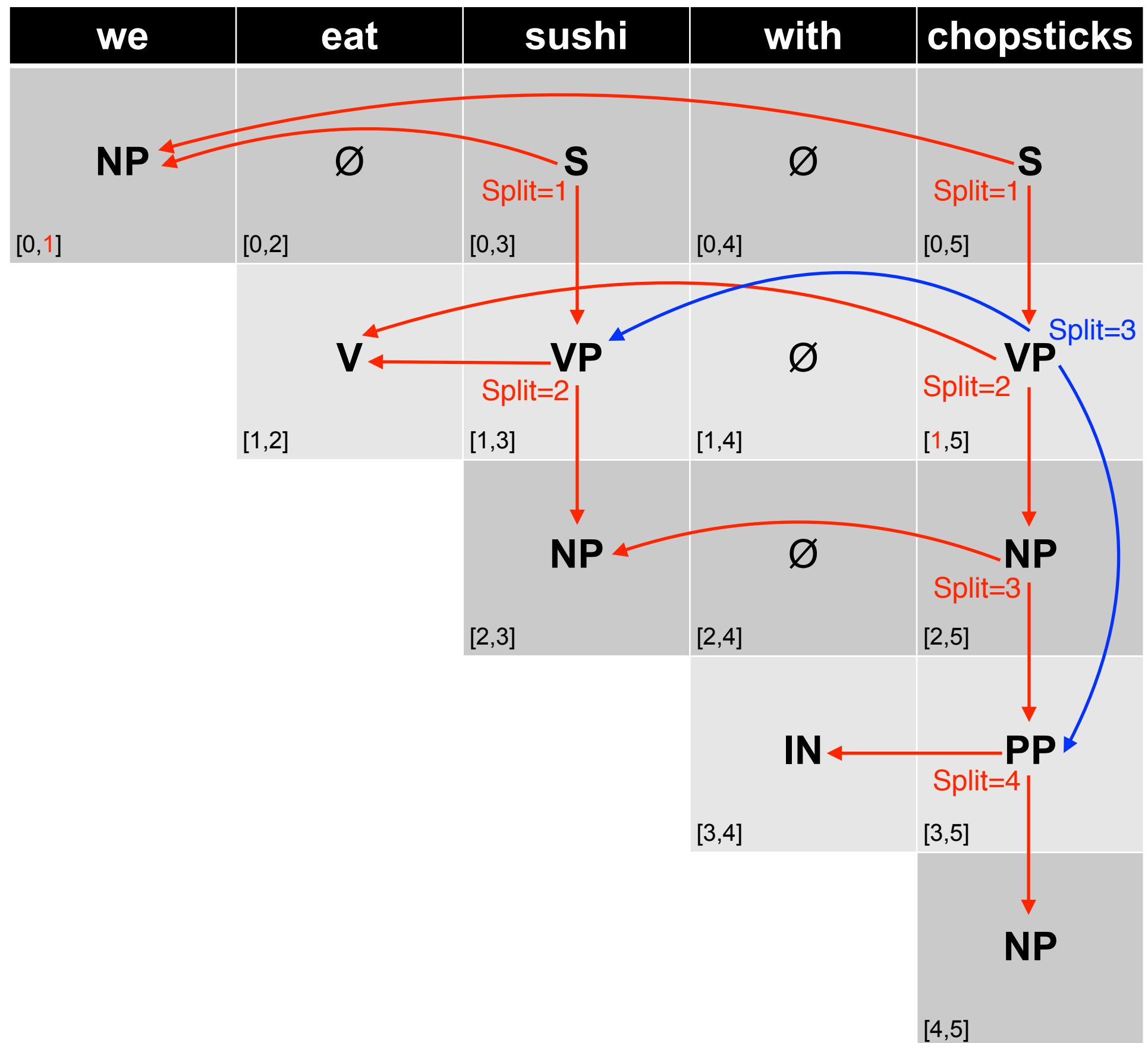
$V \rightarrow eat$



$S \rightarrow NP VP$   
 $NP \rightarrow NP PP$   
 $PP \rightarrow IN NP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow VP PP$   
 $NP \rightarrow we$   
 $NP \rightarrow sushi$   
 $NP \rightarrow chopsticks$   
 $IN \rightarrow with$   
 $V \rightarrow eat$



S → NP VP  
 NP → NP PP  
 PP → IN NP  
 VP → V NP  
 VP → VP PP  
 NP → we  
 NP → sushi  
 NP → chopsticks  
 IN → with  
 V → eat



**S → NP VP**

NP → NP PP

PP → IN NP

VP → V NP

VP → VP PP

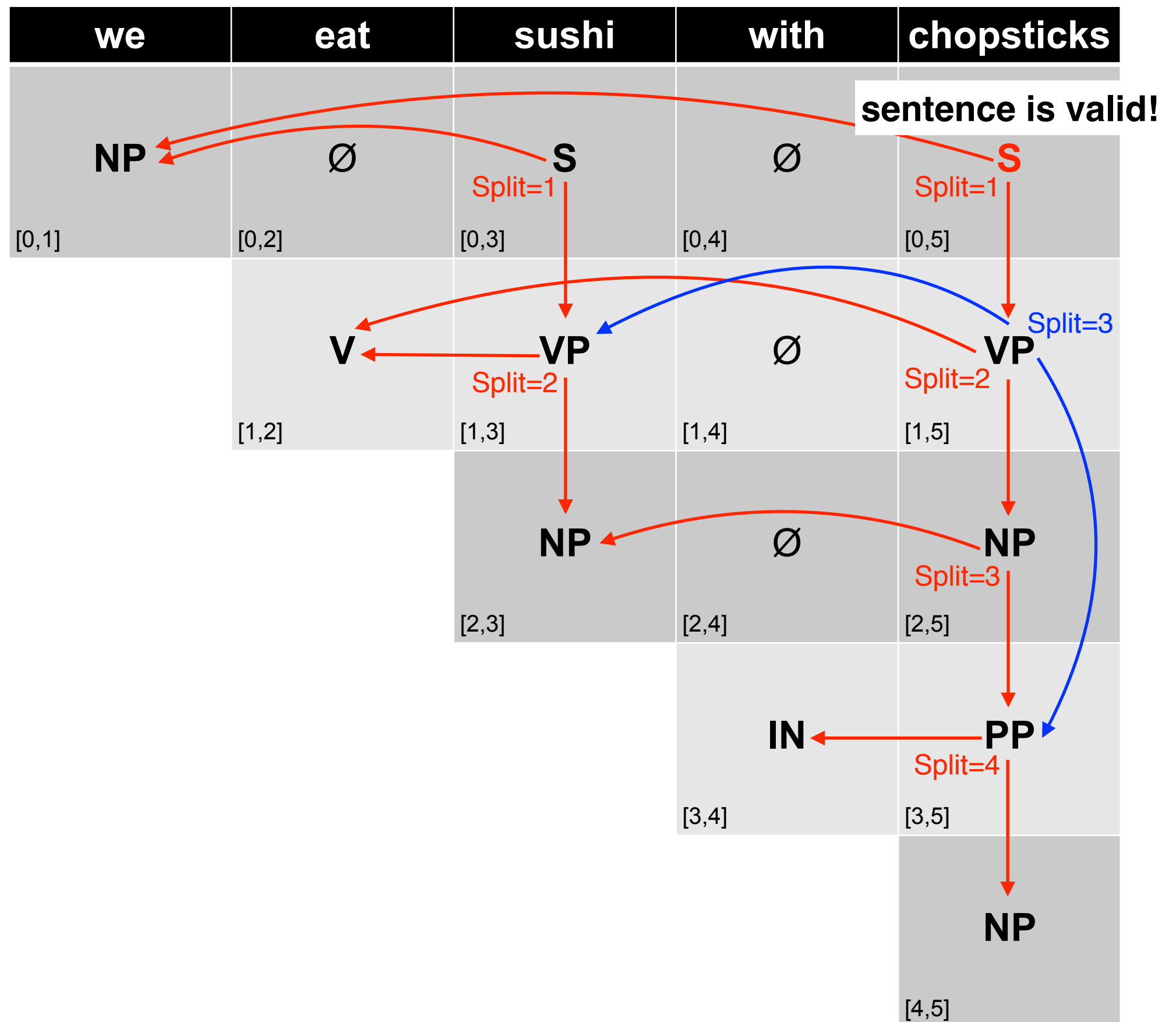
NP → we

NP → sushi

NP → chopsticks

IN → with

V → eat

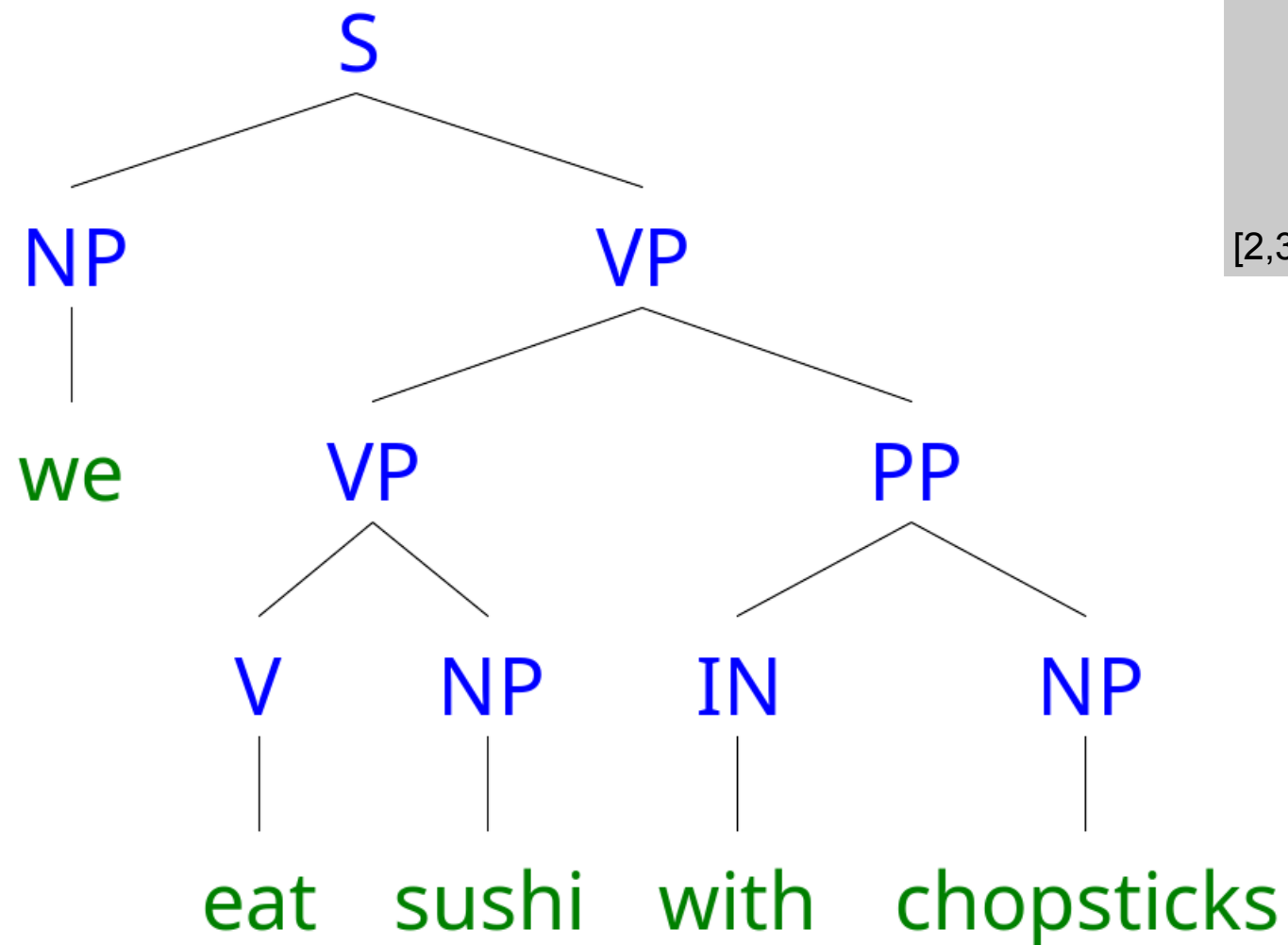
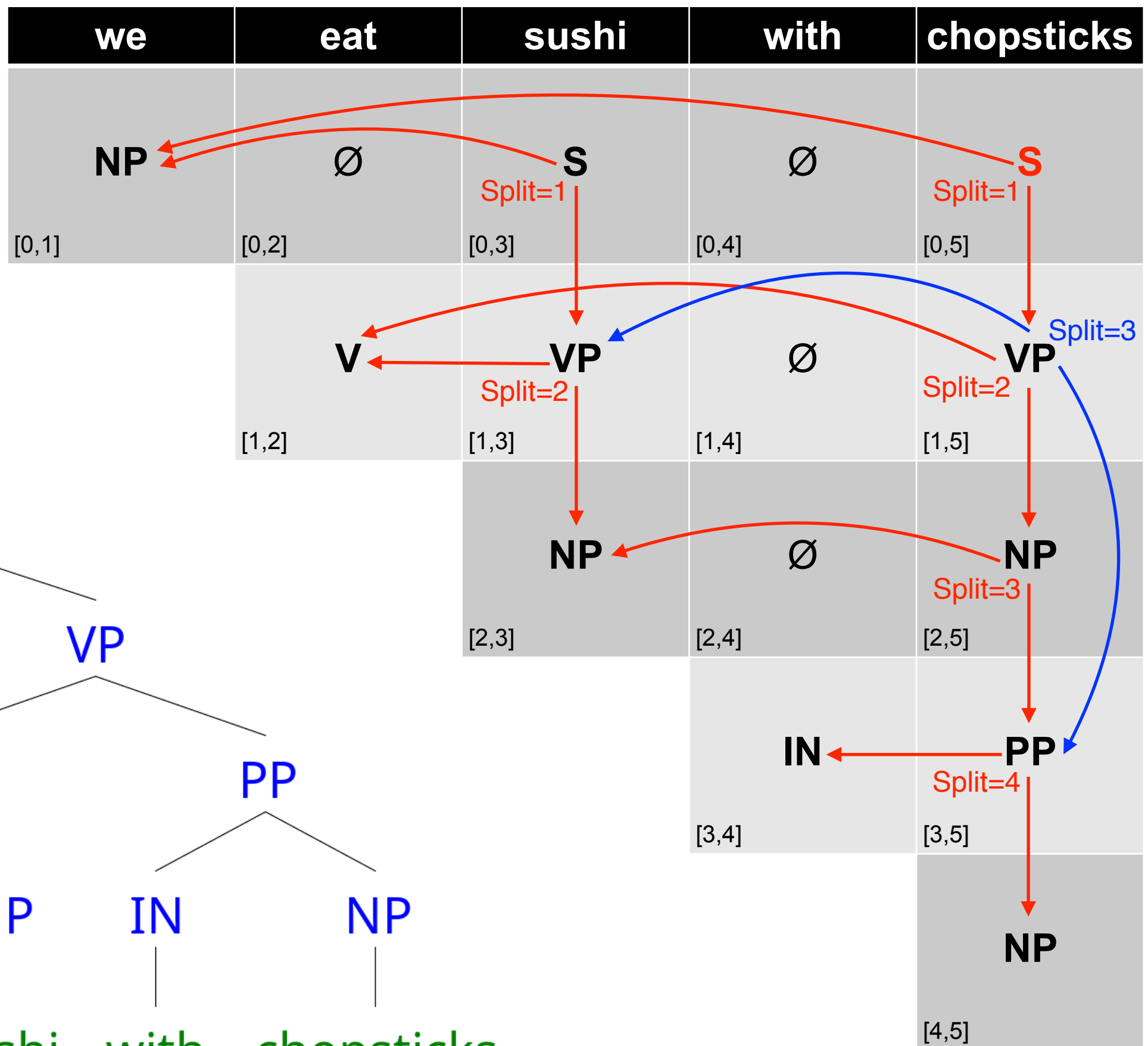


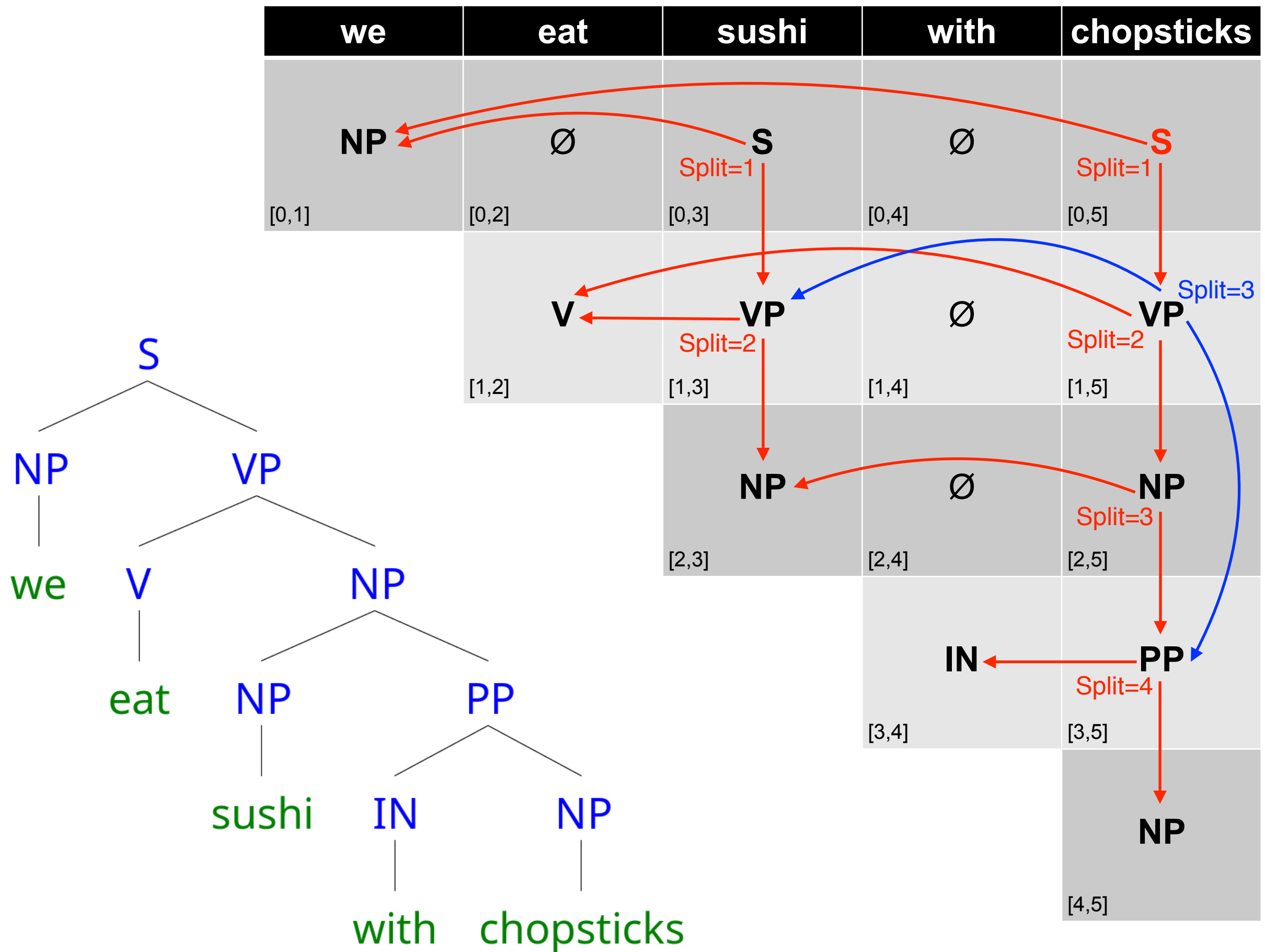
$S \rightarrow NP VP$   
 $NP \rightarrow NP PP$   
 $PP \rightarrow IN NP$   
 $VP \rightarrow V NP$   
 $VP \rightarrow VP PP$   
 $NP \rightarrow we$   
 $NP \rightarrow sushi$   
 $NP \rightarrow chopsticks$   
 $IN \rightarrow with$   
 $V \rightarrow eat$

# CYK: Retrieving the Parses

- S in the top-left corner of parse table indicates success
- To get parse(s), follow pointers back for each match







# CYK Algorithm

```
function CKY-PARSE(words, grammar) returns table

  for  $j \leftarrow$  from 1 to LENGTH(words) do
    for all  $\{A \mid A \rightarrow \text{words}[j] \in \text{grammar}\}$ 
       $\text{table}[j-1, j] \leftarrow \text{table}[j-1, j] \cup A$ 
    for  $i \leftarrow$  from  $j-2$  downto 0 do
      for  $k \leftarrow i+1$  to  $j-1$  do
        for all  $\{A \mid A \rightarrow BC \in \text{grammar} \text{ and } B \in \text{table}[i, k] \text{ and } C \in \text{table}[k, j]\}$ 
           $\text{table}[i, j] \leftarrow \text{table}[i, j] \cup A$ 
```

**Figure 12.5** The CKY algorithm.

# Representing English with CFGs

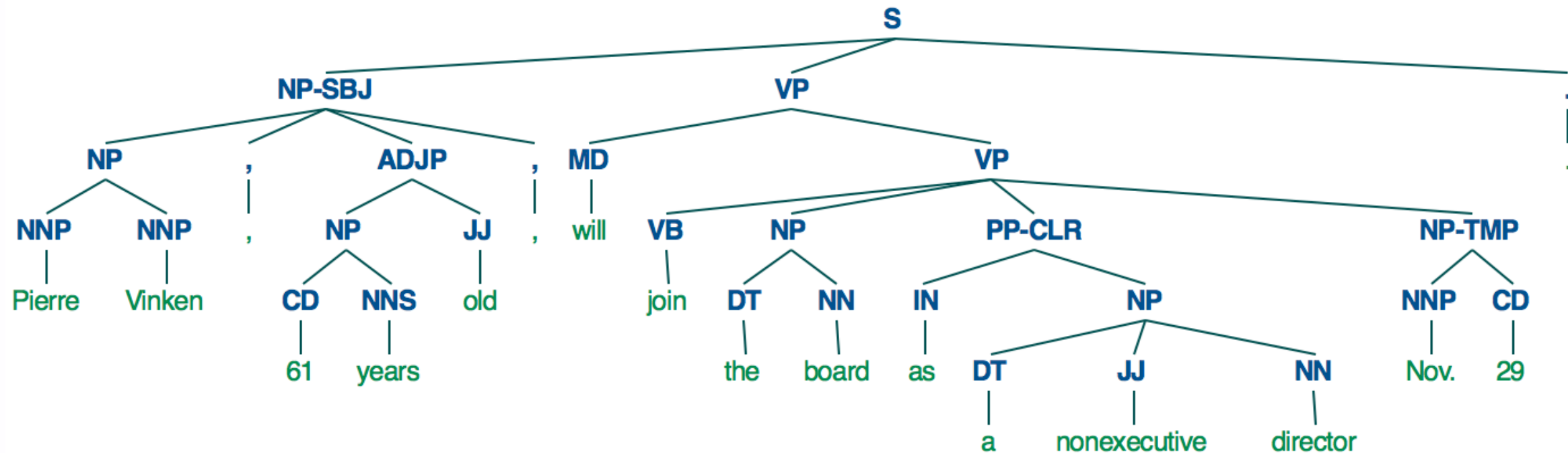
# From Toy Grammars to Real Grammars

- Toy grammars with handful of productions good for demonstration or extremely limited domains
- For real texts, we need real grammars
- Many thousands of production rules

# Key Constituents in Penn Treebank

- Sentence (S)
- Noun phrase (NP)
- Verb phrase (VP)
- Prepositional phrase (PP)
- Adjective phrase (AdjP)
- Adverbial phrase (AdvP)
- Subordinate clause (SBAR)

# Example PTB/0001



```

( (S
  (NP-SBJ
    (NP (NNP Pierre) (NNP Vinken) )
    (, ,)
    (ADJP
      (NP (CD 61) (NNS years) )
      (JJ old) )
    (, ,) )
  (VP (MD will)
    (VP (VB join)
      (NP (DT the) (NN board) )
      (PP-CLR (IN as)
        (NP (DT a) (JJ nonexecutive) (NN director) ))
      (NP-TMP (NNP Nov.) (CD 29) )))
    (. .) ))
  
```

# Basic English Sentence Structures

- Declarative sentences ( $S \rightarrow NP VP$ )
  - ▶ *The rat ate the cheese*
- Imperative sentences ( $S \rightarrow VP$ )
  - ▶ *Eat the cheese!*
- Yes/no questions ( $S \rightarrow VB NP VP$ )
  - ▶ *Did the rat eat the cheese?*
- *Wh*-subject-questions ( $S \rightarrow WH VP$ )
  - ▶ *Who ate the cheese?*
- *Wh*-object-questions ( $S \rightarrow WH VB NP VP$ )
  - ▶ *What did the rat eat?*



# English Noun Phrases

- Pre-modifiers
  - ▶ DT, CD, ADJP, NNPN, NN
  - ▶ E.g. *the two very best Philly cheese steaks*
- Post-modifiers
  - ▶ PP, VP, SBAR
  - ▶ A delivery *from Bob coming today that I don't want to miss*

NP → DT? CD? ADJP? (NNINNP)+ PP\* VP? SBAR?

NP → PRP

# Verb Phrases

- Auxiliaries

- ▶ MD, AdvP, VB, TO
- ▶ E.g. *should really have tried to wait*

VP → (MD|VB|TO) AdvP? VP

- Arguments and adjuncts

- ▶ NP, PP, SBAR, VP, AdvP
- ▶ E.g. *told him yesterday that I was ready*
- ▶ E.g. *gave John a gift for his birthday to make amends*

VP → VB NP? NP? PP\* AdvP\* VP? SBAR?

# Other Constituents

- Prepositional phrase
  - ▶  $PP \rightarrow IN\ NP$  *in the house*
- Adjective phrase
  - ▶  $AdjP \rightarrow (AdvP)\ JJ$  *really nice*
- Adverb phrase
  - ▶  $AdvP \rightarrow (AdvP)\ RB$  *not too well*
- Subordinate clause
  - ▶  $SBAR \rightarrow (IN)\ S$  *since I came here*
- Coordination
  - ▶  $NP \rightarrow NP\ CC\ NP; VP \rightarrow VP\ CC\ VP; \text{etc.}$  *Jack and Jill*
- Complex sentences
  - ▶  $S \rightarrow S\ SBAR; S \rightarrow SBAR\ S; \text{etc.}$  *if he goes, I'll go*

# A Final Word

- Context-free grammars can represent linguistic structure
- There are relatively fast dynamic programming algorithms to retrieve this structure
- But what about ambiguity?
  - ▶ Extreme ambiguity will slow down parsing
  - ▶ If multiple possible parses, which is best?

# Readings

- E18 Ch. 9.2, 10.1