

Interview Questions on Binary Search

1) facebook (Technical Round)

\hookrightarrow Not sorted manner $n \rightarrow$ size of an array



$$n = 12$$

$$6 + 7 / 2 = 6$$

array



$$\begin{aligned} \text{mid} &= \frac{6 + 11}{2} \\ &= 8 \end{aligned}$$

output = 7

as l.

✓ ↗ Linear search

$\left. \begin{array}{l} \text{for } i = 0 \text{ to } n-1 \\ \text{arr}(i) == 'inf' \\ \text{return } i \end{array} \right\} \rightarrow O(n)$

✓ 2) Modified Binary Search $\rightarrow O(\log n)$

$$\text{mid} = 0 + (11 - 0) // 2 = 0 + 5 = 5$$

if $\text{arr}(\text{mid}) == \text{int}$:

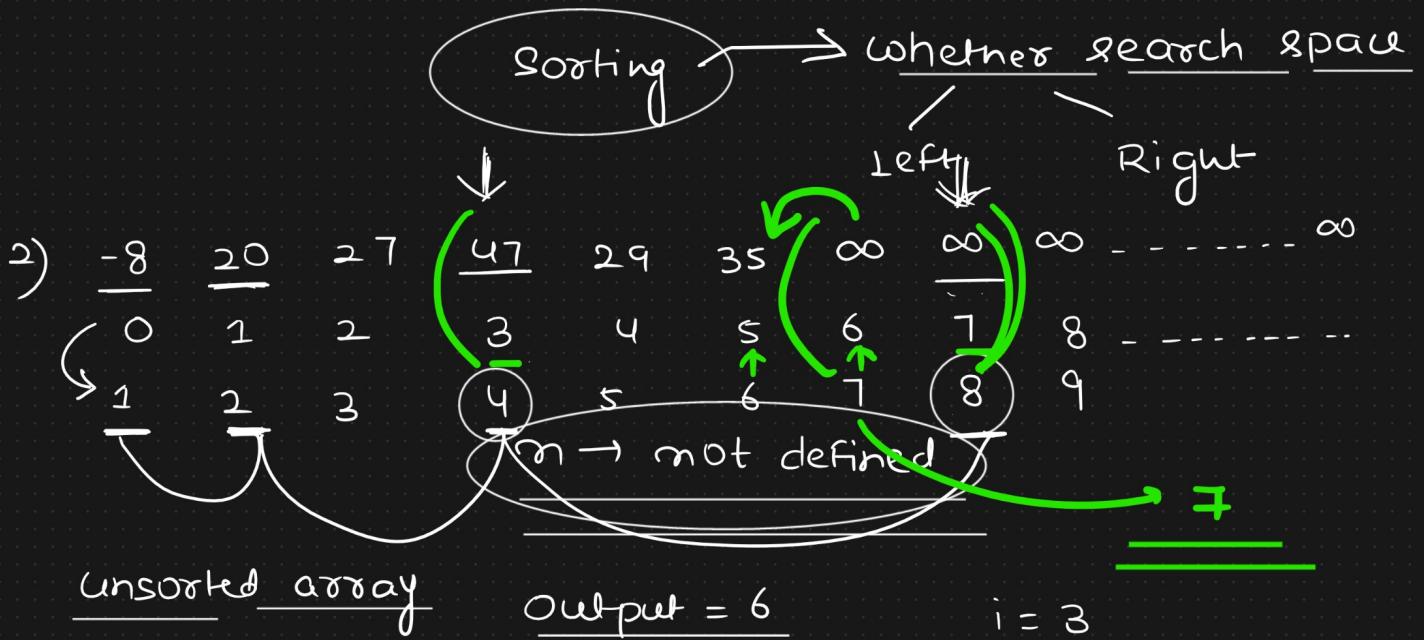
Right side \leftarrow left = mid + 1
Search space

$\text{arr}(\text{mid}) \leq \underline{\text{inf}}$

$\text{right} = \text{mid} - 1$

Note → 1) Don't judge by looking whether the array is physically sorted or not.

\hookrightarrow Logical sort



Note:
Infinite numbers

Exponential increase \rightarrow range of numbers

are only available
after all
integers.

Binary search

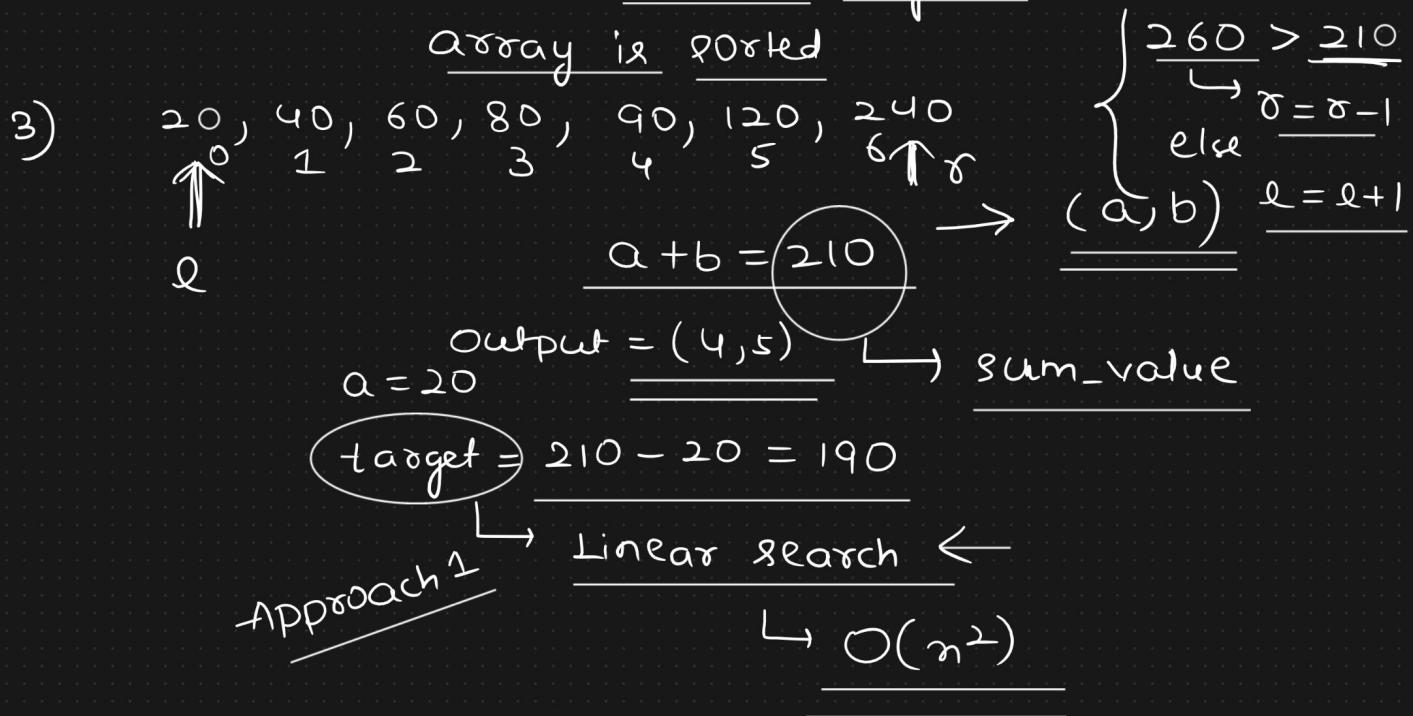
more numbers

↙

increasing search space

1 \rightarrow 2 \rightarrow 4 \rightarrow 8 \rightarrow 16

Problem Solving Skill



~~Amazon
Google~~

Approach 3

Two Pointer Approach

$l = \text{starting index}$

$r = \text{ending index}$

$O(n)$

while $l < r$:

if $\text{arr}(l) + \text{arr}(r) == \text{sum_value}$

return (l, r)

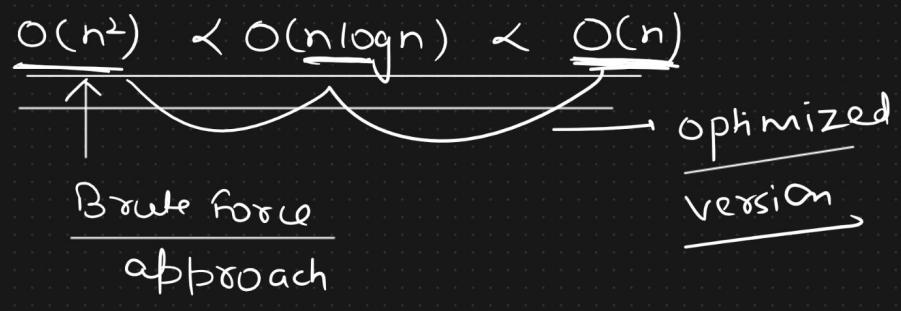
elif $\text{arr}(l) + \text{arr}(r) > \text{sum_value}$

$\text{arr}(l) + \text{arr}(r) < \text{sum_value}$

$r = r - 1$

else:

$l = l + 1$



→ Best time to buy & sell
 1 2 3 4 5 6) the stock
 4) $\text{price} = [7, 1, 5, 3, 6, 4]$
 $\max \underline{\underline{\text{profit}}} = 6 - 1 = 5$

for $i = 0$ to $n-1$
 if $\underline{\underline{\text{price}(i) < \minPrice}}$:
 $\underline{\underline{\minPrice = \text{price}(i)}}$
 elif $\text{price}(i) - \minPrice > \maxProfit$
 $\maxProfit = \text{price}(i) - \minPrice$
 return \maxProfit

5) $\frac{2D \text{ Arrays}}{0 \quad 1 \quad 2} \rightarrow (\text{Rows, columns})$

\Rightarrow 0 $\begin{bmatrix} 2 & 4 & 6 \\ 8 & 10 & 12 \\ 20 & 27 & 40 \end{bmatrix}$ \Rightarrow Matrix Computation
 \hookrightarrow 1 #rows ↓
 2 #columns ↑ $\begin{cases} \text{matrix multiplication,} \\ \text{matrix transpose} \end{cases}$
 Row number
 Column number
 Arr [1] [2]

Amazon

$m \times n$

$\begin{cases} m \rightarrow \text{no of rows} \\ n \rightarrow \text{no of columns} \end{cases}$

Constraints

- 1) Integers in each row is sorted from left to right

- 2) first integer of each row > Last integer of previous row
 $10 > 7, 23 > 20$

2D

	0	1	2	3	
0	1	3	5	7	
1	10	11	16	20	
2	23	30	34	60	

3×4

row number \downarrow_1

Column number \downarrow_1

Output \Rightarrow True

target = 3 \downarrow_1

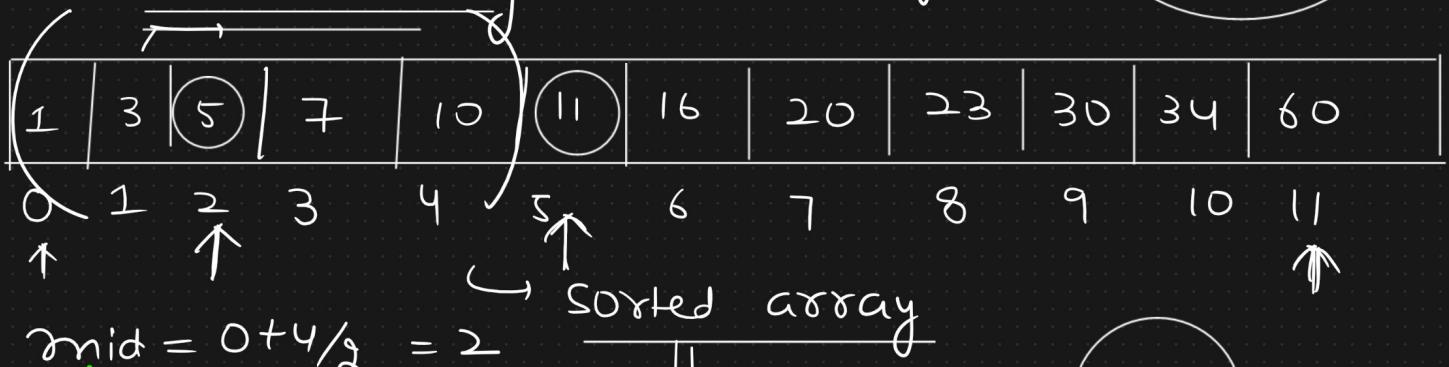
Optimised Approach

Output \Rightarrow False

target = 3

$m = 3$
 $n = 4$

Virtual array



$$m_{\text{mid}} = 0 + 1 / 2 = 0$$

$$m_{\text{mid}} = 0 + 11 / 2 = 5$$

Binary Search

$J = 11 = m \times n - 1$

$$\begin{array}{c}
 \underline{m=4} \qquad \underline{mid=5} \\
 \underline{5//4 = 1} \\
 \Rightarrow \left\{ \begin{array}{l} \underline{\text{row-number} = \underline{mid // m}} \\ \underline{\text{col-number} = \underline{mid \% n}} \end{array} \right. \\
 \underline{5 \% 4 = 1}
 \end{array}$$

$$\underline{mid = 2}$$

$$\underline{\text{row-number} = 2//4 = 0}$$

$$\underline{\text{Column-number} = 2 \% 4 = 2}$$

$$\begin{array}{c}
 \underline{mid = i + (j-i)/2} \\
 \underline{i = 0 \quad j = m * n - 1}
 \end{array}$$

Optimized approach

$$\underline{\text{Binary Search} \rightarrow \log_2(m * n)}$$

$$\Rightarrow \underline{O(\log_2(m * n))}$$

0	<u>1</u>	<u>3</u>	<u>5</u>	<u>7</u>	
1	<u>10</u>	<u>11</u>	<u>16</u>	<u>20</u>	
3	<u>23</u>	<u>30</u>	<u>34</u>	<u>60</u>	m * n

target = 60

Brute force approach $\mathcal{O}(m * n)$

$\mathcal{O}(m * n)$

for $i = 0$ to $m - 1$:
 for $j = 0$ to $n - 1$:
 if $array(i, j) == \text{target}$:
 return True
 return False