



# SMART CONTRACT AUDIT REPORT

for

## Geist Protocol



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PeckShield  
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# 1 | Introduction

Given the opportunity to review the design document and related smart contract source code of the Geist protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

## 1.1 About Geist

Geist is a decentralized non-custodial liquidity markets protocol that is developed on top of one of the largest DeFi protocols, i.e., AAVE. The protocol allows users to participate as depositors or borrowers. Depositors provide liquidity to the market to earn a passive income, while borrowers are able to borrow in an over-collateralized (perpetually) or under-collateralized (one-block liquidity) fashion. The protocol extends the original version with new features for staking-based incentivization and fee distribution.

The basic information of Geist is as follows:

Table 1.1: Basic Information of Geist

Item	Description
Name	Geist
Website	<a href="https://geist.finance/">https://geist.finance/</a>
Type	Ethereum Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	November 23, 2021

In the following, we show the Git repository of reviewed files and the commit hash value used in

this audit.

- <https://github.com/geist-finance/geist-protocol.git> (b6b13bd)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

- <https://github.com/geist-finance/geist-protocol.git> (1cfcb20)

## 1.2 About PeckShield

PeckShield Inc. [13] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email ([contact@peckshield.com](mailto:contact@peckshield.com)).

Table 1.2: Vulnerability Severity Classification

Impact	High	Critical	High	Medium
	Medium	High	Medium	Low
	Low	Medium	Low	Low
		High	Medium	Low
		Likelihood		

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on the OWASP Risk Rating Methodology [12]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Table 1.3: The Full Audit Checklist

Category	Checklist Items
Basic Coding Bugs	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
Advanced DeFi Scrutiny	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
Additional Recommendations	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a checklist of items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [11], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

## 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
<b>Configuration</b>	Weaknesses in this category are typically introduced during the configuration of the software.
<b>Data Processing Issues</b>	Weaknesses in this category are typically found in functionality that processes data.
<b>Numeric Errors</b>	Weaknesses in this category are related to improper calculation or conversion of numbers.
<b>Security Features</b>	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
<b>Time and State</b>	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
<b>Error Conditions, Return Values, Status Codes</b>	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
<b>Resource Management</b>	Weaknesses in this category are related to improper management of system resources.
<b>Behavioral Issues</b>	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
<b>Business Logic</b>	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
<b>Initialization and Cleanup</b>	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
<b>Arguments and Parameters</b>	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
<b>Expression Issues</b>	Weaknesses in this category are related to incorrectly written expressions within code.
<b>Coding Practices</b>	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.







comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.



## 2 | Findings

### 2.1 Summary

Here is a summary of our findings after analyzing the implementation of the Geist protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	0	
Medium	2	
Low	4	
Informational	1	
Undetermined	1	
Total	8	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in [Section 3](#).

## 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 2 medium-severity vulnerabilities, 4 low-severity vulnerabilities, 1 informational recommendation, and 1 undetermined issue.

Table 2.1: Key Geist Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Low	Proper answer Type in ChainlinkResponse	Coding Practice	Resolved
PVE-002	Low	Improved Oracle Status in PriceFeed::__fetchPrice()	Business Logic	Resolved
PVE-003	Low	Inconsistent Use Of IncentivesController::handleAction()	Business Logic	Resolved
PVE-004	Informational	ERC20 Compliance Of GeistToken	Coding Practice	Resolved
PVE-005	Medium	Fork-Resistant Domain Separator in AToken	Business Logic	Resolved
PVE-006	Low	Suggested Adherence Of Checks-Effects-Interactions Pattern	Time and State	Resolved
PVE-007	Undetermined	Staking Incompatibility With Deflationary Tokens	Business Logic	Confirmed
PVE-008	Medium	Trust Issue of Admin Keys	Security Features	Resolved

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

## 3 | Detailed Results

### 3.1 Proper answer Type in ChainlinkReponse

- ID: PVE-001
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: PriceFeed
- Category: Coding Practices [8]
- CWE subcategory: CWE-1099 [1]

#### Description

The Geist protocol has strengthened the oracle reliability and security with the so-called dual price oracle, i.e., it makes use of not only the Chainlink's live ETH:USD aggregator reference contract, but also bandOracle with the connection to the BandMaster contract. While reviewing the integration with the Chainlink-based oracle, we notice the internal ChainlinkResponse data structure needs to be improved.

In the following, we show the definition of the ChainlinkResponse data structure. Notice the answer member field is defined as `uint256`, not `int256`. If we examine the criteria used in `_badChainlinkResponse()` to determine whether the Chainlink response is considered bad, it explicitly considers non-positive answer. In other words, there is a type inconsistency in the answer member field.

```

49     struct ChainlinkResponse {
50         uint80 roundId;
51         uint256 answer;
52         uint256 timestamp;
53         bool success;
54         uint8 decimals;
55     }

```

Listing 3.1: The PriceFeed::ChainlinkResponse Structure

```

337     function _badChainlinkResponse(ChainlinkResponse memory _response) internal view
338         returns (bool) {
339         // Check for response call reverted

```

```

339     if (!_response.success) {return true;}
340     // Check for an invalid roundId that is 0
341     if (_response.roundId == 0) {return true;}
342     // Check for an invalid timeStamp that is 0, or in the future
343     if (_response.timeStamp == 0 || _response.timeStamp > block.timeStamp) {return
        true;}
344     // Check for non-positive price
345     if (_response.answer <= 0) {return true;}
346
347     return false;
348 }

```

Listing 3.2: PriceFeed::\_badChainlinkResponse()

**Recommendation** Resolve the type inconsistency in the `answer` member field by defining it as an `int256`, not `uint256`.

**Status** The issue has been fixed by this commit: 8702175.

## 3.2 Improved Oracle Status in PriceFeed::\_fetchPrice()

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: PriceFeed
- Category: Business Logic [9]
- CWE subcategory: CWE-837 [5]

### Description

As mentioned earlier, the Geist protocol is unique in supporting the dual price oracle, which necessitates the examination of current oracle states. In total, there are five different oracle states, i.e., `chainlinkWorking`, `usingBandChainlinkUntrusted`, `bothOraclesUntrusted`, `usingBandChainlinkFrozen`, and `usingChainlinkBandUntrusted`. While examining possible transition from the fourth state, we notice the transition logic can be revisited.

To elaborate, we show below the code snippet from the `_fetchPrice()` function. This function is designed to fetch the current price and adjust the current oracle state accordingly. Starting from the fourth state `usingBandChainlinkFrozen`, the current logic considers the conditions of `!_chainlinkIsFrozen(chainlinkResponse)` (line 263) and `_bandIsBroken(bandResponse)` (line 282) to still yield `usingBandChainlinkFrozen` as the next state, which in fact can be better adjusted as `usingChainlinkBandFrozen`.

```

248     if (status == Status.usingBandChainlinkFrozen) {
249         if (!_chainlinkIsBroken(chainlinkResponse, prevChainlinkResponse)) {
250             // If both Oracles are broken, return last good price

```

```

251         if (_bandIsBroken(bandResponse)) {
252             return (Status.bothOraclesUntrusted, lastGoodPrice);
253         }
254
255         // If Chainlink is broken, remember it and switch to using Band
256
257         if (_bandIsFrozen(bandResponse)) {return (Status.
                usingBandChainlinkUntrusted, lastGoodPrice);}
258
259         // If Band is working, return Band current price
260         return (Status.usingBandChainlinkUntrusted, bandResponse.value);
261     }
262
263     if (_chainlinkIsFrozen(chainlinkResponse)) {
264         // if Chainlink is frozen and Band is broken, remember Band broke, and
                return last good price
265         if (_bandIsBroken(bandResponse)) {
266             return (Status.usingChainlinkBandUntrusted, lastGoodPrice);
267         }
268
269         // If both are frozen, just use lastGoodPrice
270         if (_bandIsFrozen(bandResponse)) {return (Status.
                usingBandChainlinkFrozen, lastGoodPrice);}
271
272         // if Chainlink is frozen and Band is working, keep using Band (no
                status change)
273         return (Status.usingBandChainlinkFrozen, bandResponse.value);
274     }
275
276     // if Chainlink is live and Band is broken, remember Band broke, and return
                Chainlink price
277     if (_bandIsBroken(bandResponse)) {
278         return (Status.usingChainlinkBandUntrusted, chainlinkResponse.answer);
279     }
280
281     // If Chainlink is live and Band is frozen, just use last good price (no
                status change) since we have no basis for comparison
282     if (_bandIsFrozen(bandResponse)) {return (Status.usingBandChainlinkFrozen,
                lastGoodPrice);}
283
284     // If Chainlink is live and Band is working, compare prices. Switch to
                Chainlink
285     // if prices are within 5%, and return Chainlink price.
286     if (_bothOraclesSimilarPrice(chainlinkResponse, bandResponse)) {
287         return (Status.chainlinkWorking, chainlinkResponse.answer);
288     }
289
290     // Otherwise if Chainlink is live but price not within 5% of Band, distrust
                Chainlink, and return Band price
291     return (Status.usingBandChainlinkUntrusted, bandResponse.value);
292 }

```

Listing 3.3: PriceFeed::\_fetchPrice()

**Recommendation** Apply the proper state-transition logic in `_fetchPrice()` as elaborated earlier.

**Status** The issue has been confirmed.

### 3.3 Inconsistent Use Of `IncentivesController::handleAction()`

- ID: PVE-003
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: `StableDebtToken`
- Category: Business Logic [9]
- CWE subcategory: CWE-837 [5]

#### Description

The Geist protocol extends the built-in `IncentivesController` framework to engage protocol users. While reviewing the logic to integrate the incentive mechanism, we observe unnecessary inconsistency that may introduce unwanted confusion and errors.

To elaborate, we show below the `_mint()` function from both `IncentivizedERC20` and `StableDebtToken` contracts. It comes to our attention that the first contract uses the post-update balance in the invocation of `IncentivesController::handleAction()` (line 212) while the second contract uses the pre-update balance in the invocation of `IncentivesController::handleAction()` (line 413)!

```

200 function _mint(address account, uint256 amount) internal virtual {
201     require(account != address(0), 'ERC20: mint to the zero address');
202
203     _beforeTokenTransfer(address(0), account, amount);
204
205     uint256 currentTotalSupply = _totalSupply.add(amount);
206     _totalSupply = currentTotalSupply;
207
208     uint256 accountBalance = _balances[account].add(amount);
209     _balances[account] = accountBalance;
210
211     if (address(_getIncentivesController()) != address(0)) {
212         _getIncentivesController().handleAction(account, accountBalance,
213             currentTotalSupply);
214     }
215 }
```

Listing 3.4: `IncentivizedERC20::_mint()`

```

398 /**
399  * @dev Mints stable debt tokens to an user
400  * @param account The account receiving the debt tokens
401  * @param amount The amount being minted
402  * @param oldTotalSupply the total supply before the minting event
403  */
```

```

404 function _mint(
405     address account,
406     uint256 amount,
407     uint256 oldTotalSupply
408 ) internal {
409     uint256 oldAccountBalance = _balances[account];
410     _balances[account] = oldAccountBalance.add(amount);
411
412     if (address(_incentivesController) != address(0)) {
413         _incentivesController.handleAction(account, oldAccountBalance, oldTotalSupply);
414     }
415 }

```

Listing 3.5: StableDebtToken::\_mint()

**Recommendation** Be consistent in using the account balance for incentivization measurement.

**Status** The issue has been confirmed and the team clarifies the `stable lending` is disabled in Geist, which resolves the above inconsistency.

Table 3.1: Basic View-Only Functions Defined in The ERC20 Specification

Item	Description	Status
name()	Is declared as a public view function	✓
	Returns a string, for example "Tether USD"	✓
symbol()	Is declared as a public view function	✓
	Returns the symbol by which the token contract should be known, for example "USDT". It is usually 3 or 4 characters in length	✓
decimals()	Is declared as a public view function	✓
	Returns decimals, which refers to how divisible a token can be, from 0 (not at all divisible) to 18 (pretty much continuous) and even higher if required	✓
totalSupply()	Is declared as a public view function	✓
	Returns the number of total supplied tokens, including the total minted tokens (minus the total burned tokens) ever since the deployment	✓
balanceOf()	Is declared as a public view function	✓
	Anyone can query any address' balance, as all data on the blockchain is public	✓
allowance()	Is declared as a public view function	✓
	Returns the amount which the spender is still allowed to withdraw from the owner	✓



### 3.4 ERC20 Compliance Of GeistToken

---

- ID: PVE-004
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: GeistToken
- Category: Coding Practices [8]
- CWE subcategory: CWE-1126 [2]

#### Description

The GeistToken protocol is designed as the governance token, which will be disseminated to community and protocol users. In the following, we examine the ERC20 compliance of the GEIST token contract. Specifically, the ERC20 specification defines a list of API functions (and relevant events) that each token contract is expected to implement (and emit). The failure to meet these requirements means the token contract cannot be considered to be ERC20-compliant. Naturally, as part of our audit, we examine the list of API functions defined by the ERC20 specification and validate whether there exist any inconsistency or incompatibility in the implementation or the inherent business logic of the audited contract(s).

Our analysis shows that there is a minor ERC20 inconsistency or incompatibility issue. Specifically, the current implementation has defined the `decimals` state with the `uint256` type. The ERC20 specification indicates the type of `uint8` for the `decimals` state. Note that this incompatibility issue does not necessarily affect the functionality of GEIST in any negative way.

In the surrounding two tables, we outline the respective list of basic `view-only` functions (Table 3.1) and key `state-changing` functions (Table 3.2) according to the widely-adopted ERC20 specification. In addition, we perform a further examination on certain features that are permitted by the ERC20 specification or even further extended in follow-up refinements and enhancements (e.g., ERC777/ERC2222), but not required for implementation. These features are generally helpful, but may also impact or bring certain incompatibility with current DeFi protocols. Therefore, we consider it is important to highlight them as well. This list is shown in Table 3.3.

**Recommendation** Revise the GEIST implementation to ensure its ERC20-compliance.

**Status** The issue has been fixed by this commit: `de05115`.

Table 3.2: Key State-Changing Functions Defined in The ERC20 Specification

Item	Description	Status
<b>transfer()</b>	Is declared as a public function	✓
	Returns a boolean value which accurately reflects the token transfer status	✓
	Reverts if the caller does not have enough tokens to spend	✓
	Allows zero amount transfers	✓
	Emits Transfer() event when tokens are transferred successfully (include 0 amount transfers)	✓
	Reverts while transferring to zero address	✓
<b>transferFrom()</b>	Is declared as a public function	✓
	Returns a boolean value which accurately reflects the token transfer status	✓
	Reverts if the spender does not have enough token allowances to spend	✓
	Updates the spender's token allowances when tokens are transferred successfully	✓
	Reverts if the from address does not have enough tokens to spend	✓
	Allows zero amount transfers	✓
	Emits Transfer() event when tokens are transferred successfully (include 0 amount transfers)	✓
	Reverts while transferring from zero address	✓
	Reverts while transferring to zero address	✓
<b>approve()</b>	Is declared as a public function	✓
	Returns a boolean value which accurately reflects the token approval status	✓
	Emits Approval() event when tokens are approved successfully	✓
	Reverts while approving to zero address	✓
<b>Transfer() event</b>	Is emitted when tokens are transferred, including zero value transfers	✓
	Is emitted with the from address set to <i>address(0x0)</i> when new tokens are generated	✓
<b>Approval() event</b>	Is emitted on any successful call to approve()	✓

Table 3.3: Additional opt-in Features Examined in Our Audit

Feature	Description	Opt-in
<b>Deflationary</b>	Part of the tokens are burned or transferred as fee while on transfer()/transferFrom() calls	—
<b>Rebasing</b>	The balanceOf() function returns a re-based balance instead of the actual stored amount of tokens owned by the specific address	—
<b>Pausable</b>	The token contract allows the owner or privileged users to pause the token transfers and other operations	—
<b>Blacklistable</b>	The token contract allows the owner or privileged users to blacklist a specific address such that token transfers and other operations related to that address are prohibited	—
<b>Mintable</b>	The token contract allows the owner or privileged users to mint tokens to a specific address	✓
<b>Burnable</b>	The token contract allows the owner or privileged users to burn tokens of a specific address	✓

### 3.5 Fork-Resistant Domain Separator in AToken

- ID: PVE-005
- Severity: Medium
- Likelihood: Low
- Impact: High
- Target: AToken
- Category: Business Logic [9]
- CWE subcategory: CWE-841 [6]

#### Description

The various tokens in Geist are designed to strictly follow the widely-accepted ERC20 specification (Section 3.4). In the meantime, we notice the support of EIP-2612 with the permit() function that allows for approvals to be made via secp256k1 signatures. Interestingly, we notice the state variable DOMAIN\_SEPARATOR in AToken is initialized once inside the initialize() function (lines 81-89).

```

64  function initialize(
65      ILendingPool pool,
66      address treasury,
67      address underlyingAsset,
68      IAaveIncentivesController incentivesController,
69      uint8 aTokenDecimals,
70      string calldata aTokenName,
71      string calldata aTokenSymbol,
72      bytes calldata params
73  ) external override initializer {
74      uint256 chainId;
75
76      //solium-disable-next-line

```

```

77     assembly {
78         chainId := chainid()
79     }
80
81     DOMAIN_SEPARATOR = keccak256(
82         abi.encode(
83             EIP712_DOMAIN,
84             keccak256(bytes(aTokenName)),
85             keccak256(EIP712_REVISION),
86             chainId,
87             address(this)
88         )
89     );
90
91     _setName(aTokenName);
92     _setSymbol(aTokenSymbol);
93     _setDecimals(aTokenDecimals);
94
95     _pool = pool;
96     _treasury = treasury;
97     _underlyingAsset = underlyingAsset;
98     _incentivesController = incentivesController;
99
100    emit Initialized(
101        underlyingAsset,
102        address(pool),
103        treasury,
104        address(incentivesController),
105        aTokenDecimals,
106        aTokenName,
107        aTokenSymbol,
108        params
109    );
110 }

```

Listing 3.6: AToken::initialize()

The DOMAIN\_SEPARATOR is used in the permit() function and should be unique to the contract and chain in order to prevent replay attacks from other domains. However, when analyzing this permit() routine, we realize the current implementation needs to be improved by recalculating the value of DOMAIN\_SEPARATOR inside the permit() function, for the very purpose of preventing cross-chain replay attacks. Specifically, when there is a chain-level hard-fork, because of the pre-computed DOMAIN\_SEPARATOR, a valid signature for one chain could be replayed on the other.

```

336    function permit(
337        address owner,
338        address spender,
339        uint256 value,
340        uint256 deadline,
341        uint8 v,
342        bytes32 r,

```

```

343     bytes32 s
344 ) external {
345     require(owner != address(0), 'INVALID_OWNER');
346     //solium-disable-next-line
347     require(block.timestamp <= deadline, 'INVALID_EXPIRATION');
348     uint256 currentValidNonce = _nonces[owner];
349     bytes32 digest =
350         keccak256(
351             abi.encodePacked(
352                 '\x19\x01',
353                 DOMAIN_SEPARATOR,
354                 keccak256(abi.encode(PERMIT_TYPEHASH, owner, spender, value, currentValidNonce
355                                     , deadline))
356             );
357     require(owner == ecrecover(digest, v, r, s), 'INVALID_SIGNATURE');
358     _nonces[owner] = currentValidNonce.add(1);
359     _approve(owner, spender, value);
360 }

```

Listing 3.7: AToken::permit()

**Recommendation** Recalculate the value of DOMAIN\_SEPARATOR inside the permit() function.

**Status** This issue has been confirmed.

## 3.6 Suggested Adherence Of Checks-Effects-Interactions Pattern

- ID: PVE-006
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: MasterChef
- Category: Time and State [10]
- CWE subcategory: CWE-663 [4]

### Description

A common coding best practice in Solidity is the adherence of checks-effects-interactions principle. This principle is effective in mitigating a serious attack vector known as re-entrancy. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the DAO [15] exploit, and the recent Uniswap/Lendf.Me hack [14].

We notice there is an occasion where the checks-effects-interactions principle is violated. Using the `MasterChef` as an example, the `emergencyWithdraw()` function (see the code snippet below) is provided to externally call a token contract to transfer assets. However, the invocation of an external contract requires extra care in avoiding the above re-entrancy.

Apparently, the interaction with the external contract (line 297) starts before effecting the update on internal states (lines 299–300), hence violating the principle. In this particular case, if the external contract has certain hidden logic that may be capable of launching re-entrancy via the same entry function. The same issue is also applicable to other functions, including `deposit()` and `withdraw()`.

```

292 // Withdraw without caring about rewards. EMERGENCY ONLY.
293 function emergencyWithdraw(address _token) external {
294     PoolInfo storage pool = poolInfo[_token];
295     UserInfo storage user = userInfo[_token][msg.sender];
296     uint256 amount = user.amount;
297     IERC20(_token).safeTransfer(address(msg.sender), amount);
298     emit EmergencyWithdraw(_token, msg.sender, amount);
299     user.amount = 0;
300     user.rewardDebt = 0;
301     if (pool.onwardIncentives != IOnwardIncentivesController(0)) {
302         uint256 lpSupply = IERC20(_token).balanceOf(address(this));
303         try pool.onwardIncentives.handleAction(_token, msg.sender, 0, lpSupply) {}
304         catch {}
305     }

```

Listing 3.8: `MasterChef::emergencyWithdraw()`

In the meantime, we should mention that the supported tokens in the protocol do implement rather standard ERC20 interfaces and their related token contracts are not vulnerable or exploitable for re-entrancy. However, it is important to take precautions in making use of `nonReentrant` to block possible re-entrancy and the adherence of checks-effects-interactions best practice is highly recommended.

**Recommendation** Apply necessary reentrancy prevention by following the known checks-effects-interactions pattern in addition to the utilization of the `nonReentrant` modifier to block possible re-entrancy.

**Status** The issue has been fixed by this commit: [78b1414](#).

### 3.7 Staking Incompatibility With Deflationary Tokens

- ID: PVE-007
- Severity: Undetermined
- Likelihood: N/A
- Impact: N/A
- Target: MasterChef
- Category: Business Logic [9]
- CWE subcategory: CWE-841 [6]

#### Description

In the Geist protocol, the MasterChef contract is designed to take users' assets and deliver rewards depending on their share. In particular, one interface, i.e., `deposit()`, accepts asset transfer-in and records the depositor's balance. Another interface, i.e., `withdraw()`, allows the user to withdraw the asset with necessary bookkeeping under the hood. For the above two operations, i.e., `deposit()` and `withdraw()`, the contract using the `safeTransfer()/safeTransferFrom()` routines to transfer assets into or out of its pool. This routine works as expected with standard ERC20 tokens: namely the pool's internal asset balances are always consistent with actual token balances maintained in individual ERC20 token contract.

However, there exist other ERC20 tokens that may make certain customization to their ERC20 contracts. One type of these tokens is deflationary tokens that charge certain fee for every transfer. As a result, this may not meet the assumption behind asset-transferring routines. In other words, the above operations, such as `deposit()` and `withdraw()`, may introduce unexpected balance inconsistencies when comparing internal asset records with external ERC20 token contracts. Apparently, these balance inconsistencies are damaging to accurate and precise portfolio management of the pool and affects protocol-wide operation and maintenance.

Specially, if we take a look at the `_updatePool()` routine. This routine calculates `pool.accRewardPerShare` via dividing reward by `lpSupply`, where the `lpSupply` is derived from `IERC20(_token).balanceOf(address(this))` (line 212). Because the balance inconsistencies of the pool, the `lpSupply` could be 1 Wei and thus may yield a huge `pool.accRewardPerShare` as the final result, which dramatically inflates the pool's reward.

```

206 // Update reward variables of the given pool to be up-to-date.
207 function _updatePool(address _token, uint256 _totalAllocPoint) internal {
208     PoolInfo storage pool = poolInfo[_token];
209     if (block.timestamp <= pool.lastRewardTime) {
210         return;
211     }
212     uint256 lpSupply = IERC20(_token).balanceOf(address(this));
213     if (lpSupply == 0) {
214         pool.lastRewardTime = block.timestamp;
215         return;
216     }

```

```

217     uint256 duration = block.timestamp.sub(pool.lastRewardTime);
218     uint256 reward = duration.mul(rewardsPerSecond).mul(pool.allocPoint).div(
        _totalAllocPoint);
219     pool.accRewardPerShare = pool.accRewardPerShare.add(reward.mul(1e12).div(
        lpSupply));
220     pool.lastRewardTime = block.timestamp;
221 }

```

Listing 3.9: MasterChef::\_updatePool()

One mitigation is to measure the asset change right before and after the asset-transferring routines. In other words, instead of bluntly assuming the amount parameter in `safeTransfer()` or `safeTransferFrom()` will always result in full transfer, we need to ensure the increased or decreased amount in the pool before and after the `safeTransfer()` or `safeTransferFrom()` is expected and aligned well with our operation. Though these additional checks cost additional gas usage, we consider they are necessary to deal with deflationary tokens or other customized ones if their support is deemed necessary.

Another mitigation is to regulate the set of ERC20 tokens that are permitted into `Geist` for support. However, certain existing stable coins may exhibit control switches that can be dynamically exercised to convert into deflationary.

**Recommendation** Check the balance before and after the `safeTransfer()` or `safeTransferFrom()` call to ensure the book-keeping amount is accurate. An alternative solution is using non-deflationary tokens as collateral but some tokens (e.g., USDT) allow the admin to have the deflationary-like features kicked in later, which should be verified carefully.

**Status** This issue has been confirmed. The team clarifies no plan to support deflationary tokens here.

### 3.8 Trust Issue of Admin Keys

- ID: PVE-008
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple Contracts
- Category: Security Features [7]
- CWE subcategory: CWE-287 [3]

#### Description

In the `Geist` protocol, there is a privileged `owner` account that plays a critical role in governing and regulating the system-wide operations (e.g., parameter setting and price oracle adjustment). It also has the privilege to control or govern the flow of assets managed by this protocol. Our analysis



shows that the privileged account needs to be scrutinized. In the following, we examine the privileged account and the related privileged accesses in current contracts.

```

135     function setOnwardIncentives(
136         address _token,
137         IOnwardIncentivesController _incentives
138     )
139     external
140     onlyOwner
141     {
142         require(poolInfo[_token].lastRewardTime != 0);
143         poolInfo[_token].onwardIncentives = _incentives;
144     }
145
146     function setClaimReceiver(address _user, address _receiver) external {
147         require(msg.sender == _user && msg.sender == owner());
148         claimReceiver[_user] = _receiver;
149     }

```

Listing 3.10: Example Setters in ChefIncentivesController

Moreover, the LendingPoolAddressesProvider contract allows the privileged owner to configure protocol-wide contracts, including LENDING\_POOL, LENDING\_POOL\_CONFIGURATOR, POOL\_ADMIN, EMERGENCY\_ADMIN, LENDING\_POOL\_COLLATERAL\_MANAGER, PRICE\_ORACLE, and LENDING\_RATE\_ORACLE. These contracts play a variety of duties and are also considered privileged.

```

19 contract LendingPoolAddressesProvider is Ownable, ILendingPoolAddressesProvider {
20     string private _marketId;
21     mapping(bytes32 => address) private _addresses;
22
23     bytes32 private constant LENDING_POOL = 'LENDING_POOL';
24     bytes32 private constant LENDING_POOL_CONFIGURATOR = 'LENDING_POOL_CONFIGURATOR';
25     bytes32 private constant POOL_ADMIN = 'POOL_ADMIN';
26     bytes32 private constant EMERGENCY_ADMIN = 'EMERGENCY_ADMIN';
27     bytes32 private constant LENDING_POOL_COLLATERAL_MANAGER = 'COLLATERAL_MANAGER';
28     bytes32 private constant PRICE_ORACLE = 'PRICE_ORACLE';
29     bytes32 private constant LENDING_RATE_ORACLE = 'LENDING_RATE_ORACLE';
30     ...
31 }

```

Listing 3.11: The LendingPoolAddressesProvider Contract

It is worrisome if the privileged owner account is a plain EOA account. The discussion with the team confirms that the owner account is currently managed by a 2-day timelock, owned by a multisig. Note that a multi-sig account indeed greatly alleviates this concern, though it is still not perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO. In the meantime, a timelock-based mechanism can also be considered as mitigation. Moreover, it should be noted if current contracts are to be deployed behind a proxy, there is a need to properly manage the proxy-admin privileges as they fall in this trust issue as well.

**Recommendation** Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

**Status** This issue has been confirmed with the team. For the time being, it is being mitigated by a 2-day timelock (owned by a multi-sig account) to balance efficiency and timely adjustment. After the protocol becomes stable, it is expected to eventually migrate to be owned by community proposals for decentralized governance.



## 4 | Conclusion

In this audit, we have analyzed the `Geist` design and implementation. The system presents a unique, robust offering as a decentralized non-custodial money market protocol where users can participate as depositors or borrowers. The `Geist` protocol extends the original `AaveV2` with new features for staking-based incentivization and fee distribution. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and fixed.

Moreover, we need to emphasize that `Solidity`-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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