the halting problem paradox

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halting function h(m) \rightarrow \{
  true : iff m halts
  false : otherwise
}
Let A: set of all describable programs
Let U: subset of A that contains programs undecidable by h,
       cannot be mapped via function h
       up = () -> h(up) && loop_forever()
        h(up) = ???
       np = () -> !h(np) && loop_forever()
        h(np) = true and h(np) = false ???
       \therefore h(m) is also undecidable
Let D: subset of A that contains programs decidable by h,
       can be mapped via function h
                                            Α
                                 U:
                                                       D
                           {h, up, np, ...}
Presume U does not exist
We're left with D that does exist, which now should be decidable by function hd(m),
the subfunction of h(m) where m exists in D
Does hd exist?
If hd does exist, then what about
       upd = () -> hd(upd) && loop_forever()
       \therefore hd(m) is part of U, and does not exist
If hd does not exist, then how is D actually the decidable subset?
       \therefore there no subset of programs that is completely decidable, even trivially
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Neither of these conclusions are acceptable, so what went wrong?