

## CS 635 Programming Project

One of the manipulations done in converting a CFG to Chomsky Normal Form (CNF) is to replace re-writing rules with a RHS that is not exactly two non-terminals with other rules. In what follows, small letters represent terminals, capital letters from the beginning of the alphabet (like A, B, C) and anything between  $\langle \rangle$  are non-terminals. Capital letters from the end of the alphabet (like X, Y, Z) represent a symbol that may be a terminal or a non-terminal. Greek letters represent strings that may contain any mix of terminals and non-terminals.

Step 1:

For a rule  $A \rightarrow X\beta$

Split the RHS into two parts, the first symbol X, and the rest  $\beta$

For each part, if it is a single non-terminal, don't change it. Otherwise put  $\langle \rangle$  around it, ie replace it with  $\langle X \rangle$  or  $\langle \beta \rangle$ .

Example:

$A \rightarrow BC$	unchanged
$A \rightarrow aB$	replace with $A \rightarrow \langle a \rangle B$
$A \rightarrow BCD$	replace with $A \rightarrow B \langle CD \rangle$
$A \rightarrow bcd$	replace with $A \rightarrow \langle b \rangle \langle cd \rangle$

Step 2:

Now, for any new non-terminals created, for a symbol of the form  $\langle X\beta \rangle$  add a re-writing rule with  $\langle X\beta \rangle$  as the LHS and the X and  $\beta$  are treated as described above to get the RHS: For both the X and  $\beta$ , if it is a single non-terminal, don't change it. Otherwise put  $\langle \rangle$  around it, ie replace it with  $\langle X \rangle$  or  $\langle \beta \rangle$ . The two resulting symbols together for the RHS. For a symbol of the form  $\langle a \rangle$  add  $\langle a \rangle \rightarrow a$ .

Example:

for the symbol $\langle ABC \rangle$	add a rule $\langle ABC \rangle \rightarrow A \langle BC \rangle$
for the symbol $\langle aBC \rangle$	add a rule $\langle aBC \rangle \rightarrow \langle a \rangle \langle BC \rangle$
for the symbol $\langle AB \rangle$	add a rule $\langle AB \rangle \rightarrow AB$
for the symbol $\langle a \rangle$	add a rule $\langle a \rangle \rightarrow a$

Repeat step 2 for any new non-terminals created.

Complete Example:

$A \rightarrow aBcdE$

Step 1:

$A \rightarrow \langle a \rangle \langle BcdE \rangle$  creates 2 new non-terminals:  $\langle a \rangle$  and  $\langle BcdE \rangle$

Step 2:

$\langle a \rangle \rightarrow a$

$\langle BcdE \rangle \rightarrow B \langle cdE \rangle$  creates new non-terminal  $\langle cdE \rangle$

repeat step 2:

$\langle cdE \rangle \rightarrow \langle c \rangle \langle dE \rangle$  creates 2 new non-terminals:  $\langle c \rangle$  and  $\langle dE \rangle$

repeat step 2:

$\langle c \rangle \rightarrow c$

$\langle dE \rangle \rightarrow \langle d \rangle E$  creates new non-terminal  $\langle d \rangle$

repeat step 2:

$\langle d \rangle \rightarrow d$

Finished.

Write two programs, one in your choice of C++ or Java, the other in Python, to execute the above algorithm.

Input:

$A \rightarrow aBcdE$

Output:

$A \rightarrow \langle a \rangle \langle BcdE \rangle$

$\langle a \rangle \rightarrow a$

$\langle BcdE \rangle \rightarrow B \langle cdE \rangle$

$\langle cdE \rangle \rightarrow \langle c \rangle \langle dE \rangle$

$\langle c \rangle \rightarrow c$

$\langle dE \rangle \rightarrow \langle d \rangle E$

$\langle d \rangle \rightarrow d$