



PostgreSQL Query Optimization

Step by step techniques



1. What is a slow query?
2. How to chose queries to optimize?
3. What is a query plan?
4. Optimization tools
5. Optimization examples



QUERY PLAN

```
-----  
Limit  (cost=12993.17..12993.17 rows=1 width=20) (actual time=606.385..606.385 rows=1 loops=1)  
...  
Planning time: 1.236 ms  
Execution time: 607.057 ms
```

Does this query perform well enough for your system?

4



- What is your baseline?

- What is your baseline?
- 607.057 ms can be extremely fast for OLAP

- What is your baseline?
- 607.057 ms can be extremely fast for OLAP
- But $607.057 \text{ ms} * 10000$ parallel queries on OLTP?

- What is your baseline?
- 607.057 ms can be extremely fast for OLAP
- But $607.057 \text{ ms} * 10000$ parallel queries on OLTP?
- 607.057 ms on 10 y.o. SATA disks vs modern SSD

- Often it is useless to optimize all queries

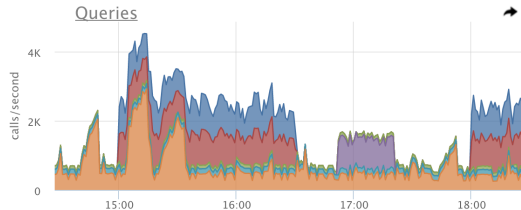
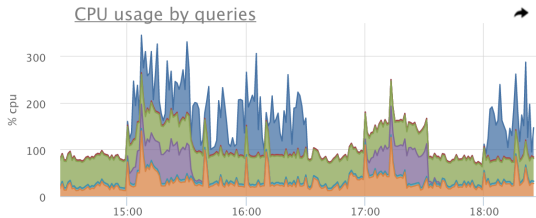
- Often it is useless to optimize all queries
- *log_min_duration_statement* = 100ms
Everything that's in the logs is due for review

- Often it is useless to optimize all queries
- *log_min_duration_statement* = 100ms
Everything that's in the logs is due for review
- *pg_stat_statements*
Lot's of useful stuff inside

- Often it is useless to optimize all queries
- *log_min_duration_statement* = 100ms
Everything that's in the logs is due for review
- *pg_stat_statements*
Lot's of useful stuff inside
- Monitoring system of choice
Hopefully it has query info accumulated and ranged

How to find the queries to optimize?

6



```
SELECT sum(total_time) AS total_time,  
       sum(blk_read_time + blk_write_time) AS io_time,  
       sum(total_time - blk_read_time - blk_write_time) AS cpu_time,  
       sum(calls) AS ncalls,  
       sum(rows) AS total_rows  
FROM pg_stat_statements  
WHERE dbid IN (SELECT oid FROM pg_database WHERE datname=current_database())
```

```
WITH ttl AS (  
    SELECT sum(total_time) AS total_time, sum(blk_read_time + blk_write_time) AS io_time,  
           sum(total_time - blk_read_time - blk_write_time) AS cpu_time,  
           sum(calls) AS ncalls, sum(rows) AS total_rows  
    FROM pg_stat_statements WHERE dbid IN (  
        SELECT oid FROM pg_database WHERE datname=current_database())  
    )  
SELECT *, (pss.total_time - pss.blk_read_time - pss.blk_write_time) / ttl.cpu_time * 100 cpu_pct  
    FROM pg_stat_statements pss, ttl  
WHERE (pss.total_time - pss.blk_read_time - pss.blk_write_time) / ttl.cpu_time >= 0.05  
ORDER BY pss.total_time - pss.blk_read_time - pss.blk_write_time DESC LIMIT 1;
```

- Lot's of metrics are possible to extract
- Requires time to come up with a good usable report
- DataEgret maintains it's report in the public domain¹

¹https://github.com/dataegret/pg-utils/blob/master/sql/global_reports/query_stat_total.sql



- Report operates with *total_time*, *io_time* and *cpu_time*, that is a difference of the first two
- Report also normalizes queries and calculates *md5* hash for faster processing
- Main part of the report includes only those entries, that (any of the conditions qualifies):
 1. used more than 1% of total CPU or total IO time
 2. returned more than 2% of all rows
 3. had been called more than 2% of all query executions
- all other queries are combined into the *other* group
- report orders queries by total time spent, longest at the top



```
total time:      19:59:57 (IO: 16.43%)
total queries:   200,609,344 (unique: 2,342)
report for all databases, version 0.9.5 @ PostgreSQL 9.6.3
tracking top 10000 queries, utilities off, logging 100ms+ queries
=====
pos:1    total time: 05:38:45 (28.2%, CPU: 30.9%, IO: 14.5%)    calls: 84,592,220 (42.17%)    avg_time: 0.24ms (IO: 8.3%)
user: all      db: all      rows: 198,391,036 (24.34%)    query:
other
=====
pos:2    total time: 04:59:15 (24.9%, CPU: 24.0%, IO: 29.9%)    calls: 5,610 (0.00%)    avg_time: 3200.60ms (IO: 19.7%)
user: postgres db: ----- rows: 5,608,185 (0.69%)    query:

WITH _deleted AS (DELETE FROM foos_2rm WHERE id IN (SELECT id FROM foos_2rm ORDER BY id LIMIT ?) RETURNING id)
DELETE FROM foos WHERE id IN (SELECT id FROM _deleted);
=====
pos:3    total time: 00:45:06 (3.8%, CPU: 2.3%, IO: 11.1%)    calls: 853,864 (0.43%)    avg_time: 3.17ms (IO: 48.6%)
user: -----_background      db: ----- rows: 164,706 (0.02%)    query:
SELECT  "foo_stats_master".* FROM "foo_stats_master" WHERE (foo_stats_master.created_at >= ?) AND (foo_stats_master.created_at < ?)
AND "foo_stats_master"."action" IN (?, ?, ?, ?) AND ("foo_stats_master"."foo_board_id" IS NOT NULL)
AND "foo_stats_master"."user_ip_inet" = ? AND "foo_stats_master"."employer_id" = ?
ORDER BY "foo_stats_master"."created_at" DESC LIMIT ?
```

So, we identified some queries to optimize

12



So, we identified some queries to optimize

12

What comes next?



- Any query can be prepended with *EXPLAIN* to see its **execution plan**
- `EXPLAIN SELECT * FROM pg_database;`
QUERY PLAN

Seq Scan on pg_database (cost=0.00..0.16 rows=6 width=271)
(1 row)

- Query goes through several stages in it's lifecycle
 - 1. Connection
 - 2. Parser
 - 3. Rewrite system
 - 4. Planner / Optimizer
 - 5. Executor ↔ [Workers]
 - 6. Send results
- Planner prepares a **plan** for executor



- It is a tree
- Nodes and operations on them
- Planner uses statistics to chose the optimal plan



```
EXPLAIN SELECT * FROM pg_database;  
               QUERY PLAN
```

Seq Scan on pg_database (cost=0.00..0.16 rows=6 width=271)
(1 row)

Seq Scan	type of node operation
on pg_database	object of node operation
cost=0.00..0.16	cost of the node
rows=6	estimated rows
width=271	average width of a row

- **Seq Scan** — sequential scan of whole relation
- **Parallel Seq Scan** — parallel sequential scan of whole relation
- **Index Scan** — targeted random IO (read index + read table)
- **Index Only Scan** — read only from index²
- **Bitmap Index Scan** — prepare a map of rows to read from relation, possibly combining maps from several indexes
- **Bitmap Heap Scan** — use map from Bitmap Index Scan and read rows from relation, *always* follows Bitmap Index Scan
- **CTE Scan** - read from Common Table Expression (*WITH Block*)
- **Function Scan** - read results, returned by a function

²https://wiki.postgresql.org/wiki/Index-only_scans



- A cost of fetching 8K block sequentially
- Cost is a relative value: a cost of 10 is 10× greater than a cost of 1

```
explain select * from posts order by id limit 5;
```

```
QUERY PLAN
```

```
-----  
Limit  (cost=0.29..0.46 rows=5 width=28)
```

```
  ->  Index Scan using posts_pkey on posts  (cost=0.29..347.29 rows=10000 width=28)  
(2 rows)
```

- A cost of fetching 8K block sequentially
- Cost is a relative value: a cost of 10 is 10× greater than a cost of 1

```
explain select * from posts order by id limit 5;
```

```
QUERY PLAN
```

```
-----  
Limit  (cost=0.29..0.46 rows=5 width=28)
```

```
  ->  Index Scan using posts_pkey on posts  (cost=0.29..347.29 rows=10000 width=28)  
(2 rows)
```

- $0.29 + (347.29 - 0.29) * 5 / 10000 = 0.4635$

- Rows × width of a root node gives a clue of a result size in bytes
- Even if the query is fast, lots of it's calls can cause a huge traffic between database and an application
- That's why *SELECT** is not a good idea



- **join** – joins data from two nodes using appropriate join method
- **sort** – various methods of sorting
- **limit** – cuts the dataset off
- **aggregate** – performs aggregation
- **hash aggregate** – groups data
- **unique** – removes duplicates from sorted datasets
- **gather** – gather data from different workers



```
EXPLAIN [ ANALYZE ] [ VERBOSE ] statement
```

```
EXPLAIN [ ( option [, ...] ) ] statement
```

- *ANALYZE* **executes** statement and shows execution details
- *VERBOSE* verbose output
- *COSTS* show plan costs
- *BUFFERS* show information about buffers operated by the query
- *TIMING* show time spent
- *SUMMARY* show totals at the end of output
- *FORMATTEXT|XML|JSON|YAML* output in selected format



```
EXPLAIN (analyze) SELECT relname,relpages,reltuples FROM pg_class WHERE reltuples>10000;
                                QUERY PLAN
```

```
-----
Seq Scan on pg_class (cost=0.00..5.55 rows=6 width=72) (actual time=0.069..0.073 rows=6 loops=1)
  Filter: (reltuples > '10000'::double precision)
    Rows Removed by Filter: 334
Planning time: 0.102 ms
Execution time: 0.087 ms
(5 rows)
```

actual time=0.069..0.073	startup and total time of node execution
rows=6	actual rows
loops=1	number of times node had been executed
Rows Removed by Filter: 334	node processing details

```
EXPLAIN (analyze, buffers) SELECT r.relname, a.attname FROM pg_class r JOIN pg_attribute a ON a.attrelid=r.oid
WHERE a.attnum>0 AND NOT attisdropped;
```

QUERY PLAN

```
-----
Hash Join  (cost=8.95..66.58 rows=1770 width=128) (actual time=0.215..2.246 rows=2039 loops=1)
  Hash Cond: (a.attrelid = r.oid)
  Buffers: shared hit=59 read=2
  I/O Timings: read=0.270
    -> Seq Scan on pg_attribute a  (cost=0.00..33.29 rows=1770 width=68) (actual time=0.009..1.148 rows=2039 loops=1)
        Filter: ((NOT attisdropped) AND (attnum > 0))
        Rows Removed by Filter: 587
        Buffers: shared hit=46 read=2
        I/O Timings: read=0.270
    -> Hash  (cost=4.70..4.70 rows=340 width=68) (actual time=0.198..0.198 rows=340 loops=1)
        Buckets: 1024  Batches: 1  Memory Usage: 42kB
        Buffers: shared hit=13
        -> Seq Scan on pg_class r  (cost=0.00..4.70 rows=340 width=68) (actual time=0.002..0.095 rows=340 loops=1)
            Buffers: shared hit=13
Planning time: 0.202 ms
Execution time: 2.554 ms
(16 rows)
```



- We know what we want in terms of performance
- We know which query to optimize
- We have all the tools (*EXPLAIN ANALYZE*)

- We know what we want in terms of performance
- We know which query to optimize
- We have all the tools (*EXPLAIN ANALYZE*)
- Now we only need to minimize the time executor spends on each node



- We know what we want in terms of performance
- We know which query to optimize
- We have all the tools (*EXPLAIN ANALYZE*)
- Now we only need to minimize the time executor spends on each node
- Or actually try to figure out what the query should do:
Never optimize a SQL-query itself, try to optimize the operation it does

```
EXPLAIN ANALYZE SELECT * FROM test WHERE val=10;
```

QUERY PLAN

```
-----  
Seq Scan on test  (cost=0.00..160.59 rows=37 width=16) (actual time=0.036..1.640 rows=18 loops=1)  
  Filter: (val = 10)  
    Rows Removed by Filter: 8900  
Planning time: 0.163 ms  
Execution time: 2.037 ms  
(5 rows)
```

```
=> create index CONCURRENTLY test_val_idx on test using btree (val);  
CREATE INDEX
```

```
=> EXPLAIN ANALYZE SELECT * FROM test WHERE val=10;
```

QUERY PLAN

```
-----  
Bitmap Heap Scan on test  (cost=4.42..41.22 rows=18 width=16) (actual time=0.041..0.062 rows=18 loops=1)  
  Recheck Cond: (val = 10)  
  Heap Blocks: exact=12  
    -> Bitmap Index Scan on test_val_idx  (cost=0.00..4.42 rows=18 width=0)  
        (actual time=0.033..0.033 rows=18 loops=1)  
          Index Cond: (val = 10)  
Planning time: 1.136 ms  
Execution time: 0.240 ms  
(7 rows)
```



```
explain analyze select distinct f1 from test_ndistinct ;
```

```
QUERY PLAN
```

```
-----  
Unique  (cost=1571431.43..1621431.49 rows=100000 width=4)
```

```
  (actual time=4791.872..7551.150 rows=90020 loops=1)
```

```
    -> Sort  (cost=1571431.43..1596431.46 rows=10000012 width=4)
```

```
      (actual time=4791.870..6893.413 rows=10000000 loops=1)
```

```
      Sort Key: f1
```

```
      Sort Method: external merge  Disk: 101648kB
```

```
      -> Seq Scan on test_ndistinct  (cost=0.00..135314.12 rows=10000012 width=4)
```

```
        (actual time=0.041..938.093 rows=10000000 loops=1)
```

```
Planning time: 0.099 ms
```

```
Execution time: 7714.701 ms
```

```
set work_mem = '8MB';
```

```
SET
```

```
explain analyze select distinct f1 from test_ndistinct ;
```

QUERY PLAN

```
-----  
HashAggregate  (cost=160314.15..161314.15 rows=100000 width=4)  
               (actual time=2371.902..2391.415 rows=90020 loops=1)
```

```
  Group Key: f1
```

```
    -> Seq Scan on test_ndistinct  (cost=0.00..135314.12 rows=10000012 width=4)  
                                         (actual time=0.093..871.619 rows=10000000 loops=1)
```

```
Planning time: 0.048 ms
```

```
Execution time: 2396.186 ms
```

1. `SELECT * FROM test WHERE id<10000`
1.2ms
2. `SELECT * FROM test WHERE id<10000 AND val IN (a list from 1 to 10)`
2.1ms
3. `SELECT * FROM test WHERE id<10000 AND val IN (a list from 1 to 100)`
6ms
4. `SELECT * FROM test WHERE id<10000 AND val IN (a list from 1 to 1000)`
38ms
5. `SELECT * FROM test WHERE id<10000 AND val IN (a list from 1 to 10000)`
380ms


```
explain analyze select * from test where id<10000 and val IN (1,...,100);  
QUERY PLAN
```

```
-----  
Index Scan using test_pkey on test  (cost=0.43..1666.85 rows=10  
width=140) (actual time=0.448..5.602 rows=16 loops=1)  
Index Cond: (id < 10000)  
Filter: (val = ANY ('1,...,100'::integer[]))  
Rows Removed by Filter: 9984
```

```
explain select count(*) from test JOIN (VALUES (1),..., (10)) AS  
v(val) USING (val) where id<10000;  
QUERY PLAN
```

```
-----  
Aggregate  (cost=497.65..497.66 rows=1 width=0)  
->  
Hash Join  
  (cost=0.69..497.65 rows=1 width=0)  
Hash Cond: (test.val = "*VALUES*".column1)  
->  Index Scan using test_pkey on test  (cost=0.43..461.22  
rows=9645 width=4)  
Index Cond: (id < 10000)  
->  Hash  (cost=0.12..0.12 rows=10 width=4)  
->  Values Scan on "*VALUES*"  (cost=0.00..0.12 rows=10  
width=4)
```

```
1. SELECT * FROM test WHERE id<10000  
1.2ms  
2. JOIN (VALUES (1),...,(10))  
1.6ms (was 2.1ms)  
3. JOIN (VALUES (1),...,(100))  
2ms (was 6ms)  
4. JOIN (VALUES (1),...,(1000))  
3.9ms (was 38ms)  
5. JOIN (VALUES (1),...,(10000))  
10ms (was 380ms)
```



```
EXPLAIN (analyze) SELECT DISTINCT author_id FROM blog_post;  
                                QUERY PLAN
```

```
-----  
Unique  (cost=0.42..32912.78 rows=1001 width=4) (actual time=0.019..347.327 rows=1001 loops=1)  
  ->  Index Only Scan using u_bp_author_ctime on blog_post  (cost=0.42..30412.72 rows=1000020 width=4)  
                                     (actual time=0.018..268.112 rows=1000000 loops=1)  
  
        Heap Fetches: 0  
Planning time: 0.068 ms  
Execution time: 347.495 ms  
(5 rows)
```

```
EXPLAIN (analyze) WITH RECURSIVE t AS (  
  -- start from least author_id -- anchor  
  (SELECT author_id AS _author_id FROM blog_post ORDER BY author_id LIMIT 1)  
  UNION ALL  
  -- find the next author_id > "current" author_id -- iterator  
  SELECT author_id AS _author_id  
    FROM t, LATERAL (SELECT author_id FROM blog_post WHERE author_id>t._author_id  
                     ORDER BY author_id LIMIT 1) AS a_id  
)  
-- return found values  
SELECT _author_id FROM t;
```

QUERY PLAN

```
-----  
CTE Scan on t (cost=52.27..54.29 rows=101 width=4) (actual time=0.017..11.176 rows=1001 loops=1)  
  CTE t  
    -> Recursive Union (cost=0.42..52.27 rows=101 width=4) (actual time=0.016..10.154 rows=1001 loops=1)  
      -> Limit (cost=0.42..0.46 rows=1 width=4) (actual time=0.015..0.015 rows=1 loops=1)  
        -> Index Only Scan using u_bp_author_ctime on blog_post (cost=0.42..30412.72 rows=1000020 width=4)  
              (actual time=0.014..0.014 rows=1 loops=1)  
              Heap Fetches: 0  
      -> Nested Loop (cost=0.42..4.98 rows=10 width=4) (actual time=0.009..0.010 rows=1 loops=1001)  
        -> WorkTable Scan on t t_1 (cost=0.00..0.20 rows=10 width=4) (actual time=0.000..0.000 rows=1 loops=1001)  
        -> Limit (cost=0.42..0.46 rows=1 width=4) (actual time=0.009..0.009 rows=1 loops=1001)  
          -> Index Only Scan using u_bp_author_ctime on blog_post blog_post_1 (cost=0.42..10973.87 rows=333340 width=4)  
                (actual time=0.009..0.009 rows=1 loops=1001)  
                Index Cond: (author_id > t_1._author_id)  
                Heap Fetches: 0  
Planning time: 0.143 ms  
Execution time: 11.301 ms  
(14 rows)
```

- NOT IN (query) instead of EXISTS
- JOIN instead IN/EXISTS
- unordered LIMIT
- ORDER BY random()



- NOT IN (query) instead of EXISTS
- JOIN instead IN/EXISTS
- unordered LIMIT
- ORDER BY random()
- **Avoid them!**



- Do not optimize all the queries - start with most critical for your **production** system
- Find your baseline
- Do not tune the query, try to figure out how to do what it does more effectively!



ik@dataegret.com

Please, leave your feedback at <https://2018.nordicpgday.org/feedback>