# **Informed Search**



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## **Objectives**

#### Students are able:

- to contrast greedy best-first search and A\* search, and
- to apply greedy best-first search and A\* search to solve a simple problem.



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#### Siswa mampu:

- untuk membedakan greedy best-first search dan A\* search, serta
- untuk menerapkan *greedy best-first search* dan *A\* search* untuk menyelesaikan persoalan sederhana.

#### **Informed Search**

- Use domain knowledge!
  - Are we getting close to the goal?
  - Use a heuristic function that estimates how close a state is to the goal
  - A heuristic does NOT have to be perfect!
  - Example of strategies:
    - 1. Greedy Best-First Search
    - 2. A\* Search
    - 3. IDA\*

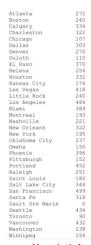


#### **Informed Search**



The distance is the straight line distance. The goal is to get to Sault Ste. Marie, so all the distances are from each city to Sault Ste. Marie.

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#### Heuristic!



## **Greedy Best-First Search**

- Evaluation function h(n) (heuristic)
- h(n) estimates the cost from n to the closest goal
- Example:  $h_{SLD}(n) = \text{straight-line distance from } n \text{ to Sault Ste. Marie}$
- Greedy search expands the node that appears to be closest to goal



## **Greedy Best-First Search**

```
function Greedy-Best-First-Search(initialState, goalTest)
returns Success or Failure: /* Cost f(n) = h(n) */

frontier = Heap.new(initialState)
explored = Set.new()

while not frontier.isEmpty():
    state = frontier.deleteMin()
    explored.add(state)

if goalTest(state):
    return Success(state)

for neighbor in state.neighbors():
    if neighbor not in frontier ∪ explored:
        frontier.insert(neighbor)
    else if neighbor in frontier:
        frontier.decreaseKey(neighbor)

return Failure
```

# Greedy Best-First Search Example The initial state:





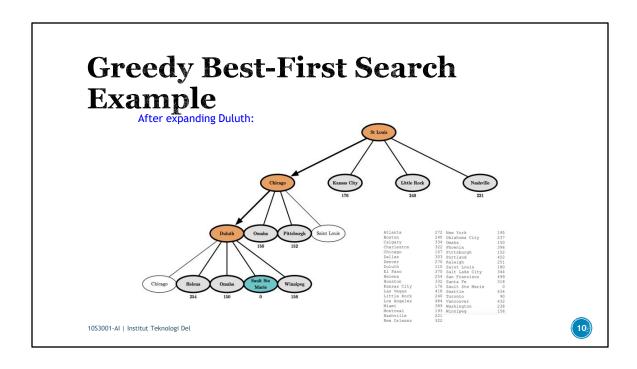
# Greedy Best-First Search Example After expanding St Louis:



Atlanta	272	New York	195
Boston	240	Oklahoma City	237
Calgary	334	Omaha	150
Charleston	322	Phoenix	396
Chicago	107	Pittsburgh	152
Dallas	303	Portland	452
Denver	270	Raleigh	251
Duluth	110	Saint Louis	180
El Paso	370	Salt Lake City	344
Helena	254	San Francisco	499
Houston	332	Santa Fe	318
Kansas City	176	Sault Ste Marie	0
Las Vegas	418	Seattle	434
Little Rock	240	Toronto	90
Los Angeles	484	Vancouver	432
Miami	389	Washington	238
Montreal	193	Winnipeg	156
Nashville	221		
New Orleans	322		



# 







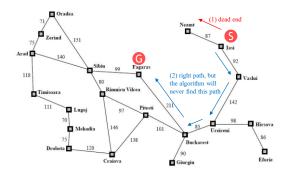
#### Greedy Best-First Search Criteria

• Complete: No

■ Time: *O*(*b*<sup>*m*</sup>)

**■ Space:** *O*(*b*<sup>*m*</sup>)

• Optimal: No



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#### A\* Search

- Minimize the total estimated solution cost
- Combines:
  - g(n): cost to reach node n
  - h(n): cost to get from n to the goal
  - $\bullet f(n) = g(n) + h(n)$

f(n) is the estimated cost of the cheapest solution through n

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#### A\* Search

```
\begin{aligned} & \textbf{function A-Star-Search}(& \text{initialState, goalTest}) \\ & & returns \textbf{Success or Failure:} \ /^* \ \text{Cost} \ f(n) = g(n) + h(n) \ ^*/ \\ & \text{frontier} = \text{Heap.new}(& \text{initialState}) \\ & \text{explored} = \text{Set.new}() \\ & \textbf{while not frontier.} & \text{isempty}(): \\ & \text{state} = & \text{frontier.} & \text{deleteMin}() \\ & \text{explored.add}(& \text{state}) \\ & \textbf{if goalTest}(& \text{state}): \\ & \text{return Success}(& \text{state}) \\ & \textbf{for neighbor in state.neighbors}(): \\ & \text{if neighbor not in frontier} \ \cup & \text{explored:} \\ & & \text{frontier.insert}(& \text{neighbor}) \\ & & \text{else if neighbor in frontier:} \\ & & \text{frontier.decreaseKey}(& \text{neighbor}) \\ & \text{return Failure} \end{aligned}
```

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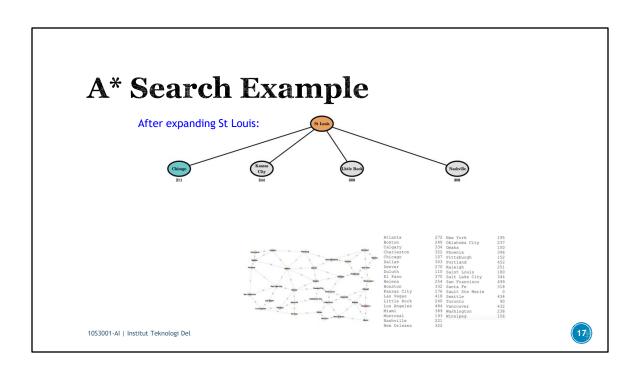
# **A\* Search Example**

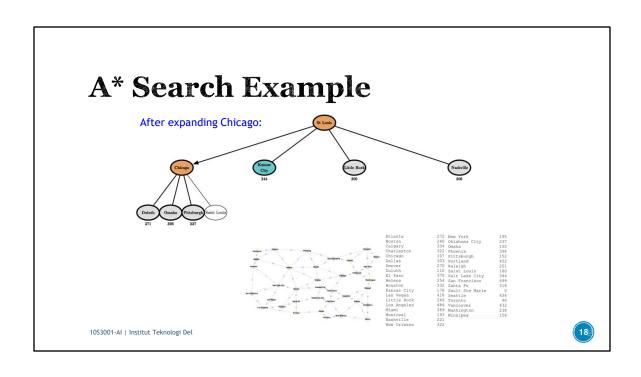
The initial state:

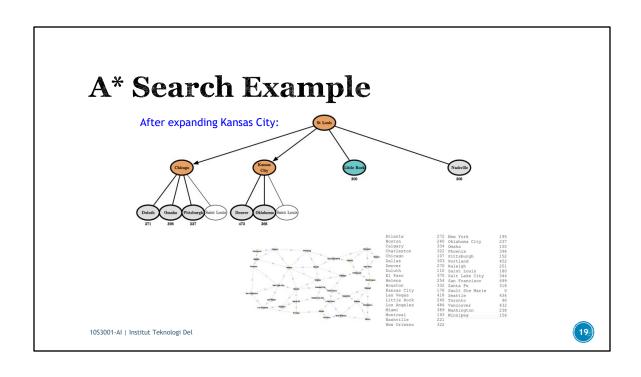


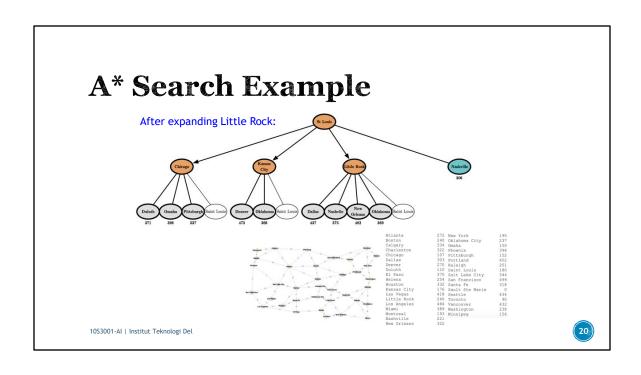


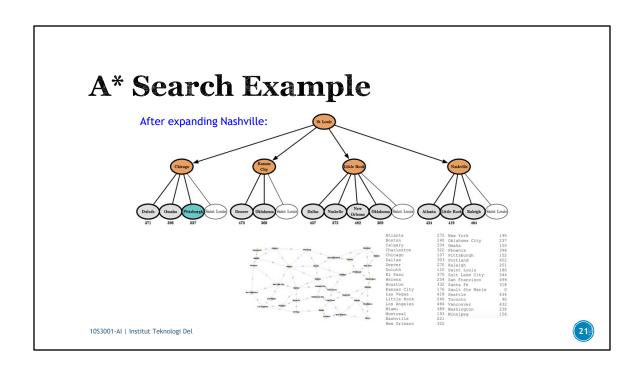


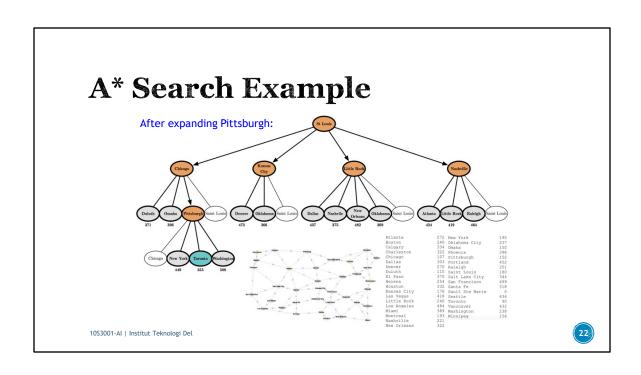


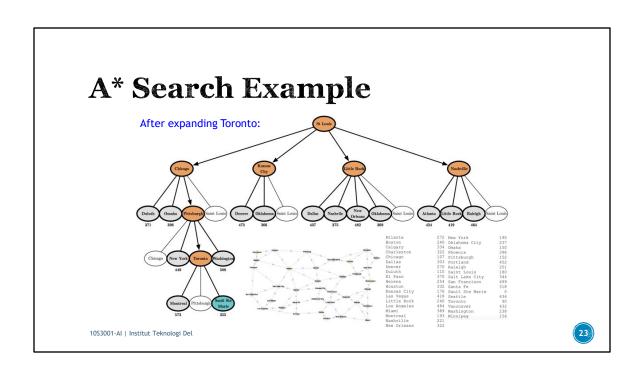


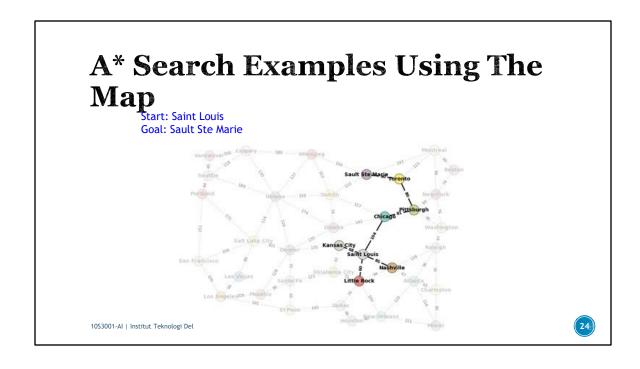


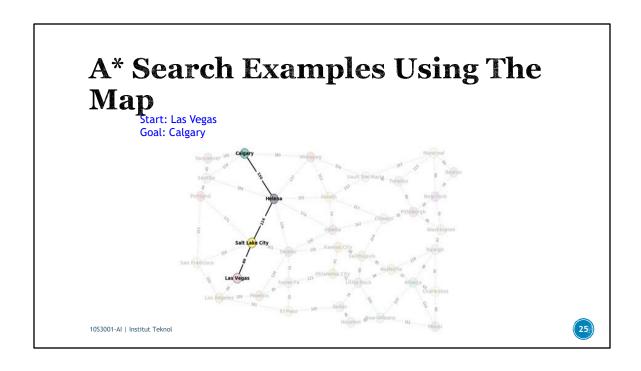












#### **Admissible Heuristics**

A good heuristic can be powerful. Only if it is of a "good quality".

A good heuristic must be admissible.

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A\* always find the optimal solution. Greedy Search sometimes hits the best solution but sometimes does not. So, A\* does a pretty good job at finding the optimal solution. However, it must rely on a good heuristic because a good heuristic can turn out to be powerful one that can lead us to the optimal solution.

So, a good heuristic is known to be powerful if it has the property of being what we call admissible. What is it?

#### **Admissible Heuristics**

- An admissible heuristic never overestimates the cost to reach the goal, that is it is optimistic.
- A heuristic h is admissible if

 $\forall$  node  $n, h(n) \leq h^*(n)$ 

where  $h^*$  is true cost to reach the goal from n.

 h<sub>SLD</sub> (used as a heuristic in the map example) is admissible because it is by definition the shortest distance between two points.

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Admissible heuristic is a heuristic that never overestimates the cost to reach the goal. That is it is optimistic, so it always has less than what it needs to reach the goal.

A heuristic h is admissible if for all node n, the value of the heuristic at the node n is actually bounded by the true cost to reach the goal from the node n.

So, in the example of the map, we use the heuristic  $h_{SLD}$  that actually is the straightest line between two cities. And this heuristic is admissible because it's always less than the true cost to reach the node.

This can be illustrated by a small example. If we have two cities A and B, the straightest distance between the two cities is actually not a real configuration. It's just an estimation of the distance between A and B. However, if you use a true path between A and B through roads, you would end up with some path that actually links the two cities, but whatever distance you take is going to be bigger than the straight line distance between A and B.

#### A\* Search Criteria

Complete: Yes

• Time: exponential

• Space: keeps every node in memory, the biggest problem

• Optimal: Yes!

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Complete because if there is a solution, you will find it.

Unfortunately, the time is still exponential, and the space is still exponential. The main reason is that it keeps all the nodes at memory to find the best possible solution of the optimal solution, which is the biggest problem. And as in the previous example, there actually lots of nodes were grey because they are all in the fringe waiting to be explored. So, there is an important use of memory to solve the problem.

It is optimal. If there is a solution, we find that it finds the best possible solution, which is a solution with the least cost in your search.

#### **A\* Optimality**

If h(n) is admissible, A\* using tree search is optimal.

#### Rationale:

- Suppose G<sub>o</sub> is the optimal goal.
- Suppose  $G_S$  is some suboptimal goal.
- Suppose n is an unexpanded node in the fringe such that n is on the shortest path to  $G_o$ .
- $f(G_s) = g(G_s)$  since  $h(G_s) = 0$  $f(G_o) = g(G_o)$  since  $h(G_o) = 0$  $f(G_s) > g(G_o)$  since  $G_s$  is suboptimal Then  $f(G_s) > f(G_o) \dots (1)$
- $h(n) \le h^*(n)$  since h is admissible

```
g(n)+h(n)\leq g(n)+h^*(n)=g(G_o)=f(G_o)
```

Then,  $f(n) \leq f(G_o) \dots (2)$ 

From (1) and (2)  $f(G_s) > f(n)$ 

so  $A^*$  will never select  $G_s$  during the search and hence  $A^*$  is optimal.

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A\* will always select the node that will lead to the optimal goal.

#### Heuristics





Goal State

- The solution is 26 steps long.
- $h_1(n)$  = number of misplaced tiles
- $h_2(n)=$  total Manhattan distance (sum of the horizontal and vertical distances).
- $h_1(n) = 8$
- Tiles 1 to 8 in the start state gives:  $h_2(n)=3+1+2+2+2+3+3+2=18$  which does not overestimate the true solution.

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Last note about heuristics. We have seen heuristics in the context of maps. Straightline distance turned out to be a good heuristic when you explore the map. In other contexts, you may need to find other kinds of heuristic.

26 steps is the true cost/optimal cost, to reach the goal states starting from the start

Both  $h_1$  and  $h_2$  never overestimate the true cost to find the solution. I prefer  $h_2$  because it's closer as a number to the real total number of steps.

## References

• S. J. Russell and P. Borvig, *Artificial Intelligence: A Modern Approach (4<sup>th</sup> Edition)*, Prentice Hall International, 2020.



