

# Logic

Logic is the \_\_\_\_\_ of math jic the rules of the \_\_\_\_\_

Def: A \_\_\_\_\_ is a \_\_\_\_\_ sentance which can be  
\_\_\_\_\_ determined to be \_\_\_\_\_ or \_\_\_\_\_

- ex) a) The integer 11 is divisible by \_\_\_\_\_  
(and non examples)  
b) The integer 11 is a  $\frac{\text{?}}{\text{?}}$  #  
 $\frac{\text{?}}{\text{?}}$

c) Is  $10^{10} \in \mathbb{Z}$

d) The integer  $10^{10}$  is bigggg

Def: Given a statement P, the negation of P is the  $\neg P$  (or  $\sim P$ ) (\text{neg})

~> It is characterized by the following —

P	
T	
F	

ex) a) P: The # II is odd

— = The # || is        odd  
= the # || is

b) P: The day of the week is Wednesday

: The day of the week is

Def: Given statements P and Q, the \_\_\_\_\_ of P and Q  
is the statement \_\_\_\_\_ denoted \_\_\_\_\_ (1 or )

• It is defined via the following    table

P	Q	
T	T	
T	F	
F	T	
F	F	

→ ie,        is  
       if        P or Q  
one of      is       

Ex) a) P: 5 is odd                   $\Rightarrow P \vee Q$  is         
Q: 10 is prime

b) Let P be any statement. Then \_\_\_\_\_ is always true.

P		
T		
F		

→ we call such statements  
(i.e. always true)

Def: Given statements P and Q the \_\_\_\_\_ of P and Q  
is the statement  
denoted \_\_\_\_\_ ( $\wedge$  land)

It is defined via the truth table

P	G
T	T

→ is only true  
when

T	F	
F	T	
F	F	

P and Q is  
\_\_\_\_\_

Def: Given statements P and Q the  
is the statement

\_\_\_\_\_ P \_\_\_\_\_ Q , denoted \_\_\_\_\_ (Implies)

• We call P the \_\_\_\_\_ and Q the \_\_\_\_\_

Its truth table is defined as

P	Q
T	T
T	F
F	T
F	F

Note that when the \_\_\_\_\_ P is false the \_\_\_\_\_ Q may be T or F, but \_\_\_\_\_ is always \_\_\_\_\_

Ex) P: It is raining

Q: I will stay at home

~ \_\_\_\_\_ : If it is raining,

Γ If it doesn't rain and I go out, did I lie?

— —  
→  $P \Rightarrow Q =$  —



Some terminology : \_\_\_\_\_ is read in diff ways

- if \_\_\_\_\_ then \_\_\_\_\_
- \_\_\_\_\_ implies \_\_\_\_\_
- \_\_\_\_\_ if \_\_\_\_\_
- \_\_\_\_\_ only if \_\_\_\_\_
- P is \_\_\_\_\_ for Q
- Q is \_\_\_\_\_ for P

all these mean



We will be considering the truth value of \_\_\_\_\_ that depend on a \_\_\_\_\_.

Def: An \_\_\_\_\_ sentence is a \_\_\_\_\_ sentence which contains \_\_\_\_\_, where each \_\_\_\_\_ can assume \_\_\_\_\_ value in a given \_\_\_\_\_, called the \_\_\_\_\_ of the \_\_\_\_\_, which becomes a statement if the \_\_\_\_\_ are replaced with specific values.

ex) For  $x \in \mathbb{R}$  consider the statement.

$$P(x) : |x| = 3$$

$$\underline{x=1} : P(1) \text{ _____}$$

$x = 2 : P(-2)$  \_\_\_\_\_

$x = 3 : P(-3)$  \_\_\_\_\_

•  $P(x) : x$  is odd for  $x \in \mathbb{Z}$

this is the \_\_\_\_\_

We can combine open sentences using \_\_\_\_\_ to form new open sentences

ex) Domain :  $x \in \mathbb{R}$      $P(x) : |x| = 3$

$Q(x) : x = -3$

$\sim \neg P(x) :$

$\neg P(1) :$  \_\_\_\_\_

$P(x) \wedge Q(x)$ :

$P(3) \wedge Q(3)$

$Q(x) \Rightarrow P(x)$ :

$P(x) \vee Q(x)$ :

Def: The implication \_\_\_\_\_ is called the \_\_\_\_\_ of  $P \Rightarrow Q$

- Given statements  $P$  and  $Q$  the \_\_\_\_\_ of  $P$  and  $Q$  is the statement

denoted \_\_\_\_\_  $\wedge$  \_\_\_\_\_  
\_\_\_\_\_ (iff)

• We read \_\_\_\_\_ as

• P \_\_\_\_\_ and only \_\_\_\_\_ Q

• P is \_\_\_\_\_ to Q

• P is \_\_\_\_\_ and \_\_\_\_\_ for Q  
 $(Q \Rightarrow P)$        $(P \Rightarrow Q)$

The truth table is given as

P	Q	$P \Rightarrow Q$	$Q \Rightarrow P$	$P \Rightarrow Q \wedge Q \Rightarrow P$
T	T			
T	F			
F	T			

F

F

→ \_\_\_\_\_ is true when \_\_\_\_\_ P and Q are either  
or \_\_\_\_\_

## Compound Statements

Def: A \_\_\_\_\_ statement is a statement consisting of  
one statement involving \_\_\_\_\_ one  
\_\_\_\_\_ connective ( $\neg, \wedge, \vee, \Rightarrow, \Leftarrow$ )

- Each statement in a \_\_\_\_\_ statement is called a  
\_\_\_\_\_ statement

ex)  $P \Leftrightarrow Q$  is a compound statement with \_\_\_\_\_

•

•

Def) 1) A compound statement is called a \_\_\_\_\_ if it is  
for all possible \_\_\_\_\_ of truth values  
for its \_\_\_\_\_ statements

2) A compound statement is called a \_\_\_\_\_ if it is  
for all possible \_\_\_\_\_ of truth values  
for its \_\_\_\_\_ statements

ex) i)  $P \vee (\neg P)$  is a u

ii)  $P \wedge (\neg P)$  is a u

$P$	$\neg P$	$P \wedge (\neg P)$

iii)  $(P \wedge (P \Rightarrow Q)) \Rightarrow Q$  is a u

$P$	$Q$	$P \Rightarrow Q$	$P \wedge (P \Rightarrow Q)$	$(P \wedge (P \Rightarrow Q)) \Rightarrow Q$
T	T	T	T	T
T	F	F	F	F
F	T	T	F	F
F	F	T	F	F

iv)  $((P \Rightarrow Q) \wedge (Q \Rightarrow R)) \Rightarrow (P \Rightarrow R)$  is a

P	Q	R	$P \Rightarrow Q$	$Q \Rightarrow R$	$(P \Rightarrow Q) \wedge (Q \Rightarrow R)$	$P \Rightarrow R$	$((P \Rightarrow Q) \wedge (Q \Rightarrow R)) \Rightarrow (P \Rightarrow R)$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	F
T	F	T	F	T	F	F	F
T	F	F	F	F	F	F	F
F	T	T	T	T	T	T	T
F	T	F	T	F	F	F	F
F	F	T	F	T	F	F	F
F	F	F	F	F	F	F	F

notice  
how  
there's  
more  
options  
now

# Logical Equivalences

Def.: Two compound statements  $R$  and  $S$  are \_\_\_\_\_ if they have the \_\_\_\_\_ for all possible combinations of \_\_\_\_\_ values for its component statements

~ We denote this as  $R \equiv S$  (equiv)

Thrm: Let  $P$  and  $Q$  be statements. Then

$$P \Rightarrow Q \quad ((\neg P) \vee Q))$$

Pf) We will build the truth table

P	Q	$\neg P$	$P \Rightarrow Q$	$(\neg P) \vee Q$
T	T			
T	F			
F	T			
F	F			

Thm: Let  $P, Q, R$  be statements. Let        and        be a        and       . Then

1) Identity laws :  $P \vee \underline{P} \wedge \underline{P} \equiv$

2) Domination laws :  $P \vee \underline{P} \wedge \underline{P} \equiv$

3) Double negation :  $\neg(\neg P) \equiv$

4) Commutative laws :  $P \vee Q \equiv$   
 $P \wedge Q \equiv$

5) Associative laws :  $(P \vee Q) \vee R \equiv$   
 $(P \wedge Q) \wedge R \equiv$

\* 6) Distributive Laws :  $P \vee (Q \wedge R) \equiv$   
 $P \wedge (Q \vee R) \equiv$

\* \* 7) De-Morgan's Laws :  $\neg(P \vee Q) \equiv$   
 $\neg(P \wedge Q) \equiv$

Pf) We will prove De-Morgan's Laws (the first part)

P	Q	$\neg P$	$\neg Q$	$P \vee Q$	$\neg(P \vee Q)$	$(\neg P) \wedge (\neg Q)$
T	T	F	F			
T	F	F	T			
F	T	T	F			
F	F	T	T			

Application: Prove that

$$\neg \neg (P \Rightarrow Q) \equiv P \wedge \neg Q$$

Pf)

2)  $((\neg Q) \Rightarrow (P \wedge \neg P)) \equiv Q$

Pf)

$$3) ((P \wedge Q) \Rightarrow (R \wedge \neg R)) \equiv P \Rightarrow Q$$