EvalML5

Syntax:

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i \in \operatorname{int} b \in \operatorname{bool} x, y \in \operatorname{Var} v \in \operatorname{Value} ::= i \mid b \mid (\mathcal{E}) \text{ [fun } x \to e \text{] } \mid (\mathcal{E}) \text{ [rec } x = \operatorname{fun } y \to e \text{] } \mid [] \mid v :: v \mathcal{E} \in \operatorname{Env} ::= \bullet \mid \mathcal{E}, x = v p \in \operatorname{Pat} ::= x \mid [] \mid p :: p \mid \_ res \in \operatorname{Res} ::= \mathcal{E} \mid F c \in \operatorname{Clauses} ::= p \to e \mid p \to e \mid c e \in \operatorname{Exp} ::= i \mid b \mid x \mid e \ op \ e \mid \text{ if } e \ \text{ then } e \ \text{ else } e \mid \text{ let } x = e \ \text{ in } e \mid [] \mid e :: e \mid \operatorname{match} \ e \ \text{ with } \ c op \in \operatorname{Prim} ::= + \mid - \mid * \mid <
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空の環境 ● (とそれに続くコンマ) は入力時には省略する.

Derivation Rules:

$$\frac{}{\mathcal{E} \vdash i \Downarrow i}$$
 (E-Int)

$$\frac{}{\mathcal{E} \vdash b \Downarrow b} \tag{E-Bool}$$

$$\frac{\mathcal{E} \vdash e_1 \Downarrow \text{true} \qquad \mathcal{E} \vdash e_2 \Downarrow v}{\mathcal{E} \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow v}$$
 (E-IFT)

$$\frac{\mathcal{E} \vdash e_1 \Downarrow \mathtt{false} \qquad \mathcal{E} \vdash e_3 \Downarrow v}{\mathcal{E} \vdash \mathtt{if} \ e_1 \ \mathtt{then} \ e_2 \ \mathtt{else} \ e_3 \Downarrow v} \tag{E-IFF}$$

$$\frac{\mathcal{E} \vdash e_1 \Downarrow i_1 \qquad \mathcal{E} \vdash e_2 \Downarrow i_2 \qquad i_1 \text{ plus } i_2 \text{ is } i_3}{\mathcal{E} \vdash e_1 + e_2 \Downarrow i_3} \tag{E-Plus}$$

$$\frac{\mathcal{E} \vdash e_1 \Downarrow i_1 \qquad \mathcal{E} \vdash e_2 \Downarrow i_2 \qquad i_1 \text{ minus } i_2 \text{ is } i_3}{\mathcal{E} \vdash e_1 \vdash e_2 \Downarrow i_3} \tag{E-Minus}$$

$$\frac{\mathcal{E} \vdash e_1 \Downarrow i_1 \qquad \mathcal{E} \vdash e_2 \Downarrow i_2 \qquad i_1 \text{ times } i_2 \text{ is } i_3}{\mathcal{E} \vdash e_1 * e_2 \Downarrow i_3} \tag{E-TIMES}$$

$$\frac{\mathcal{E} \vdash e_1 \Downarrow i_1 \qquad \mathcal{E} \vdash e_2 \Downarrow i_2 \qquad i_1 \text{ less than } i_2 \text{ is } b_3}{\mathcal{E} \vdash e_1 \lessdot e_2 \Downarrow b_3} \tag{E-LT}$$

$$\frac{(\mathcal{E}(x) = v)}{\mathcal{E} \vdash x \Downarrow v}$$
 (E-VAR)

$$\frac{\mathcal{E} \vdash e_1 \Downarrow v_1 \qquad \mathcal{E}, x = v_1 \vdash e_2 \Downarrow v}{\mathcal{E} \vdash \text{let } x = e_1 \text{ in } e_2 \Downarrow v}$$
(E-Let)

$$\frac{}{\mathcal{E} \vdash \text{fun } x \to e \Downarrow (\mathcal{E}) [\text{fun } x \to e]}$$
 (E-Fun)

$$\frac{\mathcal{E} \vdash e_1 \Downarrow (\mathcal{E}_2) \left[\text{fun } x \to e_0 \right] \qquad \mathcal{E} \vdash e_2 \Downarrow v_2 \qquad \mathcal{E}_2, x = v_2 \vdash e_0 \Downarrow v}{\mathcal{E} \vdash e_1 \mid e_2 \mid \downarrow v} \tag{E-App}$$

$$\frac{\mathcal{E}, x = (\mathcal{E}) [\text{rec } x = \text{fun } y \to e_1] \vdash e_2 \Downarrow v}{\mathcal{E} \vdash \text{let rec } x = \text{fun } y \to e_1 \text{ in } e_2 \Downarrow v}$$
 (E-Letrec)

$$\frac{}{\mathcal{E} \vdash [] \Downarrow []}$$
 (E-Nil)

$$\frac{\mathcal{E} \vdash e_1 \Downarrow v_1 \qquad \mathcal{E} \vdash e_2 \Downarrow v_2}{\mathcal{E} \vdash e_1 :: e_2 \Downarrow v_1 :: v_2}$$
 (E-Cons)

$$\frac{\mathcal{E} \vdash e_0 \Downarrow v \qquad p \text{ matches } v \text{ when } (\mathcal{E}_1) \qquad (\mathcal{E}_2 = \mathcal{E}; \mathcal{E}_1) \qquad \mathcal{E}_2 \vdash e \Downarrow v'}{\mathcal{E} \vdash \text{match } e_0 \text{ with } p \rightarrow e \Downarrow v'} \text{ (E-MATCHM1)}$$

$$\frac{\mathcal{E} \vdash e_0 \Downarrow v \qquad p \text{ matches } v \text{ when } (\mathcal{E}_1) \qquad (\mathcal{E}_2 = \mathcal{E}; \mathcal{E}_1) \qquad \mathcal{E}_2 \vdash e \Downarrow v'}{\mathcal{E} \vdash \text{match } e_0 \text{ with } p \rightarrow e \mid c \Downarrow v'} \text{ (E-MATCHM2)}$$

$$\frac{\mathcal{E} \vdash e_0 \Downarrow v \qquad p \text{ doesn't match } v \qquad \mathcal{E} \vdash \text{match } e_0 \text{ with } c \Downarrow v'}{\mathcal{E} \vdash \text{match } e_0 \text{ with } p \rightarrow e \mid c \Downarrow v'} \qquad \text{(E-MATCHN)}$$

$$rac{(i_3=i_1+i_2)}{i_1 ext{ plus } i_2 ext{ is } i_3}$$
 (B-PLUS)

$$rac{(i_3=i_1-i_2)}{i_1 ext{ minus } i_2 ext{ is } i_3}$$
 (B-MINUS)

$$rac{(i_3=i_1*i_2)}{i_1 ext{ times } i_2 ext{ is } i_3}$$
 (B-TIMES)

$$rac{(b_3=(i_1 < i_2))}{i_1 ext{ less than } i_2 ext{ is } b_3}$$
 (B-LT)