





Indian Academy of Sciences, Bengaluru Indian National Science Academy, New Delhi

The National Academy of Sciences India, Prayagraj SUMMER RESEARCH FELLOWSHIPS — 2023

Format for the four-week Report*,^,@

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Name of the candidate : Daypar	s Bhattacharya
Application Registration no. : ENGS	302
Date of joining : 22.05	2023
	- Sunit Chandran
Guide's institution : T L	T III / C
	Institute of Science, Bengaluru
Place of stay during the tenure of : Hostel prov the fellowship Guide	ided by
Own arrang	ement
Øther (Spec	ify) IASc Fellows Residency,
Darpan Bhatlacharya	
Signature of the candidate	Signature of the guide
Date: 20.06.2023	Date: 21.06.2023
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3. I receive a monthly fellowship of Rs	
my living expenses	
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	Signature of the candidate
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IASc-INSA-NASI SRFP 2023

Four-week Report

Darpan Bhattacharya

Application ID: ENGS302

Guide: Prof L. Sunil Chandran, IISc Bengaluru

19 June, 2023

Introduction

This is the four-week report of my summer research project under Prof L. Sunil Chandran at Indian Institute of Science, Bengaluru.

I have primarily worked for the past four weeks trying to learn and understand the nature and number of monochromatic perfect matchings in graphs and *n*-uniform hypergraphs, and the main topics that I have learnt and worked on over the past four weeks are summarized below.

Perfect matchings in graphs

In a graph G = (V, E), a perfect matching is a matching M of the graph such that $\forall v \in V(G)$, v is a part of some matched edge in M.

For a perfectly monochromatic colouring c of G, the number of normal colour classes containing at least one perfect matching is represented by $\mu(G,c)$. The matching index of G, denoted by $\mu(G)$ is defined as $\mu(G) = \max_{c \in \mathcal{C}(G)} \mu(G,c)$, where $\mathcal{C}(G)$ is the set of all perfectly monochromatic colourings of G.

In a graph, Bogdanov[3] proved that $\forall G$ isomorphic to K_4 , the number of disjoint perfect matchings of G = 3, and for all other graphs non-isomorphic to K_4 , the number of disjoint perfect matchings is at most 2. This proof was used in [2] to theorize that $\forall G$ isomorphic to K_4 , $\mu(G) = 3$, otherwise $\mu(G) \leq 2$.

Hypergraphs

According to [1], a hypergraph is a pair H = (V, E), where V is a finite set whose elements are called *vertices*, and E is a family of subsets of V, called *edges*.

A hypergraph H is said to be n-uniform if $\forall e \in E$, e contains precisely n vertices.

A perfect matching, \mathcal{M} of a n-uniform hypergraph H = (V, E) can be defined as a matching M of H such that $\forall v \in V$, v is part of some edge in M. It is trivial to note that for a n-uniform hypergraph to have a perfect matching, $|V| \equiv 0 \pmod{n}$.

A n-uniform hypergraph H is said to be perfectly monochromatic if the following properties hold:

1. Every perfect matching, \mathcal{M} of H is monochromatic.

2. Let \mathcal{C} be set of colour classes such that \exists a perfect matching \mathcal{M} of H such that \mathcal{M} belongs to a colour class $c \in \mathcal{C}$. Then $\forall c \in \mathcal{C}$, c must belong to exactly 1 perfect matching \mathcal{M} of H.

Now, for a perfectly monochromatic hypergraph H, we call the dimension $D = |\mathcal{C}|$ of the hypergraph to be the number of colours in it. We are intereseted in finding the bounds on D for a n-uniform hypergraph containing $k \cdot n$ vertices.

Perfect matchings in 3-uniform hypergraphs

Since we are interested in finding the bounds on D, we start with trying to find the bounds on D for 3-uniform hypergraphs. We analyzed some particular cases and the results we came up with are:

- For |V| = 3, it can be trivially observed that D = 1.
- For |V| = 6, the upper bound on D = 10. This can be verified by observing that in a complete 3-uniform hypergraph with |V| = 6, there are $\binom{6}{3} = 20$ edges. We can colour each edge and its complement with the same colour, forming a monochromatic perfect matching, and thus the number of monochromatic perfect matchings for a 3-uniform hypergraph with |V| = 6 is 20/2 = 10.
- For |V| = 9, the upper bound on D turned out to be 13. I verified it independently by writing a program which found out the value of D by combining randomization and brute force techniques.
- For |V| = 12, I developed a partly randomized program to compute the bounds on D, and the upper bound from that program turned out to be 10, which surprisingly enough, turns out to be the same value as for |V| = 6.

Besides trying to analyze and compute the bounds for D for 3-uniform hypergraphs, I have also tried implementing Algorithm 1 - To decide whether a non-trivial matching covered graph non-isomorphic to K_4 is Type 1 or Type 2 in [2], in the process of which I learnt about the Edmonds' blossom algorithm for finding the maximum number of perfect matchings in a graph, and the Edmonds-Karp algorithm for finding the maximum flow in a graph.

References

- [1] N. Alon and J. H. Spencer. *The Probabilistic Method*. Wiley Publishing, 4th edition, 2016.
- [2] L. S. Chandran and R. Gajjala. Perfect matchings and quantum physics: Progress on krenn's conjecture, 2022.
- [3] I. B. (https://mathoverflow.net/users/17581/ilya bogdanov). Graphs with only disjoint perfect matchings. MathOverflow. URL:https://mathoverflow.net/q/267013 (version: 2017-04-12).