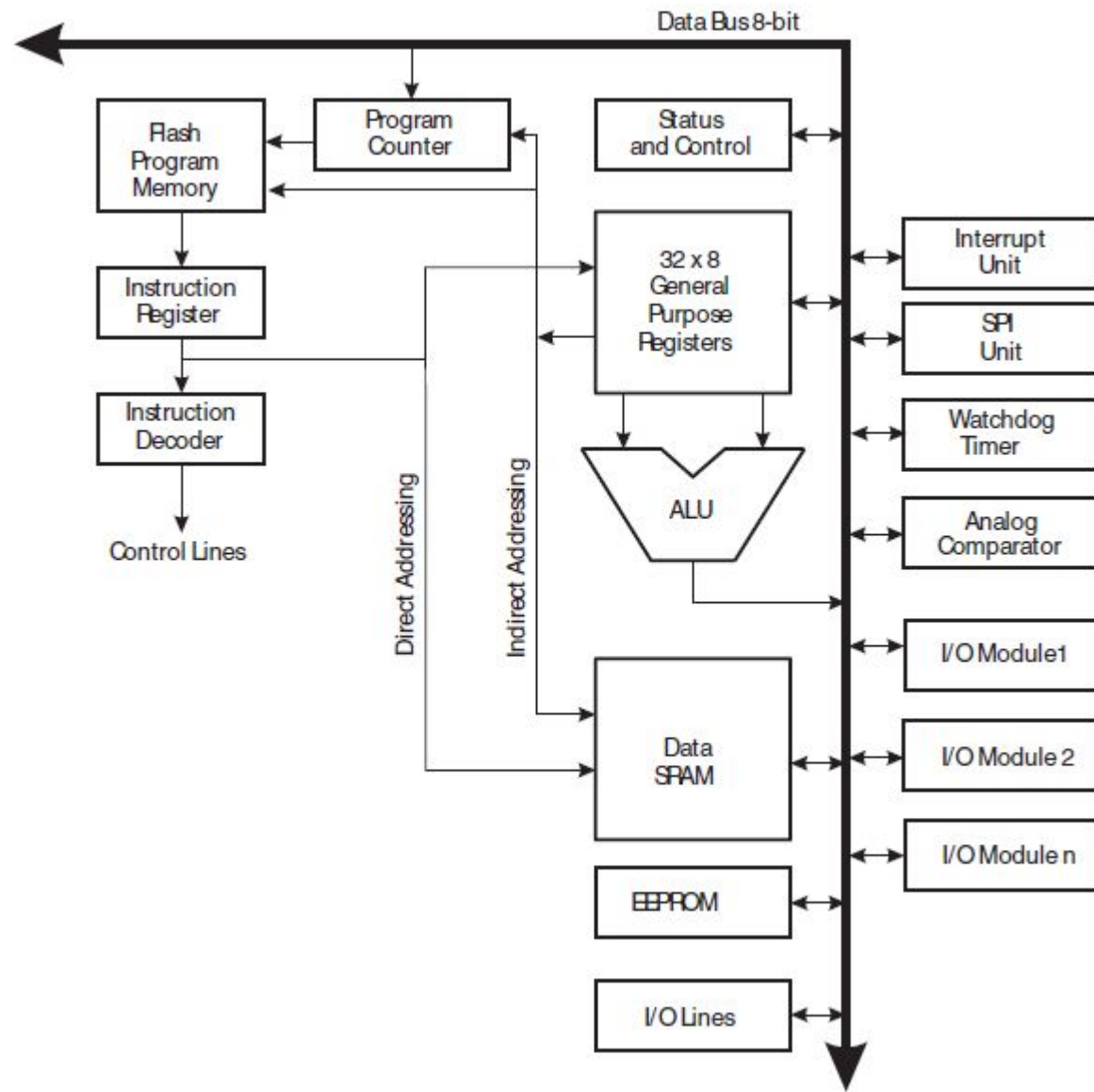


Núcleo (AVR CPU Core)

Sistemas Microcontrolados



ALU - Arithmetic Logic Unit

Mnemonics	Operands	Description	Operation	Flags	#Clocks
ARITHMETIC AND LOGIC INSTRUCTIONS					
ADD	Rd, Rr	Add two Registers	$Rd \leftarrow Rd + Rr$	Z, C, N, V, H	1
ADC	Rd, Rr	Add with Carry two Registers	$Rd \leftarrow Rd + Rr + C$	Z, C, N, V, H	1
ADIW	RdI, K	Add Immediate to Word	$Rdh:Rdl \leftarrow Rdh:Rdl + K$	Z, C, N, V, S	2
SUB	Rd, Rr	Subtract two Registers	$Rd \leftarrow Rd - Rr$	Z, C, N, V, H	1
SUBI	Rd, K	Subtract Constant from Register	$Rd \leftarrow Rd - K$	Z, C, N, V, H	1
SBC	Rd, Rr	Subtract with Carry two Registers	$Rd \leftarrow Rd - Rr - C$	Z, C, N, V, H	1
SBCI	Rd, K	Subtract with Carry Constant from Reg.	$Rd \leftarrow Rd - K - C$	Z, C, N, V, H	1
SBIW	RdI, K	Subtract Immediate from Word	$Rdh:Rdl \leftarrow Rdh:Rdl - K$	Z, C, N, V, S	2
AND	Rd, Rr	Logical AND Registers	$Rd \leftarrow Rd \bullet Rr$	Z, N, V	1
ANDI	Rd, K	Logical AND Register and Constant	$Rd \leftarrow Rd \bullet K$	Z, N, V	1
OR	Rd, Rr	Logical OR Registers	$Rd \leftarrow Rd \vee Rr$	Z, N, V	1
ORI	Rd, K	Logical OR Register and Constant	$Rd \leftarrow Rd \vee K$	Z, N, V	1
EOR	Rd, Rr	Exclusive OR Registers	$Rd \leftarrow Rd \oplus Rr$	Z, N, V	1
COM	Rd	One's Complement	$Rd \leftarrow 0xFF - Rd$	Z, C, N, V	1
NEG	Rd	Two's Complement	$Rd \leftarrow 0x00 - Rd$	Z, C, N, V, H	1
SBR	Rd, K	Set Bit(s) in Register	$Rd \leftarrow Rd \vee K$	Z, N, V	1
CBR	Rd, K	Clear Bit(s) in Register	$Rd \leftarrow Rd \bullet (0xFF - K)$	Z, N, V	1
INC	Rd	Increment	$Rd \leftarrow Rd + 1$	Z, N, V	1
DEC	Rd	Decrement	$Rd \leftarrow Rd - 1$	Z, N, V	1
TST	Rd	Test for Zero or Minus	$Rd \leftarrow Rd \bullet Rd$	Z, N, V	1

ALU - Arithmetic Logic Unit

Instruction	Op	Description	Operation	Flags	Cycles
CLR	Rd	Clear Register	$Rd \leftarrow Rd \oplus Rd$	Z, N, V	1
SER	Rd	Set Register	$Rd \leftarrow 0xFF$	None	1
MUL	Rd, Rr	Multiply Unsigned	$R1:R0 \leftarrow Rd \times Rr$	Z, C	2
MULS	Rd, Rr	Multiply Signed	$R1:R0 \leftarrow Rd \times Rr$	Z, C	2
MULSU	Rd, Rr	Multiply Signed with Unsigned	$R1:R0 \leftarrow Rd \times Rr$	Z, C	2
FMUL	Rd, Rr	Fractional Multiply Unsigned	$R1:R0 \leftarrow (Rd \times Rr) \ll 1$	Z, C	2
FMULS	Rd, Rr	Fractional Multiply Signed	$R1:R0 \leftarrow (Rd \times Rr) \ll 1$	Z, C	2
FMULSU	Rd, Rr	Fractional Multiply Signed with Unsigned	$R1:R0 \leftarrow (Rd \times Rr) \ll 1$	Z, C	2
BRANCH INSTRUCTIONS					
RJMP	k	Relative Jump	$PC \leftarrow PC + k + 1$	None	2
IJMP		Indirect Jump to (Z)	$PC \leftarrow Z$	None	2
EIJMP		Extended Indirect Jump to (Z)	$PC \leftarrow (EIND:Z)$	None	2
JMP	k	Direct Jump	$PC \leftarrow k$	None	3
RCALL	k	Relative Subroutine Call	$PC \leftarrow PC + k + 1$	None	4
ICALL		Indirect Call to (Z)	$PC \leftarrow Z$	None	4
EICALL		Extended Indirect Call to (Z)	$PC \leftarrow (EIND:Z)$	None	4
CALL	k	Direct Subroutine Call	$PC \leftarrow k$	None	5
RET		Subroutine Return	$PC \leftarrow STACK$	None	5
RETI		Interrupt Return	$PC \leftarrow STACK$	I	5
CPSE	Rd,Rr	Compare, Skip if Equal	if $(Rd = Rr)$ $PC \leftarrow PC + 2$ or 3	None	1/2/3
CP	Rd,Rr	Compare	$Rd - Rr$	Z, N, V, C, H	1
CPC	Rd,Rr	Compare with Carry	$Rd - Rr - C$	Z, N, V, C, H	1

Status Register

- ▶ O *Status Register* contém informações sobre os resultados de operações aritméticas mais recentes.
- ▶ Estas informações podem ser usadas para alterar o fluxo de programa para atender instruções condicionais.
- ▶ Portanto o *Status Register* é atualizado após cada operação da ALU
- ▶ Em muitos casos excluí-se a necessidade de utilizar instruções de comparação dedicadas resultando em maior velocidade e código mais compacto.

7.4.1 SREG – AVR Status Register

The AVR Status Register – SREG – is defined as:

Bit	7	6	5	4	3	2	1	0	
0x3F (0x5F)	I	T	H	S	V	N	Z	C	SREG
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7 – I: Global Interrupt Enable**

The Global Interrupt Enable bit must be set for the interrupts to be enabled. The individual interrupt enable control is then performed in separate control registers. If the Global Interrupt Enable Register is cleared, none of the interrupts are enabled independent of the individual interrupt enable settings. The I-bit is cleared by hardware after an interrupt has occurred, and is set by the RETI instruction to enable subsequent interrupts. The I-bit can also be set and cleared by the application with the SEI and CLI instructions, as described in the [“Instruction Set Summary” on page 404](#).

- **Bit 6 – T: Bit Copy Storage**

The Bit Copy instructions BLD (Bit Load) and BST (Bit Store) use the T-bit as source or destination for the operated bit. A bit from a register in the Register File can be copied into T by the BST instruction, and a bit in T can be copied into a bit in a register in the Register File by the BLD instruction.

BST	Rr, b	Bit Store from Register to T	$T \leftarrow Rr(b)$	T	1
BLD	Rd, b	Bit load from T to Register	$Rd(b) \leftarrow T$	None	1

- **Bit 5 – H: Half Carry Flag**

The Half Carry Flag H indicates a Half Carry in some arithmetic operations. Half Carry Is useful in BCD arithmetic. See the [“Instruction Set Summary” on page 404](#) for detailed information.

- **Bit 4 – S: Sign Bit, $S = N \oplus V$**

The S-bit is always an exclusive or between the Negative Flag N and the Two's Complement Overflow Flag V. See the [“Instruction Set Summary” on page 404](#) for detailed information.

- **Bit 3 – V: Two's Complement Overflow Flag**

The Two's Complement Overflow Flag V supports two's complement arithmetics. See the [“Instruction Set Summary” on page 404](#) for detailed information.

- **Bit 2 – N: Negative Flag**

The Negative Flag N indicates a negative result in an arithmetic or logic operation. See the [“Instruction Set Summary” on page 404](#) for detailed information.

- **Bit 1 – Z: Zero Flag**

The Zero Flag Z indicates a zero result in an arithmetic or logic operation. See the [“Instruction Set Summary” on page 404](#) for detailed information.

- **Bit 0 – C: Carry Flag**

The Carry Flag C indicates a carry in an arithmetic or logic operation. See the [“Instruction Set Summary” on page 404](#) for detailed information.

General Purpose Register File

Figure 7-2. AVR CPU General Purpose Working Registers

	7	0	Addr.	
General Purpose Working Registers	R0		0x00	
	R1		0x01	
	R2		0x02	
	...			
	R13		0x0D	
	R14		0x0E	
	R15		0x0F	
	R16		0x10	
	R17		0x11	
	...			
	R26		0x1A	X-register Low Byte
	R27		0x1B	X-register High Byte
	R28		0x1C	Y-register Low Byte
	R29		0x1D	Y-register High Byte
	R30		0x1E	Z-register Low Byte
	R31		0x1F	Z-register High Byte

Stack Pointer

- ▶ A Pilha “Stack” é usado principalmente para armazenar dados temporariamente, variáveis locais e endereço de retorno após interrupções e chamadas de sub-rotinas.
- ▶ O registrador de ponteiro de pilha “Stack Pointer Register” sempre aponta para o topo da pilha.
- ▶ Note que a pilha é implementada de forma decrescente, de posições de memória maiores para menores. Isso significa que o comando “Stack PUSH” decrementa o ponteiro de memória.

Stack Pointer

- ▶ O stack pointer deve estar configurado para o endereço 0x0200.
- ▶ O endereço inicial da Pilha é o ultimo endereço da SRAM interna.
- ▶ O Stack pointer é decrementado quando dados são empurrados na pilha pela instrução PUSH e incrementado quando os dados forem retirado com a instrução POP.