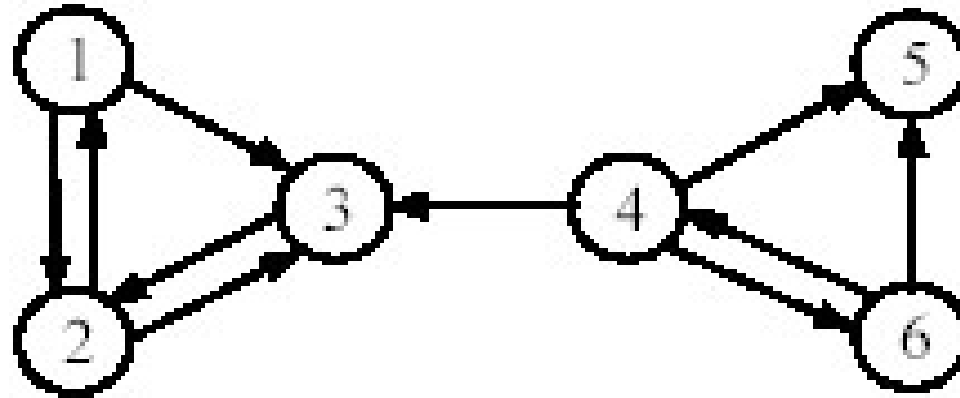


An example Web hyperlink graph



PageRank from
nodes 4,5,6 goes
to nodes 1,2,3.
At the end, 4,5,6
have PageRank
as zero

$$A = \begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 1/2 & 0 & 1/2 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1/3 & 0 & 1/3 & 1/3 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1/2 & 1/2 & 0 \end{pmatrix}$$

The final PageRank algorithm

- $(1-d)\mathbf{E}/n + d\mathbf{A}^T$ is a **stochastic matrix** (transposed). It is also **irreducible** and **aperiodic**
- If we scale Equation (25) so that $\mathbf{e}^T \mathbf{P} = n$,

$$\mathbf{P} = (1-d)\mathbf{e} + d\mathbf{A}^T \mathbf{P} \quad (27)$$

- PageRank for each page i is

$$P(i) = (1-d) + d \sum_{j=1}^n A_{ji} P(j) \quad (28)$$

The final PageRank (cont ...)

- (28) is equivalent to the formula given in the PageRank paper

$$P(i) = (1 - d) + d \sum_{(j,i) \in E} \frac{P(j)}{O_j}$$

- The parameter d is called the **damping factor** which can be set to between 0 and 1. $d = 0.85$ was used in the PageRank paper.

Compute PageRank

- Use the **power iteration** method

PageRank-Iterate(G)

$P_0 \leftarrow e/n$

$k = 1$

repeat

$P_{k+1} \leftarrow (1-d)e + dA^T P_k ;$

$k = k + 1 ;$

until $\|P_{k+1} - P_k\|_1 < \varepsilon$

return P_{k+1}

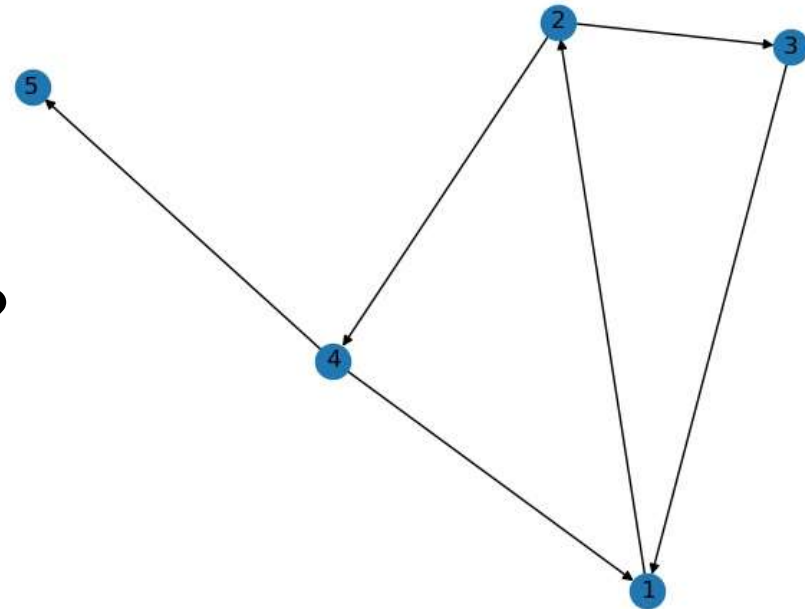
Fig. 6. The power iteration method for PageRank

Advantages of PageRank

- **Fighting spam.** A page is important if the pages pointing to it are important.
 - Since it is not easy for Web page owner to add in-links into his/her page from other important pages, it is thus not easy to influence PageRank.
- **PageRank is a global measure and is query independent.**
 - PageRank values of all the pages are computed and saved off-line rather than at the query time.
- **Criticism:** Query-independence. It could not distinguish between pages that are authoritative in general and pages that are authoritative on the query topic.

Examples of Centrality

- Consider node 4:
 - ❑ Degree centrality?
 - ❑ Closeness Centrality?
 - ❑ Betweenness Centrality?
 - ❑ Degree Prestige?
 - ❑ Proximity Prestige?
 - ❑ Rank Prestige?



Examples of Centrality

■ Consider node 4:

□ Degree centrality?

Degree centrality=2

Normalized degree centrality=1/2

□ Closeness Centrality?

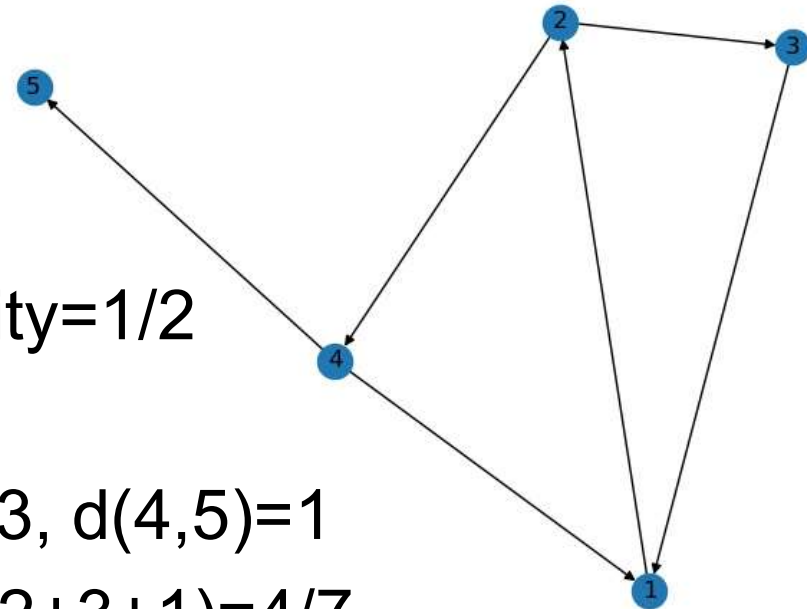
$d(4,1)=1$, $d(4,2)=2$, $d(4,3)=3$, $d(4,5)=1$

Closeness centrality= $4/(1+2+3+1)=4/7$

□ Betweenness Centrality?

Number of shortest paths that include 4=3.5
(includes 2-5, 3-5, 1-5, 2-4-1)

Normalized value= $3.5/(4*3)=0.283$



Examples of Prestige

- Consider node 4:

- Degree Prestige?

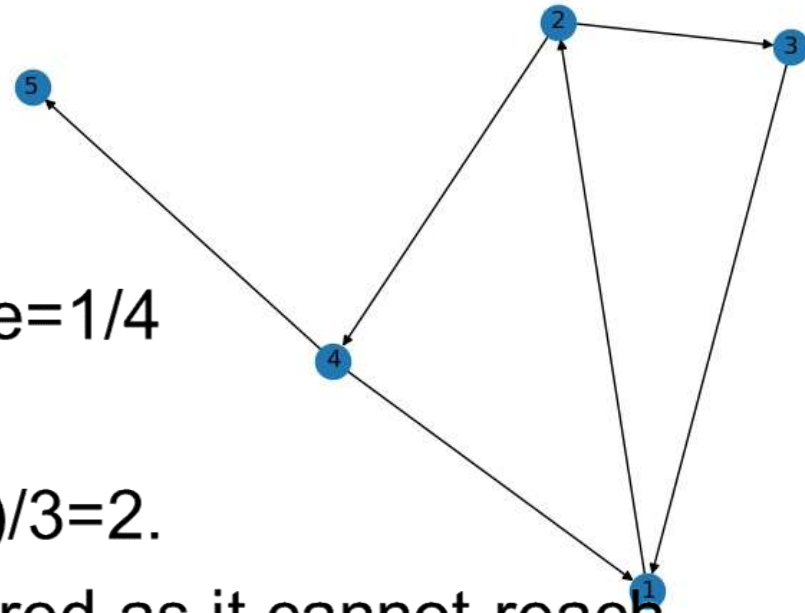
Degree prestige=1

Normalized degree prestige=1/4

- Proximity Prestige?

Proximity prestige=(2+1+3)/3=2.

Here node 5 is not considered as it cannot reach node 4.



2 / 1 2

Examples of Prestige

- Consider node 4:

- Rank Prestige?

Let $P(i)$ be Rank prestige of node i

$$P(1) = P(3) + P(4)$$

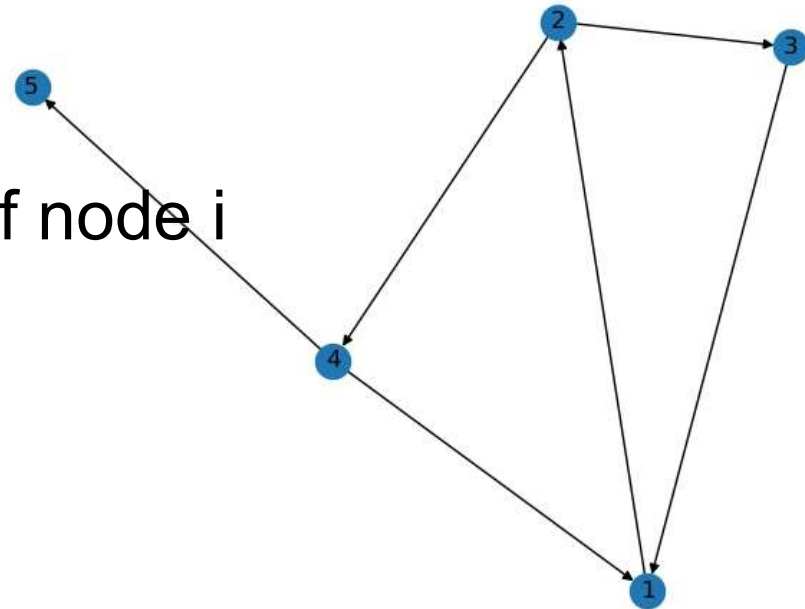
$$P(2) = P(1)$$

$$P(3) = P(2)$$

$$P(4) = P(2)$$

$$P(5) = P(4)$$

On solving, we get $P(i) = 0$ for every i .



Examples of PageRank without damping factor

- Consider node 4:

- Rank Prestige?

Let $P(i)$ be Rank prestige of node i

$$P(1) = P(3) + P(4)/2$$

$$P(2) = P(1)$$

$$P(3) = P(2)/2$$

$$P(4) = P(2)/2$$

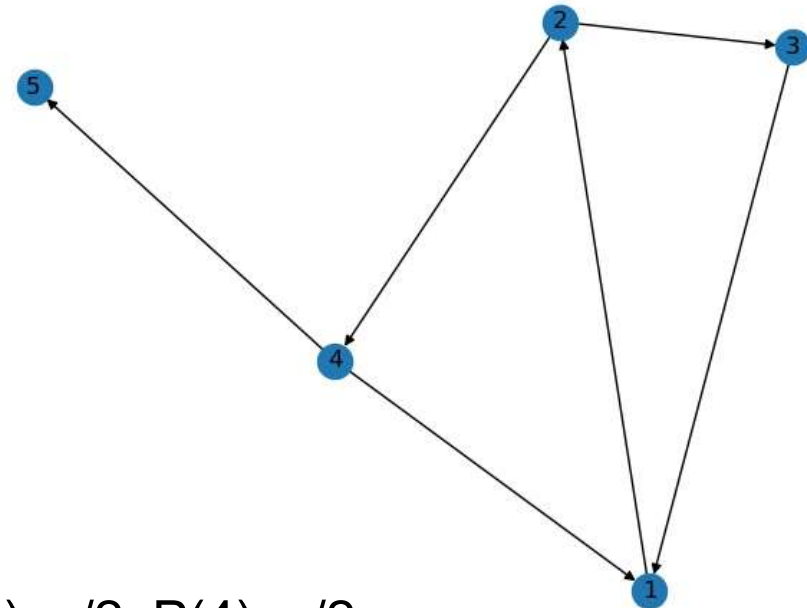
$$P(5) = P(4)/2$$

On solving, Let $P(1) = x$, $P(2) = x$, $P(3) = x/2$, $P(4) = x/2$,

$$P(5) = x/4$$

For $P(1)$,

$x = x/2 + x/4$. This works for only $x = 0$.



Examples of PageRank with damping factor say 0.8

■ Consider node 4:

□ Rank Prestige?

Let $P(i)$ be Rank prestige of node i

$$P(1) = 0.2 + 0.8(P(3) + P(4)/2)$$

$$P(2) = 0.2 + 0.8(P(1))$$

$$P(3) = 0.2 + 0.8 * (P(2)/2)$$

$$P(4) = 0.2 + 0.8 * (P(2)/2)$$

$$P(5) = 0.2 + 0.8 * (P(4)/2)$$

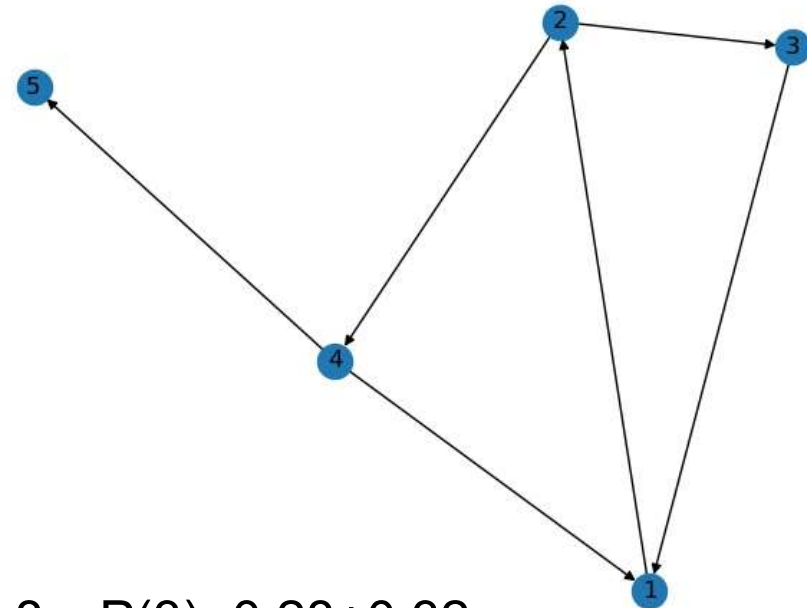
On solving, Let $P(1) = x$, $P(2) = 0.2 + 0.8x$, $P(3) = 0.28 + 0.32x$,

$$P(4) = 0.28 + 0.32x, P(5) = 0.424 + 0.256x$$

For $P(1)$,

$$x = 0.2 + 0.8(0.2 + 0.8x + 0.14 + 0.16x)$$

This works for only $x = 2.0344..$



Road map

- Introduction
- Social network analysis
- Co-citation and bibliographic coupling
- PageRank
- **HITS**
- Community Discovery
- Summary

HITS

- HITS stands for **Hypertext Induced Topic Search**.
- Unlike PageRank which is a static ranking algorithm, **HITS is search query dependent**.
- When the user issues a search query,
 - HITS first expands the list of relevant pages returned by a search engine and
 - then produces two rankings of the expanded set of pages, **authority ranking** and **hub ranking**.

Authorities and Hubs

Authority: Roughly, a authority is a page with many in-links.

- ❑ The idea is that the page may have good or authoritative content on some topic and
- ❑ thus many people trust it and link to it.

Hub: A hub is a page with many out-links.

- ❑ The page serves as an organizer of the information on a particular topic and
- ❑ points to many good authority pages on the topic.

Examples

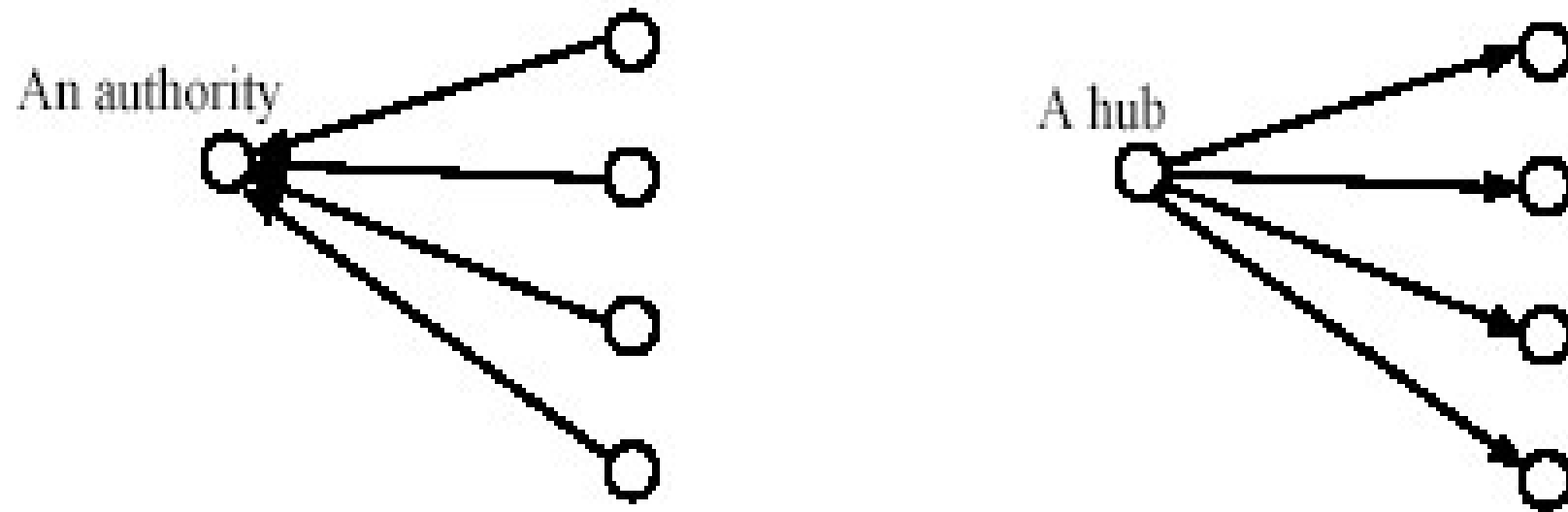


Fig. 7. An authority page and a hub page

The key idea of HITS

- A good hub points to many good authorities, and
- A good authority is pointed to by many good hubs.
- Authorities and hubs have a **mutual reinforcement relationship**. Fig. 8 shows some densely linked authorities and hubs (a **bipartite sub-graph**).

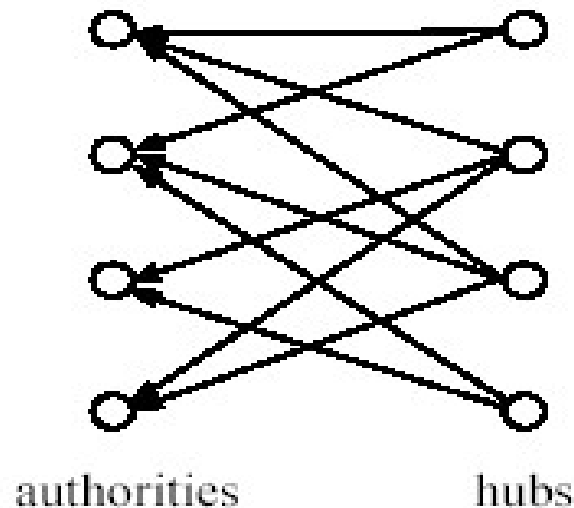


Fig. 8. A densely linked set of authorities and hubs

The HITS algorithm: Grab pages

- Given a broad search query, q , HITS collects a set of pages as follows:
 - It sends the query q to a search engine.
 - It then collects t ($t = 200$ is used in the HITS paper) highest ranked pages. This set is called the **root** set W .
 - It then grows W by including any page pointed to by a page in W and any page that points to a page in W . This gives a larger set S , **base set**.

The link graph G

- HITS works on the pages in S , and assigns every page in S an **authority score** and a **hub score**.
- Let the number of pages in S be n .
- We again use $G = (V, E)$ to denote the hyperlink graph of S .
- We use L to denote the adjacency matrix of the graph.

$$L_{ij} = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{otherwise} \end{cases}$$

The HITS algorithm

- Let the authority score of the page i be $a(i)$, and the hub score of page i be $h(i)$.
- The mutual reinforcing relationship of the two scores is represented as follows:

$$a(i) = \sum_{(j,i) \in E} h(j) \quad (31)$$

$$h(i) = \sum_{(i,j) \in E} a(j) \quad (32)$$

HITS in matrix form

- We use \mathbf{a} to denote the column vector with all the authority scores,

$$\mathbf{a} = (a(1), a(2), \dots, a(n))^T, \text{ and}$$

- use \mathbf{h} to denote the column vector with all the authority scores,

$$\mathbf{h} = (h(1), h(2), \dots, h(n))^T,$$

- Then,

$$\mathbf{a} = \mathbf{L}^T \mathbf{h} \tag{33}$$

$$\mathbf{h} = \mathbf{L} \mathbf{a} \tag{34}$$

Computation of HITS

- The computation of authority scores and hub scores is the same as the computation of the PageRank scores, using **power iteration**.
- If we use \mathbf{a}_k and \mathbf{h}_k to denote authority and hub vectors at the k th iteration, the iterations for generating the final solutions are

$$\mathbf{a}_k = L^T L \mathbf{a}_{k-1} \quad (35)$$

$$\mathbf{h}_k = L L^T \mathbf{h}_{k-1} \quad (36)$$

starting with

$$\mathbf{a}_0 = \mathbf{h}_0 = (1, 1, \dots, 1), \quad (37)$$

The algorithm

HITS-Iterate(G)

$a_0 = h_0 = (1, 1, \dots, 1);$

$k = 1$

Repeat

$a_k = L^T L a_{k-1};$

$h_k = L L^T h_{k-1};$

normalize a_k ;

normalize h_k ;

$k = k + 1;$

until a_k and h_k do not change significantly;

return a_k and h_k

Fig. 9. The HITS algorithm based on power iteration

Relationships with co-citation and bibliographic coupling

- Recall that co-citation of pages i and j , denoted by C_{ij} , is

$$C_{ij} = \sum_{k=1}^n L_{ki} L_{kj} = (\mathbf{L}^T \mathbf{L})_{ij}$$

- the authority matrix $(\mathbf{L}^T \mathbf{L})$ of HITS is the co-citation matrix \mathbf{C}

- bibliographic coupling of two pages i and j , denoted by B_{ij} is

$$B_{ij} = \sum_{k=1}^n L_{ik} L_{jk} = (\mathbf{L} \mathbf{L}^T)_{ij},$$

- the hub matrix $(\mathbf{L} \mathbf{L}^T)$ of HITS is the bibliographic coupling matrix \mathbf{B}

Strengths and weaknesses of HITS

- **Strength:** its ability to rank pages according to the query topic, which may be able to provide more relevant authority and hub pages.
- **Weaknesses:**
 - ❑ **It is easily spammed.** It is in fact quite easy to influence HITS since adding out-links in one's own page is so easy.
 - ❑ **Topic drift.** Many pages in the expanded set may not be on topic.
 - ❑ **Inefficiency at query time:** The query time evaluation is slow. Collecting the root set, expanding it and performing eigenvector computation are all expensive operations

Road map

- **Introduction**
- **Social network analysis**
- **Co-citation and bibliographic coupling**
- **PageRank**
- **HITS**
- **Community Discovery**
- **Summary**

Communities

- A community is simply a group of entities (e.g., people or organizations) that shares a common interest or is involved in an activity or event.

Definition (community): Given a finite set of **entities** $S = \{s_1, s_2, \dots, s_n\}$ of the same **type**, a **community** is a pair $C = (T, G)$, where T is the **community theme** and $G \subseteq S$ is the set of all entities in S that shares the theme T . If $s_i \in G$, s_i is said to be a **member** of the community C .

Web Pages:

- 1. Hyperlinks: A group of content creators sharing a common interest is usually interconnected through hyperlinks. That is, members in a community are more likely to be connected among themselves than outside the community.
- 2. Content words: Web pages of a community usually contain words that are related to the community theme.

Emails:

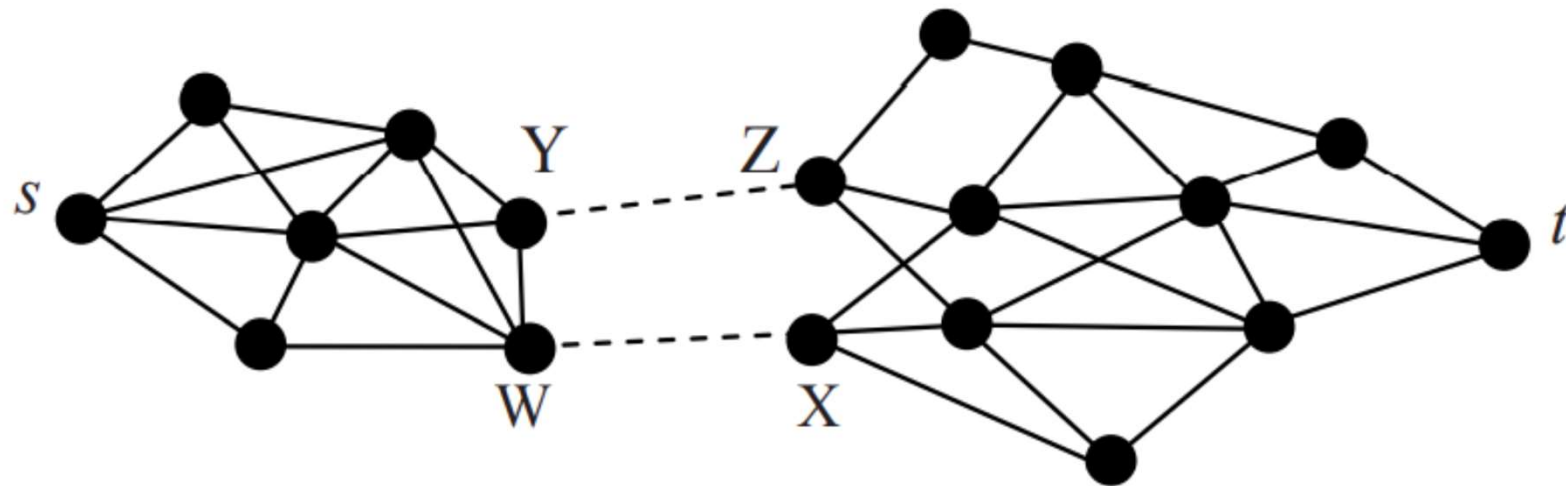
- 1. Email exchange between entities:
Members of a community are more likely to communicate with one another.
- 2. Content words: Email contents of a community also contain words related to the theme of the community.

Text documents:

- 1. Co-occurrence of entities: Members of a community are more likely to appear together in the same sentence and/or the same document.
- 2. Content words: Words in sentences indicate the community theme.

Maximum Flow Communities Detection

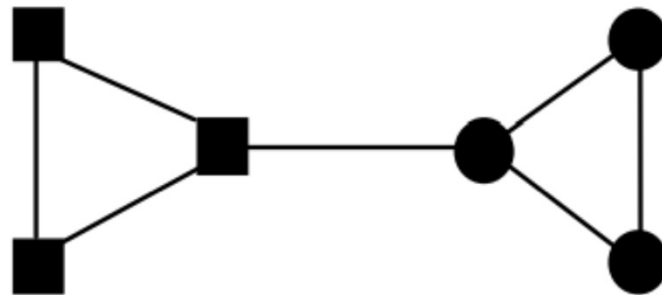
- Given a Web link graph $G = (V, E)$, a maximum flow community is defined as a collection $C \subseteq V$ of Web pages such that each member page $u \subseteq C$ has more hyperlinks (in either direction) within the community C than outside of the community $V-C$.
- Identifying a community is NP-complete graph partition problems.



- The Max Flow-Min Cut theorem of Ford and Fulkerson [26] proves that the maximum flow of a network is identical to the minimum cut that separates s and t .

Community Detection using Betweenness

- Identify Edges through which maximum shortest paths pass.
- Start removing edges till the Graph disconnects.



Road map

- **Introduction**
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Summary

- In this chapter, we introduced
 - Social network analysis, centrality and prestige
 - Co-citation and bibliographic coupling
 - PageRank, which powers Google
 - HITS
- Yahoo! and MSN have their own link-based algorithms as well, but not published.
- **Important to note:** Hyperlink based ranking is not the only algorithm used in search engines. In fact, it is combined with many **content based factors** to produce the final ranking presented to the user.

Summary

- Links can also be used to find **communities**, which are groups of content-creators or people sharing some common interests.
 - ❑ Web communities
 - ❑ Email communities
 - ❑ Named entity communities
 - Focused crawling: combining contents and links to crawl Web pages of a specific topic.
 - ❑ Follow links and
 - ❑ Use learning/classification to determine whether a page is on topic.
-