

# Econ 703 TA Note 9

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## 1 Notation

Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ . We define  $Df(x) := \nabla f(x)^\top$ . Let  $g : \mathbb{R}^n \rightarrow \mathbb{R}^k$ , where  $g(x) = \begin{bmatrix} g_1(x) \\ g_2(x) \\ \vdots \\ g_n(x) \end{bmatrix}$ . We define

$$Dg(x) := \begin{bmatrix} \nabla g_1(x)^\top \\ \nabla g_2(x)^\top \\ \vdots \\ \nabla g_n(x)^\top \end{bmatrix}.$$

## 2 Equality Constraints - Lagrange's Method

The Lagrange's method provides a way of *selecting* points that are *possibly* local optimal points when there are only equality constraints. Namely, it provides a *necessary* condition for a point to be a local optimal point.

### 2.1 First Order Condition

**Theorem 2.1 (Lagrange):** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  and  $g_i : \mathbb{R}^n \rightarrow \mathbb{R}$  for  $i = 1, \dots, k$  be  $C^1$  functions. Suppose  $x^*$  is a local maximal or minimal point of  $f$  on the set

$$\mathcal{D} = \{x : g_i(x) = 0 \text{ for all } 1 \leq i \leq k\}.$$

Write  $g = (g_1, \dots, g_k)^\top : \mathbb{R}^n \rightarrow \mathbb{R}^k$ . Suppose also that  $\text{rank}(Dg(x^*)) = k$ . Then there exists  $\lambda_1^*, \dots, \lambda_k^* \in \mathbb{R}$  such that

$$Df(x^*) + \sum_{i=1}^k \lambda_i^* Dg_i(x^*) = 0.$$

Two points should be emphasized.

- Lagrange's Theorem provides a necessary condition for local optimal points, not a sufficient condition. A point satisfying the condition may not be a local optimal point.

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\*This TA note is based on Prof. John Kennan's math camp lecture taught in 2025 at UW-Madison. All errors are mine.

- If at an optimal point  $x^*$ ,  $Dg(x^*)$  does not have rank  $k$ , then the conclusion in [Theorem 2.1](#) does not necessarily hold: there may not exist  $\lambda_1, \lambda_2, \dots, \lambda_n$  such that

$$Df(x^*) + \sum_{i=1}^k \lambda_i^* Dg_i(x^*) = 0.$$

## 2.2 Intuition

Let us consider the case when there is only one constraint:  $g(x_1, \dots, x_n) = I - p_1x_1 - \dots - p_nx_n$ , and we set the constraint  $g(x) = 0$ ,  $x \in \mathbb{R}^n$ . Charlie faces the problem:

$$\max U(x) \text{ s.t. } g(x) = 0.$$

(Suppose he is forced to spend up all of his money). At the optimal point  $x^* = (x_1^*, \dots, x_n^*)$ ,

$$\frac{MU_1}{p_1} = \frac{MU_2}{p_2} = \dots = \frac{MU_n}{p_n} \quad (1)$$

Why? Assume the contrary that at the optimal point  $x^*$ ,

$$\frac{MU_1}{p_1} > \frac{MU_2}{p_2}.$$

Suppose Charlie sells a small amount  $\epsilon$  of good 2. By doing so, he obtains  $\epsilon \cdot p_2$  units of money. This sale reduces his utility by approximately  $MU_2 \cdot \epsilon$ . He can then use the money to purchase

$$\frac{\epsilon \cdot p_2}{p_1}$$

units of good 1. The additional amount of good 1 yields an increase in utility of

$$MU_1 \cdot \frac{\epsilon \cdot p_2}{p_1}.$$

Comparing the two changes in utility, Charlie becomes strictly better off since

$$MU_1 \cdot \frac{p_2}{p_1} \cdot \epsilon > MU_2 \cdot \epsilon.$$

This contradicts the fact that  $x^*$  is the optimal point. Therefore, [Equation 1](#) must hold. Note that

$$\begin{aligned} DU(x) &= [MU_1(x) \quad MU_2(x) \quad \dots \quad MU_n(x)] \\ Dg(x) &= [-p_1 \quad -p_2 \quad \dots \quad -p_n] \end{aligned}$$

If we choose  $\lambda^* = MU_i/p_i$ , then

$$DU(x^*) + \lambda^* Dg(x^*) = 0.$$

This verifies the conclusion of [Theorem 2.1](#).

Suppose Charlie's income increases by  $\epsilon$ . He can use the additional money to buy  $\epsilon/p_i$  units of good  $i$ . This gives him an additional utility of

$$\frac{MU_1}{p_1} \epsilon = \lambda^* \epsilon.$$

Note that the additional utility he gains is the same regardless of which good he purchases.  $\lambda^*$  is Charlie's marginal utility as his income increases, and is called the **shadow price** of the constraint  $g(x)$ .

### 2.3 A Cookbook Procedure

Let an equality constraint problem be written in the following form:  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $g_i : \mathbb{R}^n \rightarrow \mathbb{R}$  for  $1 \leq i \leq k$ .

$$\max f(x) \text{ over } \mathcal{D} = \{x \in \mathbb{R}^n : g_i(x) = 0 \text{ for all } 1 \leq i \leq k\}.$$

To solve the problem, we follow the procedure:

1. Set up the Lagrangian:

$$L(x; \lambda) = f(x) + \sum_{i=1}^n \lambda_i g_i(x)$$

2. Solve  $(x, \lambda)$  that satisfies the following system of equations:

$$\begin{aligned} \frac{\partial L}{\partial x_i}(x; \lambda) &= 0 \text{ for all } i = 1, \dots, n; \\ \frac{\partial L}{\partial \lambda_j}(x; \lambda) &= 0 \text{ for all } j = 1, \dots, k. \end{aligned}$$

The solutions are called the critical points of the Lagrangian.

3. If we are maximizing  $f$ , then choose from the critical points that maximizes  $f$ . If we are minimizing  $f$ , then choose the one that minimizes  $f$ .

## 3 Inequality Constraints - Kuhn and Tucker's Method

Now we turn to a more general case: when some constraints are inequalities. Kuhn and Tucker provided a method to solve the problem:

$$\text{maximize } f \text{ over } \mathcal{D} = \{x \in \mathbb{R}^n : g_j(x) = 0 \text{ for all } 1 \leq j \leq m_1, h_i(x) \geq c_i \text{ for all } 1 \leq i \leq m_2\},$$

where  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ , each  $g_j : \mathbb{R}^n \rightarrow \mathbb{R}$  and  $h_i : \mathbb{R}^n \rightarrow \mathbb{R}$ .

### 3.1 First Order Condition

**Theorem 3.1 (Kuhn and Tucker):** Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $g_j : \mathbb{R}^n \rightarrow \mathbb{R}$  and  $h_i : \mathbb{R}^n \rightarrow \mathbb{R}$  be  $C^1$  functions,  $c_i \in \mathbb{R}$ , where  $1 \leq j \leq m_1$  and  $1 \leq i \leq m_2$ . Suppose  $x^*$  is a local maximal point of  $f$  on

$$\mathcal{D} = \{x \in \mathbb{R}^n : g_j(x) = 0 \text{ for all } 1 \leq j \leq m_1, h_i(x) \geq c_i \text{ for all } 1 \leq i \leq m_2\}.$$

Let  $E \subset \{1, \dots, l\}$  be the set of inequality constraints that are binding and write  $h_E = (h_i)_{i \in E}$ . Suppose  $D(g, h_E)$  has rank  $m_1 + |E|$ . Then there exists  $(\lambda_1^*, \dots, \lambda_{m_1}^*)$  and  $(\mu_1^*, \dots, \mu_{m_2}^*)$  such that the following conditions are met:

$$\mu_i^* \geq 0 \text{ and } \mu_i^*(h_i(x^*) - c_i) = 0 \text{ for } i = 1, \dots, m_2 \quad (\text{KT-1})$$

$$Df(x^*) + \sum_{j=1}^{m_1} \lambda_j^* Dg_j(x^*) + \sum_{i=1}^{m_2} \mu_i^* Dh_i(x^*) = 0. \quad (\text{KT 2})$$

When it is a minimization problem, KT-1 is replaced by

$$\mu_i^* \leq 0 \text{ and } \mu_i^*(h_i(x^*) - c_i) = 0 \text{ for } i = 1, \dots, m_2.$$

KT-1 is called the *complementary slackness* condition. It is called complementary because at least one of the following must be true: the constraint is binding or the multiplier is 0.

Some points worth noticing:

1. Similar to Lagrange's Theorem, Kuhn and Tucker's Theorem provides a *necessary* condition for local optimal points, not a *sufficient* condition.
2. For local optimal points that does not satisfy  $\text{rank}\{D(g, h_E)(x^*)\} = m_1 + |E|$ , the conclusion of the theorem may not hold.
3. It is recommended that one write all inequalities in terms of  $h_i(x) \geq c_i$  so that one doesn't have to think repeatedly about the sign of  $\lambda_i$  for each  $i$  when solving the question. In that way, for maximization problems,  $\lambda_i \geq 0$  for all  $i$  and, for minimization problems,  $\lambda_i \leq 0$  for all  $i$ .

### 3.2 Intuition

Charlie is maximizing his utility  $u(x_1, \dots, x_n)$  with the following constraints:

$$\begin{aligned} g(x) &= I - p_1x_1 - \dots - p_nx_n = 0, \\ h_i(x) &= x_i \geq c_i, \quad \forall 1 \leq i \leq n. \end{aligned}$$

Let  $x^* = (x_1^*, \dots, x_n^*)$ . There are two possible cases:

- (i)  $x_i^* > c_i$  for all  $1 \leq i \leq n$ : Then as we discussed in [subsection 2.2](#), at  $x^*$ , it must be

$$\frac{MU_1}{p_1} = \frac{MU_2}{p_2} = \dots = \frac{MU_n}{p_n}.$$

Choose  $\lambda^*$  to be the common ratio, and  $\mu_i^* = 0$  for all  $1 \leq i \leq n$ . Then we would have

$$Df(x^*) + \lambda^* Dg(x^*) + \sum_{i=1}^n \mu_i^* Dh_i(x) = 0.$$

- (ii) There is some  $x_i = c_i$ : Consider the case  $x_i^* > c_i$  for all  $1 \leq i \leq n-1$  while  $x_n^* = c_i$ . Then our argument in [subsection 2.2](#) no longer works: it may be the case that

$$\frac{MU_1}{p_1} = \dots = \frac{MU_{n-1}}{p_{n-1}} > \frac{MU_n}{p_n}.$$

Charlie would like to exchange  $x_n$  for other goods, but the requirement that he must keep at least  $c_n$  units of good  $n$  makes this infeasible. Write

$$\lambda^* = \frac{MU_1}{p_1} = \dots = \frac{MU_{n-1}}{p_{n-1}}, \quad \mu_n^* = \lambda^* p_n - MU_n > 0.$$

Note that

$$Dh_n(x^*) = [0, \dots, 0, 1].$$

Then we would have

$$Df(x^*) + \lambda^* Dg(x^*) + \sum_{i=1}^n \mu_i^* Dh_i(x) = 0.$$

This verifies the conclusion of [Theorem 3.1](#).

Suppose Charlie were allowed to hold only  $c_n - \epsilon$  units of good  $n$ . He could then sell  $\epsilon$  units of good  $n$  for  $\epsilon \cdot p_n$  units of money. With this money, he could purchase  $(\epsilon \cdot p_n)/p_1$  units of good 1, yielding an additional utility of  $(\epsilon \cdot p_n \cdot MU_1)/p_1 = \lambda^* p_n \epsilon$  of utility. At the same time, selling  $\epsilon$  units of good  $n$  reduces his utility by  $\epsilon \cdot MU_n$ . His net gain of utility is

$$\lambda^* p_n \epsilon - MU_n \epsilon = \mu_n^* \epsilon.$$

Hence,  $\mu_n^*$  is the marginal utility of *loosening* the constraint  $h_n(x) \geq c_n$ , and is called the **shadow price** of the constraint  $h_n(x)$ . (For a constraint  $h_i(x) \geq c_i$ , loosening it means to decrease  $c_i$ .)

This also explained why  $\mu_i \leq 0$  in minimization problems: loosening the constraint allows the minimum value to decrease.

### 3.3 A Cookbook Procedure

Let an inequality constraint problem be written in the following form:  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ ,  $g_j : \mathbb{R}^n \rightarrow \mathbb{R}$  and  $h_i : \mathbb{R}^n \rightarrow \mathbb{R}$  be  $C^1$  functions,  $c_i \in \mathbb{R}$ , where  $1 \leq j \leq m_1$  and  $1 \leq i \leq m_2$ .

$$\max f(x) \text{ over } \mathcal{D} = \{x \in \mathbb{R}^n : g_j(x) = 0 \text{ for all } 1 \leq j \leq m_1, h_i(x) \geq c_i \text{ for all } 1 \leq i \leq m_2\}.$$

We follow the procedure to find the solution.

1. Set up the Lagrangian:

$$L(x; \lambda, \mu) = f(x) + \sum_{j=1}^{m_1} \lambda_j g_j(x) + \sum_{i=1}^{m_2} \mu_i (h_i(x) - c_i).$$

2. Solve the following system of equations and inequalities:

$$\frac{\partial L}{\partial \lambda_j}(x; \lambda, \mu) = 0 \text{ for all } j = 1, \dots, m_1 \quad (\text{the equality constraints})$$

$$\frac{\partial L}{\partial \mu_i}(x; \lambda, \mu) \geq 0 \text{ for all } i = 1, \dots, m_2 \quad (\text{the inequality constraints})$$

$$\mu_i \geq 0, \quad \mu_i \frac{\partial L}{\partial \mu_i}(x; \lambda, \mu) = 0 \text{ for all } i = 1, \dots, m_2 \quad (\text{KT-1})$$

$$\frac{\partial L}{\partial x_i}(x; \lambda, \mu) = 0 \text{ for all } i = 1, \dots, n. \quad (\text{KT-2})$$

Points  $(x, \lambda)$  that solve the system are called critical points.

3. Choose the critical point that maximizes  $f$ .

For minimization problems, change (KT-1) to

$$\mu_i \leq 0, \quad \mu_i \frac{\partial L}{\partial \mu_i}(x; \lambda, \mu) = 0 \text{ for all } i = 1, \dots, m_2. \quad (\text{KT-1})$$

And of course, choose the critical point that minimizes  $f$  in the final step.